Sharp Bounds on Eigenvalues via Spectral Embedding Based on Signless Laplacians[*]

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Abstract

Using spectral embedding based on the probabilistic signless Laplacian, we obtain bounds on the spectrum of transition matrices on graphs. As a consequence, we bound return probabilities and the uniform mixing time of simple random walk on graphs. In addition, spectral embedding is used in this article to bound the spectrum of graph adjacency matrices.

Keywords: return probability, spectral embedding, signless Laplacian, spectral gap, minimum eigenvalue, eigenvalues, regular graphs.

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Contents

1   Introduction

Spectral embedding is a popular tool in modern data clustering, as summarized in [von Luxburg, 2007]. Also, spectral embedding was exploited to study graphs. In particular, [Lyons and Oveis Gharan, 2018 (henceforth cited as LOG18)] introduced spectral embedding as a new tool in analyzing reversible Markov chains (random walks on graphs). For instance, [LOG18, Theorem 4.9] gave a sharp bound on return probabilities of lazy random walk. Here, the lazy random walk on a graph G stays put at a vertex with probability $1/2$ and moves to a random uniform neighbor otherwise.

Theorem 1.1 (LOG18, Theorem 4.9). *Let G be a regular, simple, connected graph. For each vertex x of G and $t > 0$,*

$$0 \leqslant p'_t(x, x) - \frac{1}{|V(G)|} \leqslant \frac{13}{\sqrt{t}},$$

where $p'_t(x, x)$ is the probability of returning to x at step t for the lazy random walk on G starting from x. □

Now, a question arises naturally: if simple random walk, instead of lazy random walk, is considered, do we still have good bounds on return probabilities? Here, the simple random walk on a graph moves to a random neighbor uniformly.

To get some feeling, we recall the method used in the proof of the above bound in [LOG18]. Given a graph G, let P be the transition matrix of the simple random walk on G; then $P' := \frac{I+P}{2}$ is the transition matrix of the lazy random walk on G. We know that P and P' are symmetric operators on the space of functions on $V(G)$ square summable with respect to the degree sequence of G, so their spectra are both real. Denoting the spectra of P and P' as $\sigma(P)$ and $\sigma(P')$ respectively, we know that $\sigma(P')$ is closely related to $\sigma(P)$: $\lambda \in \sigma(P)$ if and only if $\frac{\lambda+1}{2} \in \sigma(P')$. Since $\sigma(P)$ is contained in the interval $[-1, 1]$, $\sigma(P')$ is contained in $[0, 1]$. Also, it is well known that the non-trivial (different from 1) spectrum of P' "governs" the "convergence rate" of lazy random walk: for instance, if the gap between the largest non-trivial eigenvalue of P' and 1 is large, then intuitively, the convergence will be faster. Thus, by the relation of $\sigma(P)$ and $\sigma(P')$, we only need to consider the spectrum of P near 1. Moreover, we notice that $\sigma(P)$ is closely related to the spectrum of the probabilistic Laplacian $\mathcal{L} := I - P$ corresponding to P. To prove the aforementioned bound in [LOG18], using the spectral embedding based on $\mathcal{L}$, [LOG18] first revealed upper bounds on the vertex spectral measure of $\mathcal{L}$. It is known that return probabilities are determined by vertex spectral measures; therefore, bounds on return probabilities can be obtained from bounds on vertex spectral measures (see [LOG18, Proposition 3.5]).

When it comes to simple random walk, the situation is different. Note that simple random walks on bipartite graphs have period two and simple random walks on non-bipartite graphs are aperiodic. For simplicity of presentation, we discuss non-bipartite case only in this introduction. To begin, we still have the intuition that the "convergence rate" of simple random walk is related to the non-trivial spectrum of P: for instance, [Diaconis and Stroock,

1991, Proposition 3] proved that the geometric convergence rate in total variation norm is determined by the maximum non-trivial eigenvalue *in absolute value* when the random walk has finitely many states. Note that $\sigma(P) \subseteq [-1, 1]$, but $\sigma(P)$ is not necessarily non-negative. Therefore, in order to deal with "convergence" of simple random walk, one also has to consider the negative spectrum of P. However, we usually study the spectral gap of the probabilistic Laplacian $\mathcal{L} = I - P$, which is related to the spectrum of P near 1 but does not provide much information about the *negative* spectrum of P.

Our solution is to consider another operator: the probabilistic signless Laplacian operator $\mathcal{Q} := I + P$. Obviously, the spectrum of P is closely related to the spectrum of $\mathcal{Q}$ by a shift of 1 unit horizontally. The spectrum of $\mathcal{Q}$ is therefore real and non-negative. This brings us some convenience: we have more tools to deal with the operator $\mathcal{Q}$, its associated quadratic form for a start. We will first consider the vertex spectral measure of $\mathcal{Q}$ using the spectral embedding based on $\mathcal{Q}$, so the negative spectrum of P is bounded. Then we may proceed to get a bound on return probabilities of simple random walk. See Theorem 3.4 for details.

1.1   Results

To give an overview of our results, in this subsection, we constrain ourselves to the case of unweighted graphs. Some notation will be needed, which will be explained in more detail in subsequent sections. Consider a locally finite, simple, connected graph $G = (V, E)$. A vertex $x \in V$ has degree $d(x)$ in G. When G is finite, we denote the eigenvalues of the transition matrix P on G as

$$-1 \leqslant \lambda_{\min}^P = \lambda_1^P \leqslant \lambda_2^P \leqslant \lambda_3^P \leqslant \cdots \leqslant \lambda_{n-1}^P < \lambda_n^P = \lambda_{\max}^P = 1,$$

where $n = |V|$. It is well known that $\lambda_{\min}^P = -1$ if and only if G is bipartite. However, for infinite graphs, it is more appropriate to consider their vertex spectral measures. For instance, denoting the probability of returning to x at step t as $p_t(x, x)$ for the simple random walk on G starting at x, then we have

$$p_t(x, x) = \int_{[0,2]} (1 - \lambda)^t \mathrm{d}\mu_x(\lambda),$$

where μ_x is the vertex spectral measure at x of $\mathcal{L}$. Denote the vertex spectral measure at x of $\mathcal{Q}$ as $\mu_x^{\mathcal{Q}}$. It is shown in Lemma 3.2 that when G is finite,

$$\sum_{x \in V(G)} \mu_x^{\mathcal{Q}}(\delta) = \left| \{ j ; \ 1 + \lambda_j^P \leqslant \delta \} \right|.$$

This enables us to count eigenvalues of P on the interval $[-1, -1 + \delta]$. Therefore, we can get lower bounds on eigenvalues of P from upper bounds of the vertex spectral measure of $\mathcal{Q}$. Write $\pi(x) := \frac{d(x)}{2|E|}$, which is 0 when G is an infinite graph.

We first consider non-bipartite graphs. In fact, when the graph is non-bipartite, the simple random walk is aperiodic, so some troubles are avoided. Our result for simple random walk on regular graphs reads as follows.

Theorem 1.2. *For a regular, non-bipartite, simple, connected graph G, we have*

$$\mu_x^{\mathcal{Q}}(\delta) \leqslant 10\sqrt{\delta}, \qquad 0 \leqslant \delta \leqslant 2,\ x \in V.$$

For each $x \in V$, simple random walk on G satisfies

$$0 \leqslant p_t(x,x) - \pi(x) \leqslant \frac{18}{\sqrt{t}} \qquad \text{for} \quad t \equiv 0 \bmod 2,$$

$$\left| p_t(x,x) - \pi(x) \right| \leqslant \frac{9}{\sqrt{t}} \qquad \text{for} \quad t \equiv 1 \bmod 2.$$

Furthermore, when G is finite, for $1 \leqslant k \leqslant n$, $\lambda_k^P \geqslant -1 + \frac{k^2}{100n^2}$.

The above result is sharp by Example 3.6 and its proof is presented in Section 3.1. Theorem 1.2 is interesting since the degree and size of the regular graph G are not involved. If we consider general finite graphs and introduce the relaxation time t_{rel}, we have the following bound on return probabilities.

Theorem 1.3. *For $t \geqslant t_{\mathrm{rel}}$, simple random walk on a non-bipartite, finite, simple, connected graph G satisfies*

$$0 \leqslant p_{2t}(x,x) - \pi(x) \leqslant \frac{8d(x)\sqrt{t_{\mathrm{rel}}+1}}{(t+1)d_{\min}},$$

$$\left| p_{2t+1}(x,x) - \pi(x) \right| \leqslant \frac{8d(x)\sqrt{t_{\mathrm{rel}}+1}}{\sqrt{(t+1)(t+2)}\,d_{\min}},$$

where $d_{\min}$ is the minimum degree of G and $t_{\mathrm{rel}} := \left(1 - |\lambda_1^P| \vee |\lambda_{n-1}^P|\right)^{-1}$.

For lazy random walk, a similar result was obtained in [Oliveira and Peres, 2019]. We use a similar method as theirs in Section 4 to prove the above theorem.

The following proposition is for graphs satisfying volume growth conditions.

Proposition 1.4. *Let G be a non-bipartite, infinite, simple, connected graph. Suppose that $c > 0$ and $D \geqslant 1$ are constants such that $\mathsf{wt}(x,r) \geqslant c(r+1)^D$ for all $r \geqslant 0$. Then for all $\delta \in (0,2)$,*

$$\mu_x^*(\delta) \leqslant Cw(x)\delta^{D/(D+1)},$$

$$\mu_x^{\mathcal{Q}}(\delta) \leqslant Cw(x)\delta^{D/(D+1)},$$

where

$$C := \frac{(D+1)^2}{c^{1/(D+1)}D^{2D/(D+1)}}.$$

Hence for all $t \geqslant 1$, simple random walk satisfies

$$p_t(x,x) \leqslant 2C'w(x)t^{-D/(D+1)} \qquad \text{for} \quad t \equiv 0 \bmod 2,$$

$$p_t(x,x) \leqslant C'w(x)t^{-D/(D+1)} \qquad \text{for} \quad t \equiv 1 \bmod 2,$$

where

$$C' := \frac{D+1}{c^{1/(D+1)}D^{(D-1)/(D+1)}}\Gamma\left(\frac{D}{D+1}\right).$$

Using the spectral embedding based on $\mathcal{Q}$, we also bound the vertex spectral measure of $\mathcal{Q}$ in another way and get bounds on the uniform mixing time.

Proposition 1.5. *For any non-bipartite, finite, simple, connected graph G, we have*

$$\mu_x^{\mathcal{Q}}(\delta) \leqslant \frac{d(x)\delta}{\mathcal{K}(G)}, \qquad \delta \in [0, 2), \ x \in V,$$

where $\mathcal{K}(G)$ is defined as

$$\min\left\{\frac{\sum_{(v,u)\in E(G)}|f(v) + f(u)|^2}{\max_{y\in V}|f(y)|^2} \ ; \ \min_{y\in V} f(y) < 0 < \max_{y\in V} f(y)\right\}$$

and satisfies $\mathcal{K}(G) \geqslant \frac{1}{\mathrm{diam}(G)+1}$. Consequently, for $1 \leqslant k \leqslant n$,

$$\lambda_k^P \geqslant -1 + \frac{k\mathcal{K}(G)}{\sum_{x\in V} d(x)} \geqslant -1 + \frac{k}{\big(\mathrm{diam}(G) + 1\big) \sum_{x\in V} d(x)}.$$

Furthermore, $\tau_\infty(1/4) \leqslant 8n^3$. If G is also regular, then we have $\tau_\infty(1/4) \leqslant 24n^2$.

Note that for a finite, simple, connected graph G, [LOG18, Proposition 4.2] implies

$$\lambda_{n-k}^P \leqslant 1 - \frac{k}{\mathcal{R}_{\mathrm{diam}}(G) \sum_{x\in V} d(x)}, \qquad 0 \leqslant k \leqslant n - 1,$$

where $\mathcal{R}_{\mathrm{diam}}(G)$ is the resistance diameter of G. This bound combined with the lower bound on eigenvalues of P in Proposition 1.5 improves [Landau and Odlyzko, 1981], which asserted that each eigenvalue λ of P that is neither 1 nor -1 satisfies

$$|\lambda| \leqslant 1 - \frac{1}{\big(\mathrm{diam}(G) + 1\big)nd_{\mathrm{max}}},$$

where d_{max} is the maximum degree of G.

The bound $\tau_\infty(1/4) = O(n^2)$ in Proposition 1.5 is sharp by the example of cycles: there is a constant $c > 0$ such that for all odd number n, simple random walk on an n-cycle satisfies $\tau_\infty(1/4) \geqslant cn^2$ (cf. [Tetali and Montenegro, 2005, Example 3.11]). See Section 6 for more details.

As a special class of regular graphs, vertex-transitive graphs are of interest, since they are intuitively "homogeneous" and especially well studied.

Theorem 1.6. *Let G be a non-bipartite, simple, connected, vertex-transitive graph with degree $d < \infty$. For each $x \in V$, $c \in (0,1)$, and $\delta \in (0,2]$, we have*

$$\mu_x^{\mathcal{Q}}(\delta) \leqslant \frac{1}{c^2 N^\#\left(\frac{\sqrt{1-c}}{\sqrt{d\delta}}\right)},$$

where $N^\#(r)$ denotes the number of vertices in a ball of radius r. In addition, if G is finite,

$$\lambda_{\mathrm{min}}^{\mathcal{Q}} \geqslant \frac{2}{d}\left(\sin\frac{\pi}{4\big(\mathrm{diam}(G) + 1\big)}\right)^2.$$

As usual, a bound on return probabilities of simple random walk follows immediately from a bound on vertex spectral measures. See Corollary 7.3 as an example.

Apart from vertex spectral measure at one single vertex, we may consider average spectral measure for finite graphs. Given a finite graph G, the average spectral measure of $\mathcal{Q}$ is defined as $\mu^{\mathcal{Q}} := \sum_{x \in V} \mu_x^{\mathcal{Q}}/n$.

Theorem 1.7. *For any non-bipartite, finite, simple, connected graph G and $\delta \in (0,2)$, we have*

$$\mu^{\mathcal{Q}}(\delta) < (4000\delta)^{1/3}.$$

Consequently, $\lambda_k^P \geqslant -1 + \frac{k^3}{4000n^3}$ for $1 \leqslant k \leqslant n$. Furthermore,

$$0 \leqslant \frac{\sum_{x \in V(G)} p_t(x,x) - 1}{n} \leqslant \frac{30}{t^{1/3}} \qquad for \quad t \equiv 0 \bmod 2,$$

$$\frac{\left|\sum_{x \in V(G)} p_t(x,x) - 1\right|}{n} \leqslant \frac{15}{t^{1/3}} \qquad for \quad t \equiv 1 \bmod 2.$$

Theorem 1.7 is treated in Section 8.

Non-bipartite graphs are considered in previous paragraphs. But the bipartite case is a bit different, because simple random walk on a bipartite graph has period two. Fortunately, by [Mohar and Woess, 1989, Theorem 4.8], the vertex spectral measures of $\mathcal{L}$ and $\mathcal{Q}$ coincide. Therefore, only the estimates for the vertex spectral measure of $\mathcal{L}$ from [LOG18] will be enough for us to get bounds on return probabilities of simple random walk. For example, we have the following result for bipartite graphs in Section 9.

Theorem 1.8. *Consider simple random walk on a regular, bipartite, simple, connected graph G. Then for each $x \in V$,*

$$0 \leqslant p_t(x,x) - 2\pi(x) \leqslant \frac{18}{\sqrt{t}} \qquad for \quad t \equiv 0 \bmod 2.$$

Our method work as well for analyzing the spectrum of the adjacency matrix A of a finite graph G. Suppose the eigenvalues of A are

$$-d_{\max} \leqslant \lambda_{\min}^A = \lambda_1^A \leqslant \lambda_2^A \leqslant \lambda_3^A \leqslant \cdots \leqslant \lambda_{n-1}^A < \lambda_n^A = \lambda_{\max}^A \leqslant d_{\max},$$

where $d_{\max}$ is the maximum degree of G.

Proposition 1.9. *Let G be a non-bipartite, finite, simple, connected graph. For $1 \leqslant k \leqslant n$,*

$$d_{\max} + \lambda_k^A \geqslant \frac{k\mathcal{K}(G)}{n} \geqslant \frac{k}{\big(\mathrm{diam}(G)+1\big)n}.$$

Proposition 1.9 improves [Alon and Sudakov, 2000, Theorem 1.1], which obtained that $d_{\max} + \lambda_1^A \geqslant \frac{1}{(\mathrm{diam}(G)+1)n}$. To get Proposition 1.9, in Section 10, we first consider the spectral embedding based on the combinatorial signless Laplacian $\Theta := D + A$, where D is the diagonal degree matrix of G. Then Proposition 1.9 follows from the bound below on the vertex spectral measure of Θ. by a linear algebra argument.

Proposition 1.10. *Let G be a non-bipartite, finite, simple, connected graph. Then for each $\delta \in [0, \lambda_{\max}^{\Theta})$ and $x \in V$, we have*

$$\mu_x^{\Theta}(\delta) \leqslant \frac{\delta}{\mathscr{K}(G)} \leqslant \big(\mathrm{diam}(G) + 1\big)\delta.$$

Suppose the eigenvalues of Θ are

$$0 \leqslant \lambda_{\min}^{\Theta} = \lambda_1^{\Theta} \leqslant \lambda_2^{\Theta} \leqslant \lambda_3^{\Theta} \leqslant \cdots \leqslant \lambda_{n-1}^{\Theta} < \lambda_n^{\Theta}.$$

The above proposition has the following corollary.

Corollary 1.11. *Let G be a non-bipartite, finite, simple, connected graph . For $1 \leqslant k \leqslant n$, we have*

$$\lambda_k^{\Theta} \geqslant \frac{k\mathscr{K}(G)}{n} \geqslant \frac{k}{\big(\mathrm{diam}(G) + 1\big)n}.$$

In fact, it is known that a graph is bipartite if and only if $\lambda_{\min}^{\Theta} = 0$; for a nonbipartite graph G, [M. Desai and V. Rao, 1994] showed that $\lambda_{\min}^{\Theta}$ measures non-bipartiteness of G.

1.2   Related Works

In the field of spectral graph theory, the combinatorial signless Laplacian has already drawn wide attention: [Cvetković and Simić, 2009, 2010a,b] are surveys on the study of the combinatorial signless Laplacian. In fact, [Alon and Sudakov, 2000, Theorem 1.1] also used the combinatorial signless Laplacian implicitly by considering its associated quadratic form. We use the method of spectral embedding, so not only the minimum eigenvalue but all eigenvalues of the graph adjacency matrices are bounded from below in Proposition 1.9.

One highlight of this article is the introduction of the probabilistic signless Laplacian $\mathcal{Q}$, enabling us to deal with the negative spectrum of the transition matrix P on a graph with the tool of spectral embedding. Indeed, the quadratic form associated to $\mathcal{Q}$ was used in [M. P. Desai and V. B. Rao, 1993; Diaconis and Stroock, 1991] to give lower bounds on the minimum eigenvalue of P. We consider the operator $\mathcal{Q}$ explicitly and exploit the spectral embedding based on it. Therefore, the entire spectrum of P is treated, rather than only the minimum eigenvalue. For instance, [Landau and Odlyzko, 1981] is improved as we discussed after Proposition 1.5.

In the literature, there are many fewer results on the negative spectrum of P than on the positive spectrum of P. Recall that in the study of the positive spectrum of P, the probabilistic Laplacian operator $\mathcal{L}$ was usually used; in particular, the spectral gap of $\mathcal{L}$ equals the gap between the largest nontrivial spectrum of P and 1. However, when one wants to study the negative spectrum of P, for instance by considering the signless probabilistic Laplacian operator $\mathcal{Q}$, it is harder: since we don't have many tools. For example, tools from electrical network theory are pretty useful in the study of $\mathcal{L}$, but they are not readily available to deal with $\mathcal{Q}$. We have to make a "detour" and adapt the existing tools.

There has been extensive study of return probability bounds of random walks, both on finite and infinite graphs. The asymptotics of large-time return probabilities correspond to the asymptotics of the spectral measures of $\mathcal{L}$ and $\mathcal{Q}$ near 0, as suggested in Section 2.3. Therefore, our bounds on the vertex spectral measure of $\mathcal{Q}$, combined with the bounds on the vertex spectral measure of $\mathcal{L}$ from [LOG18], will result in return probability bounds for simple random walks, which are parallel to the bounds for lazy random walks in [LOG18].

1.3 Structure of this Article

We review notation for graphs and introduce spectral embedding based on the signless Laplacian in Sections 2.1 and 2.2. Some preliminaries are included in Section 2.3. Return probabilities of simple random walk on regular graphs are considered in Section 3. A bound on return probabilities involving the relaxation time is presented in Section 4. Return probability bounds based on volume growth conditions are discussed in Section 5. The case of transitive graphs is treated in Section 7. We also consider average return probabilities for finite graphs in Section 8. In Section 6, we bound the uniform mixing time. Bipartite graphs are discussed in Section 9. The tool of spectral embedding is exploited to study eigenvalues of graph adjacency matrices in Section 10.

2 Graph Notation and Spectral Embedding

2.1 Graph Notation and Laplacian Operators

Let $G = \big(V(G), E(G)\big)$ be a finite or infinite, undirected, simple, connected, weighted graph. For an edge of G, say $e = (x, y) \in E(G)$, let $w(e) = w(x, y) > 0$ be its weight. We say G is *unweighted* if $w(e) = 1$ for each edge $e \in E(G)$. We assume G has weighted adjacency matrix $A(G)$, (unweighted) diameter $\mathrm{diam}(G)$, and (weighted) resistance diameter $\mathcal{R}_{\mathrm{diam}}(G)$. Also, minimum, maximum, and average degrees in G are denoted by $d_{\min}(G)$, $d_{\max}(G)$, and $d_{\mathrm{avg}}(G)$, respectively. When G is finite, we denote $\big|V(G)\big| = n(G)$. If G is understood, reference to G may be omitted.

