THE LEBESGUE UNIVERSAL COVERING PROBLEM

JOHN C. BAEZ AND KARINE BAGDASARYAN AND PHILIP GIBBS

ABSTRACT. In 1914 Lebesgue defined a 'universal covering' to be a convex subset of the plane that contains an isometric copy of any subset of diameter 1. His challenge of finding a universal covering with the least possible area has been addressed by various mathematicians: Pál, Sprague and Hansen have each created a smaller universal covering by removing regions from those known before. However, Hansen's last reduction was microsopic: he claimed to remove an area of $6 \cdot 10^{-18}$, but we show that he actually removed an area of just $8 \cdot 10^{-21}$. In the following, with the help of Greg Egan, we find a new, smaller universal covering with area less than 0.8441153. This reduces the area of the previous best universal covering by $2.2 \cdot 10^{-5}$.

1. Introduction

A 'universal covering' is a convex subset of the plane that covers a translated, reflected and/or rotated version of every subset of the plane with diameter 1. In 1914, Lebesgue [8] laid down the challenge of finding the universal covering with least area, and since then various other mathematicians have worked on it, slightly improving the results each time. The most notable of them are J. Pál [10], R. Sprague [12], and H. C. Hansen [8]. More recently Brass and Sharifi [3] have found a lower bound on the area of a universal covering, while Duff [4] showed that dropping the requirement of convexity allows for an even smaller area.

The simplest universal covering is a regular hexagon in which one can inscribe a circle of diameter 1. All the improved universal coverings so far have been constructed by removing pieces of this hexagon. Here we describe a new universal covering with less area than those known so far.

To state the problem precisely, recall that the diameter of a set of points A in a metric space is defined by

$$diam(A) = \sup\{d(x, y) : x, y \in A\}.$$

We define a **universal covering** to be a convex set of points U in the Euclidean plane \mathbb{R}^2 such that any set $A \subseteq \mathbb{R}^2$ of diameter 1 is isometric to a subset of U. If we write \cong to mean 'isometric to', we can define the collection of universal coverings to be

$$\mathcal{U} = \{U : \text{ for all } A \text{ such that } \operatorname{diam}(A) = 1 \text{ there exists } B \subseteq U \text{ such that } B \cong A\}.$$

The quest for smaller and smaller universal coverings is in part the search for this number:

$$a = \inf\{\operatorname{area}(U) : U \in \mathcal{U}\}.$$

In 1920, Pál [10] made significant progress on Lebesgue's problem. The regular hexagon in which a circle of diameter 1 can be inscribed has sides of length $1/\sqrt{3}$. Pál showed that this hexagon gives a universal covering, and thus

$$a \le \frac{\sqrt{3}}{2} = 0.86602540\dots$$

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1

Pál then took this regular hexagon and fit the largest possible regular dodecagon inside it, as shown in Figure 1. He then proved that two of the resulting corners could be removed to give a smaller universal covering. This implies that

$$a \le 2 - \frac{2}{\sqrt{3}} = 0.84529946...$$

(Here and in all that follows, when we speak of 'removing' a closed set $Y \subseteq \mathbb{R}^2$ from another closed set $X \subseteq \mathbb{R}^2$, we mean to take the set-theoretic difference X - Y and then form its closure. The universal covers we discuss are all closed.)

In 1936, Sprague went on to prove that more area could be removed from another corner of the original hexagon. This proved

$$a \le 0.844137708436$$
.

In 1992, Hansen [8] took these reductions even further by removing more area from two different corners of Pál's universal covering. He claimed that the areas that could be removed were $4 \cdot 10^{-11}$ and $6 \cdot 10^{-18}$, but we have found that his second calculation was mistaken. The actual areas removed were $3.7507 \cdot 10^{-11}$ and $8.4541 \cdot 10^{-21}$, showing

We have created a Java applet illustrating Hansen's universal covering [7]. This allows the user to zoom in and see the tiny, extremely narrow regions removed by Hansen.

Our new improved universal covering gives

$$a \le 0.8441153$$
.

This is about $2.2 \cdot 10^{-5}$ less than the previous best upper bound, due to Hansen. On the other hand, Brass and Sharifi [3] have given the best known lower bound on the area of a universal covering. They did so by using various shapes such as a circle, a pentagon, and a triangle of diameter 1 in a certain alignment. They concluded that

$$a \ge 0.832$$
.

