

Wednesday 27 April 2022 14:00–15:30 BST

Duration: 1 hour 30 minutes Additional time: 30 minutes Timed exam – fixed start time

DEGREES OF MSci, MEng, BEng, BSc, MA and MA (Social Sciences)

ALGORITHMS AND DATA STRUCTURES 2 COMPSCI2007

Answer all 5 questions

This examination paper is an open book, online assessment and is worth a total of 60 marks

1. Algorithm **F** takes as input an array of integers A and two indices l, r of A. Assume A contains both negative and positive integers and no duplicate elements. The algorithm is described by the following pseudocode:

```
1: \mathbf{F}(A, l, r)
 2:
       x := 0
 3:
       y := 0
       if l = r
 4:
          x := A[l]
 5:
 6:
          y := A[l]
       else if l+1=r
 7:
             if A[l] < A[r]
 8:
                x := A[r]
 9:
10:
                y := A[l]
             else
11:
                x := A[l]
12:
                y := A[r]
13:
       else
14:
          q := l + (r - l)/2
                                    // integer division
15:
          (xl,yl) := \mathbf{F}(A,l,q)
16:
17:
          (xr, yr) := \mathbf{F}(A, q+1, r)
          if xl > xr
18:
             x := xl
19:
          else
20:
21:
             x := xr
          if yl < yr
22:
             y := yl
23:
24:
25:
             y := yr
       return (x, y)
26:
```

(a) Briefly explain what algorithm **F** implements.

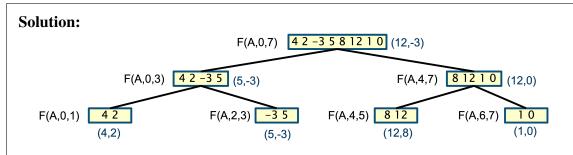
Solution: This algorithm finds the maximum (1 point) and minimum (1 point) elements in the subarray A[l..r] (1 point). The maximum and minimum values are stored in variables x and y (2 points), respectively.

(b) What is the output of $\mathbf{F}(A,0,7)$, where A = [4,2,-3,5,8,12,1,0]?

```
Solution: The output is the pair (12, -3).
```

(c) Draw the recursion tree for $\mathbf{F}(A,0,7)$. [3]

[5]



1 point for correct tree structure, 1 point for correct input values, and 1 point for correct return values.

(d) Is **F** an in-place algorithm? Justify your answer.

[2]

Solution: Yes, as no auxiliary data structure is allocated during the execution of the algorithm. Only the following integer auxiliary variables are used: x, y, q, xl, yl, xr, and yr.

(e) Identify and briefly explain with simple language the Divide, Conquer and Combine steps in Algorithm F. [3]

Solution: The three steps are as follows:

Divide: input array A is divided into two equal parts around the mid-value.

Conquer: recursively find the maximum and minimum of both the left and right subarrays.

Combine: compare the maximum and minimum of both sub-arrays to find the maximum and minimum of the whole array.

(f) Write the recurrence equation for **F** and solve it to compute its running time using Big-O notation. [5]

Solution: The recurrence equation is

$$T(1) = T(2) = O(1)$$

 $T(n) = T(n/2) + T(n/2) + O(1)$

Solving with the iterative method we obtain

$$T(n) = 2T(n/2) + c_1$$

$$= 4T(n/4) + 2c_1$$

$$= ...$$

$$= 2^k T(n/2^k) + kc_1$$

$$= 2^{\log n} T(1) + (\log n)c_1$$

$$= nc_2 + (\log n)c_1$$

$$= O(n)$$

1 point for base case, 1 point for step case, 1 point for correct iteration up to k, 1 point for correct substitution with $\log n$, 1 point for correct solution.

2. Apply the Master Theorem to give tight asymptotic bounds for the following recurrences. In your answers, briefly explain which case applies and why.

(a)
$$T(n) = 2T(n/2) + n^3$$
 [3]

Solution: This is a divide-and-conquer recurrence with a=2, b=2, $f(n)=n^3$, and c=3. Since $c>\log_b a$, that is $3>\log_2 2=1$, case 3 of the master theorem applies, and $T(n)=\Theta(n^3)$.

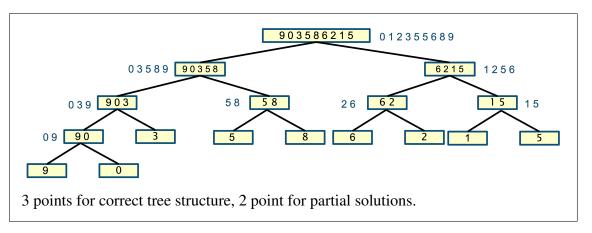
(b)
$$T(n) = 16T(n/4) + n^2$$
 [3]

Solution: This is another divide-and-conquer recurrence with a=16, b=4, $f(n)=n^2$, and c=2. Since $c=\log_b a$, that is $2=\log_4 16=2\log_4 4=2$, case 2 of the master theorem applies, and $T(n)=\Theta(n^2\log n)$.