For $x \in V(G)$, we use standard graph notation $N(x) \coloneqq \big\{ y \in V(G) \, ; \, (x, y) \in E(G) \big\}$ to denote the collection of all neighbors of x. Throughout this article, we require that G is *locally finite*, i.e., $\sum_{y \in N(x)} w(x, y) < \infty$ for each $x \in V(G)$. We say that $w(x) \coloneqq \sum_{y \in N(x)} w(x, y)$ is the weight of $x \in V(G)$ in G, and the weight of a vertex subset $S \subseteq V$ in G is $\mathsf{wt}(S; G) \coloneqq \sum_{x \in S} w(x)$. If G is unweighted, $w(x)$ equals the degree $d(x)$ of x. Again, when G is understood, the reference to G may be omitted. In particular, when G is vertex-transitive, all vertices have the same weight, denoted by w. Write $\pi(x) = \frac{w(x)}{\mathsf{wt}(V(G))}$. For a vertex $x \in V(G)$ and $r \geqslant 0$, let

$$B(x, r; G) = \big\{ y \in V(G) \, ; \, \mathsf{dist}(x, y) \leqslant r \big\},$$

where dist is the distance on G. Set

$$\mathsf{wt}(x, r; G) \coloneqq \mathsf{wt}\big(B(x, r; G); G\big)$$

We have a series of useful operators: the probability transition matrix P on G, the probabilistic Laplacian $\mathcal{L} = I - P$, and the probabilistic signless Laplacian $\mathcal{Q} = I + P$. We know that P, $\mathcal{L}$, and $\mathcal{Q}$ are bounded self-adjoint operators on Hilbert space $\ell^2(V, w)$. Recall that we write $\ell^2(V, w)$ for the (real or complex) Hilbert space of functions $f \colon V \to \mathbb{R}$ or $\mathbb{C}$ with inner product

$$\langle f, g \rangle_w := \sum_{x \in V} w(x) f(x) \overline{g(x)}$$

and squared norm $\|f\|_w^2 := \langle f, f \rangle_w$. We reserve $\langle \cdot, \cdot \rangle$ and $\| \cdot \|$ for the standard inner product and norm on $\mathbb{R}^k$ ($k \in \mathbb{N}$) and $\ell^2(V)$, and also for the norm on a generic Hilbert space. For a vertex $x \in V$, we use $\mathbf{1}_x$ to denote the indicator vector of x:

$$\mathbf{1}_x(y) := \begin{cases} 1 & \text{if } y = x, \\ 0 & \text{otherwise.} \end{cases}$$

We also write $\mathbf{e}_x := \frac{\mathbf{1}_x}{\sqrt{w(x)}}$. Note that $\mathbf{e}_x \in \ell^2(V, w)$ is of unit norm: $\|\mathbf{e}_x\|_w = 1$.

Obviously, whether G is finite or infinite, the spectrum of P and $\mathcal{Q}$ are related by a shift of 1 unit horizontally. Denote the resolution of identity for $\mathcal{L}$ as $I_\mathcal{L}$. As in [LOG18], the vertex spectral measure of $\mathcal{L}$ at $x \in V$ is defined by

$$\mu_x(\delta) := \big\langle I_\mathcal{L}\big([0, \delta]\big) \mathbf{e}_x, \mathbf{e}_x \big\rangle_w = \big\langle I_\mathcal{L}\big([0, \delta]\big) \mathbf{1}_x, \mathbf{1}_x \big\rangle, \qquad 0 \leqslant \delta \leqslant 2.$$

For convenience, we will also use

$$\mu_x^*(\delta) := \big\langle I_\mathcal{L}\big((0, \delta]\big) \mathbf{e}_x, \mathbf{e}_x \big\rangle_w = \big\langle I_\mathcal{L}\big((0, \delta]\big) \mathbf{1}_x, \mathbf{1}_x \big\rangle, \qquad 0 \leqslant \delta \leqslant 2, \ x \in V.$$

It is easy to see that $\mu_x(\delta) = \mu_x^*(\delta) + \pi(x)$ for $0 \leqslant \delta \leqslant 2$ (cf. [LOG18, Section 3.1]).

When G is finite, the spectra of P and $\mathcal{Q}$ consist of eigenvalues (point spectrum) only:

- Denote the eigenvalues of the transition matrix P on G as

$$-1 \leqslant \lambda_{\min}^P = \lambda_1^P \leqslant \lambda_2^P \leqslant \lambda_3^P \leqslant \cdots \leqslant \lambda_{n-1}^P < \lambda_n^P = \lambda_{\max}^P = 1.$$

- Denote the eigenvalues of the probabilistic signless Laplacian $\mathcal{Q}$ as

$$0 \leqslant \lambda_{\min}^\mathcal{Q} = \lambda_1^\mathcal{Q} \leqslant \lambda_2^\mathcal{Q} \leqslant \lambda_3^\mathcal{Q} \leqslant \cdots \leqslant \lambda_{n-1}^\mathcal{Q} < \lambda_n^\mathcal{Q} = \lambda_{\max}^\mathcal{Q} = 2.$$

2.2 Spectral Embedding Based on the Signless Laplacian

The spectral embedding based on $\mathcal{L}$ is introduced in [LOG18, Section 3.4] as a powerful tool in analyzing random walk on graphs. In this subsection, we will introduce the spectral embedding based on $\mathcal{Q}$ in parallel. Denote the resolution of identity for $\mathcal{Q}$ as $I_\mathcal{Q}$, with vertex spectral measure at x

$$\mu_x^\mathcal{Q}(\delta) := \big\langle I_\mathcal{Q}(\delta) \mathbf{e}_x, \mathbf{e}_x \big\rangle_w, \qquad 0 \leqslant \delta \leqslant 2, \ x \in V,$$

where $I_\mathcal{Q}(\delta) := I_\mathcal{Q}\big([0, \delta]\big)$.

Lemma 2.1. *Let $f \in \ell^2(V, w)$. We have*

$$\langle \mathcal{Q}f, f \rangle_w = \sum_{(v,u) \in E(G)} w(v, u)\big|f(v) + f(u)\big|^2.$$

For $\delta \in [0, 2]$ and $f \in \mathrm{img}\big(I_{\mathcal{Q}}(\delta)\big)$,

$$\langle \mathcal{Q}f, f \rangle_w \leqslant \delta\|f\|_w^2 = \delta \sum_{v \in V} w(v)\big|f(v)\big|^2.$$

Proof. We need prove the first assertion only. First, we claim that

$$\sum_{v,u \in V} w(v, u)|f(v)||f(v) + f(u)| < \infty.$$

In fact, we have

$$\sum_{v,u \in V} w(v, u)|f(v)||f(v) + f(u)| \leqslant \sum_{v,u \in V} w(v, u)\big(|f(v)|^2 + |f(v)||f(u)|\big)$$

$$= \sum_{v,u \in V} w(v, u)|f(v)|^2 + \sum_{v,u \in V} w(v, u)|f(v)||f(u)|.$$

By the Cauchy–Schwarz inequality, we may proceed and get

$$\sum_{v,u \in V} w(v, u)|f(v)||f(v) + f(u)| \leqslant \sum_{v,u \in V} w(v, u)|f(v)|^2$$

$$+ \Big(\sum_{v,u \in V} w(v, u)|f(v)|^2\Big)^{1/2}\Big(\sum_{v,u \in V} w(v, u)|f(u)|^2\Big)^{1/2}$$

$$= \sum_{v,u \in V} w(v, u)|f(v)|^2 + \sum_{v,u \in V} w(v, u)|f(v)|^2$$

$$= 2\sum_{v,u \in V} w(v, u)|f(v)|^2 = 2\sum_{v \in V}|f(v)|^2 \sum_{u \in V} w(v, u)$$

$$= 2\sum_{v \in V}|f(v)|^2 w(v) = 2\|f\|_w < \infty.$$

The claim is proved. By this claim, using Fubini's theorem, we have

$$\sum_{v,u \in V} w(v, u)f(v)\overline{\big(f(v) + f(u)\big)} = \sum_{v \in V} f(v)\overline{\sum_{u \in V} w(v, u)\big(f(v) + f(u)\big)}$$

$$= \sum_{v \in V} w(v)f(v)\overline{\big(f(v) + \sum_{u \in V} p(v, u)f(u)\big)}$$

$$= \sum_{v \in V} w(v)f(v)\overline{\big((I + P)f\big)(v)} = \langle f, (I + P)f \rangle_w$$

$$= \langle f, \mathcal{Q}f \rangle_w.$$

By interchanging u and v, we have

$$\sum_{v,u \in V} w(v,u) f(u) \overline{\big(f(v) + f(u)\big)} = \langle f, \mathcal{Q}f \rangle_w.$$

Therefore,

$$\sum_{(v,u) \in E(G)} w(v,u) \big|f(v) + f(u)\big|^2 = \frac{1}{2} \sum_{v,u \in V} w(v,u) \big|f(v) + f(u)\big|^2$$

$$= \frac{1}{2}\Big(\sum_{v,u \in V} w(v,u) f(v) \overline{\big(f(v) + f(u)\big)} + \sum_{v,u \in V} w(v,u) f(u) \overline{\big(f(v) + f(u)\big)}\Big)$$

$$= \tfrac{1}{2}\big(\langle f, \mathcal{Q}f \rangle_w + \langle f, \mathcal{Q}f \rangle_w\big)$$

$$= \langle f, \mathcal{Q}f \rangle_w. \qquad \square$$

Corollary 2.2. *If G is connected and non-bipartite, then $\mathrm{img}\big(I_{\mathcal{Q}}(\{0\})\big)$ contains only the zero function. Therefore in this case, $I_{\mathcal{Q}}(\delta) = I_{\mathcal{Q}}\big([0,\delta]\big) = I_{\mathcal{Q}}\big((0,\delta]\big)$. It also follows that 0 is not an eigenvalue of $\mathcal{Q}$ if G is non-bipartite.*

Proof. Assume $f \in \mathrm{img}\big(I_{\mathcal{Q}}(\{0\})\big)$. We see that

$$\sum_{(v,u) \in E(G)} w(v,u) \big|f(v) + f(u)\big|^2 = \langle \mathcal{Q}f, f \rangle_w \leqslant 0\|f\|_w^2 = 0.$$

Therefore, for $(v,u) \in E(G)$, we have $f(v) = -f(u)$. Since G is non-bipartite, there is an odd cycle $C = v_1 v_2 \cdots v_s v_1$, where $s \equiv 1 \bmod 2$. Thus, $f(v_1) = -f(v_1)$, implying that $f(v_1) = 0$. By connectedness, $f(v) = 0$ for all $v \in V$. This corollary is proved. $\qquad \square$

For a fixed $\delta \in [0,2]$, we define spectral embedding $F^{\mathcal{Q}}$ based on $\mathcal{Q}$ as

$$F^{\mathcal{Q}} \colon V(G) \to \ell^2(V(G), w)$$

$$x \mapsto F_x^{\mathcal{Q}} := \frac{I_{\mathcal{Q}}(\delta)\mathbf{e}_x}{\sqrt{w(x)}} = \frac{I_{\mathcal{Q}}(\delta)\mathbf{1}_x}{w(x)}.$$

It is clear that for each $x \in V$, $F_x^{\mathcal{Q}} \in \ell^2(V, w)$ is a real-valued function on V.

Lemma 2.3. *For each finite or infinite graph G and $x \in V(G)$,*

$$\big\|F_x^{\mathcal{Q}}\big\|_w^2 = F_x^{\mathcal{Q}}(x) = \frac{\mu_x^{\mathcal{Q}}(\delta)}{w(x)}. \qquad \square$$

Lemma 2.4. *If $\mu_x^{\mathcal{Q}}(\delta) > 0$, define $f \colon V \to \mathbb{C}$ as $f := \frac{F_x^{\mathcal{Q}}}{\|F_x^{\mathcal{Q}}\|_w}$. Then*

(1) $\|f\|_w = 1$;

(2) $f(x) = \sqrt{\mu_x^{\mathcal{Q}}(\delta)/w(x)}$;

(3) $f \in \mathrm{img}\big(I_{\mathcal{Q}}(\delta)\big)$. $\qquad \square$

We need only mimic the proofs of [LOG18, Lemmas 3.11 and 3.12] to prove Lemmas 2.3 and 2.4.

2.3 Some Preliminaries

Lemma 4.5 in [LOG18] gave an estimate for diameter of regular graphs. Using a similar argument, we get the following enhanced Lemma 2.5.

Lemma 2.5. *Let G be a finite, simple, connected graph. We have* $\mathrm{diam}(G) \leqslant \frac{3n}{d_{\min}+1} - 1$.

Proof. Suppose that u_0 and $u_{\mathrm{diam}(G)}$ are two vertices of G with

$$\mathsf{dist}\big(u_0, u_{\mathrm{diam}(G)}\big) = \mathrm{diam}(G),$$

and $\mathcal{P} = u_0 u_1 \dots u_{\mathrm{diam}(G)}$ is a shortest path joining these two vertices. Let

$$S = \{u_0, u_3, \dots, u_{\lfloor \mathrm{diam}(G) \rfloor_{(3)}}\},$$

where $\lfloor \mathrm{diam}(G) \rfloor_{(3)}$ is the largest multiple of 3 that is smaller than or equal to $\mathrm{diam}(G)$. It is easy to see that

$$|S| = \frac{\lceil \mathrm{diam}(G) + 1 \rceil^{(3)}}{3},$$

where $\lceil \mathrm{diam}(G)+1 \rceil^{(3)}$ is the smallest multiple of 3 that is greater than or equal to $\mathrm{diam}(G)+1$. Because $\mathcal{P}$ is a shortest path, no vertex in S is adjacent to any other vertex in S, and no pair of vertices in S have any common neighbors. Hence,

$$n \geqslant |S| + d_{\min}|S| = |S|(d_{\min} + 1)$$
$$= \frac{\lceil \mathrm{diam}(G) + 1 \rceil^{(3)}}{3}(d_{\min} + 1) \geqslant \frac{\mathrm{diam}(G) + 1}{3}(d_{\min} + 1).$$

Thus,

$$\mathrm{diam}(G) \leqslant \frac{3n}{d_{\min} + 1} - 1. \qquad \square$$

The following two results will be useful when we are dealing with return probabilities in subsequent sections.

Lemma 2.6. *Let t be a positive integer and η be an increasing and right-continuous function on $[0, 2]$ with $\eta(0) = 0$. Then we have*

$$t \int_0^2 \eta(\lambda)(1 - \lambda)^{t-1} \, d\lambda = (-1)^{t+1}\eta(2) + \int_{(0,2]} (1 - \lambda)^t \, d\eta(\lambda).$$

Proof. Using integration by parts, we have

$$t \int_0^2 \eta(\lambda)(1 - \lambda)^{t-1} \, d\lambda = -\int_0^2 \eta(\lambda) \, d(1 - \lambda)^t$$
$$= -\eta(\lambda)(1 - \lambda)^t \big|_0^2 + \int_{(0,2]} (1 - \lambda)^t \, d\eta(\lambda)$$
$$= (-1)^{t+1}\eta(2) + \int_{(0,2]} (1 - \lambda)^t \, d\eta(\lambda). \qquad \square$$

Lemma 2.7. *Consider simple random walk on a non-bipartite, simple, connected, weighted graph G.*

(1) *Let φ be an increasing and right-continuous function on $[0, 2]$. Assume further that φ satisfies the following conditions:*

$$\varphi(0) = 0 = \mu_x^*(0), \qquad \varphi(2) = 1 - \pi(x) = \mu_x^*(2),$$
$$\mu_x^*(\lambda) \leqslant \varphi(\lambda) \quad for \quad \lambda \in [0, 1),$$
$$\mu_x^*(\lambda) \geqslant \varphi(\lambda) \quad for \quad \lambda \in [1, 2].$$

Then for $t \equiv 0 \bmod 2$,

$$\pi(x) \leqslant p_t(x, x) \leqslant \pi(x) + \int_{(0,2]} (1 - \lambda)^t \, \mathrm{d}\varphi(\lambda).$$

(2) *Let ψ_1 and ψ_2 be increasing and right-continuous functions on $[0, 2]$. Assume further that ψ_1 and ψ_2 satisfy the following conditions:*

$$\psi_1(0) = \psi_2(0) = 0 = \mu_x^*(0),$$
$$\psi_1(2) = \psi_1(2) = 1 - \pi(x) = \mu_x^*(2),$$
$$\mu_x^*(\lambda) \leqslant \psi_2(\lambda) \quad for \quad \lambda \in [0, 2],$$
$$\mu_x^*(\lambda) \geqslant \psi_1(\lambda) \quad for \quad \lambda \in [0, 2].$$

Then for $t \equiv 1 \bmod 2$, we have

$$\pi(x) + \int_{(0,2]} (1 - \lambda)^t \, \mathrm{d}\psi_1(\lambda) \leqslant p_t(x, x) \leqslant \pi(x) + \int_{(0,2]} (1 - \lambda)^t \, \mathrm{d}\psi_2(\lambda).$$

Proof. a) We consider a non-negative integer t in this part. Since $P = I - \mathcal{L}$, we have

$$p_t(x, x) = \left\langle (I - \mathcal{L})^t \mathbf{1}_x, \mathbf{1}_x \right\rangle.$$

Symbolic calculus gives

$$(I - \mathcal{L})^t = \int_{[0,2]} (1 - \lambda)^t \, I_{\mathcal{L}}(\mathrm{d}\lambda).$$

Therefore, by the definition of the vertex spectral measure of $\mathcal{L}$, we have

$$p_t(x, x) = \int_{[0,2]} (1 - \lambda)^t \, \mathrm{d}\left\langle I_{\mathcal{L}}(\lambda) \mathbf{1}_x, \mathbf{1}_x \right\rangle = \int_{[0,2]} (1 - \lambda)^t \, \mu_x(\mathrm{d}\lambda).$$

It follows that

$$p_t(x, x) = \pi(x) + \int_{[0,2]} (1 - \lambda)^t \, \mu_x^*(\mathrm{d}\lambda). \tag{1}$$

Furthermore, using integration by parts, we have

$$p_t(x,x) = \pi(x) + \int_{[0,2]} (1-\lambda)^t \, \mu_x^*(\mathrm{d}\lambda)$$

$$= \pi(x) + (1-\lambda)^t \mu_x^*(\lambda)\big|_0^2 - \int_0^2 \mu_x^*(\lambda)\,\mathrm{d}(1-\lambda)^t$$

$$= \pi(x) + (-1)^t \mu_x^*(2) + t \int_0^2 \mu_x^*(\lambda)(1-\lambda)^{t-1}\,\mathrm{d}\lambda.$$

b) When $t \equiv 0 \bmod 2$, because $(1-\lambda)^t$ is always non-negative, by Eq. (1), $p_t(x,x) \geqslant \pi(x)$. On the other hand, by the result in part a),

$$p_t(x,x) = \pi(x) + \big(1 - \pi(x)\big) + t \int_0^2 \mu_x^*(\lambda)(1-\lambda)^{t-1}\,\mathrm{d}\lambda$$

$$\leqslant 1 + t \int_0^2 \varphi(\lambda)(1-\lambda)^{t-1}\,\mathrm{d}\lambda$$

$$= 1 - \big(1 - \pi(x)\big) + \int_{(0,2]} (1-\lambda)^t \,\mathrm{d}\varphi(\lambda)$$

$$= \pi(x) + \int_{(0,2]} (1-\lambda)^t \,\mathrm{d}\varphi(\lambda),$$

where the second equality is by Lemma 2.6. The first assertion is proved.

c) When $t \equiv 1 \bmod 2$, by the result in part a), we have

$$p_t(x,x) = 2\pi(x) - 1 + t \int_0^2 \mu_x^*(\lambda)(1-\lambda)^{t-1}\,\mathrm{d}\lambda.$$

Therefore, by Lemma 2.6,

$$p_t(x,x) \geqslant 2\pi(x) - 1 + \big(1 - \pi(x)\big) + t \int_0^2 \psi_1(\lambda)(1-\lambda)^{t-1}\,\mathrm{d}\lambda$$

$$= \pi(x) + \int_{(0,2]} (1-\lambda)^t \,\mathrm{d}\psi_1(\lambda).$$

Similarly,

$$p_t(x,x) \leqslant \pi(x) + \int_{(0,2]} (1-\lambda)^t \,\mathrm{d}\psi_2(\lambda). \qquad \square$$

3 Return Probability on Regular Graphs

In this section, we consider regular graphs.

3.1 Estimate of Spectral Measure and Convergence Rate

Theorem 3.1. *Let G be a non-bipartite, connected, regular, unweighted, simple graph. For each $x \in V(G)$, we have*

$$\mu_x^{\mathcal{Q}}(\delta) \leqslant 10\sqrt{\delta}, \qquad 0 \leqslant \delta \leqslant 2.$$

Proof. Assume G is d-regular. Now fixing a vertex $x \in V(G)$, we define f as in Lemma 2.4.

a) If $f(y) \geqslant 0$ for all $y \in V(G)$, then

$$10\sqrt{\delta} \geqslant \delta \geqslant \langle \mathcal{Q}f, f \rangle_w = \sum_{(v,u) \in E(G)} w(v,u) \big| f(v) + f(u) \big|^2$$

$$\geqslant \sum_{y \in N(x)} w(x,y) \big| f(x) + f(y) \big|^2 \geqslant w(x) \big| f(x) \big|^2$$

$$= \mu_x^{\mathcal{Q}}(\delta).$$

b) We may now assume without loss of generality that $\{v \in V(G) \,;\; f(v) < 0\}$ is non-empty. Let $S := \{s \in V \,;\; f(s) > 0\}$ and $T := \{t \in V \,;\; f(t) < 0\}$. By our assumptions, both S and T are non-empty. Recall that for each edge $e \in E(G)$, we have $w(e) = 1$ because G is assumed to be unweighted. However, we do the following construction of an auxiliary graph G' for general weighted graphs because this will also be useful subsequently in the proof of Proposition 5.1:

(1) The vertex set V' of G' includes V. Also, if $(u,v) = e \in E(G)$ and $u, v \in S$, we introduce two vertices $u^{(e)}$ and $v^{(e)}$ in V'. Similarly, if $(u,v) = e \in E(G)$ and $u, v \in T$, we introduce two vertices $u^{(e)}$ and $v^{(e)}$ in V'.

(2) Construct the edge set E' of G' and their weights: If $e = (u,v) \in E(G)$ with u and v both in S, we introduce three edges $(u, u^{(e)})$, $(u^{(e)}, v^{(e)})$, and $(v^{(e)}, v)$ in E'; if $e = (u,v) \in E(G)$ with u and v both in T, we introduce three edges $(u, u^{(e)})$, $(u^{(e)}, v^{(e)})$, and $(v^{(e)}, v)$ in E'. Suppose $e = (u,v)$ has weight $w(u,v)$ in G. For edges introduced above, we assign $w(u,v)$ as their weights w' in G', i.e.,

$$w'(u, u^{(e)}) = w'(u^{(e)}, v^{(e)}) = w'(v^{(e)}, v) = w(u,v).$$

If $e = (u,v) \in E(G)$ is not of the aforementioned forms, introduce one edge (u,v) in E' and set its weight $w'(u,w)$ in G' as $w(u,v)$.

It is obvious that G' is connected. By the above construction, for $v \in V \subseteq V'$, v has the same weight in G and in G'. For $v \in V'$, denote its weight in G' as $w'(v)$. Since G is unweighted in the current setup, G' is also unweighted. So for a vertex v' of G', $w'(v')$ equals the degree of v' in G'.

c) Define a function $g \colon V' \to \mathbb{R}$:

$$g(v) := \begin{cases} \bigl|f(v)\bigr|, & v \in S \cup T, \\ 0, & \text{otherwise.} \end{cases}$$

It follows that

$$g(x) = f(x) = \sqrt{\mu_x^{\mathcal{Q}}(\delta)/w(x)},$$

$$\sum_{v \in V'} \bigl|g(v)\bigr|^2 w'(v) = \sum_{v \in V} \bigl|f(v)\bigr|^2 w(v) = 1,$$

$$\sum_{(v_1,v_2) \in E'} w'(v_1, v_2)\bigl|g(v_1) - g(v_2)\bigr|^2 \leqslant \sum_{(v_1,v_2) \in E} w(v_1, v_2)\bigl|f(v_1) + f(v_2)\bigr|^2.$$

We claim that $g(v') = 0$ for some $v' \in V'$. Otherwise, if $V' \setminus (S \cup T) = \varnothing$, all edges in G would be between S and T, contradicting the assumption that G is non-bipartite.

Set $B := \bigl\{y \in V'; \ \bigl|g(y) - g(x)\bigr| \leqslant \frac{g(x)}{2}\bigr\}$. It follows that $B \subseteq S \cup T \subsetneq V'$. Since G' is connected, there exists a shortest path $\mathcal{P}$ in G' from x to $V' \setminus B$.

d) If $|\mathcal{P}| = 1$ and at least half of the neighbors of x are outside of B, then we have

$$\delta \geqslant \langle \mathcal{Q}f, f \rangle_w = \sum_{(v_1,v_2) \in E} \bigl|f(v_1) + f(v_2)\bigr|^2$$

$$\geqslant \sum_{(v_1,v_2) \in E'} \bigl|g(v_1) - g(v_2)\bigr|^2$$

$$\geqslant \frac{d}{2} \cdot \frac{1}{4}\bigl|g(x)\bigr|^2 = \frac{d}{2} \cdot \frac{\mu_x^{\mathcal{Q}}(\delta)}{4w(x)}$$

$$= \frac{\mu_x^{\mathcal{Q}}(\delta)}{8}.$$

Therefore, when $\delta \leqslant 1$, $\mu_x^{\mathcal{Q}}(\delta) \leqslant 8\delta \leqslant 8\sqrt{\delta} \leqslant 10\sqrt{\delta}$; when $\delta > 1$, $\mu_x^{\mathcal{Q}}(\delta) \leqslant 1 < \delta < 10\sqrt{\delta}$. The result holds in this case.

e) If $|\mathcal{P}| = 1$ and at least half of the neighbors of x are inside of B, then the neighbors of x inside of B are in $S \cup T$ so of degree d in G'. Thus,

$$\mathsf{wt}(B; G') \geqslant \frac{d}{2} \cdot d = \frac{d^2}{2}|\mathcal{P}| \geqslant \frac{d^2}{6}|\mathcal{P}|.$$

If $|\mathcal{P}| \geqslant 2$, we claim that $\mathsf{wt}(B; G') \geqslant \frac{d^2}{6}|\mathcal{P}|$ still holds. To justify this claim, we assume that $\mathcal{P} = u_0 u_1 \cdots u_r u_{r+1}$, where $r = |P| - 1$. Consider $\widetilde{B} := \bigl\{y \in V'; \ \mathsf{dist}(x, y; G') \leqslant r\bigr\} \subseteq B \subseteq S \cup T$. Since $\mathcal{P}$ is a shortest path, $u_0, u_1, \ldots, u_{r-2}$, and u_{r-1} are all in $S \cup T$, and are therefore of degree d (in both G and G'). Setting $K := \{u_0, u_3, \ldots, u_{\lfloor r-1 \rfloor_{(3)}}\}$, we have

$$|K| = \frac{\lceil r \rceil^{(3)}}{3},$$

where $\lceil r \rceil^{(3)}$ is the smallest multiple of 3 that is greater than or equal to r. By a similar argument as in the proof of Lemma 2.5, it follows that

$$|\widetilde{B}| \geqslant (r+1) + |K|(d-2),$$

where $r+1$ is counting the vertices $\{u_0, u_1, \ldots, u_r\}$, and $|K|(d-2)$ is counting the neighbors of K that are not on $\mathcal{P}$. Hence, we get that

$$\begin{aligned}
\mathsf{wt}(B; G') \geqslant \mathsf{wt}(\widetilde{B}; G') &= d|\widetilde{B}| \\
&\geqslant d \cdot \big((r+1) + |K|(d-2)\big) \\
&\geqslant d \cdot \big((r+1) + \frac{r}{3}(d-2)\big) \\
&\geqslant \frac{d^2 r}{3} = \frac{d^2}{3}(|\mathcal{P}| - 1) \\
&\geqslant \frac{d^2}{6}|\mathcal{P}|.
\end{aligned}$$

So our claim holds.

Therefore, we may assume that $\mathsf{wt}(B; G') \geqslant \frac{d^2}{6}|\mathcal{P}|$.

f) Note that $\sum_{v \in V'} |g(v)|^2 w'(v) = 1$. It is easy to get

$$\mathsf{wt}(B; G') \leqslant \frac{1}{\left|\frac{1}{2}g(x)\right|^2} = \frac{4w(x)}{\mu_x^{\mathcal{Q}}(\delta)} = \frac{4d}{\mu_x^{\mathcal{Q}}(\delta)}.$$

Therefore,

$$\frac{4d}{\mu_x^{\mathcal{Q}}(\delta)} \geqslant \mathsf{wt}(B, G') \geqslant \frac{d^2}{6}|\mathcal{P}|.$$

Thus,

$$|\mathcal{P}| \leqslant \frac{24}{d\mu_x^{\mathcal{Q}}(\delta)}. \tag{2}$$

We note that

$$\begin{aligned}
\delta \geqslant \langle \mathcal{Q}f, f \rangle_w &= \sum_{(v_1, v_2) \in E} \big|f(v_1) + f(v_2)\big|^2 \\
&\geqslant \sum_{(v_1, v_2) \in E'} \big|g(v_1) - g(v_2)\big|^2 \geqslant \sum_{i=0}^{|\mathcal{P}|-1} \big|g(u_i) - g(u_{i+1})\big|^2 \\
&\geqslant \frac{1}{|\mathcal{P}|}\left(\sum_{i=0}^{|\mathcal{P}|-1} \big|g(u_i) - g(u_{i+1})\big|\right)^2 \geqslant \frac{1}{|\mathcal{P}|}\big|g(u_0) - g(u_{|\mathcal{P}|})\big|^2 \\
&\geqslant \frac{1}{|\mathcal{P}|}\frac{|g(x)|^2}{4} = \frac{\mu_x^{\mathcal{Q}}(\delta)}{4|\mathcal{P}|w(x)} \\
&= \frac{\mu_x^{\mathcal{Q}}(\delta)}{4d|\mathcal{P}|}.
\end{aligned}$$

Proceeding further by Eq. (2), we have

$$\delta \geqslant \frac{\mu_x^{\mathcal{Q}}(\delta)}{4d|\mathcal{P}|} \geqslant \frac{\mu_x^{\mathcal{Q}}(\delta)}{4d\frac{24}{d\mu_x^{\mathcal{Q}}(\delta)}} = \frac{\left(\mu_x^{\mathcal{Q}}(\delta)\right)^2}{96}.$$

Therefore,

$$\mu_x^{\mathcal{Q}}(\delta) \leqslant \sqrt{96\delta}, \qquad 0 \leqslant \delta \leqslant 2. \qquad \square$$

For Theorem 3.1, we remark that the same upper bound holds with an extra factor of the maximum degree over the minimum degree for general unweighted graphs; it is similar for the return probability bound in Theorem 3.4. The upper bound in Theorem 3.1 could be easily used to get lower bounds on the eigenvalues of P. To this end, we need Lemma 3.2:

Lemma 3.2. *Let G be a finite, connected, weighted graph. We have*

$$\sum_{x \in V} \mu_x^{\mathcal{Q}}(\delta) = \left|\{j \,;\, \lambda_j^{\mathcal{Q}} \leqslant \delta\}\right|.$$

Proof. Recall that the eigenvalues of $\mathcal{Q}$ are

$$0 \leqslant \lambda_{\min}^{\mathcal{Q}} = \lambda_1^{\mathcal{Q}} \leqslant \lambda_2^{\mathcal{Q}} \leqslant \lambda_3^{\mathcal{Q}} \leqslant \cdots \leqslant \lambda_{n-1}^{\mathcal{Q}} < \lambda_n^{\mathcal{Q}} = \lambda_{\max}^{\mathcal{Q}} = 2.$$

Let $h_1, h_2, \ldots, h_n$ be an orthonormal basis of $\ell^2(V, w)$ such that

$$\mathcal{Q}h_j = \lambda_j^{\mathcal{Q}} h_j, \qquad 1 \leqslant j \leqslant n.$$

It follows that

$$\sum_{x \in V} |\langle h_j, \mathbf{e}_x \rangle_w|^2 = \|h_j\|_w^2 = 1.$$

Therefore,

$$\sum_{x \in V} \mu_x^{\mathcal{Q}}(\delta) = \sum_{x \in V} \left\langle I_{\mathcal{Q}}(\delta)\mathbf{e}_x, \mathbf{e}_x \right\rangle_w = \sum_{x \in V} \|I_{\mathcal{Q}}(\delta)\mathbf{e}_x\|_w^2$$

$$= \sum_{x \in V} \left\| \sum_{j\,;\, \lambda_j^{\mathcal{Q}} \leqslant \delta} \langle \mathbf{e}_x, h_j \rangle_w h_j \right\|^2 = \sum_{x \in V} \sum_{j\,;\, \lambda_j^{\mathcal{Q}} \leqslant \delta} |\langle \mathbf{e}_x, h_j \rangle_w|^2$$

$$= \sum_{j\,;\, \lambda_j^{\mathcal{Q}} \leqslant \delta} \sum_{x \in V} |\langle \mathbf{e}_x, h_j \rangle_w|^2 = \sum_{j\,;\, \lambda_j^{\mathcal{Q}} \leqslant \delta} \|h_j\|_w^2 = \sum_{j\,;\, \lambda_j^{\mathcal{Q}} \leqslant \delta} 1$$

$$= \left|\{j \,;\, \lambda_j^{\mathcal{Q}} \leqslant \delta\}\right|. \qquad \square$$

Corollary 3.3. *Let G be a regular, non-bipartite, finite, simple, connected, unweighted graph. For $1 \leqslant k \leqslant n$, we have $\lambda_k^{\mathcal{Q}} \geqslant \frac{k^2}{100n^2}$. Therefore, $\lambda_k^{P} \geqslant -1 + \frac{k^2}{100n^2}$.*

Proof. By Lemma 3.2 and Theorem 3.1,

$$\left|\{j \,;\, \lambda_j^{\mathcal{Q}} \leqslant \delta\}\right| = \sum_{x \in V} \mu_x^{\mathcal{Q}}(\delta) = 10n\sqrt{\delta}.$$

Therefore, if $10n\sqrt{\delta} < k$, $\left|\{j \,;\, \lambda_j^{\mathcal{Q}} \leqslant \delta\}\right| < k$. In other words, $\lambda_k^{\mathcal{Q}} \geqslant \frac{k^2}{100n^2}$. $\square$

Our main interest is to get bounds on return probabilities of simple random walk on regular graphs.

Theorem 3.4. *Let G be a regular, non-bipartite, simple, connected, unweighted graph. For each $x \in V(G)$, simple random walk on G satisfies*

$$0 \leqslant p_t(x,x) - \pi(x) \leqslant \frac{18}{\sqrt{t}} \qquad for \quad t \equiv 0 \bmod 2,$$

$$\left| p_t(x,x) - \pi(x) \right| \leqslant \frac{9}{\sqrt{t}} \qquad for \quad t \equiv 1 \bmod 2.$$

In the proof of Theorem 3.4, we will need the following elementary calculation.

Lemma 3.5. *For $t > 0$, we have*

$$\int_0^1 (1-\lambda)^t \frac{1}{\sqrt{\lambda}} \, d\lambda \leqslant \frac{1.8}{\sqrt{t}}.$$

Proof. We have

$$\int_0^1 (1-\lambda)^t \frac{1}{\sqrt{\lambda}} \, d\lambda \leqslant \int_0^1 \exp\{-\lambda t\} \lambda^{-1/2} \, d\lambda \leqslant \int_0^\infty \exp\{-\lambda t\} \lambda^{-1/2} \, d\lambda.$$

Taking a change of variable $\lambda t = s$, we get that

$$\int_0^\infty e^{-s}(s/t)^{-1/2} \, d(s/t) \leqslant \frac{1}{\sqrt{t}} \int_0^\infty e^{-s} s^{-1/2} \, ds$$

$$= \frac{\Gamma(1/2)}{\sqrt{t}} \leqslant \frac{1.8}{\sqrt{t}}.$$

Hence, the inequality follows immediately. $\qquad\square$

Proof of Theorem 3.4. a) Set

$$\varphi(\lambda) := \begin{cases} 10\sqrt{\lambda} & \text{if } \lambda \geqslant 0 \text{ and } 10\sqrt{\lambda} \leqslant \mu_x^*(1), \\ \mu_x^*(1) & \text{for intermediate values of } \lambda, \\ \mu_x^*(2) - 10\sqrt{2-\lambda} & \text{if } \lambda \leqslant 2 \text{ and } \mu_x^*(2) - 10\sqrt{2-\lambda} \geqslant \mu_x^*(1). \end{cases}$$

We claim that the function φ defined above satisfies the conditions in Lemma 2.7(1). In fact, it is known that $\mu_x^*(\lambda) < 10\sqrt{\lambda}$ from [LOG18, Theorem 4.9]. On the other hand, by Theorem 3.1,

$$\mu_x^*(\lambda) = \left\langle I_{\mathcal{L}}\big((0,\lambda]\big) \mathbf{e}_x, \mathbf{e}_x \right\rangle_w$$
$$= \left\langle I_{\mathcal{L}}\big((0,2]\big) \mathbf{e}_x, \mathbf{e}_x \right\rangle_w - \left\langle I_{\mathcal{L}}\big((\lambda,2]\big) \mathbf{e}_x, \mathbf{e}_x \right\rangle_w$$
$$= 1 - \left\langle I_{\mathcal{L}}(0) \mathbf{e}_x, \mathbf{e}_x \right\rangle_w - \left\langle I_{\mathcal{L}}\big((\lambda,2]\big) \mathbf{e}_x, \mathbf{e}_x \right\rangle_w$$
$$\geqslant 1 - \left\langle I_{\mathcal{L}}(0) \mathbf{e}_x, \mathbf{e}_x \right\rangle_w - \left\langle I_{\mathcal{L}}\big([\lambda,2]\big) \mathbf{e}_x, \mathbf{e}_x \right\rangle_w$$
$$= 1 - \pi(x) - \left\langle I_{\mathcal{Q}}(2-\lambda) \mathbf{e}_x, \mathbf{e}_x \right\rangle_w$$
$$\geqslant 1 - \pi(x) - 10\sqrt{2-\lambda}$$
$$= \mu_x^*(2) - 10\sqrt{2-\lambda}.$$

The claim is now proved.

Therefore, for $t \equiv 0 \bmod 2$, by Lemma 2.7,

$$\pi(x) \leqslant p_t(x,x) \leqslant \pi(x) + \int_0^2 (1-\lambda)^t \varphi'(\lambda)\, \mathrm{d}\lambda$$

$$\leqslant \pi(x) + 2\int_0^1 (1-\lambda)^t \frac{5}{\sqrt{\lambda}}\, \mathrm{d}\lambda$$

$$= \pi(x) + 10\int_0^1 (1-\lambda)^t \frac{1}{\sqrt{\lambda}}\, \mathrm{d}\lambda$$

$$\leqslant \pi(x) + \frac{18}{\sqrt{t}}.$$

The first assertion is proved.

b) Set

$$\psi_1(\lambda) := \big(\mu_x^*(2) - 10\sqrt{2-\lambda}\,\big) \vee 0, \qquad \lambda \in [0,2],$$

and

$$\psi_2(\lambda) := (10\sqrt{\lambda}\,) \wedge \mu_x^*(2), \qquad \lambda \in [0,2].$$

Then they satisfy the conditions of Lemma 2.7(2). Consequently, we have

$$\pi(x) + \int_0^2 (1-\lambda)^t \psi_1'(\lambda)\, \mathrm{d}\lambda \leqslant p_t(x,x) \leqslant \pi(x) + \int_0^2 (1-\lambda)^t \psi_2'(\lambda)\, \mathrm{d}\lambda.$$

By some elementary calculation, we get that

$$\pi(x) - \int_0^1 (1-\lambda)^t \frac{5}{\sqrt{\lambda}}\, \mathrm{d}\lambda \leqslant p_t(x,x) \leqslant \pi(x) + \int_0^1 (1-\lambda)^t \frac{5}{\sqrt{\lambda}}\, \mathrm{d}\lambda.$$

The second assertion follows immediately from Lemma 3.5. $\qquad\qquad\square$

3.2   Sharpness and Spectral Radius

The order of $\frac{1}{\sqrt{t}}$ in Theorem 3.4 is sharp. We show this by the following Example 3.6.

Example 3.6. *We will exploit a local limit theorem. Consider an unweighted graph G with $V(G) = \mathbb{Z}$: $(i,j) \in E(G)$ if and only if $0 < |i - j| \leqslant 2$. Obviously, G is non-bipartite, connected, and 4-regular. We also consider an i.i.d. sequence $\{\xi_k\}_{k=1}^\infty$ of random variables that are uniform on $\{-2, -1, 1, 2\}$. Then for the simple random walk on G,*

$$p_t(0,0) = \mathbf{P}\Big[\sum_{k=1}^t \xi_k = 0\Big].$$

Using the notation in [Davis and McDonald, 1995, Theorem 1.1], we have

$$a_t = 0, \qquad t \geqslant 1,$$

$$b_t = \sqrt{\frac{5t}{2}}, \qquad t \geqslant 1.$$

Therefore, by [Davis and McDonald, 1995, Theorem 1.1], we have

$$\left| \sqrt{\frac{5t}{2}} \, \mathbf{P}\Big[\sum_{k=1}^{t} \xi_k = 0\Big] - \frac{1}{\sqrt{2\pi}} \right| \to 0, \qquad as \ t \to \infty.$$

In other words,

$$p_t(0,0) \sim \frac{1}{\sqrt{5\pi t}}, \qquad as \ t \to \infty.$$

Hence, the sharpness is demonstrated. □

In Section 3.1, we focused on the negative spectrum of P by considering $\mathcal{Q} = I + P$. As a comparison, in [LOG18, Theorem 4.9], $\mathcal{L} = I - P$ is exploited and so essentially the positive spectrum of P is focused on. Recalling that the spectrum $\sigma(P)$ of P, as an operator on $\ell^2(V, w)$, is contained in $[-1, 1]$, we set

$$\gamma_- := \inf \sigma(P) - (-1),$$
$$\gamma_+ := 1 - \sup\big(\sigma(P) \setminus \{1\}\big).$$

It is obvious that γ_- and γ_+ are both non-negative. But which is larger between γ_- and γ_+? When G is finite, it depends; but when G is infinite, we have the following simple fact:

Proposition 3.7. *Let $G = (V, E)$ be a connected, weighted, infinite, locally finite graph, with $\mathsf{wt}(V) = \infty$. Then we have $\gamma_+ \leqslant \gamma_-$. In other words, the spectral radius of P is achieved by the positive spectrum.*

Proof. To begin, since $\mathsf{wt}(V) = \infty$, we see that the constant function is not in $\ell^2(V, w)$ and 1 is not in the point spectrum of P.

For an arbitrarily small number $\varepsilon > 0$, there exists $f \in \ell^2(V, w)$, with $\|f\|_w = 1$, satisfying

$$-1 + \gamma_- + \varepsilon \geqslant \langle Pf, f\rangle_w = \sum_{x \in V} w(x)(Pf)(x)\overline{f(x)} = \sum_{x \in V} w(x)\Big(\sum_{y \in N(x)} p(x,y)f(y)\Big)\overline{f(x)}.$$

Hence, we have

$$1 - \gamma_- - \varepsilon \leqslant |\langle Pf, f\rangle_w| = \left| \sum_{x \in V} w(x)\Big(\sum_{y \in N(x)} p(x,y)f(y)\Big)\overline{f(x)} \right|$$
$$\leqslant \sum_{x \in V} w(x)\Big(\sum_{y \in N(x)} p(x,y)|f(y)|\Big)|f(x)|$$
$$= \langle P|f|, |f|\rangle_w.$$

Since 1 is not in the point spectrum of P, we further have

$$\langle P|f|, |f|\rangle_w \leqslant 1 - \gamma_+.$$

Hence, $1 - (\gamma_- + \varepsilon) \leqslant 1 - \gamma_+$. It follows that

$$\gamma_+ \leqslant \gamma_- + \varepsilon.$$

Because $\varepsilon > 0$ is arbitrary, we get $\gamma_+ \leqslant \gamma_-$. □

Remark. A proof of Proposition 3.7 could also be found in [Mao and Song, 2013], but our proof is much shorter. Proposition 3.7 is stated in the language of graphs, but essentially it also applies to irreducible, non-ergodic, symmetrizable Markov chains with invariant measure, π. Note that the condition of "non-ergodic" means the total mass of π is infinity.

4　A Slight Improvement of Return Probability Bound

In this section, we consider finite graphs only. We use essentially a similar method to [Oliveira and Peres, 2019]; yet the negative spectrum of P is also considered. Let G be a non-bipartite, finite, simple, connected, unweighted graph. Then $\lambda_{\min}^P = \lambda_1^P > -1$. Set $\Lambda := |\lambda_1^P| \vee |\lambda_{n-1}^P|$, $t_{\mathrm{rel}} := (1 - \Lambda)^{-1}$, $t' := \lceil t_{\mathrm{rel}} \rceil - 1$.

