

Black box computes an unknown periodic function $f(x)$
find the period

classically : hard problem

However Quantum: easy *

let N be a positive integer ($N \geq 2$)

. Modulo N multiplication

let $G_N = \{a : \gcd(a, N) = 1 \text{ and } 1 \leq a < N\}$
a & N are relatively prime

G_N is a group under multiplication
modulo N

$|G_N| = \text{order of group} = \phi(N) = \text{Euler's totient function}$

if $N = p_1^{n_1} p_2^{n_2} \cdots p_k^{n_k}$ = prime number decomp.
of N

$$\phi(N) = p_1^{n_1-1} (p_1 - 1) \cdot p_2^{n_2-1} (p_2 - 1) \cdots$$

• Order of a is smallest number m such that $a^m \equiv 1 \pmod{N}$

$\rightarrow m$ divides $\phi(N)$

\rightarrow if a is relatively prime to N
then $a^{\phi(N)} \equiv 1 \pmod{N}$

$\hookrightarrow f(x) \equiv a^x \pmod{N}$

$f: \mathbb{Z} \rightarrow G_N$

$$\boxed{f(m) = 1} \Rightarrow f(x) = f(x+m)$$
$$a^x = a^x \cdot \underbrace{a^m}_1 = a^{x+m}$$

Period of $f(x)$ is m = the order of a

Problem: Given N find the prime factors of N .

Simplest case: $N = pq$ = product of two primes

Question: N is known. Find p (or q).

Classically: very hard.

Example: Try every prime number between $2, \dots, \sqrt{N}$

Shor's Algorithm:

Pick a random a .

Find the period of $f(x) = a^x \pmod{N}$

$$\begin{aligned}
 m: \text{ we know } m \mid \phi(N) &= (p-1)(q-1) \\
 &= pq - p - q + 1 \\
 &= N - (p+q) + 1
 \end{aligned}$$

Find a few more such orders

$$\left. \begin{array}{l} a \rightarrow m \\ a' \rightarrow m' \\ a'' \rightarrow m'' \end{array} \right\} \text{find least common multiple} \left. \begin{array}{l} \\ \\ \end{array} \right\} \text{get } \phi(N)$$

$$\left. \begin{array}{l} \text{You now know } N = pq \\ N+1 - \phi(N) = p+q \end{array} \right\} \text{find } p \text{ and } q$$

