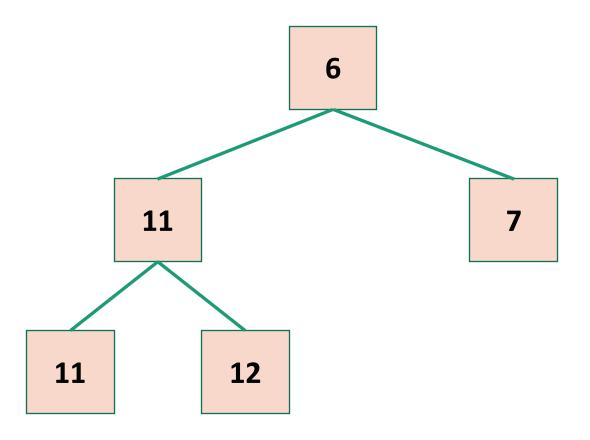
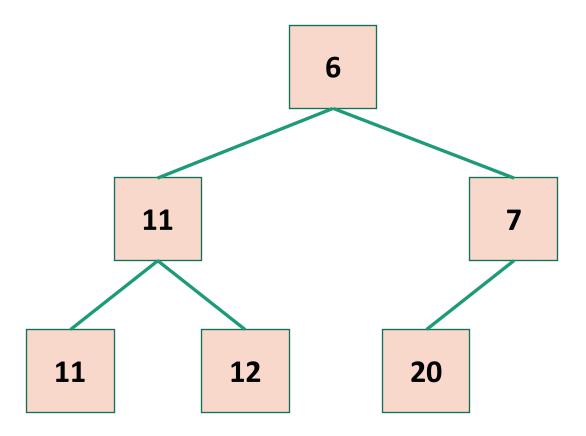


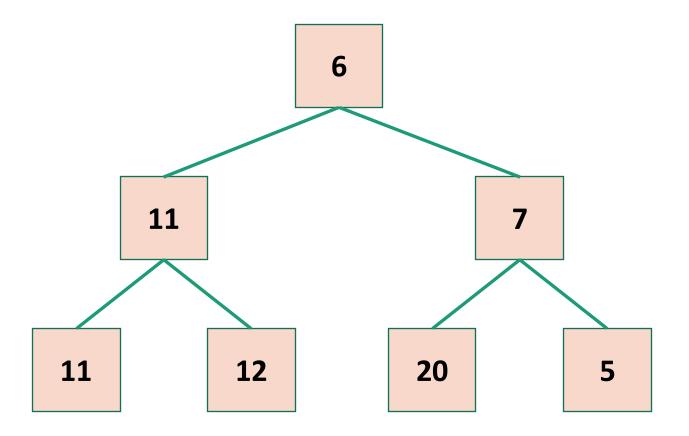
Priority Queue and Heap

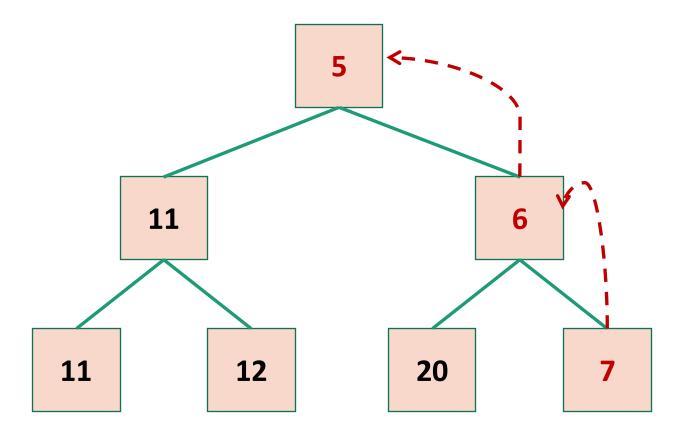
Some slides were created by Dr. Jianfeng Ren. Edited by Heshan Du