Thus, at present we know that

$$0.832 < a < 0.8441153$$
.

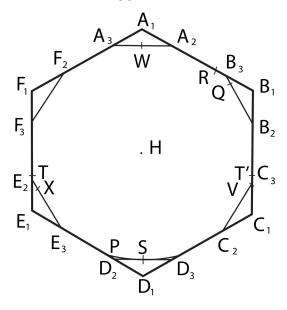
In their book Old and New Unsolved Problems in Plane Geometry and Number Theory, Klee and Wagon [9] wrote:

...it does seem safe to guess that progress on [this problem], which has been painfully slow in the past, may be even more painfully slow in the future.

However, our reduction of Hansen's universal covering is about a million times greater than Hansen's reduction of Sprague's. The use of computers played an important role. Given the gap between the area of our universal covering and Brass and Sharifi's lower bound, we can hope for even faster progress in the decades to come!

2. Hansen's Calculation

Let us recall Hansen's universal covering [8] and correct his calculation of its area.



1. Figure similar to the one used by Hansen

Figure 1 shows a regular hexagon $A_1B_1C_1D_1E_1F_1$ in which can be inscribed a circle of diameter 1. Inside this is a regular dodecagon, $A_3A_2B_3B_2C_3C_2D_3D_2E_3E_2F_3F_2$. Pál showed that one can take the hexagon, remove the triangles $C_1C_2C_3$ and $E_1E_2E_3$, and still obtain a universal covering. We generalize his proof in Reduction 1 of Section 3.

Later, Sprague [12] constructed a smaller universal covering. To do this he showed that near A_1 the region outside the circle of radius 1 centered at E_3 could be removed, as well as the region outside the circle of radius 1 centered at C_2 . We generalize this argument in Reduction 2.

Building on Sprague's work, Hansen constructed a still smaller universal covering. First he removed a tiny region XE_2T , almost invisible in this diagram. Then he removed a much smaller region $T'C_3V$. Each of these regions is a thin sliver bounded by two line segments and an arc, as shown in our Java applet [7].

To compute the areas of these regions, Hansen needed to determine certain distances x_0, \ldots, x_4 . The points and arcs he used, and these distances, are defined as follows:

- The distance from W, the midpoint of A_2A_3 , to the dodecagon corner A_3 is x_0 .
- The circle of radius 1 centered at W is tangent to D_2D_3 at S and hits the dodecagon at P. The distance from S to the dodecagon corner D_2 can be seen to equal x_0 . The distance from D_2 to P is x_1 .
- The circle of radius 1 centered at P is tangent to B_3A_2 at R and hits the dodecagon at Q. The distance from R to the dodecagon corner B_3 can be seen to equal x_1 . The distance from B_3 to Q is x_2 .
- The circle of radius 1 centered at Q is tangent to E_2E_3 at X and hits the dodecagon at T. The distance from X to the dodecagon corner E_2 can be seen to equal x_2 . The distance from E_2 to T is x_3 .
- The circle of radius 1 centered at T is tangent to C_3B_2 at T' and hits the dodecagon at V. The distance from T' to the dodecagon corner C_3 can be seen to equal x_3 . The distance from C_3 to V is x_4 .

Here are the two regions Hansen removed:

- Hansen's first region XE_2T is bounded by the line segment XE_2 , the much shorter line segment E_2T , and the arc XT that is part of the circle of radius 1 centered at Q. He claimed this first region has an area $4 \cdot 10^{-11}$. This calculation is approximately right.
- Hansen's second region $T'C_3V$ is bounded by the line segment $T'C_3$, the much shorter line segment C_3V , and the arc T'V that is part of the circle of radius 1 centered at T. He claimed this second region had an area of $6 \cdot 10^{-18}$. This calculation is far from correct.

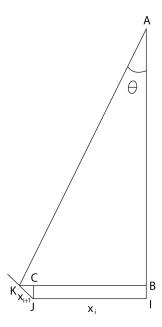
In order to find the correct areas, we first calculate x_0 and then recursively calculate the other distances x_i .