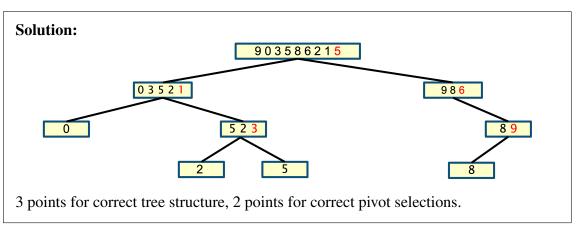
3. Draw the recursion tree computed when sorting array A = [9,0,3,5,8,6,2,1,5] with the following two algorithms:

(a)
$$MERGE-SORT(A, 0, 8)$$
. [5]

Solution:



(b) QUICKSORT(A,0,8). Assume the right-most element is selected as pivot when partitioning. [5]

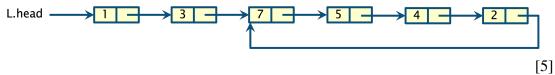


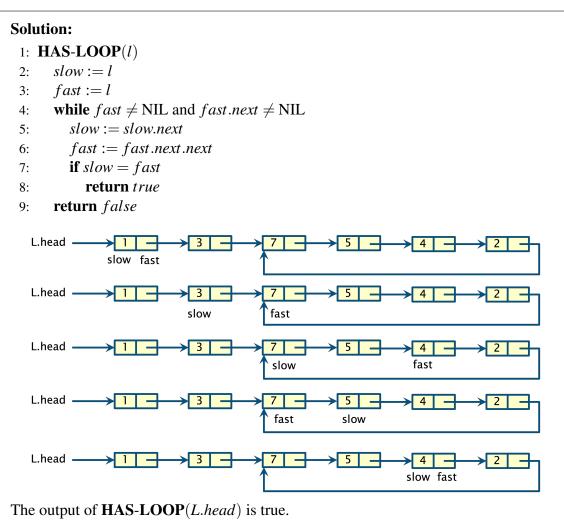
4. (a) Consider a bounded queue implemented by array Q[0..4]. Illustrate the result of each operation in the sequence $\mathbf{ENQUEUE}(Q,2)$, $\mathbf{DEQUEUE}(Q)$, $\mathbf{ENQUEUE}(Q,6)$, $\mathbf{ENQUEUE}(Q,0)$, and $\mathbf{DEQUEUE}(Q)$. Assume the initial configuration is Q = [*,1,7,5,*], with Q.head = 1 and Q.tail = 4. [5]

Solution:

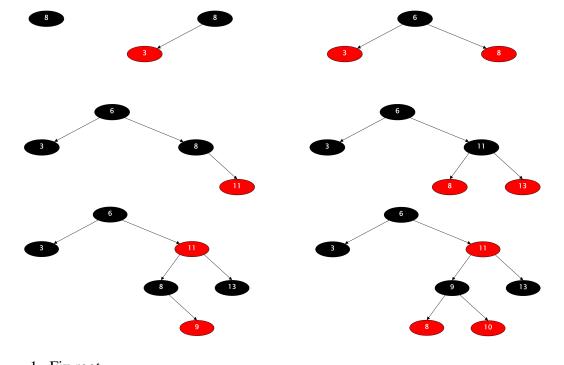
- **ENQUEUE**(Q,2), Q = [*,1,7,5,2], Q.head = 1, Q.tail = 0;
- **DEQUEUE**(Q), return 1, Q = [*, *, 7, 5, 2], Q.head = 2;
- **ENQUEUE**(Q,6), Q = [6,*,7,5,2], Q.head = 2, Q.tail = 1;
- **ENQUEUE**(Q,0), error "overflow";
- **DEQUEUE**(Q), return 7, Q = [6, *, *, 5, 2], Q.head = 3.
- (b) Given the head of a singly-linked list, define using pseudocode an algorithm that detects

if the list has a cycle. Use an iterator with slow and fast pointers. Illustrate how it operates on the following input:





5. (a) Insert keys 8,3,6,11,13,9,10 (in this order) in the empty red-black tree *T* and show at each step how *T* is updated. Explicitly mention at each step which operations are performed (e.g. rotations, recolourings, etc). [7]



- 1. Fix root
- 2. –
- 3. Left-rotation on 3, right-rotation on 6, recolouring 6 and 8
- 4. Recolouring 3 and 8, fix root
- 5. Left-rotation on 8, recolouring 11 and 8
- 6. Recolouring 8, 13, and 11
- 7. Left-rotation on 8, recolouring 8 and 9
- (b) Insert keys 12,7,15,4,6,14,13 (in this order) in the empty B-tree T with t=2 (i.e. each node can store at most 3 keys) and show at each step how T is updated. Explicitly mention at each step which operations are performed. Assume insertion implements a one-pass strategy. [7]

Solution:

