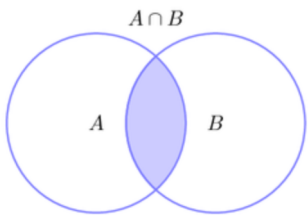
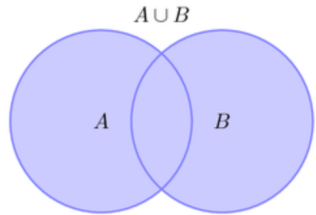
A Framework of Private Set Operations





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Shandong University

joint work with Min Zhang, Cong Zhang, Minglang Dong and Weiran Liu

Outline

Outline

Privacy Preserving Computation

国家重大战略

国务院《关于构建要素市场化的意见》
《十四五规划和 2035 年远景目标纲要》

数据是新型生产要素 ⇝ 激活数据要素潜能

数据保护需求

数据泄露事件频发，损失难以估量
三法五典出台

严格保护数据安全 ⇝ 数据流动性降低

Gartner 2021: 变革型前沿技术 ⇒ 破局的关键、数字经济的安全底座

高级密码方案

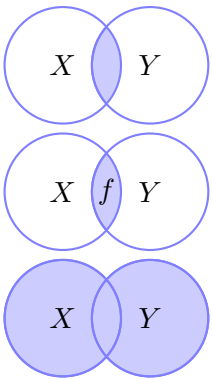
零知识证明

安全多方计算

打破数据孤岛

释放数据价值

Private Set Operations (high frequency and high valuable)



PSI $= X \cap Y$

$$\text{PCSI} = \begin{cases} |X \cap Y| & \text{cardinality} \\ |X \cap Y|, \sum_{x_i \in X \cap Y} v_i & \text{cardinality-sum} \\ f(X \cap Y) & \text{general computation} \end{cases}$$

PSU $= X \cup Y$

Wide Applications of PSO

PSI

- privacy-preserving location sharing
- private contact discovery
- DNA testing and pattern matching

PCSI

- measuring the effectiveness of online advertising

PSU

- IP blacklist and vulnerability data aggregation
- private DB supporting full join
- private-ID

SOTA of PSO

PSI has been extensively studied in the last two decades

- balanced setting: [KKRT16, CM20, RR22] achieves linear complexity, and almost as efficient as insecure hash protocol
- unbalanced setting: [CLR17, CHLR18, CMdG+21] achieves sub-linear complexity of large set

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concretely $20\times$ slower in timing and $30\times$ more communication than PSI

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PSU

- [KS05, Fri07, HN10, KRTW19, JSZ+22] have superlinear complexity
- [DC17, ZCL+23] achieve linear complexity, but not strict (comm. complexity additionally depends on statistical parameter $\lambda \approx 40$)

concretely $20\times$ slower in timing and $25\times$ more communication than PSI

Motivation

Different approaches are used for different private set operations $\rightsquigarrow$ require much more engineering effort and maintaining cost

- Goal: a unified framework of PSO

[1] [GMR+21] presented a PSO framework from permuted characteristic. However, its oblivious shuffle functionality is not necessary for PSO, and incurs superlinear complexity.

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There exists huge efficiency gap between PSI and other PSO protocols

- Goal: efficient instantiations to close the gap[1]

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After $\approx$ 40 years, DH-PSI [Mea86] is still the most easily understood and implemented one among numerous PSI protocols. Surprisingly, no counterpart is known in the PSU setting yet. Existing protocols are very complicated.

- Goal: build DDH-based PSU protocol as simple as DH-PSI

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- Goal: build DDH-based PSU protocol as simple as DH-PSI

Is there a central building block that enables a unified framework for PSO?
How to give instantiations with optimal asymptotic complexity and good concrete efficiency?
Can the DDH assumption strike back with efficient PSU protocol?

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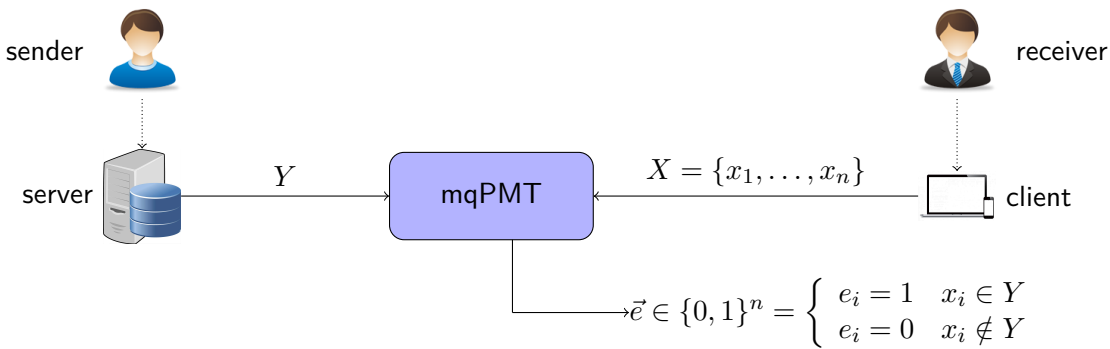
Outline

Start Point: multi-query Private Membership Test (mqPMT) underlying PSI



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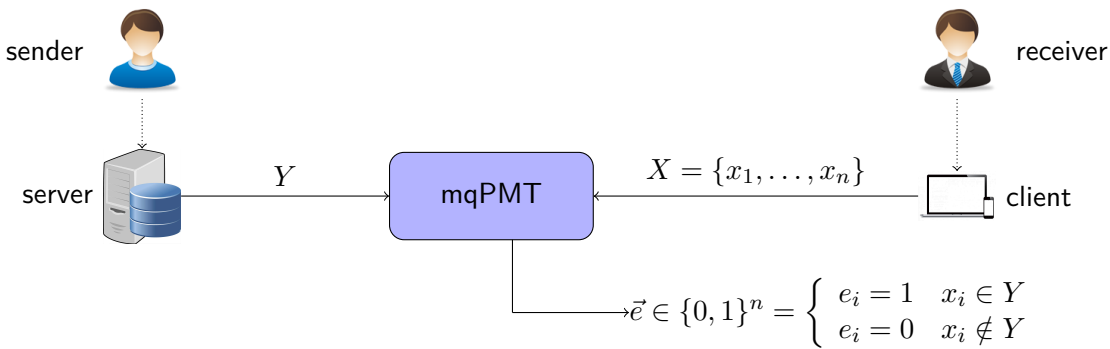
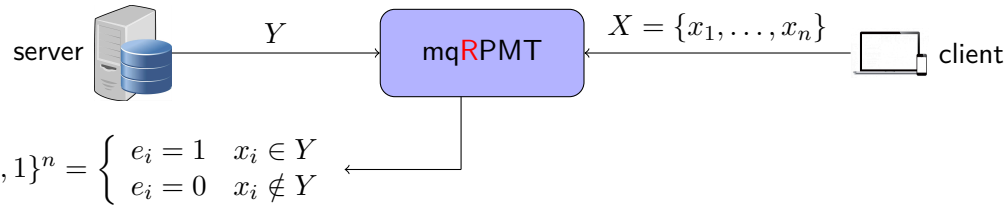


- Problem: the client learns both x_i and e_i, a.k.a. the intersection $\rightsquigarrow$ not suitable for protocols that should hide intersection, such as PCSI and PSU.

The core protocol: multi-query Reverse Private Membership Test (mqRPMT)



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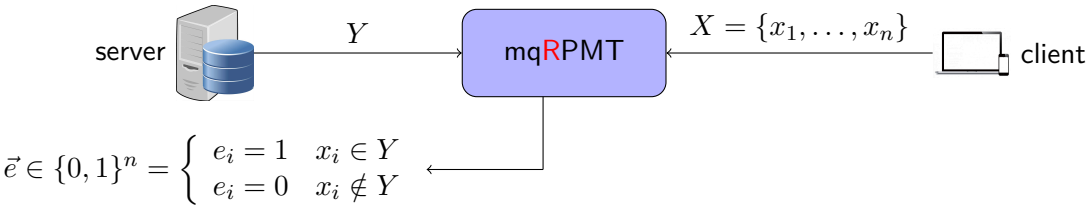
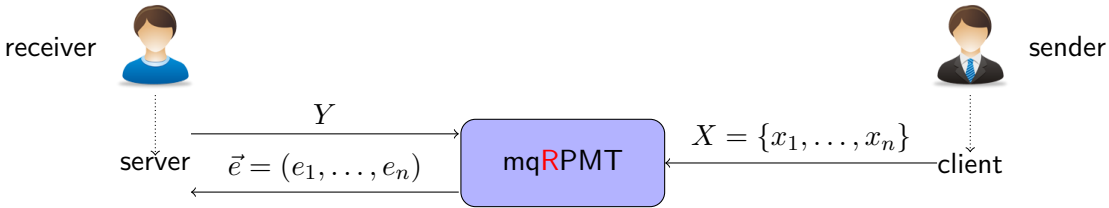
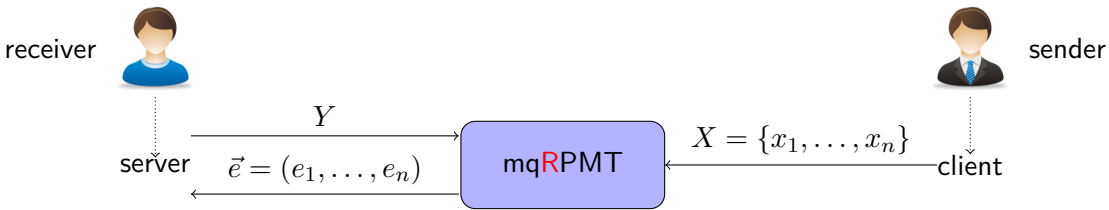


- The server learns e_i, while the client learns x_i, a.k.a. the information of intersection is shared between the two parties $\rightsquigarrow$ suitable for all PSO protocols

PSO from mqRPMT

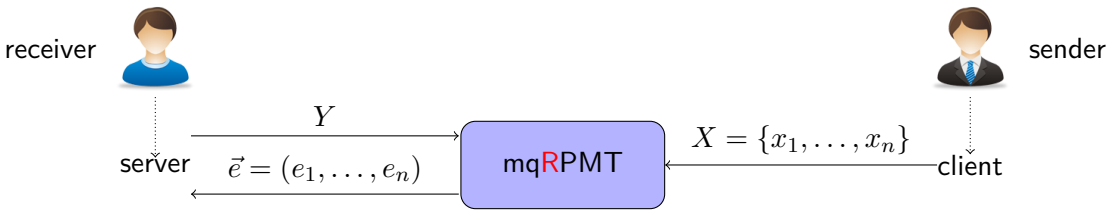


PSO from mqRPMT

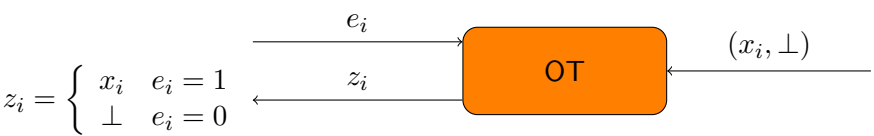


directly yields PSI-card: $|X \cap Y|$ is the Hamming weight of $\vec{e}$

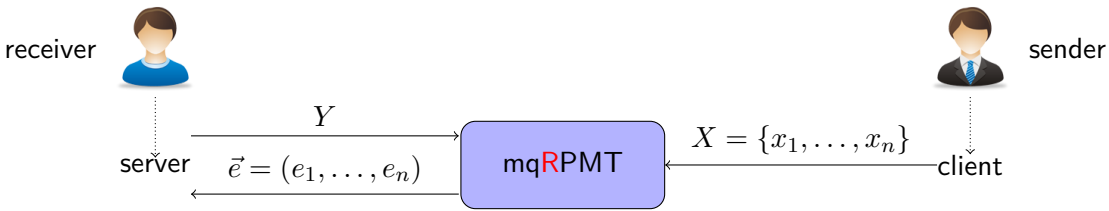
PSO from mqRPMT



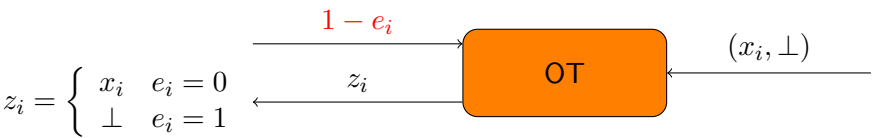
yields PSI coupled with OT: receiver obtains $X \cap Y$



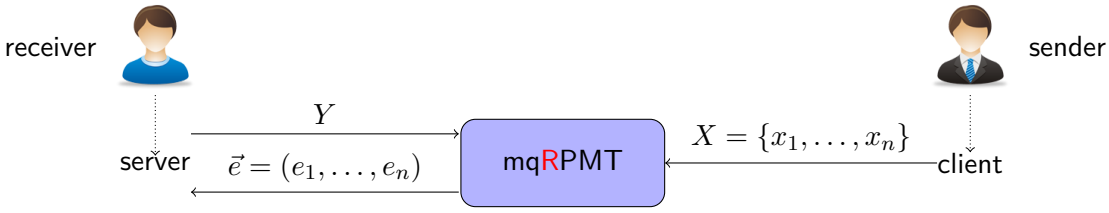
PSO from mqRPMT



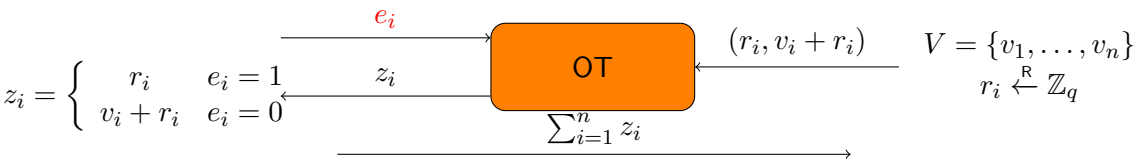
yields PSU coupled with OT (flipping $\vec{e}$): receiver obtains $X - Y$



PSO from mqRPMT



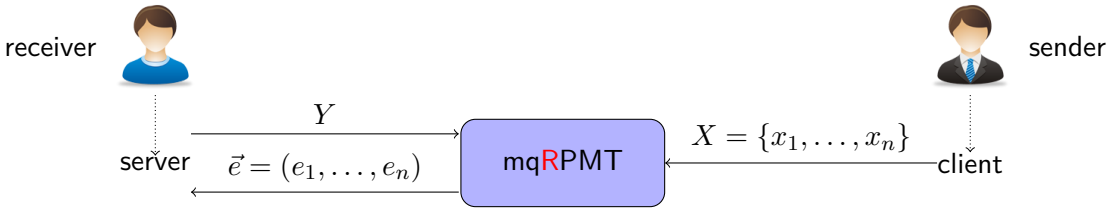
yields PSI-card-sum coupled with OT and masking trick



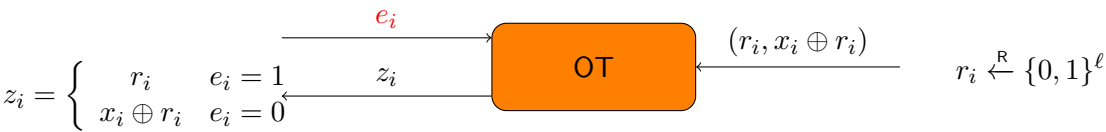
receiver obtains $|X \cap Y|$

sender obtains $\sum_{x_i \in Y} v_i = \sum_{i=1}^{n} z_i - \sum_{i=1}^{n} r_i$

PSO from mqRPMT



yields PSI-card-secret-share coupled with OT and masking trick



receiver obtains $|X \cap Y|$ and z_i

sender has $x_i \oplus r_i$

Outline

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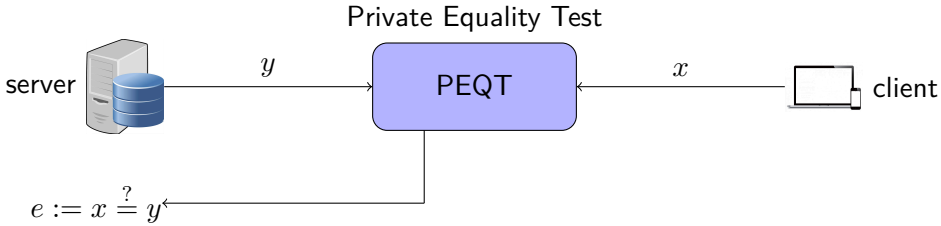
Starting Point: PEQT



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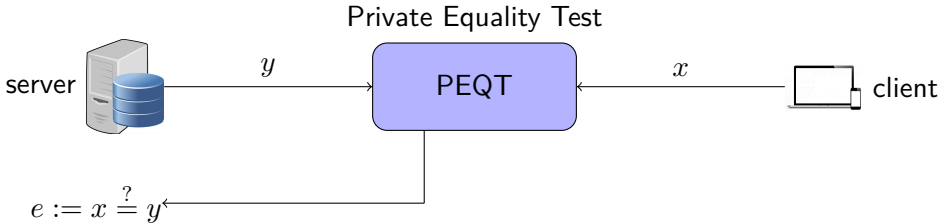
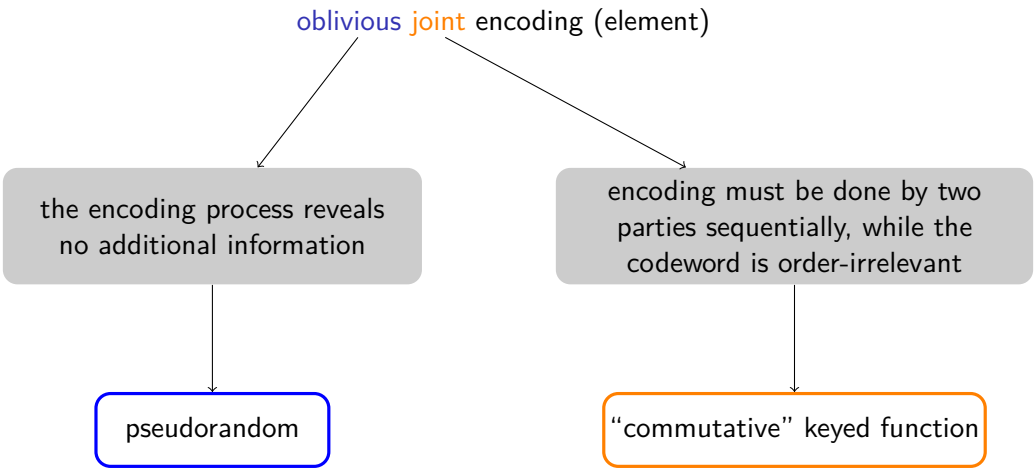


Observation: PEQT is not only an extreme case of mqPMT, but also an extreme case of mqRPMT

Goal: build PEQT amenable to extension:

$$y \rightsquigarrow Y = \{y_1, \ldots, y_m\},\ x \rightsquigarrow X = \{x_1, \ldots, x_n\},\ e \rightsquigarrow \vec{e} = (e_1, \ldots, e_n)$$

High-level Idea



Commutative Weak PRF

We first formally define two standard properties for keyed functions.

Composable. For a family of keyed functions $F : K \times D \to R$, F is 2-composable if $R \subseteq D$ (special case $R = D$) $\rightsquigarrow F_{k_1}(F_{k_2}(\cdot))$ is well-defined.

Commutative. A family of composable keyed functions is commutative if:

$$\forall k_1, k_2 \in K, \forall x \in D : F_{k_1}(F_{k_2}(x)) = F_{k_2}(F_{k_1}(x))$$

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Definition 1 (Commutative Weak PRF)

$F : K \times D \to D$ is cwPRF if it satisfies weak pseudorandomness ($k \xleftarrow{R} K$, $x \xleftarrow{R} X$) and commutative property simultaneously. When F is a permutation, we say F is cwPRP.

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Why merely weak pseudorandomness?

Standard pseudorandomness denies commutativity. Consider the following attack:

- $\mathcal{A}$ picks $k' \xleftarrow{R} K$, $x \xleftarrow{R} D$, queries the real-or-random oracle at point $F_{k'}(x)$ and x, receiving y' and y. $\mathcal{A}$ then outputs '1' iff $F_{k'}(y) = y'$.

Construction of cwPRF

Construction (DDH-based cwPRF)

- Setup(1^κ): runs GroupGen$(1^\kappa) \to (\mathbb{G}, g, p)$, output $pp = (\mathbb{G}, g, p)$ which defines
$$F : \mathbb{Z}_p \times D \to \mathbb{G} \text{ as } F_k(x) := x^k$$
- KeyGen(pp): outputs $k \xleftarrow{R} \mathbb{Z}_p$.
- Eval(k, x): on input $k \in \mathbb{Z}_p$ and $x \in D$, outputs x^k.

DDH assumption $\Rightarrow$ weak pseudorandomness

Commutativity: $\forall k_1, k_2 \in K$ and $\forall x \in D$: $F_{k_1}(F_{k_2}(x)) = x^{k_1 k_2} = F_{k_2}(F_{k_1}(x))$

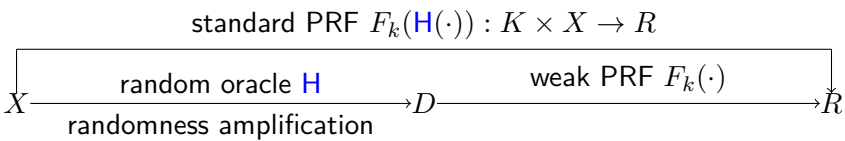
cwPRF is the "right" cryptographic abstraction of the classic DH function.

Randomness Amplification

But what we need for mqRPMT is standard pseudorandomness.

Solution: hash-then-evaluate

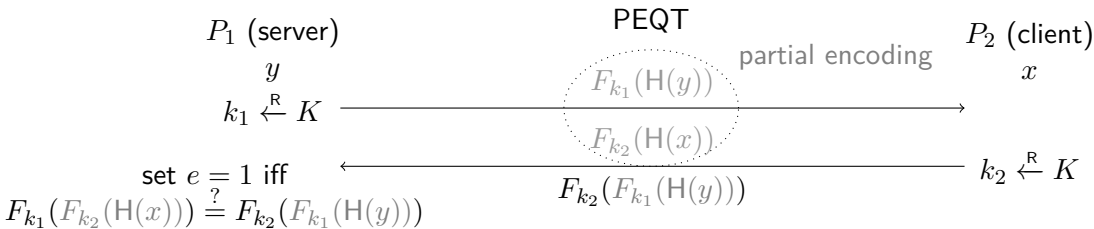
- Domain extension: handle arbitrary domain $X = \{0,1\}^*$
- Randomness amplification: weak $\rightsquigarrow$ standard