Let S be the midpoint of D_2D_3 . Recall that x_0 is the distance from S to D_2 , while x_1 is the distance from D_2 to P. We compute these distances using some elementary geometry. To do this we must keep in mind some basic facts. The length of each side of the hexagon is $\frac{1}{\sqrt{3}}$. The circle inscribed in the hexagon has radius $\frac{1}{2}$. The arc from S to P has a radius of 1 and is tangent to SD_2 at S.

To compute the distance x_0 we first need to determine the distance SD_1 . Let H be the center of the circle, as in Figure 1. The distance SD_1 is the distance HD_1 minus the distance HS, that is, $\frac{1}{\sqrt{3}} - \frac{1}{2}$. Since D_2SD_1 is a 30–60–90 triangle, the length $x_0 = D_2S$ is found $\sqrt{3}$ times SD_1 . Thus,

$$x_0 = 1 - \frac{\sqrt{3}}{2}$$
.

Knowing x_0 we can then calculate x_1 . More generally, knowing x_i we can calculate x_{i+1} . The diagram in Figure 2 can be used to solve for x_{i+1} , and further to calculate the area of the region IJK bounded by the line IJ, the line IK, and the arc through I and K that is part of a circle of radius 1 centered at A. The two regions Hansen removed, namely XE_2T and $T'C_3V$, are examples of this type of region IJK.



2. Diagram used to solve for x_i and x_{i+1}

In Figure 2 we let AI and AK be line segments of length 1. We let AI be perpendicular to JI and KB and parallel to CJ. Let x_i be the distance from J to I and x_{i+1} be the distance from K to J; BI = CJ. KI is an arc of a circle centered at A with radius 1. We set $\angle IJK = 150^{\circ}$, the interior angle of a regular dodecagon.

Note that $\angle IJC = 90^{\circ}$ and $\angle CJK = 60^{\circ}$. The triangle CJK is a right triangle with $\angle CKJ = 30^{\circ}$. We thus obtain

$$BI = CJ = \frac{x_{i+1}}{2}, \quad CK = \frac{\sqrt{3}x_{i+1}}{2}, \quad BC = x_i.$$

Then by applying the Pythagorean theorem to the triangle ABK, we see

$$1 = (AK)^2 = (BK)^2 + (AB)^2 = (CK + BC)^2 + (1 - BI)^2 = \left(\frac{\sqrt{3}x_{i+1}}{2} + x_i\right)^2 + \left(1 - \frac{x_{i+1}}{2}\right)^2.$$

This gives a quadratic relationship between x_i and x_{i+1} which can be solved for x_{i+1} :

$$x_{i+1} = \frac{1 - \sqrt{3}x_i - \sqrt{1 - 2\sqrt{3}x_i - x_i^2}}{2}.$$

To avoid rounding errors we multiply the numerator and denominator by $1-\sqrt{3}x_i+\sqrt{1-2\sqrt{3}x_i-x_i^2}$. This leaves us with

$$x_{i+1} = \frac{2x_i^2}{1 - \sqrt{3}x_i + \sqrt{1 - 2\sqrt{3}x_i - x_i^2}}.$$

Let a_i be the area of the region IJK bounded by the line segments IJ, JK and the arc centered at A going through I and K. To compute a_i , we first calculate the area of the triangle IJK and then subtract the area of the region between the line segment IK and the arc through I and K. The area of the triangle IJK is

$$\frac{IJ \cdot CJ}{2} = \frac{x_i x_{i+1}}{4}.$$

The area of the other region is

$$\frac{R^2}{2}(\theta - \sin \theta),$$

where θ is the angle $\angle BAK$ and R is the radius of the circle, namely AK, which is 1. We determine θ using the formula $\theta = 2\sin^{-1}(d/2R)$, where d is the length of the chord IK. We thus obtain

$$a_i = \frac{x_i x_{i+1}}{4} - \frac{(\theta - \sin \theta)}{2}$$

where $\theta = 2\sin^{-1}(\frac{d}{2})$ and

$$d = \sqrt{\frac{x_{i+1}^2}{4} + \left(x_i + \frac{\sqrt{3}}{2}x_{i+1}\right)^2}.$$