Theorem 4.1. *Let G be a non-bipartite, finite, simple, connected, unweighted graph. For $t \geqslant t_{\mathrm{rel}}$, simple random walk on G satisfies*

$$0 \leqslant p_{2t}(x, x) - \pi(x) \leqslant \frac{8 d(x) \sqrt{t_{\mathrm{rel}} + 1}}{(t + 1) d_{\min}},$$

$$\left| p_{2t+1}(x, x) - \pi(x) \right| \leqslant \frac{8 d(x) \sqrt{t_{\mathrm{rel}} + 1}}{\sqrt{(t + 1)(t + 2)}\, d_{\min}}.$$

To prove Theorem 4.1, we need some preparation. Recall that the hitting time of $A \subseteq V$ is $\tau_A := \inf\{t \geqslant 0\,;\ X_t \in A\}$ and Green's function is

$$g_t(x, y) := \sum_{s=0}^{t} p_t(x, y), \qquad t \geqslant 0,\ x, y \in V.$$

Lemma 4.2. *For $t \geqslant t'$ and $x \in V$, we have*

$$0 \leqslant p_{2t}(x, x) - \pi(x) \leqslant \frac{1}{1 - \mathrm{e}^{-2}} \cdot \frac{g_{2t'}(x, x)}{t + 1}.$$

Proof. By Eq. (1), $p_{2t}(x, x) - \pi(x)$ is nonnegative and decreasing in t. Therefore, we have

$$0 \leqslant p_{2t}(x, x) - \pi(x) \leqslant \frac{1}{t + 1} \left(\sum_{s=0}^{t} p_{2s}(x, x) - (t + 1)\pi(x) \right). \tag{3}$$

In addition,

$$\sum_{s=0}^{t} p_{2s}(x, x) - (t + 1)\pi(x) = \sum_{s=0}^{t} \int_{(-1,1)} \lambda^{2s} \left\| I_P(\mathrm{d}\lambda)\mathbf{e}_x \right\|_w^2 = \int_{(-1,1)} \left(\sum_{s=0}^{t} \lambda^{2s} \right) \left\| I_P(\mathrm{d}\lambda)\mathbf{e}_x \right\|_w^2$$

$$= \int_{(-1,1)} \frac{1 - \lambda^{2t+2}}{1 - \lambda^2} \left\| I_P(\mathrm{d}\lambda)\mathbf{e}_x \right\|_w^2$$

$$\leqslant \int_{(-1,1)} \frac{1}{1 - \lambda^2} \left\| I_P(\mathrm{d}\lambda)\mathbf{e}_x \right\|_w^2, \tag{4}$$

where I_P is the resolution of identity for P.

On the other hand, since for $1 \leqslant i \leqslant n$,

$$(\lambda_i^P)^{2t'+2} \leqslant \Lambda^{2t'+2} \leqslant \Lambda^{2/(1-\Lambda)} \leqslant e^{-2},$$

we have

$$
\begin{aligned}
\int_{(-1,1)} \frac{1 - e^{-2}}{1 - \lambda^2} \, \|I_P(d\lambda)\mathbf{e}_x\|_w^2 &\leqslant \int_{(-1,1)} \frac{1 - \lambda^{2t'+2}}{1 - \lambda^2} \, \|I_P(d\lambda)\mathbf{e}_x\|_w^2 \\
&\leqslant \sum_{s=0}^{t'} p_{2s}(x,x) - (t'+1)\pi(x) \leqslant \sum_{s=0}^{t'} p_{2s}(x,x) \\
&\leqslant g_{2t'}(x,x).
\end{aligned}
\tag{5}
$$

By Eqs. (3) to (5),

$$
\begin{aligned}
0 \leqslant p_{2t}(x,x) - \pi(x) &\leqslant \frac{1}{t+1}\left(\sum_{s=0}^{t} p_{2s}(x,x) - (t+1)\pi(x)\right) \\
&\leqslant \frac{1}{t+1} \int_{(-1,1)} \frac{1}{1 - \lambda^2} \, \|I_P(d\lambda)\mathbf{e}_x\|_w^2 \\
&\leqslant \frac{1}{(t+1)(1 - e^{-2})} \int_{(-1,1)} \frac{1 - e^{-2}}{1 - \lambda^2} \, \|I_P(d\lambda)\mathbf{e}_x\|_w^2 \\
&\leqslant \frac{1}{(t+1)(1 - e^{-2})} g_{2t'}(x,x).
\end{aligned}
$$

Lemma 4.2 is proved. $\qquad\square$

Proposition 4.3. *We have*

$$t_{\mathrm{rel}} \leqslant \max_{x,y\in V}\big(\mathbf{E}_x[\tau_y] + \mathbf{E}_y[\tau_x]\big) \leqslant 3n^2 \frac{d_{\mathrm{avg}}}{d_{\mathrm{min}}} - 2.$$

Proof. The first inequality follows from [Levin and Peres, 2017, Lemma 12.17]. For the second inequality, by [Lyons and Peres, 2017, Corollary 2.21] and a simple network reduction argument, we have

$$\max_{x,y\in V}\big(\mathbf{E}_x[\tau_y] + \mathbf{E}_y[\tau_x]\big) \leqslant \mathrm{diam}(G)\,\mathsf{wt}(V).$$

Since Lemma 2.5 gives $\mathrm{diam}(G) \leqslant \frac{3n}{d_{\mathrm{min}}} - 1$, we may proceed and get

$$
\begin{aligned}
\max_{x,y\in V}\big(\mathbf{E}_x[\tau_y] + \mathbf{E}_y[\tau_x]\big) &\leqslant \left(\frac{3n}{d_{\mathrm{min}}} - 1\right)\mathsf{wt}(V) \\
&= \frac{3n\,\mathsf{wt}(V)}{d_{\mathrm{min}}} - \mathsf{wt}(V) \\
&\leqslant 3n^2 \frac{d_{\mathrm{avg}}}{d_{\mathrm{min}}} - 2. \qquad\square
\end{aligned}
$$

For non-empty $A \subsetneq V$ and $x \in V \setminus A$, define $G(x, x; A) \coloneqq \mathbf{E}_x\big[\sum_{s=0}^{\tau_A - 1} \mathbf{1}_{\{X_s \notin A\}}\big]$.

Proposition 4.4. *We have*
$$\frac{G(x, x; A)}{\pi(x)} \leqslant \frac{9}{2}\Big(\frac{d_{\text{avg}}n}{d_{\min}}\Big)^2\big(1 - \pi(A)\big).$$

Proof. By [Lyons and Peres, 2017, Eq. (2.5)] and network reduction, we have
$$\frac{G(x, x; A)}{\pi(x)} \leqslant \mathsf{wt}(V)\mathcal{R}_{\text{eff}}(x \leftrightarrow A) = d_{\text{avg}}n\mathcal{R}_{\text{eff}}(x \leftrightarrow A),$$

where $\mathcal{R}_{\text{eff}}(x \leftrightarrow A)$ is the effective resistance between x and A. Then we need only follow the proof of [Oliveira and Peres, 2019, Proposition 3.2] to get
$$\mathcal{R}_{\text{eff}}(x \leftrightarrow A) \leqslant \frac{9}{2}\frac{d_{\text{avg}}n}{d_{\min}^2}\big(1 - \pi(A)\big). \qquad \square$$

Now fix $x \in V$. For $\alpha > 1$, let
$$A_\alpha \coloneqq \big\{y \in V \; ; \; g_{2t'}(y, x) \leqslant \alpha\pi(x)(2t' + 1)\big\}.$$

We claim that $A_\alpha \neq \varnothing$ for $\alpha > 1$. In fact,
$$\begin{aligned}
1 - \pi(A_\alpha) = \sum_{y \notin A_\alpha} \pi(y) &< \sum_{y \notin A_\alpha} \pi(y)\frac{g_{2t'}(y, x)}{\alpha\pi(x)(2t' + 1)} \\
&= \sum_{y \notin A_\alpha} \frac{\pi(x)g_{2t'}(x, y)}{\alpha\pi(x)(2t' + 1)} \leqslant \sum_{y \in V} \frac{g_{2t'}(x, y)}{\alpha(2t' + 1)} \\
&= \frac{1}{\alpha} < 1.
\end{aligned} \qquad (6)$$

So for $\alpha > 1$, A_α is non-empty.

Lemma 4.5. *For $x \in V$,*
$$\frac{g_{2t'}(x, x)}{\pi(x)} \leqslant \frac{9d_{\text{avg}}n}{2d_{\min}}\sqrt{2t' + 1}.$$

Proof. a) Set $\alpha_0 \coloneqq \frac{3d_{\text{avg}}n}{d_{\min}} \cdot \frac{1}{\sqrt{2t' + 1}}$. We claim $\alpha_0 > 1$. In fact, Proposition 4.3 ensures
$$\lceil t_{\text{rel}}\rceil \leqslant 3\frac{d_{\text{avg}}}{d_{\min}}n^2.$$

Therefore, we have
$$\begin{aligned}
\alpha_0 = \frac{3d_{\text{avg}}n}{d_{\min}} &\cdot \frac{1}{\sqrt{2\lceil t_{\text{rel}}\rceil - 1}} \\
&\geqslant \frac{3d_{\text{avg}}n}{d_{\min}} \cdot \frac{1}{\sqrt{\frac{6d_{\text{avg}}n^2}{d_{\min}}}} = \frac{3}{\sqrt{6}} \cdot \sqrt{\frac{d_{\text{avg}}}{d_{\min}}} \\
&> 1.
\end{aligned}$$

As a consequence, A_{α_0} is non-empty.

b) If $x \in A_{\alpha_0}$, by the definition of A_{α_0},

$$\frac{g_{2t'}(x,x)}{\pi(x)} \leqslant \alpha_0(2t'+1) = \frac{3d_{\mathrm{avg}}n}{d_{\min}}\sqrt{2t'+1}.$$

c) If $x \notin A_{\alpha_0}$, by the strong Markov property, Proposition 4.4, and the definition of A_{α_0},

$$\begin{aligned}
\frac{g_{2t'}(x,x)}{\pi(x)} &\leqslant \frac{G(x,x;A_{\alpha_0})}{\pi(x)} + \mathbf{E}_x\left[\frac{g_{2t'}(X_{\tau_{A_{\alpha_0}}},x)}{\pi(x)}\right] \\
&\leqslant \frac{9}{2}\cdot\left(\frac{d_{\mathrm{avg}}n}{d_{\min}}\right)^2(1-\pi(A_{\alpha_0})) + \alpha_0(2t'+1) \\
&\leqslant \frac{9}{2}\cdot\left(\frac{d_{\mathrm{avg}}n}{d_{\min}}\right)^2\frac{1}{\alpha_0} + \alpha_0(2t'+1) \\
&= (\tfrac{3}{2}+3)\sqrt{2t'+1}\frac{d_{\mathrm{avg}}n}{d_{\min}} \\
&= \frac{9d_{\mathrm{avg}}n}{2d_{\min}}\sqrt{2t'+1},
\end{aligned}$$

where the third inequality is by Eq. (6). $\qquad\square$

We are now in position to prove Theorem 4.1.

Proof of Theorem 4.1. a) Recalling that $t' = \lceil t_{\mathrm{rel}}\rceil - 1$, we have $t_{\mathrm{rel}} \geqslant t'$. Therefore, $t \geqslant t_{\mathrm{rel}}$ implies $t \geqslant t'$.

b) By Lemmas 4.2 and 4.5, for $t \geqslant t_{\mathrm{rel}}$ and $x \in V$,

$$\begin{aligned}
0 \leqslant p_{2t}(x,x) - \pi(x) &\leqslant \frac{1}{1-\mathrm{e}^{-2}}\cdot\frac{g_{2t'}(x,x)}{t+1} \\
&\leqslant \frac{9d_{\mathrm{avg}}n\pi(x)}{2(1-\mathrm{e}^{-2})(t+1)d_{\min}}\sqrt{2t'+1} = \frac{9d(x)}{2(1-\mathrm{e}^{-2})(t+1)d_{\min}}\sqrt{2\lceil t_{\mathrm{rel}}\rceil - 1} \\
&\leqslant \frac{9d(x)}{2(1-\mathrm{e}^{-2})(t+1)d_{\min}}\sqrt{2(t_{\mathrm{rel}}+1)} \leqslant \frac{8d(x)}{d_{\min}}\frac{\sqrt{t_{\mathrm{rel}}+1}}{t+1}.
\end{aligned}$$

c) Our calculation in part b) implies that for $t \geqslant t_{\mathrm{rel}}$,

$$p_{2t}(x,x) - \pi(x) = \int_{(-1,1)}\lambda^{2t}\|I_P(\mathrm{d}\lambda)\mathbf{e}_x\|_w^2 \leqslant \frac{8d(x)\sqrt{t_{\mathrm{rel}}+1}}{(t+1)d_{\min}}.$$

Therefore, when $t \geqslant t_{\mathrm{rel}}$, we have

$$\begin{aligned}
\left|p_{2t+1}(x,x) - \pi(x)\right| &= \left|\int_{(-1,1)}\lambda^{2t+1}\|I_P(\mathrm{d}\lambda)\mathbf{e}_x\|_w^2\right| \leqslant \int_{(-1,1)}|\lambda|^{2t+1}\|I_P(\mathrm{d}\lambda)\mathbf{e}_x\|_w^2 \\
&= \int_{(-1,1)}|\lambda|^t|\lambda|^{t+1}\|I_P(\mathrm{d}\lambda)\mathbf{e}_x\|_w^2.
\end{aligned}$$

Hence, the Cauchy–Schwarz inequality gives

$$\left|p_{2t+1}(x,x) - \pi(x)\right| \leqslant \sqrt{\int_{(-1,1)} |\lambda|^{2t} \, \|I_P(\mathrm{d}\lambda)\mathbf{e}_x\|_w^2} \sqrt{\int_{(-1,1)} |\lambda|^{2t+2} \, \|I_P(\mathrm{d}\lambda)\mathbf{e}_x\|_w^2}$$

$$\leqslant \frac{8d(x)\sqrt{t_{\mathrm{rel}} + 1}}{\sqrt{(t+1)(t+2)}\, d_{\min}}. \qquad \square$$

5   Volume Growth Conditions

For lazy random walk, [LOG18, Corollaries 4.10 and 4.11] presented return probability bounds depending on volume growth conditions. We have parallel results for simple random walk. Let us begin with the following proposition, which is comparable to the first assertion of [LOG18, Proposition 4.7].

Proposition 5.1. *Let G be a non-bipartite, finite or infinite, simple, connected, weighted graph with weight at least 1 for each edge. For each vertex $x \in V(G)$, $\delta \in (0,2)$, $\alpha \in (0,1)$, and $r \geqslant 0$, $\mathsf{wt}(x,r) > \frac{w(x)}{(1-\alpha)^2 \mu_x^{\mathcal{Q}}(\delta)}$ implies that $\mu_x^{\mathcal{Q}}(\delta) \leqslant \frac{\delta w(x)}{\alpha^2} r$.*

Proof. We may assume $\mu_x^{\mathcal{Q}}(\delta) > 0$. Fixing a vertex $x \in V(G)$, we define f as in Lemma 2.4.

a) If $f(y) \geqslant 0$ for all $y \in V(G)$, then

$$\delta \geqslant \langle \mathcal{Q}f, f \rangle_w = \sum_{(v,u) \in E(G)} w(v,u)\big|f(v) + f(u)\big|^2 \geqslant \big|f(x)\big|^2 = \frac{\mu_x^{\mathcal{Q}}(\delta)}{w(x)}.$$

Note that $\mathsf{wt}(x,r) > \frac{w(x)}{(1-\alpha)^2 \mu_x^{\mathcal{Q}}(\delta)}$ implies $r \geqslant 1$. Therefore,

$$\mu_x^{\mathcal{Q}}(\delta) \leqslant \delta w(x) \leqslant \frac{\delta w(x)}{\alpha^2} r.$$

b) From now on, we may assume that $T := \big\{t \in V(G);\ f(t) < 0\big\}$ is non-empty. We construct G' and g as in the proof of Theorem 3.1 and use notations there. Set

$$B^{(\alpha)} := \big\{y \in V';\ |g(y) - g(x)| \leqslant \alpha g(x)\big\}.$$

It follows that $B^{(\alpha)} \subseteq S \cup T \subsetneq V'$. Since G' is connected, there exists a shortest path

$\mathcal{P} = u_0 u_1 \cdots u_{|\mathcal{P}|}$ in G' from x to $V' \setminus B^{(\alpha)}$ with $u_0 = x$. Hence,

$$\delta \geqslant \langle \mathcal{Q}f, f \rangle_w = \sum_{(v_1, v_2) \in E(G)} w(v_1, v_2) |f(v_1) + f(v_2)|^2$$

$$\geqslant \sum_{(v_1, v_2) \in E'} w'(v_1, v_2) |g(v_1) - g(v_2)|^2 \geqslant \sum_{i=0}^{|\mathcal{P}|-1} |g(u_i) - g(u_{i+1})|^2$$

$$\geqslant \frac{1}{|\mathcal{P}|} \left(\sum_{i=0}^{|\mathcal{P}|-1} |g(u_i) - g(u_{i+1})| \right)^2 \geqslant \frac{1}{|\mathcal{P}|} |g(u_0) - g(u_{|\mathcal{P}|})|^2$$

$$> \frac{\alpha^2 |g(x)|^2}{|\mathcal{P}|} = \frac{\alpha^2 \mu_x^{\mathcal{Q}}(\delta)}{w(x)|\mathcal{P}|}.$$

c) We claim that

$$\{ y \in V' ; \ \mathsf{dist}(y, x; G') \leqslant |\mathcal{P}| - 1 \} = \{ y \in V ; \ \mathsf{dist}(y, x; G) \leqslant |\mathcal{P}| - 1 \}. \tag{7}$$

In fact,

$$\{ y \in V' ; \ \mathsf{dist}(y, x; G') \leqslant |\mathcal{P}| - 1 \} \subseteq B^{(\alpha)} \subseteq S \cup T \subseteq V.$$

By the construction of G', $\mathsf{dist}(y_1, y_2; G') \geqslant \mathsf{dist}(y_1, y_2; G)$ for $y_1, y_2 \in V$. Therefore,

$$\{ y \in V' ; \ \mathsf{dist}(y, x; G') \leqslant |\mathcal{P}| - 1 \} \subseteq \{ y \in V ; \ \mathsf{dist}(y, x; G) \leqslant |\mathcal{P}| - 1 \}.$$

Suppose the above inclusion is strict. Then there will be a vertex $\widetilde{v} \in V(G)$ that is not in the left-hand side of Eq. (7), and a path $\widetilde{\mathcal{P}} = \widetilde{u}_0 \widetilde{u}_1 \cdots \widetilde{u}_{|\widetilde{\mathcal{P}}|}$ in G with $|\widetilde{\mathcal{P}}| \leqslant |\mathcal{P}| - 1$, $\widetilde{u}_0 = x$, and $\widetilde{u}_{|\widetilde{\mathcal{P}}|} = \widetilde{v}$. Set

$$L := \left\{ 0 \leqslant i \leqslant |\widetilde{\mathcal{P}}| - 1 ; \ \{\widetilde{u}_i, \widetilde{u}_{i+1}\} \subseteq S \text{ or } \{\widetilde{u}_i, \widetilde{u}_{i+1}\} \subseteq T \right\}.$$

Then L must be non-empty, otherwise $\widetilde{v} \in V(G)$ would be in the left-hand side of Eq. (7). Pick the smallest number i^* in L and write $e = (\widetilde{u}_{i^*}, \widetilde{u}_{i^*+1})$. By the construction of G' and g, we have $g(\widetilde{u}_{i^*}^{(e)}) = 0$. Thus, $\widetilde{u}_0 \widetilde{u}_1 \cdots \widetilde{u}_{i^*} \widetilde{u}_{i^*}^{(e)}$ is a path in G' linking x and $V' \setminus B^{(\alpha)}$, whose length is $i^* + 1 \leqslant |\widetilde{\mathcal{P}}| < |\mathcal{P}|$. This is a contradiction. Therefore, Eq. (7) holds.

d) If $\mathsf{wt}(x, r; G) > \frac{w(x)}{(1-\alpha)^2 \mu_x^{\mathcal{Q}}(\delta)}$, we must have $r \geqslant |\mathcal{P}|$. In fact, we notice that

$$\mathsf{wt}(B^{(\alpha)}; G')(1-\alpha)^2 g^2(x) \leqslant \sum_{y \in B^{(\alpha)}} |g(y)|^2 w'(y) \leqslant \sum_{y \in V'} |g(y)|^2 w'(y) = 1.$$

Therefore,

$$\mathsf{wt}(x, |\mathcal{P}| - 1; G') \leqslant \mathsf{wt}(B^{(\alpha)}; G') \leqslant \frac{1}{(1-\alpha)^2 g^2(x)} = \frac{w(x)}{(1-\alpha)^2 \mu_x^{\mathcal{Q}}(\delta)}.$$

Hence, Eq. (7) gives

$$\mathsf{wt}(x, |\mathcal{P}| - 1; G) = \mathsf{wt}(x, |\mathcal{P}| - 1; G') \leqslant \frac{w(x)}{(1 - \alpha)^2 \mu_x^{\mathcal{Q}}(\delta)}.$$

Consequently, $\mathsf{wt}(x, r) > \frac{w(x)}{(1-\alpha)^2 \mu_x^{\mathcal{Q}}(\delta)}$ implies $r \geqslant |\mathcal{P}|$.

e) By the results in b) and d), we have

$$\delta \geqslant \frac{\alpha^2 \mu_x^{\mathcal{Q}}(\delta)}{w(x)|\mathcal{P}|} \geqslant \frac{\alpha^2 \mu_x^{\mathcal{Q}}(\delta)}{w(x)r}.$$

Hence, we arrive at the conclusion that $\mu_x^{\mathcal{Q}}(\delta) \leqslant \frac{\delta w(x)}{\alpha^2} r$. $\qquad\square$

Corollary 5.2. *Let G be a non-bipartite, finite or infinite, simple, connected, weighted graph with weight at least 1 for each edge. For each vertex $x \in V(G)$ and $\delta < \frac{1}{r\,\mathsf{wt}(x,r)}$, we have $\mu_x^{\mathcal{Q}}(\delta) \leqslant \frac{4w(x)}{\mathsf{wt}(x,r)}$.* $\qquad\square$

Corollary 5.2 is comparable to the second assertion of [LOG18, Proposition 4.7]. To get a proof, one need only exploit Proposition 5.1 and mimic the proof for the second assertion of [LOG18, Proposition 4.7].

Now we present return probability bounds based on volume growth conditions.