* RSA public key cryptosystem

Rivest, Shamir, Adleman

if a is relatively prime to N

$$a^{\phi(N)} \equiv 1 \pmod{N}$$

$$\rightarrow a^{\phi(N)+1} \equiv a \pmod{N}$$

If you know the prime factors of N ,
then you can find $\phi(N)$

$$N = pq \rightarrow \phi(N) = (p-1)(q-1)$$

find x that divides $\phi(N)+1$

" y such that $xy = \phi(N)+1$

$$a^{xy} \equiv a \pmod{N}$$

$$(a^x)^y \equiv a \pmod{N}$$

$$f: a \longrightarrow a^x$$

$$g: b \longrightarrow b^y$$

f and g are inverses of each other

$$g(f(a)) = a \pmod{N}$$

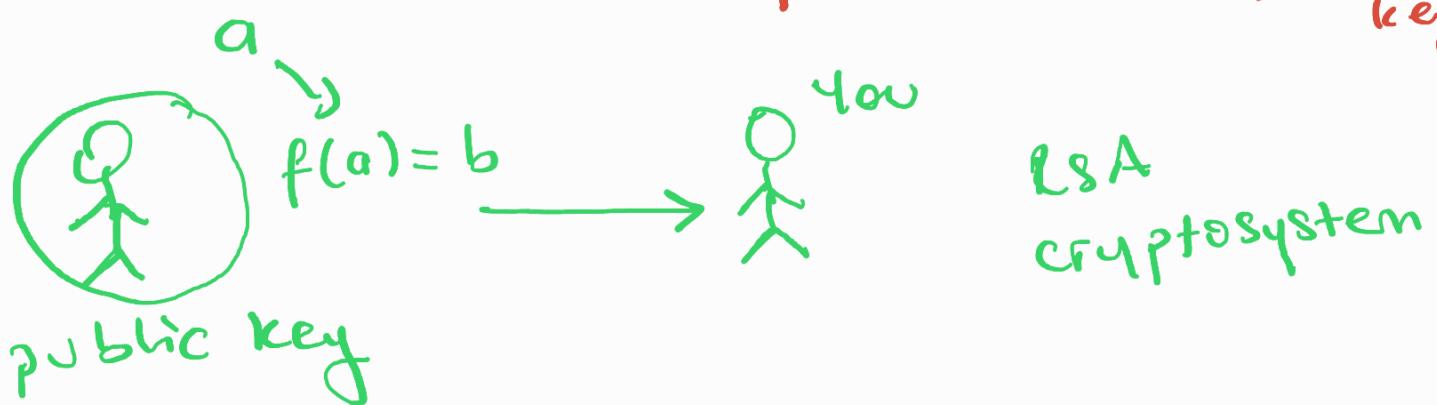
$$f(g(b)) = b \pmod{N}$$

$f: a \longrightarrow b = a^x \pmod{N}$ encoding (key= x, N)
function

$g: b \longrightarrow a = b^y \pmod{N}$ decoding (key= y, N)
function

Public key: x, N
 Private key: y, N, p, q

 } since factorization
 } is difficult →
 knowing public key
 doesn't give you any
 info about the private
 key.



* Entanglement

\textcircled{A} \textcircled{B}

$$|\Psi\rangle_{AB} = |\phi\rangle_A \otimes |x\rangle_B : \text{product state}$$

entangle if it is not a product state

$$|\Psi\rangle_{AB} = \frac{1}{\sqrt{2}} (|00\rangle + |11\rangle)$$

For two qubits A, B

$$|\Psi\rangle_{AB} = \sum_{i,j=0}^1 c_{ij} |i\rangle_A \otimes |j\rangle_B$$

$$c = \begin{bmatrix} c_{00} & c_{01} \\ c_{10} & c_{11} \end{bmatrix}$$

$|\Psi\rangle$ is entangled $\Leftrightarrow \det c \neq 0$

$|\Psi\rangle$ is product $\Leftrightarrow \det c = 0$

let A and B be any systems

\mathcal{H}_A = Hilbert space of A = n dimensional

\mathcal{H}_B = Hilbert space of B = m dimensional

let $|\Psi\rangle \in \mathcal{H}_A \otimes \mathcal{H}_B$

$|\Psi\rangle_{AB}$ = a state of the composite system A, B

→ let $|\alpha_1\rangle, \dots, |\alpha_N\rangle$ be an orthonormal basis (orb) of \mathcal{H}_A

$|\beta_1\rangle, \dots, |\beta_M\rangle$ " "

" " of \mathcal{H}_B



Expand $|\Psi\rangle = \sum_{i=1}^n \sum_{j=1}^m c_{ij} |i\rangle \otimes |\beta_j\rangle$

$$C = \begin{bmatrix} c_{11} & \dots & c_{1m} \\ \vdots & & \vdots \\ c_{n1} & \dots & c_{nm} \end{bmatrix}$$

is an $n \times m$ matrix

Singular value decomposition of C

$$C = UDV$$

U : $n \times n$ unitary
 D : $n \times m$ diagonal
 V : $n \times m$ unitary

$$D = \begin{bmatrix} * & & & & \\ * & \ddots & & & \\ & \ddots & \ddots & & \\ 0 & \dots & 0 & \ddots & \\ & & & & 0 \end{bmatrix} \quad \left\{ \text{only the } D_{ii} \text{ entries can be non-zero} \right.$$

$$C = UDV$$

$$C^+ = V^+ D^+ U^+$$

$$CC^+ = U \underbrace{D D^+}_{\text{diagonal } n \times n} V^+ \quad \left\{ \begin{array}{l} V \text{ diagonalizes } CC^+ \\ " \quad " \quad C^+ C \end{array} \right.$$

Expand

$$|\Psi\rangle = \sum_{i=1}^n \sum_{j=1}^m c_{ij} |\alpha_i\rangle \otimes |\beta_j\rangle$$

$$= \sum_{\substack{i,j,k,l \\ i,j \neq k,l}} U_{ik} D_{kl} V_{lj} |\alpha_i\rangle \otimes |\beta_j\rangle$$

Define $|\tilde{\alpha}_k\rangle = \sum_i U_{ik} |\alpha_i\rangle$

$|\tilde{\alpha}_1\rangle, \dots, |\tilde{\alpha}_N\rangle$ a new arb for H_A .

Define $|\tilde{\beta}_e\rangle = \sum_j V_{ej} |\beta_j\rangle$

$|\tilde{\beta}_1\rangle, \dots, |\tilde{\beta}_M\rangle$ is a new arb for H_B

$$|\Psi\rangle = \sum_{k,l} D_{kl} \underbrace{|\tilde{\alpha}_k\rangle}_{d_k \delta_{kl}} \otimes |\tilde{\beta}_e\rangle$$

$$|\Psi\rangle = \sum_k d_k |\tilde{\alpha}_k\rangle \otimes |\tilde{\beta}_k\rangle$$

Schmidt decomposition \Rightarrow