Table 1 shows the lengths x_i and areas a_i . We see that the area of the first region Hansen removed, namely XE_2T , is $a_2\approx 3.7507\cdot 10^{-11}$. This is close to Hansen's claim of $a_2=4\cdot 10^{-11}$. However, we see that the area of the second region Hansen removed, namely $T'C_3V$, is $a_3\approx 8.4541\cdot 10^{-21}$. This is significantly smaller than Hansen's claim of $a_3=6\cdot 10^{-18}$.

i	x_i	a_i
0	$1.339745962156 \cdot 10^{-1}$	$4.952913815765 \cdot 10^{-4}$
1	$2.413116066646 \cdot 10^{-2}$	$2.418850424555 \cdot 10^{-6}$
2	$6.080990483915 \cdot 10^{-4}$	$3.750723412843 \cdot 10^{-11}$
3	$3.701744790810 \cdot 10^{-7}$	$8.454119457933 \cdot 10^{-21}$
4	$1.370292328207 \cdot 10^{-13}$	$4.288332272809 \cdot 10^{-40}$

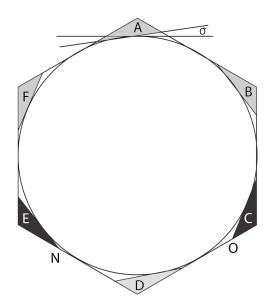
Table 1: Lengths x_i and areas a_i in Hansen's construction

3. An improved universal covering

Theorem 1. There exists a universal covering with area less than or equal to 0.8441153.

Proof. We begin with the regular hexagon in Figure 1. We shall modify this universal covering in the following ways:

Reduction 1.



3. Reduction with slant angle σ

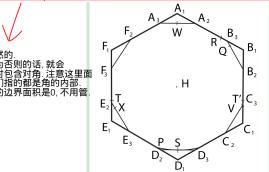
Start with a regular hexagon \mathcal{H} with a circle of diameter 1 inscribed inside of it. Pál proved that this is a universal covering [10]. Let \mathcal{H}' be a regular hexagon of the same size centered at the same point, rotated counterclockwise by an angle of $30 + \sigma$ degrees.

When $\sigma = 0$, the case considered by Pál, the intersection $\mathcal{H} \cap \mathcal{H}'$ is the largest possible regular dodecagon contained in the original hexagon \mathcal{H} . As shown in Figure 1, this dodecagon has corners $A_3A_2B_3B_2C_3C_2D_3D_2E_3E_2F_3F_2$. Pál proved that if we remove the triangles $C_1C_2C_3$ and $E_1E_2E_3$ from the hexagon \mathcal{H} , the remaining set is still a universal covering.

We generalize this to the case where σ is some other small angle—say, less than 10° . Later we shall determine that the optimal value of σ for our purposes is about 1.29° . Note that \mathcal{H} with \mathcal{H}' removed is the union of six triangles A, B, C, D, E, F shown in Figure 3. Let \mathcal{K} be \mathcal{H} with C and E removed. We claim that \mathcal{K} is a universal covering. 我们定义尽是正六边形去掉C, E

To prove this, we generalize an argument due to to Pál [10]. Imagine two parallel lines each tangent to the circle, containing two opposite edges of the rotated hexagon \mathcal{H}' . Any pair of points inside \mathcal{H} and on opposite sides of \mathcal{H}' have a distance of at least 1 from each other, since the diameter of the circle is 1. Consider the six triangles A, B, C, D, E, F. An isometric copy of any set of diameter 1 can be fit inside \mathcal{H} . Since the distance between any two triangles that are opposite each other is 1, this set cannot simultaneously have points in the interior of two opposite triangles. Thus, this set can only have a nonempty intersection with three adjacent triangles or three nonadjacent triangles.

因此被覆盖的那个长度为1的集合我们叫他set他不可能同时在A和D的内部.



说的意思是比如 面积A和面积D之间 的距离大于等于1 opposi te指的是 平行线的外侧 In the first case we can rotate the set so that the only triangles it intersects are F, A and B. In the second case we can rotate it so that the only triangles it intersects are F, B and D. In either case, it fails to intersect the triangles E and C. Thus, the set K consisting of the hexagon H with these two triangles removed is a universal covering.