Corollary 5.3. *Let G be a non-bipartite, infinite, simple, connected, weighted graph with weight at least 1 for each edge. Suppose that $c > 0$ and $D \geqslant 1$ are constants such that $\mathsf{wt}(x, r) \geqslant c(r + 1)^D$ for all $r \geqslant 0$. Then for all $\delta \in (0, 2)$,*

$$\mu_x^*(\delta) \leqslant Cw(x)\delta^{D/(D+1)},$$
$$\mu_x^{\mathcal{Q}}(\delta) \leqslant Cw(x)\delta^{D/(D+1)},$$

where

$$C := \frac{(D + 1)^2}{c^{1/(D+1)}D^{2D/(D+1)}}.$$

Hence for all $t \geqslant 1$, simple random walk satisfies

$$p_t(x, x) \leqslant 2C'w(x)t^{-D/(D+1)} \qquad for \quad t \equiv 0 \bmod 2,$$
$$p_t(x, x) \leqslant C'w(x)t^{-D/(D+1)} \qquad for \quad t \equiv 1 \bmod 2,$$

where

$$C' := \frac{D + 1}{c^{1/(D+1)}D^{(D-1)/(D+1)}}\Gamma\Big(\frac{D}{D + 1}\Big).$$

Proof. The bound on μ_x^* is proved in [LOG18, Corollary 4.10]. A similar argument can be used to prove the bound on $\mu_x^{\mathcal{Q}}$: in lieu of [LOG18, Eq. (4.4)], we use Corollary 5.2. To prove the bound on return probabilities, we set

$$\varphi(\lambda) := \begin{cases} Cw(x)\lambda^{D/(D+1)} & \text{if } \lambda \in [0, 1) \text{ and } Cw(x)\lambda^{D/(D+1)} \leqslant \mu_x^*(1), \\ \mu_x^*(1) & \text{for intermediate values of } \lambda, \\ \mu_x^*(2) - Cw(x)(2 - \lambda)^{D/(D+1)} & \text{if } \lambda \in [1, 2) \text{ and } Cw(x)(2 - \lambda)^{D/(D+1)} \geqslant \mu_x^*(1), \end{cases}$$

and

$$\psi_2(\lambda) := \begin{cases} Cw(x)\lambda^{D/(D+1)} \wedge \mu_x^*(2) & \text{if } \lambda \in [0,1), \\ \mu_x^*(2) & \text{if } \lambda \in [1,2]. \end{cases}$$

It is easy to see that φ and ψ_2 satisfy the conditions in Lemma 2.7. Therefore, by Lemma 2.7 and some elementary calculation, we have

$$p_t(x,x) \leqslant \frac{2Cw(x)D}{D+1} \int_0^1 \lambda^{-1/(D+1)}(1-\lambda)^t \, \mathrm{d}\lambda \qquad \text{for} \quad t \equiv 0 \bmod 2,$$

$$p_t(x,x) \leqslant \frac{Cw(x)D}{D+1} \int_0^1 \lambda^{-1/(D+1)}(1-\lambda)^t \, \mathrm{d}\lambda \qquad \text{for} \quad t \equiv 1 \bmod 2.$$

But we have

$$\int_0^1 \lambda^{-1/(D+1)}(1-\lambda)^t \, \mathrm{d}\lambda \leqslant \int_0^1 \lambda^{-1/(D+1)}\mathrm{e}^{-\lambda t} \, \mathrm{d}\lambda \leqslant \int_0^\infty \lambda^{-1/(D+1)}\mathrm{e}^{-\lambda t} \, \mathrm{d}\lambda$$

$$= \int_0^\infty (s/t)^{-1/(D+1)}\mathrm{e}^{-s} \, \mathrm{d}(s/t) = t^{-D/(D+1)} \int_0^\infty s^{-1/(D+1)}\mathrm{e}^{-s} \, \mathrm{d}s$$

$$= t^{-D/(D+1)}\Gamma\left(\frac{D}{D+1}\right),$$

where we introduce a change of variable $\lambda t = s$ to get the first equality. Hence, the bound on return probabilities is proved. $\qquad \square$

Corollary 5.3 is comparable to [LOG18, Corollary 4.10]. Using a similar method, one may get an analogue of [LOG18, Corollarie 4.11]. Details are omitted.

6 Mixing Time Bound

We are concerned in this section with mixing time bounds, which are based on the bounds on vertex spectral measures in Lemma 6.3. As a preparation, we introduce the following graph parameter.

Definition 6.1. *Let G be a finite and weighted graph. We call*

$$\min\left\{ \frac{\sum_{(v,u)\in E(G)} w(v,u)|f(v)+f(u)|^2}{\max_{y\in V}|f(y)|^2} \; ; \; \min_{y\in V} f(y) < 0 < \max_{y\in V} f(y) \right\}$$

*the **energy efficiency** of G and denote it as $\mathcal{K}(G)$.*

For bipartite graphs, it is easy to see that their energy efficiency must be 0. For non-bipartite graphs, we bound their energy efficiency in the following lemma.

Lemma 6.2. *Let G be a connected and weighted graph, with weight at least 1 for each edge.*

(1) *Assume that f is a function on $V(G)$ and $\mathcal{P} = z_0 z_1 \cdots z_{|\mathcal{P}|}$ is an edge-simple path. Then*

$$\sum_{(v,u)\in E(G)} w(v,u)\big|f(v) + f(u)\big|^2 \geqslant \frac{1}{|P|}\big(f(z_0) - (-1)^{|\mathcal{P}|}f(z_{|\mathcal{P}|})\big)^2.$$

(2) *If G is also non-bipartite, finite, and simple, then $\mathscr{K}(G) \geqslant \frac{1}{\operatorname{diam}(G)+1}$.*

Proof. a) Note that

$$\sum_{(v,u)\in E(G)} w(v,u)\big|f(v) + f(u)\big|^2 \geqslant \sum_{i=0}^{|P|-1}\big|f(z_i) + f(z_{i+1})\big|^2$$

By the Cauchy–Schwarz inequality, we have

$$\sum_{i=0}^{|P|-1}\big|f(z_i) + f(z_{i+1})\big|^2 = \sum_{i=0}^{|P|-1}\big|(-1)^i f(z_i) - (-1)^{i+1} f(z_{i+1})\big|^2$$

$$\geqslant \frac{1}{|P|}\left(\sum_{i=0}^{|P|-1}\big((-1)^i f(z_i) - (-1)^{i+1} f(z_{i+1})\big)\right)^2$$

$$= \frac{1}{|P|}\big(f(z_0) - (-1)^{|\mathcal{P}|}f(z_{|\mathcal{P}|})\big)^2.$$

The first assertion is proved.

b) Now we deal with the second assertion. Take f satisfying the constraints in Definition 6.1 such that

$$\mathscr{K}(G) = \frac{\sum_{(v,u)\in E(G)} w(v,u)|f(v) + f(u)|^2}{\max_{y\in V}|f(y)|^2}.$$

Assume that $|f|$ attains its maximum at x, then

$$\mathscr{K}(G) = \frac{\sum_{(v,u)\in E(G)} w(v,u)|f(v) + f(u)|^2}{|f(x)|^2}.$$

We may assume $f(x) > 0$ without loss of generality. Set

$$S := \big\{s \in V ;\ f(s) \geqslant 0\big\}, \qquad T := \big\{t \in V ;\ f(t) < 0\big\}.$$

By the assumptions, S and T are both non-empty. Because G is non-bipartite, there exists an edge $(s_1, s_2) \in E$ with $s_1, s_2 \in S$, or an edge $(t_1, t_2) \in E$ with $t_1, t_2 \in T$.

c) If there is an edge $(s_1, s_2) \in E$ with $s_1, s_2 \in S$, let $\widehat{\mathcal{P}}_0$ be a shortest path from x to $\{s_1, s_2\}$. Without loss of generality, we may assume $\widehat{\mathcal{P}}_0$ is from x to s_1. If $|\widehat{\mathcal{P}}_0|$ is odd, we set $\widehat{\mathcal{P}} := \widehat{\mathcal{P}}_0$; if $|\widehat{\mathcal{P}}_0|$ is even, we set $\widehat{\mathcal{P}} := \widehat{\mathcal{P}}_0.(s_1, s_2)$, the concatenation of $\widehat{\mathcal{P}}_0$ and the edge (s_1, s_2). Hence, $\widehat{\mathcal{P}}$ is a path of odd length in any case. Assume

$$\widehat{\mathcal{P}} = u_0 u_1 \cdots u_k,$$

with $u_0 = x$. Then k is odd and $k \leqslant \operatorname{diam}(G) + 1$. Hence, by the first assertion,

$$\sum_{(v,u) \in E(G)} w(v,u) \big| f(v) + f(u) \big|^2 \geqslant \frac{1}{k} \big(f(u_0) + f(u_k) \big)^2 \geqslant \frac{1}{k} \big| f(x) \big|^2 \geqslant \frac{\big| f(x) \big|^2}{\operatorname{diam}(G) + 1}.$$

Therefore, in this case,

$$\mathscr{K}(G) \geqslant \frac{1}{\operatorname{diam}(G) + 1}.$$

d) If there is an edge $(t_1, t_2) \in E$ with $t_1, t_2 \in T$, let $\overline{\mathcal{P}}_0$ be a shortest path from x to $\{t_1, t_2\}$. Without loss of generality, we may assume $\overline{\mathcal{P}}_0$ is from x to t_1. If $|\overline{\mathcal{P}}_0|$ is even, we set $\overline{\mathcal{P}} := \overline{\mathcal{P}}_0$; if $|\overline{\mathcal{P}}_0|$ is odd, we set $\overline{\mathcal{P}} := \overline{\mathcal{P}}_0.(t_1, t_2)$. Hence, $\overline{\mathcal{P}}$ is a path of even length in any case. Assume

$$\overline{\mathcal{P}} = v_0 v_1 \cdots v_\ell,$$

with $v_0 = x$. Then ℓ is even and $\ell \leqslant \operatorname{diam}(G) + 1$. Hence, by the first assertion,

$$\sum_{(v,u) \in E(G)} w(v,u) \big| f(v) + f(u) \big|^2 \geqslant \frac{1}{\ell} \big(f(v_0) - f(v_\ell) \big)^2 > \frac{1}{\ell} \big| f(x) \big|^2 \geqslant \frac{\big| f(x) \big|^2}{\operatorname{diam}(G) + 1}.$$

Therefore, in this case,

$$\mathscr{K}(G) \geqslant \frac{1}{\operatorname{diam}(G) + 1}. \qquad \square$$

Now, we present our bounds on vertex spectral measures.

Lemma 6.3. *Let G be a finite, simple, connected, weighted graph, with weight at least 1 for each edge.*

(1) *For $\delta \in [0, 2)$ and $x \in V$, $\mu_x^*(\delta) + \pi(x) \leqslant \mathcal{R}_{\operatorname{diam}}(G) w(x) \delta$.*

(2) *If G is also non-bipartite, for $\delta \in [0, 2)$ and $x \in V$, $\mu_x^{\mathcal{Q}}(\delta) \leqslant \frac{w(x)\delta}{\mathscr{K}(G)}$. Moreover, by Lemma 6.2, we have $\mu_x^{\mathcal{Q}}(\delta) \leqslant \big(\operatorname{diam}(G) + 1 \big) w(x) \delta$ for $\delta \in [0, 2)$.*

Proof. The first assertion follows immediately from [LOG18, Proposition 4.2]. We now deal with the second assertion. Fix a vertex $x \in V$ and define f as in Lemma 2.4.

When $\delta \in [0, \lambda_{\min}^{\mathcal{Q}})$, $\mu_x^{\mathcal{Q}}(\delta) = 0$ by definition. So the inequality holds automatically.

When $\lambda_{\min}^{\mathcal{Q}} \leqslant \delta < \lambda_{\max}^{\mathcal{Q}} = 2$, we know that f is orthogonal to the eigenspace of $\mathcal{Q}$ corresponding to $\lambda_{\max}^{\mathcal{Q}} = 2$, which is spanned by $(1, 1, \ldots, 1)$, the constant vector. Therefore, by Definition 6.1, we have

$$\delta \geqslant \langle f, \mathcal{Q}f \rangle_w = \sum_{(v,u) \in E(G)} w(v,u) |f(v) + f(u)|^2 \geqslant f(x)^2 \mathscr{K}(G) = \frac{\mu_x^{\mathcal{Q}}(\delta) \mathscr{K}(G)}{w(x)}. \qquad \square$$

By Lemma 6.2(2), Lemma 6.3(2), and Lemma 3.2, we have the following corollary.

Corollary 6.4. *For any non-bipartite, finite, simple, connected, weighted graph G with weight at least 1 for each edge, we have*

$$\lambda_k^{\mathcal{Q}} \geqslant \frac{k\mathcal{K}(G)}{\mathsf{wt}(V)} \geqslant \frac{k}{\big(\mathrm{diam}(G)+1\big)\mathsf{wt}(V)}$$

for $1 \leqslant k \leqslant n$. Therefore,

$$\lambda_k^P \geqslant -1 + \frac{k\mathcal{K}(G)}{\mathsf{wt}(V)} \geqslant -1 + \frac{k}{\big(\mathrm{diam}(G)+1\big)\mathsf{wt}(V)}. \qquad \square$$

As mentioned in Section 1.1, Corollary 6.4, combined with [LOG18, Proposition 4.2], improves the result in [Landau and Odlyzko, 1981].

To get a mixing time bound, we first give a bound on return probabilities using Lemma 6.3.

Corollary 6.5. *Let G be a non-bipartite, finite, simple, connected, unweighted graph, and $t \equiv 0 \bmod 2$. Then we have*

$$0 \leqslant \frac{p_t(x,x) - \pi(x)}{\pi(x)} \leqslant \frac{2\big(\mathrm{diam}(G)+1\big)\mathsf{wt}(V)}{t}.$$

Proof. Set

$$\widetilde{\varphi}(\lambda) := \begin{cases} \mathcal{R}_{\mathrm{diam}}(G)w(x)\lambda & \text{if } \lambda \in [0,1) \text{ and } \mathcal{R}_{\mathrm{diam}}(G)w(x)\lambda \leqslant \mu_x^*(1), \\ \mu_x^*(1) & \text{for intermediate values of } \lambda, \\ \mu_x^*(2) - \big(\mathrm{diam}(G)+1\big)w(x)(2-\lambda) & \begin{aligned}&\text{if } \lambda \in [1,2] \text{ and} \\ &\mu_x^*(2) - \big(\mathrm{diam}(G)+1\big)w(x)(2-\lambda) \geqslant \mu_x^*(1).\end{aligned} \end{cases}$$

By Lemma 6.3, using the argument as in part a) of the proof of Theorem 3.4, the function $\widetilde{\varphi}$ defined above satisfies the conditions in Lemma 2.7(1). Therefore, for $t \equiv 0 \bmod 2$,

$$0 \leqslant p_t(x,x) - \pi(x) \leqslant \int_0^2 (1-\lambda)^t \widetilde{\varphi}'(\lambda)\,\mathrm{d}\lambda$$

$$\leqslant \mathcal{R}_{\mathrm{diam}}(G)w(x)\int_0^1 (1-\lambda)^t\,\mathrm{d}\lambda + \big(\mathrm{diam}(G)+1\big)w(x)\int_0^1 (1-\lambda)^t\,\mathrm{d}\lambda$$

$$\leqslant 2\big(\mathrm{diam}(G)+1\big)w(x)\int_0^1 (1-\lambda)^t\,\mathrm{d}\lambda.$$

Hence,

$$\frac{p_t(x,x) - \pi(x)}{\pi(x)} \leqslant 2\big(\mathrm{diam}(G)+1\big)\frac{w(x)}{\pi(x)}\int_0^1 (1-\lambda)^t\,\mathrm{d}\lambda.$$

But we know $\frac{w(x)}{\pi(x)} = \mathsf{wt}(V)$ and

$$\int_0^1 (1-\lambda)^t\,\mathrm{d}\lambda \leqslant \int_0^1 \exp\{-\lambda t\}\,\mathrm{d}\lambda \leqslant \int_0^\infty \exp\{-\lambda t\}\,\mathrm{d}\lambda = \frac{1}{t}.$$

Therefore,

$$\frac{p_t(x,x) - \pi(x)}{\pi(x)} \leqslant \frac{2\big(\mathrm{diam}(G)+1\big)\mathsf{wt}(V)}{t}. \qquad \square$$

We are now almost in position to give the following mixing time bound.

Corollary 6.6. *For a non-bipartite, finite, simple, connected, unweighted graph G, the uniform mixing time of the simple random walk on G satisfies*

$$\tau_\infty(1/4) \leqslant 8n^3.$$

If we further assume that G is regular, then $\tau_\infty(1/4) \leqslant 24n^2$.

In the proof of Corollary 6.6, the following lemma will be useful.

Lemma 6.7. *For a reversible Markov chain with stationary distribution π, we have*

$$\left| \frac{p_{2t}(x, y) - \pi(y)}{\pi(y)} \right| \leqslant \sqrt{\frac{p_{2t}(x, x) - \pi(x)}{\pi(x)}} \sqrt{\frac{p_{2t}(y, y) - \pi(y)}{\pi(y)}}.$$

In addition, $\max_{x,y} \left| \frac{p_t(x,y) - \pi(y)}{\pi(y)} \right|$ is decreasing in t.

Proof. One may refer to the proof of [LOG18, Proposition A.1]. The monotonicity is mentioned in [Lyons and Oveis Gharan, 2021]. $\square$

Proof of Corollary 6.6. a) Combining Lemma 6.7 and Corollary 6.5, for $t \equiv 0 \bmod 2$, we have

$$\max_{x,y} \left| \frac{p_t(x, y) - \pi(y)}{\pi(y)} \right| \leqslant \max_x \frac{p_t(x, x) - \pi(x)}{\pi(x)} \leqslant \frac{2\big(\operatorname{diam}(G) + 1\big)\operatorname{wt}(V)}{t}.$$

In addition, we notice that $\operatorname{wt}(V) \leqslant n(n-1)$. Hence, the second assertion of Lemma 6.7 ensures

$$\tau_\infty(1/4) \leqslant \left\lceil 8\big(\operatorname{diam}(G) + 1\big)\operatorname{wt}(V) \right\rceil \leqslant \left\lceil 8n \cdot n(n-1) \right\rceil \leqslant 8n^3.$$

b) To prove the second assertion, we assume G is d-regular. By Lemma 2.5, we have

$$\operatorname{diam}(G) + 1 \leqslant \frac{3n}{d+1}.$$

In addition, $\operatorname{wt}(V) = nd$. By the second assertion of Lemma 6.7,

$$\tau_\infty(1/4) \leqslant \left\lceil 8\big(\operatorname{diam}(G) + 1\big)\operatorname{wt}(V) \right\rceil \leqslant \left\lceil 8 \cdot \frac{3n}{d+1} \cdot nd \right\rceil \leqslant 24n^2. \qquad \square$$

7 Transitive Case

We consider vertex-transitive graphs in this section.

7.1 Estimate of Spectral Measure

Given a vertex-transitive graph G, set

$$N^{\#}(r) := \{v \in V(G); \ \mathsf{dist}(v, x) \leqslant r\}.$$

This value does not depend on the choice of x since G is vertex-transitive. Furthermore, as in [LOG18, Lemma 6.4], for every two vertices $x, y \in V$, it is easy to see that $\|F_x^{\mathcal{Q}}\|_w = \|F_y^{\mathcal{Q}}\|_w$.