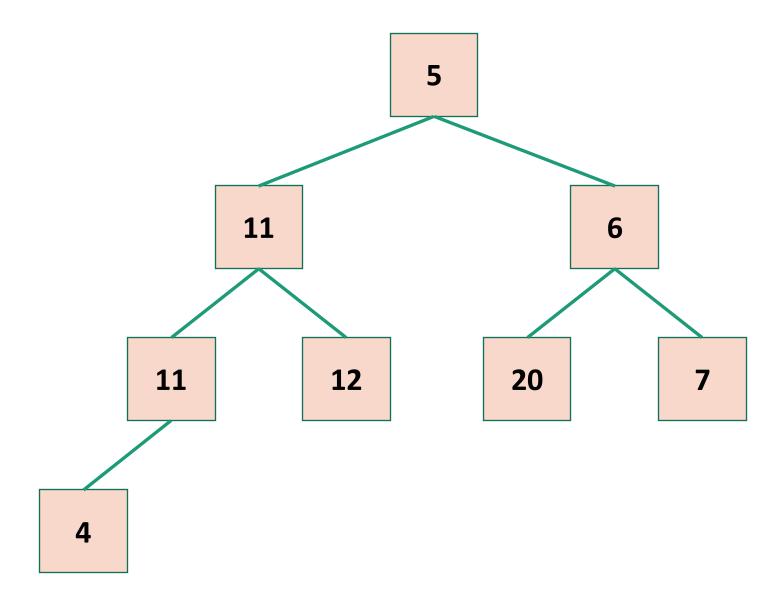
Exercise 1: insert 20, 5, 4

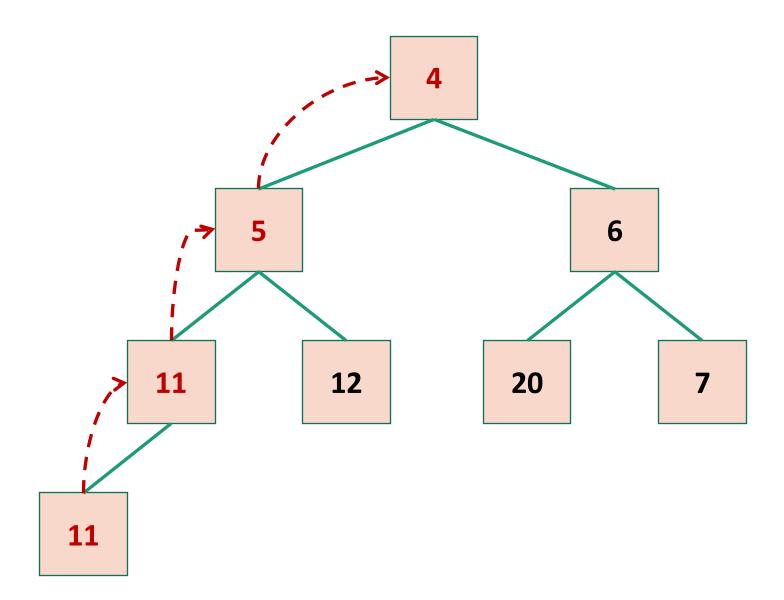




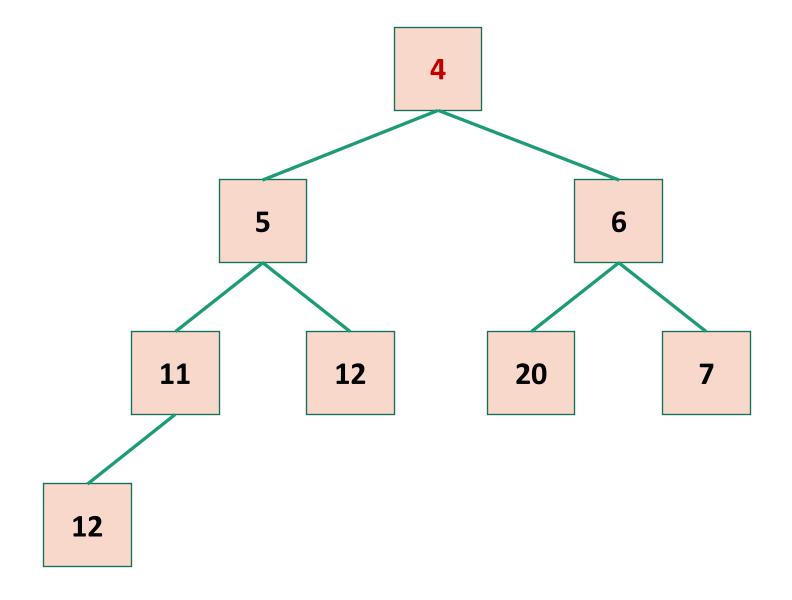




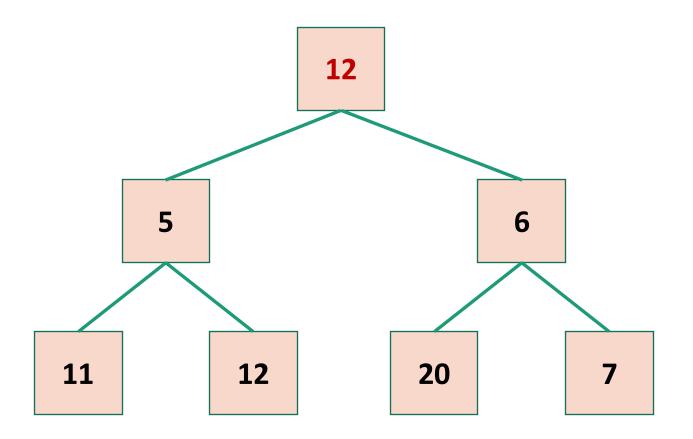




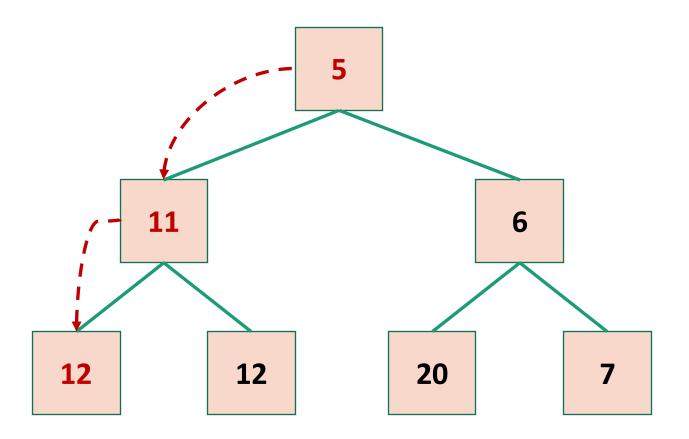
Exercise 2: remove 4



remove 4



remove 4



Exercise 3: heap sort

要使用 in-place heap sort 将序列 [3, 6, 9, 2, 5, 8] 排序为非递增顺序 (从大到小)

 Sort the following sequence in non-increasing order using in-place heap sort:

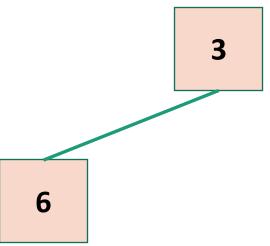
[3 6 9 2 5 8]

- Original sequence: [3 6 9 2 5 8]
- Array: [3 6 9 2 5 8]

3

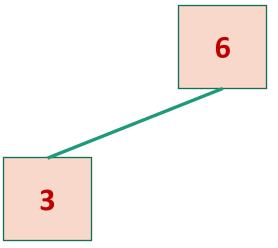
• Original sequence: [3 6 9 2 5 8]

• Array: [36 9 2 5 8]



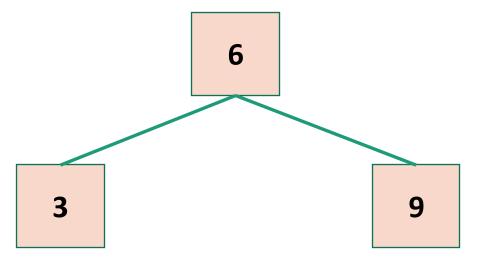
• Original sequence: [3 6 9 2 5 8]

• Array: [<u>6 3</u> 9 2 5 8]