Reduction 2. Recall that a **curve of constant width** is a convex subset of the plane whose **width**, defined as the perpendicular distance between two distinct parallel lines each having at least one point in common with the sets boundary but none with the set's interior, is independent of the direction of these lines. Vrećica [11] has shown that any subset of the plane with diameter 1 can be extended to a curve of constant width 1. Thus, a set will be a universal covering if it contains an isometric copy of every curve of constant width 1.

curve应该是 一个集合的意思

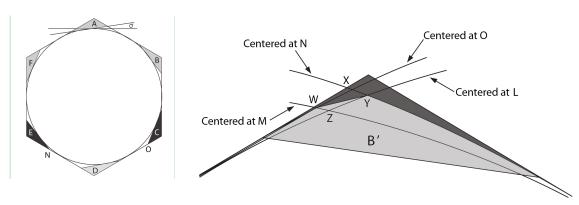
Consider any curve of constant width 1 inside K. It must touch each of the hexagon's sides at a unique point. To see this, note that if it does not touch one of the sides, it would have width less than 1. On the other hand, if it touched a side at two points, then it would have width greater than 1.

因为这个边跟他的对边距离是1,如果一条边有2个焦点,那么2个焦点跟对面边的距离必然有一个大于1.

The curve must touch the side of \mathcal{H} running from D to C at some point to the left of O, the corner of the removed triangle C. This implies that all points near A outside an arc of radius 1 centered on O can be removed from the closure of \mathcal{K} , obtaining a smaller universal covering.

Similarly, the curve must touch the side of the hexagon running from E to D somewhere to the right of N, the corner of the removed triangle E. Thus, all points near A outside an arc of radius 1 centered on N can also be removed. The remaining universal covering then has a vertex at the point X where these two arcs meet; for a closeup see Figure 4.

This reduces the area of the universal covering, but not enough to make it smaller for any nonzero value of σ than Sprague's universal covering, which is the case $\sigma = 0$. So, we must go further.



容易证明 以o点为圆心 花园是以do 这条边上任意一点花园! A相交的最大情况. 所以i 个切割是对的. 切割的部 一定无用.

三相邻

4. Reduction closeup 综上, reduction2的切割是图4中centered at o 和centered at N 2个曲线上方全切掉!

Reduction 3. For the final reduction we refer to Figures 4 and 5. Figure 4 shows a region WXY bounded by two arcs WX and XY and a straight line segment WY defined as follows:

- The arc through W and X is the arc of radius 1 centered at O, which is a point where the rotated hexagon \mathcal{H}' intersects the original hexagon \mathcal{H} .
- X is the point of intersection of the arc of radius 1 centered at O and the arc of radius 1 centered on N, which is another point where the rotated hexagon intersects the original hexagon.

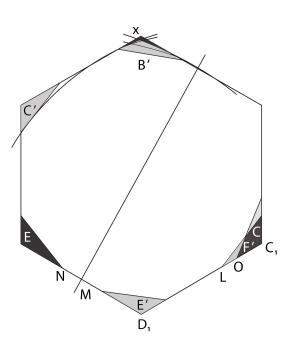
看图5.

- W is the intersection of the arc centered at O with the arc of radius 1 centered on M, the midpoint of the edge E_1D_1 of the original hexagon.
- Y is the intersection of the arc with radius 1 centered at N and the arc with radius 1 centered at L, which is another point where the rotated hexagon intersects the original hexagon.

We claim that the univeral covering considered in Reduction 2 can be further reduced by removing the region WXY for a specific angle σ that we will specify. (In fact, we could remove the whole of region WXYZ outside the arcs WZ and ZY and still be left with a set that contains an isometric copy of every curve of constant width 1. However, this set would not be convex, since it contains Wand Y but not points on the line segment between W and Y. Since a universal covering is required to be convex, we only remove the smaller region WXY, one of whose sides is the straight segment WY.)

To prove our claim, consider Figure 5. This shows an axis of symmetry of the hexagon \mathcal{H} going through a point M, the midpoint of the side D_1E_1 . When the triangles A, B, C, D, E, F are reflected reduction3 about this axis they are mapped to new triangles A', B', C', D', E', F', some of which are shown in gray in this figure.