Lemma 7.1. *When G is vertex-transitive, for each $x \in V$,*

$$\mathcal{Q}F_x^{\mathcal{Q}} = \sum_{v \in N(x)} \frac{w(x, v)}{w}(F_x^{\mathcal{Q}} + F_v^{\mathcal{Q}}).$$

Proof. By the definition of $F_x^{\mathcal{Q}}$,

$$\mathcal{Q}F_x^{\mathcal{Q}} = \frac{\mathcal{Q}I_{\mathcal{Q}}(\delta)\mathbf{1}_x}{w} = \frac{I_{\mathcal{Q}}(\delta)\mathcal{Q}\mathbf{1}_x}{w}.$$

However, by direct calculation,

$$\mathcal{Q}\mathbf{1}_x = \sum_{v \in N(x)} \frac{w(x, v)}{w}(\mathbf{1}_x + \mathbf{1}_v).$$

Thus, we have

$$\begin{aligned}
\mathcal{Q}F_x^{\mathcal{Q}} &= \frac{1}{w}I_{\mathcal{Q}}(\delta) \sum_{v \in N(x)} \frac{w(x, v)}{w}(\mathbf{1}_x + \mathbf{1}_v) \\
&= \sum_{v \in N(x)} \frac{w(x, v)}{w}\left(\frac{I_{\mathcal{Q}}(\delta)\mathbf{1}_x}{w} + \frac{I_{\mathcal{Q}}(\delta)\mathbf{1}_v}{w}\right) \\
&= \sum_{v \in N(x)} \frac{w(x, v)}{w}(F_x^{\mathcal{Q}} + F_v^{\mathcal{Q}}). \qquad \square
\end{aligned}$$

Theorem 7.2. *Let G be a vertex-transitive, non-bipartite, simple, connected, weighted graph with weight at least 1 for each edge. For each $x \in V(G)$, $c \in (0, 1)$, and $\delta \in \left(0, \frac{2}{w}\right]$, we have*

$$\mu_x^{\mathcal{Q}}(\delta) \leqslant \frac{1}{c^2 N^{\#}\left(\frac{\arcsin\sqrt{(1-c)/2}}{\arcsin\sqrt{w\delta/2}}\right)}.$$

Furthermore, for $\delta \in (0, 2]$,

$$\mu_x^{\mathcal{Q}}(\delta) \leqslant \frac{1}{c^2 N^{\#}\left(\frac{\sqrt{1-c}}{\sqrt{w\delta}}\right)}.$$

Proof. Fix a vertex $x \in V$.

a) Consider the spectral embedding $\{F_v^{\mathcal{Q}} ; \; v \in V\}$ based on $\mathcal{Q}$. Let

$$\rho := \|F_x^{\mathcal{Q}}\|_w, \qquad \beta_0 := \max_{v \in N(x)} \|F_x^{\mathcal{Q}} + F_v^{\mathcal{Q}}\|_w.$$

We have

$$\delta \geqslant \frac{1}{\|F_x^{\mathcal{Q}}\|_w^2} \langle \mathcal{Q} F_x^{\mathcal{Q}}, F_x^{\mathcal{Q}} \rangle_w = \frac{1}{\|F_x^{\mathcal{Q}}\|_w^2} \left\langle \sum_{v \in N(x)} \frac{w(x,v)}{w} (F_x^{\mathcal{Q}} + F_v^{\mathcal{Q}}), F_x^{\mathcal{Q}} \right\rangle_w$$

$$= \frac{1}{\|F_x^{\mathcal{Q}}\|_w^2} \sum_{v \in N(x)} \frac{w(x,v)}{w} \left(\|F_x^{\mathcal{Q}}\|_w^2 + \langle F_x^{\mathcal{Q}}, F_v^{\mathcal{Q}} \rangle_w \right)$$

$$= \frac{1}{\|F_x^{\mathcal{Q}}\|_w^2} \sum_{v \in N(x)} \frac{w(x,v)}{w} \frac{\|F_x^{\mathcal{Q}} + F_v^{\mathcal{Q}}\|_w^2}{2}$$

$$\geqslant \frac{1}{2w\rho^2} \max_{v \in N(x)} \|F_x^{\mathcal{Q}} + F_v^{\mathcal{Q}}\|_w^2 = \frac{1}{2w\rho^2} \beta_0^2.$$

In other words, $\frac{\beta_0}{2\rho} \leqslant \sqrt{w\delta/2}$. Since we assumed that $\delta \in (0, \frac{2}{w}]$, by the monotonicity of arcsin function on the interval $[0, 1]$, we have

$$2 \arcsin \frac{\beta_0}{2\rho} \leqslant 2 \arcsin \sqrt{w\delta/2}.$$

b) Set $B_0 := B\left(x, \frac{\arcsin \sqrt{(1-c)/2}}{\arcsin \sqrt{(w\delta)/2}}\right)$. For $v \in B_0$, let $\mathcal{P} = u_0 u_1 \cdots u_{r-1} u_r$ be a shortest path joining x and v, with $u_0 = x$ and $u_r = v$. Define

$$\widetilde{F}_i := (-1)^i F_{u_i}^{\mathcal{Q}}, \qquad i = 0, 1, 2, \ldots, r.$$

Then $r \leqslant \frac{\arcsin \sqrt{(1-c)/2}}{\arcsin \sqrt{(w\delta)/2}}$, and

$$\left\| \widetilde{F}_i - \widetilde{F}_{i+1} \right\|_w = \left\| F_{u_i}^{\mathcal{Q}} + F_{u_{i+1}}^{\mathcal{Q}} \right\|_w, \qquad i = 0, 1, 2, \ldots, r-1.$$

Write $\theta(h_1, h_2)$ for the angle between $h_1, h_2 \in \ell^2(V, w)$. Then using the sphere metric, we get

$$\theta\left(\widetilde{F}_0, \widetilde{F}_r\right) \leqslant \sum_{i=0}^{r-1} \theta\left(\widetilde{F}_i, \widetilde{F}_{i+1}\right)$$

$$= \sum_{i=0}^{r-1} 2 \arcsin \frac{\left\| \widetilde{F}_i - \widetilde{F}_{i+1} \right\|_w}{2\rho} = \sum_{i=0}^{r-1} 2 \arcsin \frac{\left\| F_{u_i}^{\mathcal{Q}} + F_{u_{i+1}}^{\mathcal{Q}} \right\|_w}{2\rho}$$

$$\leqslant \sum_{i=0}^{r-1} 2 \arcsin \frac{\beta_0}{2\rho} \leqslant 2r \arcsin \sqrt{w\delta/2}$$

$$\leqslant 2 \frac{\arcsin \sqrt{(1-c)/2}}{\arcsin \sqrt{w\delta/2}} \arcsin \sqrt{w\delta/2} = 2 \arcsin \sqrt{(1-c)/2}$$

$$= \arccos c.$$

(8)

Thus,
$$\left|\cos\theta\left(F_x^{\mathcal{Q}}, F_v^{\mathcal{Q}}\right)\right| = \cos\theta\left(\widetilde{F}_0, \widetilde{F}_r\right) \geqslant c.$$

In summary, if $v \in B_0 = B\left(x, \frac{\arcsin\sqrt{(1-c)/2}}{\arcsin\sqrt{(w\delta)/2}}\right)$,
$$\left|\cos\theta\left(F_x^{\mathcal{Q}}, F_v^{\mathcal{Q}}\right)\right| \geqslant c.$$

c) By the above discussion,
$$
\begin{aligned}
\rho^2 = \|F_x^{\mathcal{Q}}\|_w^2 &= \sum_{v \in V} w\left|F_x^{\mathcal{Q}}(v)\right|^2 \\
&= \sum_{v \in V} w\left|\langle F_x^{\mathcal{Q}}, F_v^{\mathcal{Q}}\rangle_w\right|^2 \geqslant w\sum_{v \in B_0}\left|\langle F_x^{\mathcal{Q}}, F_v^{\mathcal{Q}}\rangle_w\right|^2 \\
&= w\sum_{v \in B_0}\left|\cos\theta\left(F_x^{\mathcal{Q}}, F_v^{\mathcal{Q}}\right)\right|^2\|F_x^{\mathcal{Q}}\|_w^2\|F_v^{\mathcal{Q}}\|_w^2 \\
&\geqslant w\rho^4 c^2 N^\#\left(\frac{\arcsin\sqrt{(1-c)/2}}{\arcsin\sqrt{w\delta/2}}\right).
\end{aligned}
$$

Thus, we have
$$\rho^2 \leqslant \frac{1}{wc^2 N^\#\left(\frac{\arcsin\sqrt{(1-c)/2}}{\arcsin\sqrt{w\delta/2}}\right)}.$$

Now, by Lemma 2.3,
$$\mu_x^{\mathcal{Q}}(\delta) = w\|F_x^{\mathcal{Q}}\|_w^2 = w\rho^2 \leqslant \frac{1}{c^2 N^\#\left(\frac{\arcsin\sqrt{(1-c)/2}}{\arcsin\sqrt{w\delta/2}}\right)}.$$

The first assertion is proved.

d) Now we turn to the second assertion.

If $(1-c) < w\delta$, then
$$N^\#\left(\frac{\sqrt{1-c}}{\sqrt{w\delta}}\right) = 1.$$

Therefore,
$$\frac{1}{c^2 N^\#\left(\frac{\sqrt{1-c}}{\sqrt{w\delta}}\right)} > 1.$$

But $\mu_x^{\mathcal{Q}}(\delta) \leqslant 1$. So the second assertion holds in this case.

If $(1-c) \geqslant w\delta$, then $\sqrt{w\delta/2} < 1$. Thus, $\arcsin\sqrt{w\delta/2}$ is defined. Note that $\frac{x}{\arcsin x}$ is decreasing for $x \in (0,1)$. Therefore,
$$\frac{\sqrt{(1-c)/2}}{\arcsin\sqrt{(1-c)/2}} \leqslant \frac{\sqrt{w\delta/2}}{\arcsin\sqrt{w\delta/2}}.$$

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Consequently,

$$\frac{\sqrt{1-c}}{\sqrt{w\delta}} \leqslant \frac{\arcsin\sqrt{(1-c)/2}}{\arcsin\sqrt{w\delta/2}}.$$

Hence, the second assertion follows from the first one immediately. □

By the bounds on vertex spectral measures in Theorem 7.2 and [LOG18, Theorem 6.1], using the method as in Theorem 3.4, we may get a similar result as [LOG18, Corollary 6.6].

Corollary 7.3. *Let G be a finite or infinite, vertex-transitive, d-regular, non-bipartite, simple, connected, unweighted graph with at least polynomial growth rate $N^{\#}(r) \geqslant Cr^D$, where $C > 0$ and $D \geqslant 1$ are constants and $0 \leqslant r \leqslant \mathrm{diam}(G)$. Then for each $x \in V$ and $t > 0$,*

$$0 \leqslant p_t(x,x) - \pi(x) \leqslant 2\widetilde{C}t^{-D/2} \qquad for \quad t \equiv 0 \bmod 2,$$

$$\left|p_t(x,x) - \pi(x)\right| \leqslant \widetilde{C}t^{-D/2} \qquad for \quad t \equiv 1 \bmod 2,$$

where $\widetilde{C} = \frac{(D+4)^{D/2+2}d^{D/2}}{32CD^{D/2-1}}\Gamma\left(\frac{D}{2}\right)$. □

The proof of this corollary is omitted. Note that the rate $t^{-D/2}$ here is the correct decay rate for the simple random walk on $\mathbb{Z}^D$, $D \in \mathbb{N}$. An analogue of [LOG18, Corollary 6.7] could also be written down easily.

7.2 Minimum Eigenvalue

For a connected, finite graph G and its signless probabilistic Laplacian $\mathcal{Q} = I + P$, we know that 0 is an eigenvalue of $\mathcal{Q}$ if and only if G is bipartite from Corollary 2.2.

Theorem 7.4. *Let G be a vertex-transitive, non-bipartite, finite, simple, connected, weighted graph with weight at least 1 for each edge. Then*

$$\lambda_{\min}^{\mathcal{Q}} \geqslant \frac{2}{w}\left(\sin\frac{\pi}{4\left(\mathrm{diam}(G)+1\right)}\right)^2.$$

Proof. We may assume $\frac{\lambda_{\min}^{\mathcal{Q}}}{2} < 1$: otherwise, the inequality is trivial. Consider the spectral embedding based on $\mathcal{Q}$ with $\delta = \lambda_{\min}^{\mathcal{Q}}$. Fix $x \in V$ and let

$$S := \left\{s \in V ;\ f(s) \geqslant 0\right\}, \qquad T := \left\{t \in V ;\ f(t) < 0\right\}.$$

It is easy to see that both S and T are non-empty. Because G is assumed to be non-bipartite, there exists an edge $(s_1, s_2) \in E$ with $s_1, s_2 \in S$, or an edge $(t_1, t_2) \in E$ with $t_1, t_2 \in T$.

a) If there is an edge $(s_1, s_2) \in E$ with $s_1, s_2 \in S$, let $\widehat{\mathcal{P}}_0$ be a shortest path from x to $\{s_1, s_2\}$. Without loss of generality, we may assume $\widehat{\mathcal{P}}_0$ is from x to s_1. If $|\widehat{\mathcal{P}}_0|$ is odd, we set $\widehat{\mathcal{P}} := \widehat{\mathcal{P}}_0$; if $|\widehat{\mathcal{P}}_0|$ is even, we set $\widehat{\mathcal{P}} := \widehat{\mathcal{P}}_0.(s_1, s_2)$, the concatenation of $\widehat{\mathcal{P}}_0$ and the edge (s_1, s_2). Hence, $\widehat{\mathcal{P}}$ is a path of odd length in any case. Assume

$$\widehat{\mathcal{P}} = u_0 u_1 \cdots u_k,$$

with $u_0 = x$. We define

$$\widehat{F}_i := (-1)^i F_{u_i}^{\mathcal{Q}}, \qquad i = 0, 1, 2, \dots, k.$$

Then $\widehat{F}_0 = F_x^{\mathcal{Q}}$ and

$$\begin{aligned}\langle \widehat{F}_0, \widehat{F}_k \rangle_w &= (-1)^k \langle F_x^{\mathcal{Q}}, F_{u_k}^{\mathcal{Q}} \rangle_w = -\langle F_x^{\mathcal{Q}}, F_{u_k}^{\mathcal{Q}} \rangle_w = -\langle I_{\mathcal{Q}}(\delta)\mathbf{1}_x/w, I_{\mathcal{Q}}(\delta)\mathbf{1}_{u_k}/w \rangle_w \\ &= -\langle I_{\mathcal{Q}}(\delta)\mathbf{1}_x/w, \mathbf{1}_{u_k}/w \rangle_w = -\langle I_{\mathcal{Q}}(\delta)\mathbf{1}_x/w, \mathbf{1}_{u_k} \rangle = -\langle F_x^{\mathcal{Q}}, \mathbf{1}_{u_k} \rangle \\ &= -F_x^{\mathcal{Q}}(u_k) \leqslant 0.\end{aligned}$$

Hence,

$$\frac{\pi}{2} \leqslant \theta(\widehat{F}_0, \widehat{F}_k). \tag{9}$$

In addition, using the sphere metric as in Eq. (8), we have

$$\theta(\widehat{F}_0, \widehat{F}_k) \leqslant 2k \arcsin \sqrt{w \lambda_{\min}^{\mathcal{Q}}/2} \leqslant 2(\operatorname{diam}(G) + 1) \arcsin \sqrt{w \lambda_{\min}^{\mathcal{Q}}/2}. \tag{10}$$

Combining Eqs. (9) and (10),

$$\frac{\pi}{2} \leqslant 2(\operatorname{diam}(G) + 1) \arcsin \sqrt{w \lambda_{\min}^{\mathcal{Q}}/2}.$$

So in this case,

$$\lambda_{\min}^{\mathcal{Q}} \geqslant \frac{2}{w} \left(\sin \frac{\pi}{4(\operatorname{diam}(G) + 1)} \right)^2.$$

b) If there is an edge $(t_1, t_2) \in E$ with $t_1, t_2 \in T$, let $\overline{\mathcal{P}}_0$ be a shortest path from x to $\{t_1, t_2\}$. Without loss of generality, we may assume $\overline{\mathcal{P}}_0$ is from x to t_1. If $|\overline{\mathcal{P}}_0|$ is even, we set $\overline{\mathcal{P}} := \overline{\mathcal{P}}_0$; if $|\overline{\mathcal{P}}_0|$ is odd, we set $\overline{\mathcal{P}} := \overline{\mathcal{P}}_0.(t_1, t_2)$. Thus, $\overline{\mathcal{P}}$ is a path of even length in any case. Assume

$$\overline{\mathcal{P}} = v_0 v_1 \cdots v_\ell,$$

with $v_0 = x$, and define

$$\overline{F}_i := (-1)^i F_{v_i}^{\mathcal{Q}}, \qquad i = 0, 1, 2, \dots, \ell.$$

Then $\overline{F}_0 = F_x^{\mathcal{Q}}$ and

$$\langle \overline{F}_0, \overline{F}_\ell \rangle_w = (-1)^\ell \langle F_x^{\mathcal{Q}}, F_{v_\ell}^{\mathcal{Q}} \rangle_w = F_x^{\mathcal{Q}}(v_\ell) < 0.$$

Hence, we have

$$\frac{\pi}{2} < \theta(\overline{F}_0, \overline{F}_\ell). \tag{11}$$

In addition, using the sphere metric as in Eq. (8) again, we have

$$\theta(\overline{F}_0, \overline{F}_r) \leqslant 2\ell \arcsin \sqrt{w \lambda_{\min}^{\mathcal{Q}}/2} \leqslant 2(\operatorname{diam}(G) + 1) \arcsin \sqrt{w \lambda_{\min}^{\mathcal{Q}}/2}. \tag{12}$$

Combining Eqs. (11) and (12), we get

$$\frac{\pi}{2} < 2\big(\mathrm{diam}(G) + 1\big) \arcsin \sqrt{w\lambda^{\mathcal{Q}}_{\min}/2}.$$

So in this case,

$$\lambda^{\mathcal{Q}}_{\min} > \frac{2}{w}\left(\sin \frac{\pi}{4\big(\mathrm{diam}(G) + 1\big)}\right)^2. \qquad \square$$

8 Average Return Probability

In this section, we deal with the average spectral measure of $\mathcal{Q}$ and average return probabilities for simple random walk on generic non-bipartite, simple, finite, connected graphs. Our method is inspired by [LOG18, Section 5] and it requires more dexterity here.

8.1 Estimate of Average Spectral Measure

Define average spectral measure of $\mathcal{Q}$ as

$$\mu^{\mathcal{Q}}(\delta) := \frac{1}{n} \sum_{x \in V} \mu^{\mathcal{Q}}_x(\delta).$$

We also write $\mu^{\mathcal{Q}}_S(\delta) := \sum_{x \in S} \mu^{\mathcal{Q}}_x(\delta)$ for $S \subseteq V$.

Theorem 8.1. *For any non-bipartite, finite, simple, connected, unweighted graph G and $\delta \in (0, 2)$, we have*

$$\mu^{\mathcal{Q}}(\delta) < (4000\delta)^{1/3}.$$

Corollary 8.2. *For any non-bipartite, finite, simple, connected, unweighted graph G, we have $\lambda^{\mathcal{Q}}_k \geqslant \frac{k^3}{4000n^3}$ for $1 \leqslant k \leqslant n$. Therefore, $\lambda^{P}_k \geqslant -1 + \frac{k^3}{4000n^3}$.*

Proof. By the same argument as in the proof of Corollary 3.3, this corollary follows easily from Lemma 3.2 and Theorem 8.1. $\qquad \square$

To prove Theorem 8.1, we need some preparation.
Let $m := \left\lfloor \frac{n}{2}\mu^{\mathcal{Q}}(\delta) \right\rfloor + 1$. For each $x \in V$, set

$$\mathscr{R}(x) := \left\{y \in V\,;\, \|F^{\mathcal{Q}}_x - F^{\mathcal{Q}}_y\|_w \leqslant \frac{\|F^{\mathcal{Q}}_x\|_w}{4} \text{ or } \|F^{\mathcal{Q}}_x + F^{\mathcal{Q}}_y\|_w \leqslant \frac{\|F^{\mathcal{Q}}_x\|_w}{4}\right\}.$$

We will use Algorithm 1 to get some useful sets.
For each $x \in V(G)$, set

$$\widetilde{N}(x) := \left\{y \in N(x)\,;\, \|F^{\mathcal{Q}}_y + F^{\mathcal{Q}}_x\|_w \leqslant \frac{\|F^{\mathcal{Q}}_x\|_w}{9}\right\}.$$

Let $\widetilde{T}(x)$ be the star formed by x and $\widetilde{N}(x)$. We pick $T(x)$ as a maximal (with respect to inclusion) connected bipartite graph including $\widetilde{T}(x)$ such that

$$\forall y \in V\big(T(x)\big) \quad \left\|F^{\mathcal{Q}}_x - (-1)^{\mathsf{dist}(x,y;T(x))} F^{\mathcal{Q}}_y\right\|_w \leqslant \frac{\|F^{\mathcal{Q}}_x\|_w}{9}.$$

Algorithm 1 Set-Selection(α)

Let $S_0 \leftarrow V$.
for $i = 1 \rightarrow m$ **do**
   Choose a vertex x_i in S_{i-1} that maximizes $\mu_{x_i}^{\mathcal{Q}}(\delta)$.
   Let $S_i \leftarrow S_{i-1} \setminus \mathscr{R}(x_i)$.
end for
return $\mathscr{R}(x_1), \mathscr{R}(x_2), \ldots, \mathscr{R}(x_m)$.

Lemma 8.3. *Let G be a non-bipartite, finite, simple, connected, unweighted graph and $\delta \geqslant \lambda_{\min}^{\mathcal{Q}}$.*

(0) *For $i = 1, 2, \ldots, m$, we have $\mu_{\mathscr{R}(x_i)}^{\mathcal{Q}}(\delta) \leqslant \frac{4}{3}$.*

(1) *For $i = 1, 2, \ldots, m$, we have $\mu_{x_i}^{\mathcal{Q}}(\delta) \geqslant \frac{\mu^{\mathcal{Q}}(\delta)}{3}$. In addition, Algorithm 1 is well designed: each S_{i-1} is non-empty and x_i could be chosen for $i = 1, 2, \ldots, m$. Therefore, Algorithm 1 is not stopping before $i = m$.*

(2) *For $1 \leqslant i < j \leqslant m$, $V(T(x_i)) \cap V(T(x_j)) = \varnothing$.*

Proof. (0) Since $\delta \geqslant \lambda_{\min}^{\mathcal{Q}}$, we must have $\|F_{x_i}^{\mathcal{Q}}\|_w > 0$. By the proof of [LOG18, Lemma 3.13], for any two non-zero vectors in a Hilbert space $\mathcal{H}$,

$$\Re\left\langle \frac{f}{\|f\|_{\mathcal{H}}}, \frac{g}{\|g\|_{\mathcal{H}}} \right\rangle_{\mathcal{H}} \geqslant 1 - \frac{2\|f - g\|_{\mathcal{H}}^2}{\|f\|_{\mathcal{H}}^2}, \tag{13}$$

where $\Re z$ is the real part of a complex number z. Also, for $y \in \mathscr{R}(x_i)$, by the construction of $\mathscr{R}(x_i)$, we have an integer $\sigma_i(y)$ such that

$$\left\| F_{x_i}^{\mathcal{Q}} - (-1)^{\sigma_i(y)} F_y^{\mathcal{Q}} \right\|_w \leqslant \frac{\|F_{x_i}^{\mathcal{Q}}\|_w}{4}.$$

This implies $\|F_y^{\mathcal{Q}}\|_w > 0$ for $y \in \mathscr{R}(x_i)$.

Therefore,

$$
\begin{aligned}
1 &= \sum_{y \in \mathscr{R}(x_i)} w(y) \|F_y^{\mathcal{Q}}\|_w^2 \left| \left\langle \frac{F_y^{\mathcal{Q}}}{\|F_y^{\mathcal{Q}}\|_w}, \frac{F_{x_i}^{\mathcal{Q}}}{\|F_{x_i}^{\mathcal{Q}}\|_w} \right\rangle_w \right|^2 \\
&= \sum_{y \in \mathscr{R}(x_i)} \mu_y^{\mathcal{Q}}(\delta) \left| \left\langle \frac{(-1)^{\sigma_i(y)} F_y^{\mathcal{Q}}}{\|(-1)^{\sigma_i(y)} F_y^{\mathcal{Q}}\|_w}, \frac{F_{x_i}^{\mathcal{Q}}}{\|F_{x_i}^{\mathcal{Q}}\|_w} \right\rangle_w \right|^2 \\
&\geqslant \sum_{y \in \mathscr{R}(x_i)} \mu_y^{\mathcal{Q}}(\delta) \left(1 - \frac{2\|F_{x_i}^{\mathcal{Q}} - (-1)^{\sigma_i(y)} F_y^{\mathcal{Q}}\|_w^2}{\|F_{x_i}^{\mathcal{Q}}\|_w^2} \right)^2 \\
&\geqslant \mu_{\mathscr{R}(x_i)}^{\mathcal{Q}}(\delta) \left(1 - \frac{2}{16} \right)^2 = \frac{49}{64} \mu_{\mathscr{R}(x_i)}^{\mathcal{Q}}(\delta),
\end{aligned}
$$

where we use Eq. (13) to get the first inequality.