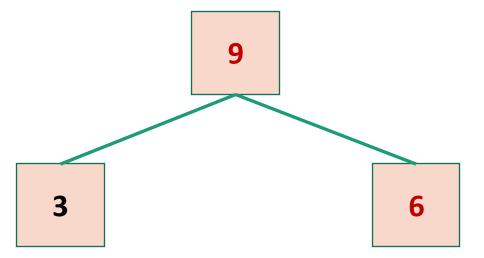
• Original sequence: [3 6 9 2 5 8]

• Array: [<u>6 3 9</u> 2 5 8]

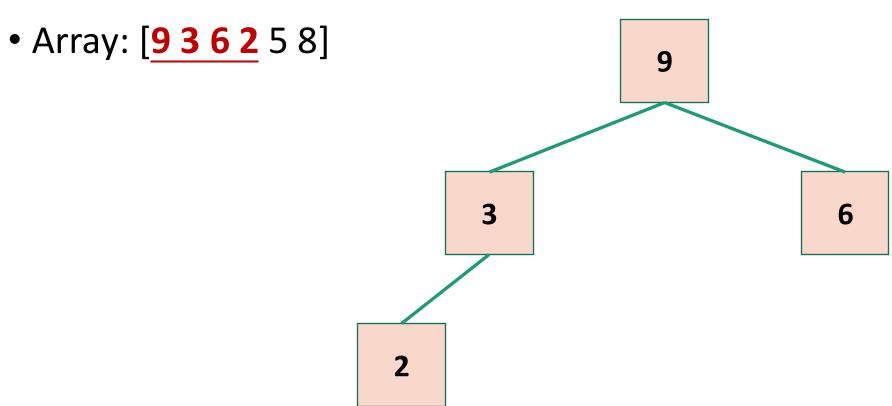


• Original sequence: [3 6 9 2 5 8]

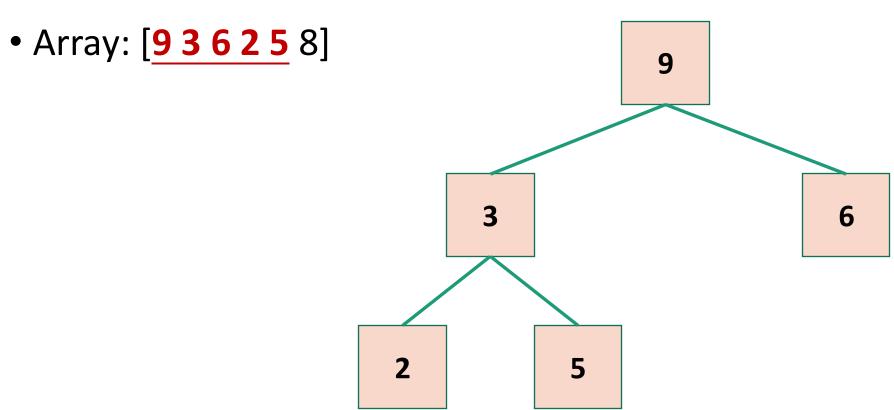
• Array: [936 258]



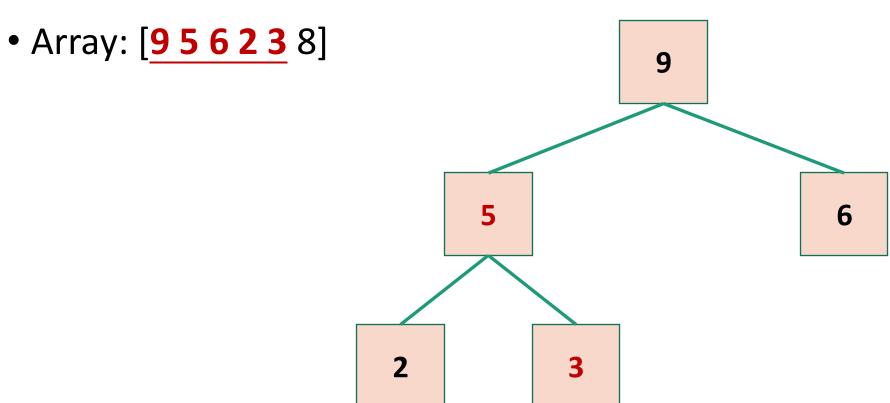
• Original sequence: [3 6 9 2 5 8]