开始证明



5. Reduction with symmetry

To prove our claim, it suffices to show that any curve of constant width 1 can be positioned inside the hexagon with the corners removed at E and C in such a way that no point enters the region WXY. To show this, we separately consider three cases. Case (1) is where the curve of constant width 1 enters the interior of triangle E'. Case (2) is where the curve does not enter the interior of E' but it does enter the interior of region C'. Case (3) is where the curve does not enter the interior of E' or C'.

Case (1): If a curve of constant width 1 enters the interior of E' then it cannot enter the interior of B' since all points inside B' are at a distance greater than 1 from points inside E'. It is therefore sufficient to show that all points in the region WXY are inside B'. This can be checked by calculation for the specific slant angle σ to be used.

Case (2): If a curve \mathcal{Q} of constant width 1 enters the interior of C' then it cannot enter the interior of the region F'. It must therefore touch the side D_1C_1 between D_1 and L. It follows that no point inside \mathcal{Q} can lie outside the arc of radius 1 centered at L within the triangle A.

Therefore, no point of \mathcal{Q} will lie inside the region WXY provided the interior of this region lies outside the arc of radius 1 centered at L. This condition holds if no points on the segment WY lie inside the arc of radius 1 centered at L. This in turn follows from $\angle WYL$ being greater than a right angle. We shall check this by a calculation for the specific angle σ we use.

Case (3): If a curve \mathcal{Q} of constant width 1 does not enter the interior of E' or C' then it can be reflected about the axis through M and the center of the hexagon, giving a curve \mathcal{Q}' also of constant width 1. This reflection maps C' to C and E' to E, so \mathcal{Q}' will lie inside the hexagon with the two corners C and E removed. Recall from Reduction 1 that \mathcal{Q} has no points in C or E. Thus \mathcal{Q}' has no points in C' or E'.

Each of the curves Q and Q' touch the side of the hexagon between N and N' at a single point. If Q touches this side between N and M then Q' will touch this side between M and N', and vice versa. Since reflections are an allowed isometry, we are free to use the curve in either position Q or Q', so we choose to use whichever of the two touches the side between M and N'. No point inside the curve will then lie both outside the arc of radius 1 centered at M and to the left of the line of reflection.

Therefore, no point of the chosen curve will lie inside the region WXY provided that no point in the interior of this region lies inside the arc of radius 1 centered at M. This condition holds if no points on the segment WY lie inside the arc of radius 1 centered at M. This is turn follows from $\angle MWY$ being greater than a right angle. We shall check this by a calculation for the specific angle σ we use.

Our computation of the area of the resulting universal covering using Java is available online [1], along with the output [2]. The output lists various choices of σ , followed by the corresponding values of the area of the would-be universal covering, together with the angles $\angle WYL$ and $\angle MWY$. We find that when $\sigma = 1.3^{\circ}$, the area of the universal covering is

0.8441153768593765

within the accuracy provided by double-precision floating-point arithmetic. For this value of σ , $\angle WYL$ is approximately 90.00593° and $\angle MWY$ is approximately 122.9277°. We could obtain a smaller area if σ were smaller, but if were too small the constraint that $\angle WYL$ exceed a right angle would be violated.

Greg Egan has checked our calculations using Mathematica with high-precision arithmetic (2000 digits working precision). He found that an angle

 $\sigma = 1.294389444703601012^{\circ}$

still obeys the necessary constraints, giving a universal covering with area

0.844115297128419059...

Egan has kindly made his Mathematica notebook available online [5], along with a printout that is readable without this software [6].

With Egan's help, we thus claim to have proved that

$$a \le 0.8441153.$$

Acknowledgements. We thank Greg Egan for catching some mistakes in this paper, and especially for verifying and improving our result. We also thank both referees for spotting and helping us fix various errors.

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DEPARTMENT OF MATHEMATICS, UNIVERSITY OF CALIFORNIA, RIVERSIDE CA 92521, USA E-mail address: baez@math.ucr.edu, karine.bagdasaryan@ucr.edu

6 Welbeck Drive Langdon Hills, Basildon, Essex, SS16 6BU, UK $E\text{-}mail\ address:\ philegibbs@gmail.com}$