It follows that

$$\mu^{\mathcal{Q}}_{\mathscr{R}(x_i)}(\delta) \leqslant \frac{64}{49} < \frac{4}{3}.$$

(1) We know $\mu^{\mathcal{Q}}_{S_0}(\delta) = \mu^{\mathcal{Q}}_V(\delta) = n\mu^{\mathcal{Q}}(\delta)$. In addition, by assertion (0), the total spectral measure of removed vertices in each iteration of the for loop in Algorithm 1 is at most $\frac{4}{3}$. We have, for each $i \leqslant m$,

$$\mu^{\mathcal{Q}}_{S_{i-1}}(\delta) \geqslant n\mu^{\mathcal{Q}}(\delta) - \frac{4}{3}(m-1) \geqslant \frac{n\mu^{\mathcal{Q}}(\delta)}{3},$$

where the last inequality holds by the definition of m. This implies $S_{i-1} \neq \varnothing$ and x_i could be chosen in Algorithm 1 for $i = 1, 2, \ldots, m$. In other words, the algorithm is well designed and is not stopping before $i = m$.

Furthermore, since x_i has the largest vertex spectral measure in S_{i-1} for $i = 1, 2, \ldots, m$, we have

$$\mu^{\mathcal{Q}}_{x_i}(\delta) \geqslant \frac{\mu^{\mathcal{Q}}_{S_{i-1}}(\delta)}{n} \geqslant \frac{\mu^{\mathcal{Q}}_{S_{m-1}}(\delta)}{n} \geqslant \frac{\mu^{\mathcal{Q}}(\delta)}{3}.$$

(2) Suppose that some vertex y lies in $V\big(T(x_i)\big) \cap V\big(T(x_j)\big)$. We deal with assertion (2) in the following two cases. For ease of notation, we write $\tau_k := \mathsf{dist}(x_k, y; T(x_k))$ for $k = i, j$. By the construction of $T(x_i)$ and $T(x_j)$,

$$\left\| F^{\mathcal{Q}}_{x_k} - (-1)^{\tau_k} F^{\mathcal{Q}}_y \right\|_w \leqslant \frac{\| F^{\mathcal{Q}}_{x_k} \|_w}{9}, \qquad k = i, j.$$

<u>Case (a)</u>: $\| F^{\mathcal{Q}}_{x_j} \|_w > \frac{5}{4} \| F^{\mathcal{Q}}_{x_i} \|_w$. In this case, we have

$$\begin{aligned}
\left\| F^{\mathcal{Q}}_y \right\|_w &= \left\| (-1)^{\tau_j} F^{\mathcal{Q}}_y \right\|_w \\
&\geqslant \frac{8}{9} \| F^{\mathcal{Q}}_{x_j} \|_w > \frac{8}{9} \cdot \frac{5}{4} \| F^{\mathcal{Q}}_{x_i} \|_w \\
&= \frac{10}{9} \| F^{\mathcal{Q}}_{x_i} \|_w.
\end{aligned}$$

Thus, we have

$$\left\| (-1)^{\tau_i} F^{\mathcal{Q}}_y \right\|_w = \left\| F^{\mathcal{Q}}_y \right\|_w > \frac{10}{9} \| F^{\mathcal{Q}}_{x_i} \|_w.$$

This contradicts our assumption that $y \in V\big(T(x_i)\big)$.

<u>Case (b)</u>: $\| F^{\mathcal{Q}}_{x_j} \|_w \leqslant \frac{5}{4} \| F^{\mathcal{Q}}_{x_i} \|_w$. In this case, we have

$$\begin{aligned}
\left\| F^{\mathcal{Q}}_{x_i} - (-1)^{\tau_i + \tau_j} F^{\mathcal{Q}}_{x_j} \right\|_w &= \left\| F^{\mathcal{Q}}_{x_i} - (-1)^{\tau_i} F^{\mathcal{Q}}_y + (-1)^{\tau_i} F^{\mathcal{Q}}_y - (-1)^{\tau_i + \tau_j} F^{\mathcal{Q}}_{x_j} \right\|_w \\
&\leqslant \left\| F^{\mathcal{Q}}_{x_i} - (-1)^{\tau_i} F^{\mathcal{Q}}_y \right\|_w + \left\| (-1)^{\tau_i} F^{\mathcal{Q}}_y - (-1)^{\tau_i + \tau_j} F^{\mathcal{Q}}_{x_j} \right\|_w \\
&= \left\| F^{\mathcal{Q}}_{x_i} - (-1)^{\tau_i} F^{\mathcal{Q}}_y \right\|_w + \left\| F^{\mathcal{Q}}_{x_i} - (-1)^{\tau_j} F^{\mathcal{Q}}_y \right\|_w \\
&\leqslant \frac{\| F^{\mathcal{Q}}_{x_i} \|_w}{9} + \frac{\| F^{\mathcal{Q}}_{x_j} \|_w}{9} \leqslant \left(\frac{1}{9} + \frac{1}{9} \cdot \frac{5}{4} \right) \| F^{\mathcal{Q}}_{x_i} \|_w \\
&= \frac{\| F^{\mathcal{Q}}_{x_i} \|_w}{4}.
\end{aligned}$$

Thus, $x_j \in \mathscr{R}(x_i)$. But this is impossible.

Therefore, we may conclude that $V(T(x_i)) \cap V(T(x_j)) = \varnothing$. $\qquad \square$

For a subset $E' \subseteq E(G)$ and a mapping $F \colon V \to \mathcal{H}$ to a Hilbert space $\mathcal{H}$, we define energy

$$\widetilde{\mathcal{E}}(E') := \sum_{(x,y) \in E'} \|F_x + F_y\|_{\mathcal{H}}^2.$$

For $S \subseteq V(G)$, let $E(S)$ be the collection of edges that are incident with the vertices in S. We define the energy of S as $\widetilde{\mathcal{E}}(S) := \widetilde{\mathcal{E}}\big(E(S)\big)$.

Lemma 8.4. *Assume that the graph G is a non-bipartite, finite, simple, connected, un-weighted graph. When $\delta \geqslant \lambda_{\min}^{\mathcal{Q}}$, for $1 \leqslant i \leqslant n$,*

$$\widetilde{\mathcal{E}}\big(V(T(x_i))\big) > \frac{\mu^{\mathcal{Q}}(\delta)}{250\big|V(T(x_i))\big|^2}.$$

Proof. a) We first consider the case $w(x_i) > \big|V(T(x_i))\big|$. Recall that $w(x_i)$ equals the degree of x_i, since we are considering unweighted graphs. In this case,

$$
\begin{aligned}
\widetilde{\mathcal{E}}\big(V(T(x_i))\big) &\geqslant \sum_{y \in N(x_i) \backslash V(T(x_i))} \|F_{x_i}^{\mathcal{Q}} + F_y^{\mathcal{Q}}\|_w^2 \\
&> \big|N(x_i) \setminus V(T(x_i))\big| \cdot \frac{\|F_{x_i}^{\mathcal{Q}}\|_w^2}{81} \\
&= \big|N(x_i) \setminus V(T(x_i))\big| \cdot \frac{\mu_{x_i}^{\mathcal{Q}}(\delta)}{81 w(x_i)} \\
&> \big(w(x_i) - \big|V(T(x_i))\big| + 1\big) \frac{\mu^{\mathcal{Q}}(\delta)}{250 w(x_i)} \\
&\geqslant \frac{\mu^{\mathcal{Q}}(\delta)}{250\big|V(T(x_i))\big|^2},
\end{aligned}
$$

where the second inequality holds because $V(T(x_i)) \supseteq \widetilde{N}(x_i)$, and the third inequality holds thanks to Lemma 8.3(1).

b) We now turn to the case $w(x_i) \leqslant \big|V(T(x_i))\big|$. In this case, we claim that there must be a path $z_0 z_1 \cdots z_\ell$, such that $\ell \leqslant \big|V(T(x_i))\big|$, $z_0 = x_i$, and

$$\big\|F_{x_i}^{\mathcal{Q}} - (-1)^\ell F_{z_\ell}^{\mathcal{Q}}\big\|_w > \frac{\|F_{x_i}^{\mathcal{Q}}\|_w}{9}.$$

In fact, if $V\big(T(x_i)\big)$ is a proper subset of $V(G)$, there is a vertex $y \in V(G) \setminus V(T(x_i))$ with $\mathsf{dist}\big(y, T(x_i); G\big) = 1$. Let $z_0 z_1 \cdots z_\ell$ be a path joining $z_0 := x_i$ and $z_\ell := y$, with $z_0 z_1 \cdots z_{\ell-1}$ being a path in $T(x_i)$. Because $T(x_i)$ is assumed maximal and $y \notin V\big(T(x_i)\big)$, we must have

$$\big\|F_{x_i}^{\mathcal{Q}} - (-1)^\ell F_{z_\ell}^{\mathcal{Q}}\big\|_w > \frac{\|F_{x_i}^{\mathcal{Q}}\|_w}{9}.$$

If $V\big(T(x_i)\big) = V(G)$, since $T(x_i)$ is assumed to be a maximal connected bipartite graph and G is a connected non-bipartite graph, there must be an edge $e = (u, v)$, whose addition to $T(x_i)$ results in an odd cycle. To be specific, we have

$$\mathsf{dist}\big(x_i, u; T(x_i)\big) \equiv \mathsf{dist}\big(x_i, v; T(x_i)\big) \quad \bmod 2.$$

Write $\tau := \mathsf{dist}\big(x_i, u; T(x_i)\big)$. Then

$$\left\| F_{x_i}^{\mathcal{Q}} - (-1)^\tau F_u^{\mathcal{Q}} \right\|_w \leqslant \frac{\| F_{x_i}^{\mathcal{Q}} \|_w}{9}.$$

This implies

$$\left\| F_{x_i}^{\mathcal{Q}} - (-1)^{\tau+1} F_u^{\mathcal{Q}} \right\|_w > \frac{\| F_{x_i}^{\mathcal{Q}} \|_w}{9}. \tag{14}$$

In fact, if otherwise $\left\| F_{x_i}^{\mathcal{Q}} - (-1)^{\tau+1} F_u^{\mathcal{Q}} \right\|_w \leqslant \frac{\| F_{x_i}^{\mathcal{Q}} \|_w}{9}$, then

$$\left\| 2 F_{x_i}^{\mathcal{Q}} \right\|_w = \left\| F_{x_i}^{\mathcal{Q}} - (-1)^\tau F_u^{\mathcal{Q}} + F_{x_i}^{\mathcal{Q}} - (-1)^{\tau+1} F_u^{\mathcal{Q}} \right\|_w < \left(\frac{1}{9} + \frac{1}{9} \right) \left\| F_{x_i}^{\mathcal{Q}} \right\|_w.$$

Therefore, $\left\| F_{x_i}^{\mathcal{Q}} \right\|_w = 0$. But $\left\| F_{x_i}^{\mathcal{Q}} \right\|_w \neq 0$ for $\delta \geqslant \lambda_{\min}^{\mathcal{Q}}$.

Now pick a shortest path $z_0 z_1 \cdots z_{\ell-1}$ joining x_i and v in $T(x_i)$, where $z_0 = x_i$ and $z_{\ell-1} = v$. Let $z_\ell = u$. Then $z_0 z_1 \cdots z_{\ell-1} z_\ell$ is a path joining x_i and u. Hence, by our assumption,

$$\ell - 1 \equiv \tau \quad \bmod 2.$$

Therefore, Eq. (14) can be written as

$$\left\| F_{x_i}^{\mathcal{Q}} - (-1)^\ell F_{z_\ell}^{\mathcal{Q}} \right\|_w > \frac{\| F_{x_i}^{\mathcal{Q}} \|_w}{9}.$$

Our claim is proved. By the claim, we have

$$\widetilde{\mathcal{E}}\big(V(T(x_i))\big) \geqslant \sum_{k=0}^{\ell-1} \| F_{z_k}^{\mathcal{Q}} + F_{z_{k+1}}^{\mathcal{Q}} \|_w^2 = \sum_{k=0}^{\ell-1} \| (-1)^k F_{z_k}^{\mathcal{Q}} - (-1)^{k+1} F_{z_{k+1}}^{\mathcal{Q}} \|_w^2$$

$$\geqslant \frac{1}{\ell} \Big(\sum_{k=0}^{\ell-1} \| (-1)^k F_{z_k}^{\mathcal{Q}} - (-1)^{k+1} F_{z_{k+1}}^{\mathcal{Q}} \|_w \Big)^2 \geqslant \frac{1}{\ell} \| F_{z_0}^{\mathcal{Q}} - (-1)^\ell F_{z_\ell}^{\mathcal{Q}} \|_w^2$$

$$> \frac{1}{\ell} \cdot \frac{1}{81} \| F_{x_i}^{\mathcal{Q}} \|_w^2 \geqslant \frac{1}{|V(T(x_i))|} \cdot \frac{\| F_{x_i}^{\mathcal{Q}} \|_w^2}{81}$$

$$= \frac{1}{|V(T(x_i))|} \cdot \frac{\mu_{x_i}^{\mathcal{Q}}(\delta)}{81 w(x_i)} > \frac{1}{|V(T(x_i))|} \cdot \frac{\mu^{\mathcal{Q}}(\delta)}{250 w(x_i)}$$

$$\geqslant \frac{\mu^{\mathcal{Q}}(\delta)}{250 |V(T(x_i))|^2},$$

where the sixth inequality follows from Lemma 8.3(1), and the last inequality holds thanks to our assumption that $w(x_i) \leqslant |V(T(x_i))|$. The proof is complete. $\qquad \square$

Proof of Theorem 8.1. To begin, we claim that

$$\delta \geqslant \frac{\sum_{(y,z)\in E(G)}\|F_y^{\mathcal{Q}} + F_z^{\mathcal{Q}}\|_w^2}{\sum_{x\in V}\|F_x^{\mathcal{Q}}\|_w^2 w(x)}.$$

This can be proved exactly in the same way as [LOG18, Lemma 3.14]. In fact, for all $x, y \in V$,

$$F_x^{\mathcal{Q}}(y) = \langle I_{\mathcal{Q}}(\delta)\mathbf{1}_x/w(x), \mathbf{1}_y \rangle = \langle I_{\mathcal{Q}}(\delta)\mathbf{1}_x/w(x), \mathbf{1}_y/w(y) \rangle_w = F_y^{\mathcal{Q}}(x).$$

Therefore, using Lemma 2.1, we have

$$\sum_{(y,z)\in E(G)} w(y,z)\|F_y^{\mathcal{Q}} + F_z^{\mathcal{Q}}\|_w^2 = \sum_{(y,z)\in E(G)} w(y,z)\sum_{x\in V} w(x)\big|\big(F_y^{\mathcal{Q}} + F_z^{\mathcal{Q}}\big)(x)\big|^2$$
$$= \sum_{(y,z)\in E(G)} w(y,z)\sum_{x\in V} w(x)\big|F_x^{\mathcal{Q}}(y) + F_x^{\mathcal{Q}}(z)\big|^2.$$

Proceeding further, we get that

$$\sum_{(y,z)\in E(G)} w(y,z)\|F_y^{\mathcal{Q}} + F_z^{\mathcal{Q}}\|_w^2 = \sum_{x\in V} w(x) \sum_{(y,z)\in E(G)} w(y,z)\big|F_x^{\mathcal{Q}}(y) + F_x^{\mathcal{Q}}(z)\big|^2$$
$$= \sum_{x\in V} w(x)\langle F_x^{\mathcal{Q}}, \mathcal{Q}F_x^{\mathcal{Q}}\rangle_w$$
$$\leqslant \sum_{x\in V} \delta\, w(x)\|F_x^{\mathcal{Q}}\|_w^2.$$

The claim is thus proved.

Now we assume $\delta \geqslant \lambda_{\min}^{\mathcal{Q}}$ without loss of generality. Since for different i and j,

$$V(T(x_i)) \cap V(T(x_j)) = \varnothing,$$

we have $\sum_{i=1}^m |V(T(x_i))| \leqslant n$. Therefore,

$$\delta \geqslant \frac{\sum_{(x,y)\in E(G)}\|F_x^{\mathcal{Q}} + F_y^{\mathcal{Q}}\|_w^2}{\sum_{y\in V}\|F_y^{\mathcal{Q}}\|_w^2 w(y)} \geqslant \frac{1}{n\mu^{\mathcal{Q}}(\delta)} \cdot \frac{1}{2}\sum_{i=1}^m \widetilde{\mathcal{E}}\big(V\big(T(x_i)\big)\big)$$
$$> \frac{1}{2n\mu^{\mathcal{Q}}(\delta)}\sum_{i=1}^m \frac{\mu^{\mathcal{Q}}(\delta)}{250\big|V(T(x_i))\big|^2} \geqslant \frac{m^3}{500n^3} \geqslant \frac{\mu^{\mathcal{Q}}(\delta)^3}{4000},$$

where the second inequality follows by Lemma 2.3 and the fact that each edge is counted in at most two sets $T(x_i)$ for energy, the fourth inequality follows by convexity of the function $s \mapsto 1/s^2$, and the last inequality holds because $m \geqslant n\mu^{\mathcal{Q}}(\delta)/2$. The proof of Theorem 8.1 is complete. $\square$

8.2 Average Return Probability

Using Theorem 8.1, we can bound average return probabilities. Lemma 8.5 will be useful in the proofs of Theorems 8.6 and 8.7.

Lemma 8.5. *For $t > 0$, we have*

$$\frac{\sqrt[3]{4000}}{3} \int_0^1 (1 - \lambda)^t \lambda^{-2/3} \, d\lambda \leqslant \frac{15}{t^{1/3}}.$$

Proof. We have

$$\int_0^1 (1 - \lambda)^t \lambda^{-2/3} \, d\lambda \leqslant \int_0^\infty \exp\{-\lambda t\} \lambda^{-2/3} \, d\lambda.$$

Taking $\lambda t = s$, we get

$$\int_0^\infty \exp\{-\lambda t\} \lambda^{-2/3} \, d\lambda = \int_0^\infty e^{-s}(s/t)^{-2/3} \, d(s/t)$$
$$= \frac{1}{t^{1/3}} \int_0^\infty e^{-s} s^{-2/3} \, ds = \frac{\Gamma(1/3)}{t^{1/3}}.$$

Therefore,

$$\frac{\sqrt[3]{4000}}{3} \int_0^1 (1 - \lambda)^t \lambda^{-2/3} \, d\lambda \leqslant \frac{\sqrt[3]{4000}\,\Gamma(1/3)}{3t^{1/3}} \leqslant \frac{15}{t^{1/3}}. \qquad \square$$

Theorem 8.6. *Let G be a non-bipartite, finite, simple, connected, unweighted graph. Then for $t \equiv 0 \bmod 2$, we have*

$$0 \leqslant \frac{\sum_{x \in V(G)} p_t(x, x) - 1}{n} \leqslant \frac{30}{t^{1/3}}.$$

Proof. Set

$$\Phi(\lambda) := \begin{cases} n\sqrt[3]{4000\lambda} & \text{if } \lambda \geqslant 0 \text{ and } n\sqrt[3]{4000\lambda} \leqslant n\mu^*(1), \\ n\mu^*(1) & \text{for intermediate values of } \lambda, \\ n - 1 - n\sqrt[3]{4000(2 - \lambda)} & \text{if } \lambda \leqslant 2 \text{ and } n - 1 - n\sqrt[3]{4000(2 - \lambda)} \geqslant n\mu^*(1). \end{cases}$$

Then as in the proof of Theorem 3.4, by Theorem 8.1 and [LOG18, Theorem 5.1],

$$\Phi(0) = 0, \qquad \Phi(2) = n - 1,$$
$$n\mu^*(\lambda) \leqslant \Phi(\lambda) \qquad \text{for} \quad \lambda \in [0, 1],$$
$$n\mu^*(\lambda) \geqslant \Phi(\lambda) \qquad \text{for} \quad \lambda \in [1, 2].$$

Therefore, by our calculation in part a) of the proof of Lemma 2.7, for $t \equiv 0 \bmod 2$,

$$\sum_{x \in V} p_t(x, x) = \sum_{x \in V} \left(\pi(x) + \left(1 - \pi(x) \right) \right) + t \int_0^2 \left(\sum_{x \in V} \mu_x^*(\lambda) \right) (1 - \lambda)^{t-1} \, d\lambda$$

$$= n + t \int_0^2 n \mu^*(\lambda)(1 - \lambda)^{t-1} \, d\lambda \leqslant n + t \int_0^2 \Phi(\lambda)(1 - \lambda)^{t-1} \, d\lambda$$

$$= n - (n - 1) + \int_0^2 (1 - \lambda)^t \Phi'(\lambda) \, d\lambda = 1 + \int_0^2 (1 - \lambda)^t \Phi'(\lambda) \, d\lambda$$

$$\leqslant 1 + 2 \int_0^1 (1 - \lambda)^t \frac{n \sqrt[3]{4000}}{3} \lambda^{-2/3} \, d\lambda$$

$$\leqslant 1 + \frac{30n}{t^{1/3}},$$

where we are using Lemma 2.6 to get the third equality. Therefore,

$$\frac{\sum_{x \in V(G)} p_t(x, x) - 1}{n} \leqslant \frac{30}{t^{1/3}}.$$

On the other hand, for all $x \in V$ and even t, $p_t(x, x) \geqslant \pi(x)$. Hence, we have

$$\frac{\sum_{x \in V(G)} p_t(x, x) - 1}{n} \geqslant 0. \qquad \square$$

Theorem 8.7. *Let G be a non-bipartite, finite, simple, connected, unweighted graph. Then for $t \equiv 1 \bmod 2$,*

$$\frac{\left| \sum_{x \in V(G)} p_t(x, x) - 1 \right|}{n} \leqslant \frac{15}{t^{1/3}}.$$

Proof. a) Let $t \equiv 1 \bmod 2$. By our calculation in part a) of the proof of Lemma 2.7,

$$\sum_{x \in V} p_t(x, x) = \sum_{x \in V} \left(\pi(x) - \left(1 - \pi(x) \right) \right) + t \int_0^2 \left(\sum_{x \in V} \mu_x^*(\lambda) \right) (1 - \lambda)^{t-1} \, d\lambda$$

$$= 2 - n + t \int_0^2 n \mu^*(\lambda)(1 - \lambda)^{t-1} \, d\lambda.$$

b) Set

$$\Psi_1(\lambda) := \left(n - 1 - n \sqrt[3]{4000(2 - \lambda)} \right) \vee 0.$$

Then by Theorem 8.1,

$$\Psi_1(0) = 0, \qquad \Psi_1(2) = n - 1,$$

$$\sum_{x \in V} \mu_x^*(\lambda) \geqslant \Psi_1(\lambda) \qquad \text{for } \lambda \in [0, 2].$$

By our calculation in part a),

$$\sum_{x \in V} p_t(x,x) \geqslant 2 - n + t \int_0^2 \Psi_1(\lambda)(1-\lambda)^{t-1} \, d\lambda$$

$$= 2 - n + (n-1) + \int_0^2 (1-\lambda)^t \Psi_1'(\lambda) \, d\lambda = 1 + \int_0^2 (1-\lambda)^t \Psi_1'(\lambda) \, d\lambda$$

$$\geqslant 1 - \int_0^1 (1-\lambda)^t \frac{n\sqrt[3]{4000}}{3} \lambda^{-2/3} \, d\lambda \geqslant 1 - \frac{15n}{t^{1/3}},$$

where the first equality holds thanks to Lemma 2.6.

c) Set

$$\Psi_2(\lambda) := \left(n\sqrt[3]{4000\lambda} \right) \wedge (n-1).$$

Then by [LOG18, Theorem 5.1],

$$\Psi_2(0) = 0, \qquad \Psi_2(2) = n - 1,$$

$$\sum_{x \in V} \mu_x^*(\lambda) \leqslant \Psi_2(\lambda) \qquad \text{for } \lambda \in [0,2].$$

Using a similar argument as in b), we get

$$\sum_{x \in V} p_t(x,x) \leqslant 1 + \frac{15n}{t^{1/3}}. \qquad \square$$

8.3 Sum of Eigenvalue Powers in Absolute Value

Since similar matrices have the same trace, we have

$$\frac{\sum_{x \in V(G)} p_t(x,x) - 1}{n} = \frac{\sum_{i=1}^n (\lambda_i^P)^t - 1}{n} = \frac{1}{n} \sum_{i=1}^{n-1} (\lambda_i^P)^t.$$

Therefore, when $t \geqslant 2$ is even, Theorem 8.6 gives

$$\frac{1}{n} \sum_{i=1}^{n-1} |\lambda_i^P|^t = \frac{\sum_{x \in V(G)} p_t(x,x) - 1}{n} \leqslant \frac{30}{t^{1/3}}. \tag{15}$$

For odd t, we have

$$\left| \sum_{x \in V(G)} p_t(x,x) - 1 \right| = \left| \sum_{i=1}^{n-1} (\lambda_i^P)^t \right| \leqslant \sum_{i=1}^{n-1} |\lambda_i^P|^t.$$

Theorem 8.7 will not give a bound on $\sum_{i=1}^{n-1} |\lambda_i^P|^t$ directly. But we can still make a detour and bound $\sum_{i=1}^{n-1} |\lambda_i^P|^t$ by Theorem 8.6 as follows.