• Original sequence: [3 6 9 2 5 8]



• Original sequence: [3 6 9 2 5 8]



• Original sequence: [3 6 9 2 5 8]

• Array: [9 5 6 2 3 8]

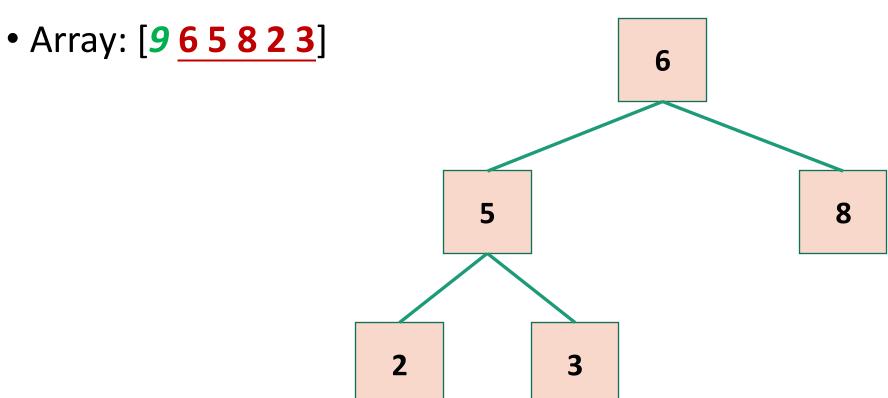
• Original sequence: [3 6 9 2 5 8]

• Array: [9 5 8 2 3 6]

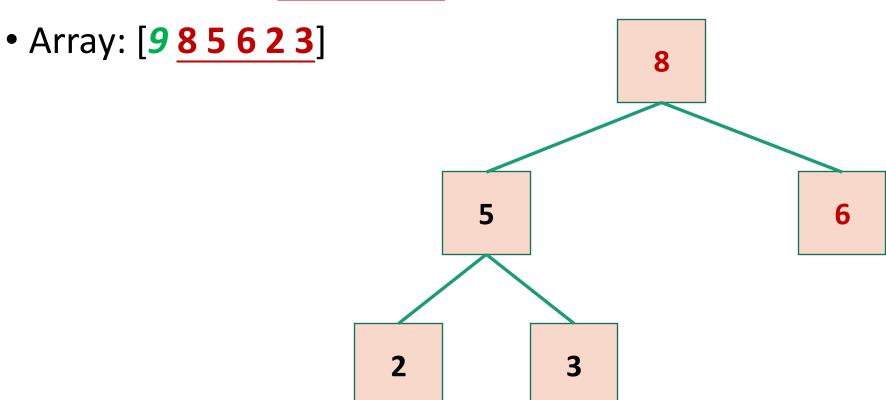
• Original heap: [958236]

• Array: [9 5 8 2 3 6]

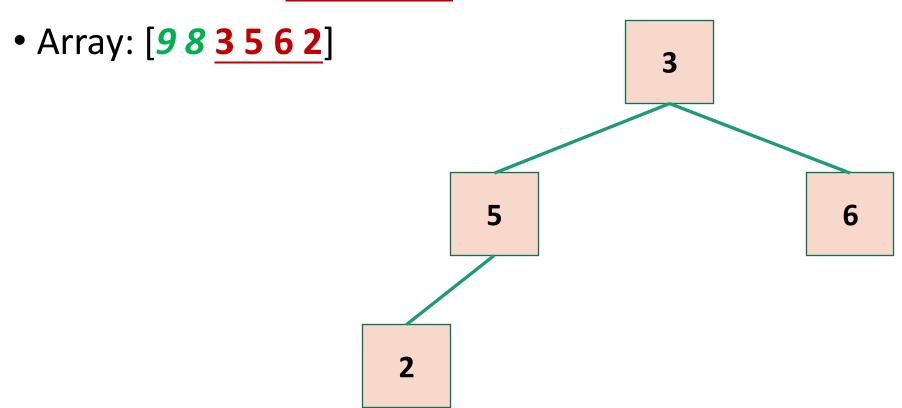
• Original heap: [958236]



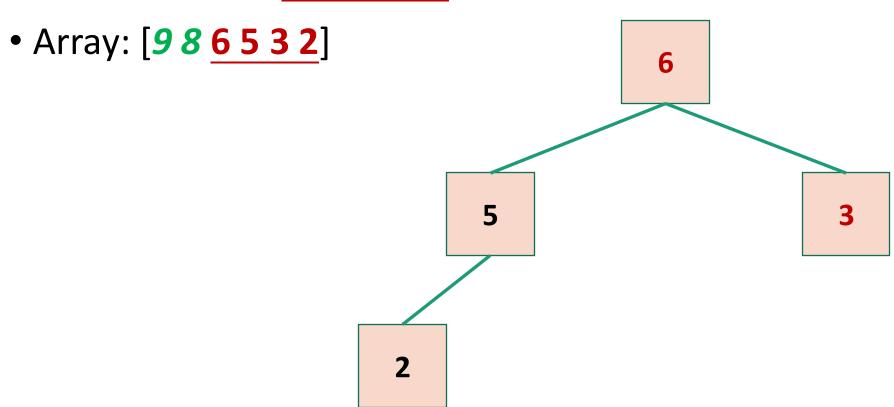
• Original heap: [958236]



• Original heap: [958236]

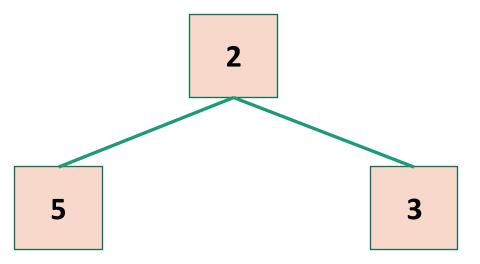


• Original heap: [9 5 8 2 3 6]



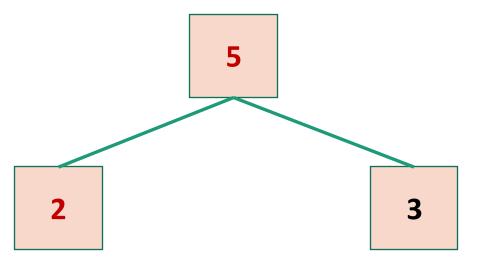
• Original heap: [9 5 8 2 3 6]

• Array: [986253]



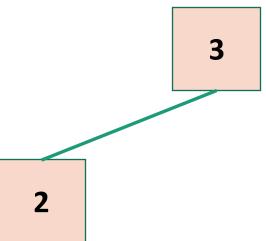
• Original heap: [9 5 8 2 3 6]

• Array: [986523]



• Original heap: [9 5 8 2 3 6]

• Array: [986532]



- Original heap: [9 5 8 2 3 6]
- Array: [986532]

2

- Original heap: [958236]
- Array: [986532]

Exercise 4

Illustrate the execution of the selection-sort algorithm on the following input sequence:

(22, 15, 36, 44, 10, 3, 9, 13, 29, 25).

每轮从未排序部分中选择最小值放到当前起始位置。总共进行 n-1 轮比较。

Exercise 4: Solution

22 15	36	44	10	3			29	25
3 15	36	44	10	22	9	13	29	25
39	36	44	10	22	15	13	29	25
39	10	44	36	22	15	13	29	25
3 9	10	13	36	22	15	44	29	25
3 9	10	13	15	22	36	44	29	25
3 9	10	13	15	22	36	44	29	25
3 9	10	13	15	22	25	44	29	36
3 9	10	13	15	22	25	29	44	36
3 9	10	13	15	22	25	29	36	44

Exercise 5

Illustrate the execution of the insertion-sort algorithm on the following input sequence:

(22, 15, 36, 44, 10, 3, 9, 13, 29, 25).