Proposition 8.8. *Let G be a non-bipartite, finite, simple, connected, unweighted graph. We have*

$$\frac{1}{n}\sum_{i=1}^{n-1}|\lambda_i^P|^t \leqslant \frac{30}{t^{1/3}} \qquad for \quad t \equiv 0 \bmod 2 \ and \ t \geqslant 2,$$

$$\frac{1}{n}\sum_{i=1}^{n-1}|\lambda_i^P|^t \leqslant \frac{30}{(t^2-1)^{1/6}} \qquad for \quad t \equiv 1 \bmod 2 \ and \ t \geqslant 3.$$

The bound in Theorem 8.7 is $\frac{15}{t^{1/3}}$; we get the bound $\frac{30}{(t^2-1)^{1/6}}$ for odd t in Proposition 8.8. These two bounds are not far apart: they are of the same order in t.

Proof of Proposition 8.8. The assertion for even t is nothing but Eq. (15). We now consider odd t. When $t \geqslant 3$ is odd, by Eq. (15),

$$\sum_{i=1}^{n-1}|\lambda_i^P|^t = \sum_{i=1}^{n-1}\big(|\lambda_i^P|^{\frac{t-1}{2}}|\lambda_i^P|^{\frac{t+1}{2}}\big)^{1/2} \leqslant \Big(\sum_{i=1}^{n-1}|\lambda_i^P|^{t-1}\Big)^{1/2}\Big(\sum_{i=1}^{n-1}|\lambda_i^P|^{t+1}\Big)^{1/2}$$

$$\leqslant \Big(\frac{30n}{\sqrt[3]{t-1}}\Big)^{1/2}\Big(\frac{30n}{\sqrt[3]{t+1}}\Big)^{1/2}$$

$$= \frac{30n}{(t^2-1)^{1/6}},$$

where the first inequality holds thanks to the Cauchy–Schwarz inequality. $\qquad\square$

9 Bipartite Case

Now we deal with bipartite graphs. On bipartite graphs, simple random walk has period two. We have the following result on return probabilities, whose proof uses only the bound on the vertex spectral measure of $\mathcal{L}$ from [LOG18, Theorem 4.9].

Theorem 9.1. *Consider a regular, bipartite, simple, connected, unweighted graph G. For each $x \in V$, simple random walk on G satisfies*

$$0 \leqslant p_t(x,x) - 2\pi(x) \leqslant \frac{18}{\sqrt{t}} \qquad for \quad t \equiv 0 \bmod 2.$$

To prove Theorem 9.1, we don't need to get a bound on the vertex spectral measure of $\mathcal{Q}$ as we did in Section 3, because the following lemma gives us a relation between the vertex spectral measures of $\mathcal{L}$ and $\mathcal{Q}$.

Lemma 9.2 (Mohar and Woess, 1989, Theorem 4.8)**.** *If G is bipartite, the spectrum of P is symmetric with respect to zero. For each $x \in V$, the vertex spectral measure of P, $\big\langle I_P(\mathrm{d}\delta)\mathbf{e}_x, \mathbf{e}_x\big\rangle_w$, is symmetric with respect to 0. As a consequence, if G is bipartite, for each $x \in V$ and $\delta \in [0,2]$, $\mu_x(\delta) = \mu_x^{\mathcal{Q}}(\delta)$.* $\qquad\square$

Proof of Theorem 9.1. a) Set

$$\mu_x^{\#}(\lambda) := \begin{cases} \mu_x^*(\lambda) & \text{if } 0 \leqslant \lambda < 2, \\ 1 - 2\pi(x) & \text{if } \lambda = 2, \end{cases}$$

and

$$\varphi^{\#}(\lambda) := \begin{cases} 10\sqrt{\lambda} & \text{if } \lambda \geqslant 0 \text{ and } 10\sqrt{\lambda} \leqslant \mu_x^{\#}(1), \\ \mu_x^{\#}(1) & \text{for intermediate values of } \lambda, \\ \mu_x^{\#}(2) - 10\sqrt{2 - \lambda} & \text{if } \lambda \leqslant 2 \text{ and } \mu_x^{\#}(2) - 10\sqrt{2 - \lambda} \geqslant \mu_x^{\#}(1). \end{cases}$$

Then $\varphi^{\#}(0) = \mu_x^{\#}(0) = 0$, $\varphi^{\#}(2) = \mu_x^{\#}(2) = 1 - 2\pi(x)$. In addition, by Lemma 9.2 and [LOG18, Theorem 4.9], we have

$$\mu_x^{\#}(\lambda) \leqslant \varphi^{\#}(\lambda) \qquad \text{for} \quad \lambda \in [0, 1],$$
$$\mu_x^{\#}(\lambda) \geqslant \varphi^{\#}(\lambda) \qquad \text{for} \quad \lambda \in [1, 2].$$

b) Now let t be a positive even number. We have

$$p_t(x, x) = \int_{[0,2]} (1 - \lambda)^t \, \mu_x(\mathrm{d}\lambda)$$
$$= 2\pi(x) + \int_{[0,2]} (1 - \lambda)^t \, \mu_x^{\#}(\mathrm{d}\lambda). \tag{16}$$

Hence, we see that

$$p_t(x, x) - 2\pi(x) = \int_{[0,2]} (1 - \lambda)^t \, \mu_x^{\#}(\mathrm{d}\lambda) \geqslant 0.$$

c) On the other hand,

$$p_t(x, x) = 2\pi(x) + \int_{[0,2]} (1 - \lambda)^t \, \mu_x^{\#}(\mathrm{d}\lambda)$$

$$= 2\pi(x) + (1 - \lambda)^t \mu_x^{\#}(\lambda)\big|_0^2 - \int_0^2 \mu_x^{\#}(\lambda) \, \mathrm{d}(1 - \lambda)^t$$

$$= 2\pi(x) + \big(1 - 2\pi(x)\big) + t \int_0^2 \mu_x^{\#}(\lambda)(1 - \lambda)^{t-1} \, \mathrm{d}\lambda$$

$$\leqslant 1 + t \int_0^2 \varphi^{\#}(\lambda)(1 - \lambda)^{t-1} \, \mathrm{d}\lambda$$

$$= 1 - \varphi^{\#}(2) + \int_0^2 (1 - \lambda)^t (\varphi^{\#})'(\lambda) \, \mathrm{d}\lambda$$

$$= 2\pi(x) + \int_0^2 (1 - \lambda)^t (\varphi^{\#})'(\lambda) \, \mathrm{d}\lambda$$

$$\leqslant 2\pi(x) + 2 \int_0^1 \frac{5(1 - \lambda)^t}{\sqrt{\lambda}} \, \mathrm{d}\lambda = 2\pi(x) + 10 \int_0^1 \frac{(1 - \lambda)^t}{\sqrt{\lambda}} \, \mathrm{d}\lambda$$

$$\leqslant 2\pi(x) + \frac{18}{\sqrt{t}},$$

where the fourth equality follows from Lemma 2.6, and we use Lemma 3.5 to get the last inequality.

d) Summing up the above discussion, we get, for $x \in V$ and $t \equiv 0 \bmod 2$,

$$0 \leqslant p_t(x,x) - 2\pi(x) \leqslant \frac{18}{\sqrt{t}}. \qquad \square$$

Theorem 9.1 is an example of treating simple random walk on bipartite graphs. Using the same method, one may get parallel results to what we had in previous sections. For instance, the following Corollary 9.3 is parallel to Corollary 7.3.

Corollary 9.3. *Let G be a finite or infinite, vertex-transitive, d-regular, bipartite, simple, connected, unweighted graph with at least polynomial growth rate $N^{\#}(r) \geqslant Cr^D$, where $C > 0$ and $D \geqslant 1$ are constants and $0 \leqslant r \leqslant \mathrm{diam}(G)$. Then for each $x \in V$ and $t \equiv 0 \bmod 2$,*

$$0 \leqslant p_t(x,x) - \pi(x) \leqslant 2\widetilde{C}t^{-D/2},$$

where $\widetilde{C} = \frac{(D+4)^{D/2+2}d^{D/2}}{32CD^{D/2-1}}\Gamma\left(\frac{D}{2}\right)$. $\qquad \square$

To prove Corollary 9.3, one need only follow our argument in the proof of Theorem 9.1 and use the bound on the vertex spectral measure of $\mathcal{L}$ obtained in [LOG18, Theorem 6.1]. Details are omitted.

10   Combinatorial Signless Laplacian

We used spectral embedding to deal with random walk on graphs in previous sections. In fact, this tool is also powerful in analyzing the spectrum of graph adjacency matrices.

Assume that G is a weighted finite graph. Let $\ell^2(V)$ be the Hilbert space of functions $f \colon V \to \mathbb{R}$ or $\mathbb{C}$ with inner product

$$\langle f, g \rangle := \sum_{x \in V} f(x)\overline{g(x)}$$

and squared norm $\|f\|^2 := \langle f, f \rangle$. Let W be the diagonal weight matrix of the graph G: $W := \mathsf{diag}(w(x)\,;\ x \in V)$. Then it is easy to see that the combinatorial signless Laplacian $\Theta := W + A$ is a bounded self-adjoint operator on $\ell^2(V)$. We denote the resolution of identity for Θ as I_Θ and define the vertex spectral measure of Θ at $x \in V$ as

$$\mu_x^\Theta(\delta) := \langle I_\Theta\big([0,\delta]\big)\mathbf{1}_x, \mathbf{1}_x \rangle, \qquad \delta \geqslant 0.$$

For ease of notation, we also write $I_\Theta(\delta) := I_\Theta\big([0,\delta]\big)$ for $\delta \geqslant 0$.

Lemma 10.1. *For $f \in \ell^2(V)$, we have*

$$\langle \Theta f, f \rangle = \sum_{(v,u) \in E(G)} w(v,u)\big|f(v) + f(u)\big|^2.$$

Therefore, the spectrum of Θ is non-negative. Moreover, if $f \in \mathrm{img}\big(I_\Theta(\delta)\big)$ for some $\delta \geqslant 0$, then $\langle \Theta f, f \rangle \leqslant \delta\|f\|^2$.

Proof. We need to prove the first assertion only. Notice that

$$\sum_{v,u\in V} w(v,u)f(v)\overline{\big(f(v)+f(u)\big)} = \sum_{v\in V} f(v)\overline{\sum_{u\in V} w(v,u)f(v)} + \sum_{v\in V} f(v)\overline{\sum_{u\in V} w(v,u)f(u)}$$

$$= \sum_{v\in V} f(v)\overline{w(v)f(v)} + \sum_{v\in V} f(v)\overline{(Af)(v)}$$

$$= \langle f, Wf\rangle + \langle f, Af\rangle$$

$$= \langle f, \Theta f\rangle.$$

By interchanging u and v, we have

$$\sum_{v,u\in V} w(v,u)f(u)\overline{\big(f(v)+f(u)\big)} = \langle f, \Theta f\rangle.$$

Therefore,

$$\sum_{(v,u)\in E(G)} w(v,u)\big|f(v)+f(u)\big|^2$$

$$= \frac{1}{2}\left(\sum_{v,u\in V} w(v,u)f(v)\overline{\big(f(v)+f(u)\big)} + \sum_{v,u\in V} w(v,u)f(u)\overline{\big(f(v)+f(u)\big)}\right)$$

$$= \tfrac{1}{2}\big(\langle f, \Theta f\rangle + \langle f, \Theta f\rangle\big)$$

$$= \langle f, \Theta f\rangle. \qquad \square$$

For $\delta \geqslant 0$, we define the spectral embedding based on Θ as

$$F^{\Theta}: V(G) \to \ell^2(V(G))$$

$$x \mapsto F_x^{\Theta} := I_{\Theta}(\delta)\mathbf{1}_x.$$

It is clear that F_x^{Θ} is a real-valued function on V for each $x \in V$.

Lemma 10.2. *For each finite graph G and $x \in V(G)$,*

$$\big\|F_x^{\Theta}\big\|^2 = F_x^{\Theta}(x) = \mu_x^{\Theta}(\delta).$$

Proof. Since $I_{\Theta}(\delta)$ is a self-adjoint projection operator on $\ell^2(V)$, we see that

$$\big\|F_x^{\Theta}\big\|^2 = \langle F_x^{\Theta}, F_x^{\Theta}\rangle = \langle I_{\Theta}(\delta)\mathbf{1}_x, I_{\Theta}(\delta)\mathbf{1}_x\rangle = \langle I_{\Theta}(\delta)\mathbf{1}_x, \mathbf{1}_x\rangle.$$

By the definition of the vertex spectral measure of Θ, we see that $\langle I_{\Theta}(\delta)\mathbf{1}_x, \mathbf{1}_x\rangle = \mu_x^{\Theta}(\delta)$. Moreover, since $I_{\Theta}(\delta)\mathbf{1}_x = F_x^{\Theta}$, we have $\langle I_{\Theta}(\delta)\mathbf{1}_x, \mathbf{1}_x\rangle = F_x^{\Theta}(x)$. $\qquad \square$

Lemma 10.3. *If $\mu_x^{\Theta}(\delta) > 0$, define $f: V \to \mathbb{C}$ as $f := \frac{F_x^{\Theta}}{\|F_x^{\Theta}\|}$. We have*

(1) $\|f\| = 1$;

(2) $f(x) = \sqrt{\mu_x^\Theta(\delta)}$;

(3) $f \in \mathrm{img}\big(I_\Theta(\delta)\big)$.

Proof. The first and third assertions are obvious. As for the second assertion, by Lemma 10.2,

$$f(x) = \frac{F_x^\Theta(x)}{\|F_x^\Theta\|} = \frac{\mu_x^\Theta(\delta)}{\sqrt{\mu_x^\Theta(\delta)}} = \sqrt{\mu_x^\Theta(\delta)}. \qquad \square$$

We are now in position to present bounds on vertex spectral measures. Denote the eigenvalues of the combinatorial signless Laplacian Θ on G as

$$0 \leqslant \lambda_{\min}^\Theta = \lambda_1^\Theta \leqslant \lambda_2^\Theta \leqslant \lambda_3^\Theta \leqslant \cdots \leqslant \lambda_{n-1}^\Theta < \lambda_n^\Theta = \lambda_{\max}^\Theta.$$

Recall that $\mathscr{K}(G)$ is defined in Definition 6.1; Lemma 6.2 shows that $\mathscr{K}(G) \geqslant \frac{1}{\mathrm{diam}(G)+1}$ when G is non-bipartite.

Proposition 10.4. *Let G be a non-bipartite, finite, simple, connected, weighted graph with weight at least 1 for each edge. Then for each $\delta \in [0, \lambda_{\max}^\Theta)$ and $x \in V$, we have*

$$\mu_x^\Theta(\delta) \leqslant \frac{\delta}{\mathscr{K}(G)} \leqslant \big(\mathrm{diam}(G) + 1\big)\delta.$$

Proof. This proof is a mimic of the proof of Lemma 6.3. Fixing a vertex $x \in V$, we define f as in Lemma 10.3.

When $0 \leqslant \delta < \lambda_{\min}^\Theta$, $\mu_x^\Theta(\delta) = 0$ by definition. So the inequality holds automatically.

When $\lambda_{\min}^\Theta \leqslant \delta < \lambda_{\max}^\Theta$, we know that f is orthogonal to the eigenspace of Θ corresponding to $\lambda_{\max}^\Theta$, which is spanned by a positive vector according to the Perron–Frobenius theorem. By Definition 6.1, we have

$$\delta \geqslant \langle f, \Theta f \rangle = \sum_{(v,u) \in E(G)} w(v,u)|f(v) + f(u)|^2 \geqslant f(x)^2 \mathscr{K}(G) = \mu_x^\Theta(\delta)\mathscr{K}(G). \qquad \square$$

In order to get a lower bound on eigenvalues of Θ from Proposition 10.4, we need the following Lemma 10.5.

Lemma 10.5. *Let G be a finite, connected, weighted graph. We have*

$$\sum_{x \in V} \mu_x^\Theta(\delta) = \big|\{j \,;\, \lambda_j^\Theta \leqslant \delta\}\big|.$$

Proof. Let $h_1, h_2, \ldots, h_n$ be an orthonormal basis of $\ell^2(V)$ such that

$$\Theta h_j = \lambda_j^\Theta h_j, \qquad 1 \leqslant j \leqslant n.$$

It follows that

$$\sum_{x \in V} |\langle h_j, \mathbf{1}_x \rangle|^2 = \|h_j\|^2 = 1.$$

Therefore,

$$\sum_{x \in V} \mu_x^\Theta(\delta) = \sum_{x \in V} \big\langle I_\Theta(\delta)\mathbf{1}_x, \mathbf{1}_x \big\rangle.$$

This is the trace of $I_\Theta(\delta)$, which equals the dimension of its image. $\qquad \square$

Corollary 10.6. *Let G be a non-bipartite, finite, simple, connected, weighted graph with weight at least 1 for each edge. For $1 \leqslant k \leqslant n$, we have*

$$\lambda_k^\Theta \geqslant \frac{k \mathscr{K}(G)}{n} \geqslant \frac{k}{\big(\mathrm{diam}(G) + 1\big)n}.$$

Proof. By Proposition 10.4,

$$\sum_{x \in V} \mu_x^\Theta(\delta) \leqslant \sum_{x \in V} \frac{\delta}{\mathscr{K}(G)} = \frac{n\delta}{\mathscr{K}(G)}.$$

Hence, Lemma 10.5 gives

$$\big|\{j ;\ \lambda_j^\Theta \leqslant \delta\}\big| \leqslant \frac{n\delta}{\mathscr{K}(G)}.$$

Therefore, we must have $\lambda_k^\Theta \geqslant \frac{k\mathscr{K}(G)}{n}$. $\qquad\square$

For the adjacency matrix A of G, we denote its eigenvalues as

$$-w_{\max} \leqslant \lambda_{\min}^A = \lambda_1^A \leqslant \lambda_2^A \leqslant \lambda_3^A \leqslant \cdots \leqslant \lambda_{n-1}^A < \lambda_n^A = \lambda_{\max}^A \leqslant w_{\max},$$

where $w_{\max} := \max_{x \in V} w(x)$. The following Corollary 10.7 improves [Alon and Sudakov, 2000, Theorem 1.1], which obtained that $d_{\max} + \lambda_1^A \geqslant \frac{1}{(\mathrm{diam}(G)+1)n}$ for unweighted graphs.

Corollary 10.7. *Let G be a non-bipartite, finite, simple, connected, weighted graph with weight at least 1 for each edge. For $1 \leqslant k \leqslant n$, we have*

$$w_{\max} + \lambda_k^A \geqslant \frac{k \mathscr{K}(G)}{n} \geqslant \frac{k}{\big(\mathrm{diam}(G) + 1\big)n}.$$

Proof. Let X_1 be the linear subspace of $\ell^2(V)$ spanned by the eigenvectors of A corresponding to λ_1^A, λ_2^A, ..., λ_k^A, and X_2 be the linear subspace of $\ell^2(V)$ spanned by the eigenvectors of Θ corresponding to λ_k^Θ, λ_{k+1}^Θ, ..., λ_n^Θ. Then

$$\dim X_1 \geqslant k, \quad \dim X_2 \geqslant n - k + 1.$$

It follows that

$$\dim X_1 + \dim X_2 > n = \dim \ell^2(V).$$

Therefore, the intersection of X_1 and X_2 contains a non-zero vector h of unit norm. Hence,

$$\frac{k \mathscr{K}(G)}{n} \leqslant \lambda_k^\Theta \leqslant (\Theta h, h) = (Wh, h) + (Ah, h) \leqslant w_{\max} + \lambda_k^A. \qquad\square$$

Remark. The combinatorial signless Laplacian is also related to line graphs. Let $\rho_1 \geqslant \rho_2 \geqslant \cdots \geqslant \rho_r$ be the positive combinatorial signless Laplacian eigenvalues of G. Then by [Brouwer and Haemers, 2012, Proposition 1.4.1], the eigenvalues of the line graph of G are

$$\theta_i = \rho_i - 2, \quad i = 1, 2, \ldots, r,$$
$$\theta_i = -2, \quad i = r+1, r+2, \ldots, |E(G)|.$$

Using this relation and Corollary 10.6, one may do some quick analysis on the spectrum of the line graph of G.

Acknowledgments

Abbreviations

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