从左往右依次将元素插入到左侧已经排好序的子序列中。 每一步都保持左侧序列有序。

步骤	已排序部分 (左侧)	当前插入元素	操作说明
初始	_	_	[22, 15, 36, 44, 10, 3, 9, 13, 29, 25]
1	[22]	15	插入前面: → [15, 22]
2	[15, 22]	36	保持原位: → [15, 22, 36]
3	[15, 22, 36]	44	保持原位: → [15, 22, 36, 44]
4	[15, 22, 36, 44]	10	插入最前面: → [10, 15, 22, 36, 44]
(5)	[10, 15, 22, 36, 44]	3	插入最前面: → [3, 10, 15, 22, 36, 44]
6	[3, 10, 15, 22, 36, 44]	9	插入 10 前面: → [3, 9, 10, 15, 22, 36, 44]
7	[3, 9, 10, 15, 22, 36, 44]	13	插入 15 前面: → [3, 9, 10, 13, 15, 22, 36, 44]
8	[3, 9, 10, 13, 15, 22, 36, 44]	29	插入 36 前面: → [3, 9, 10, 13, 15, 22, 29, 36, 44]
9	[3, 9, 10, 13, 15, 22, 29, 36, 44]	25	插入 29 前面: → [3, 9, 10, 13, 15, 22, 25, 29, 36, 44]

Exercise 5: Solution

22	15	36	44	10	3	9	13	29	25
15	22	36	44	10	3	9	13	29	25
15	22	36	44	10	3	9	13	29	25
10	15	22	36	44	3	9	13	29	25
3	10	15	22	36	44	9	13	29	25
3	9	10	15	22	36	44	13	29	25
3	9	10	13	15	22	36	44	29	25
3	9	10	13	15	22	29	36	44	25
3	9	10	13	15	22	25	29	36	44

Exercise 6

Show that the sum $\sum_{i=1}^{n} \log i$, appearing in the analysis of heap-sort, is $O(n \log n)$.

$$\sum_{i=1}^n \log i \approx \int_1^n \log x \, dx$$

计算积分:

$$\int_{1}^{n} \log x \, dx = x \log x - x \mid_{1}^{n} = n \log n - n + 1$$

因此:

$$\sum_{i=1}^n \log i = \Theta(n \log n)$$

从而我们有:

$$\sum_{i=1}^{n} \log i = O(n \log n)$$

Exercise 6: Solution

Proof: To show $\sum_{i=1}^{n} \log i$ is $O(n \log n)$, by the definition of O, we need to show that there exist a positive real constant c and a positive integer n_0 such that for all $n \geq n_0$, $\sum_{i=1}^{n} \log i \leq cn \log n$.

$$\sum_{i=1}^{n} \log i = \log 1 + \log 2 + \dots + \log n$$

$$\leq n \log n$$

for every $n \ge 1$. Let c = 1, $n_0 = 1$. By the definition of O, we have $\sum_{i=1}^{n} \log i$ is $O(n \log n)$.

Exercise 7

Show that the sum $\sum_{i=1}^{n} \log i$, appearing in the analysis of heap-sort, is $\Omega(n \log n)$.

$$\sum_{i=1}^n \log i \ge \sum_{i=\lceil n/2 \rceil}^n \log i \ge \sum_{i=\lceil n/2 \rceil}^n \log \frac{n}{2} = \left(\frac{n}{2}\right) \log \frac{n}{2}$$

计算这个下界:

$$\left(\frac{n}{2}\right)\log\frac{n}{2} = \frac{n}{2}(\log n - \log 2) = \frac{n}{2}\log n - \frac{n}{2}\log 2 = \Omega(n\log n)$$

Exercise 7: Solution

Proof: To show $\sum_{i=1}^{n} \log i$ is $\Omega(n \log n)$, by the definition of Ω , we need to show that there exist a positive real constant c and a positive integer n_0 such that for all $n \geq n_0$, $\sum_{i=1}^{n} \log i \geq c n \log n$.

$$\sum_{i=1}^{n} \log i = \log 1 + \log 2 + \dots + \log n$$

$$\geq \frac{n}{2} \log \frac{n}{2} \quad \text{(consider the last } \frac{n}{2} \text{ terms)}$$

$$= \frac{n}{2} \log n - \frac{n}{2}$$

Let $c = \frac{1}{4}$. By solving $\frac{n}{2} \log n - \frac{n}{2} \ge cn \log n$, we have $n \ge 4$. So we let $n_0 = 4$. By the definition of Ω , we have $\sum_{i=1}^{n} \log i$ is $\Omega(n \log n)$.