Commutativity still holds w.r.t. H

$$F_{k_1}(F_{k_2}(\mathsf{H}(x))) = F_{k_2}(F_{k_1}(\mathsf{H}(x)))$$

mqRPMT from cwPRF



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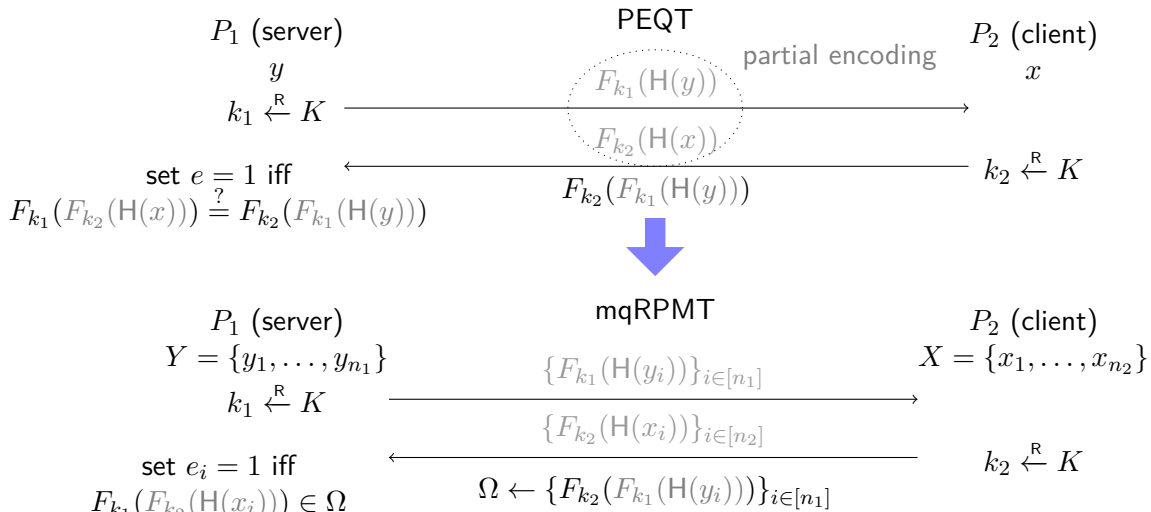

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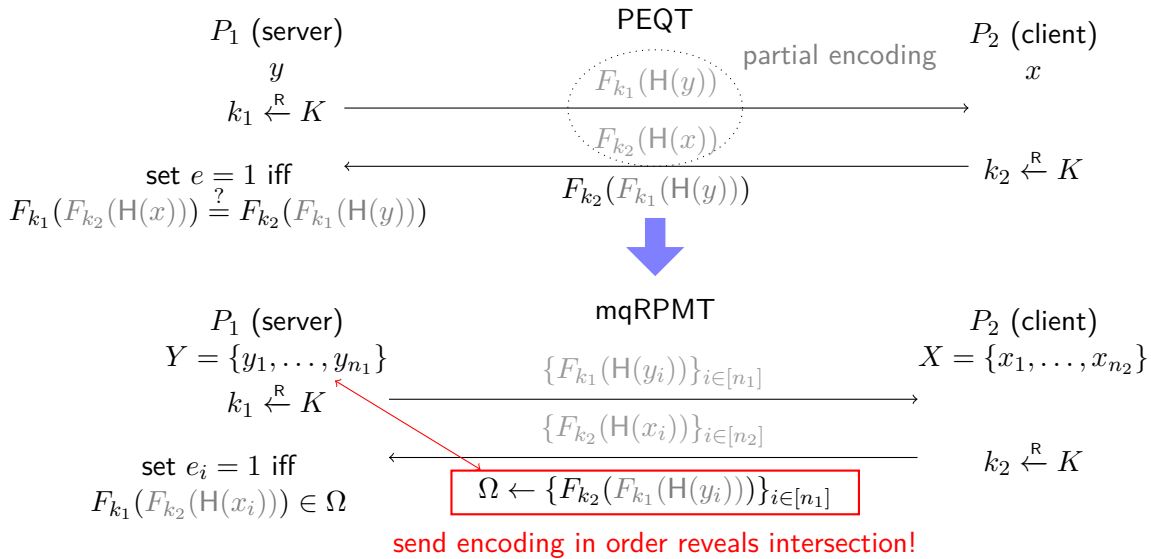

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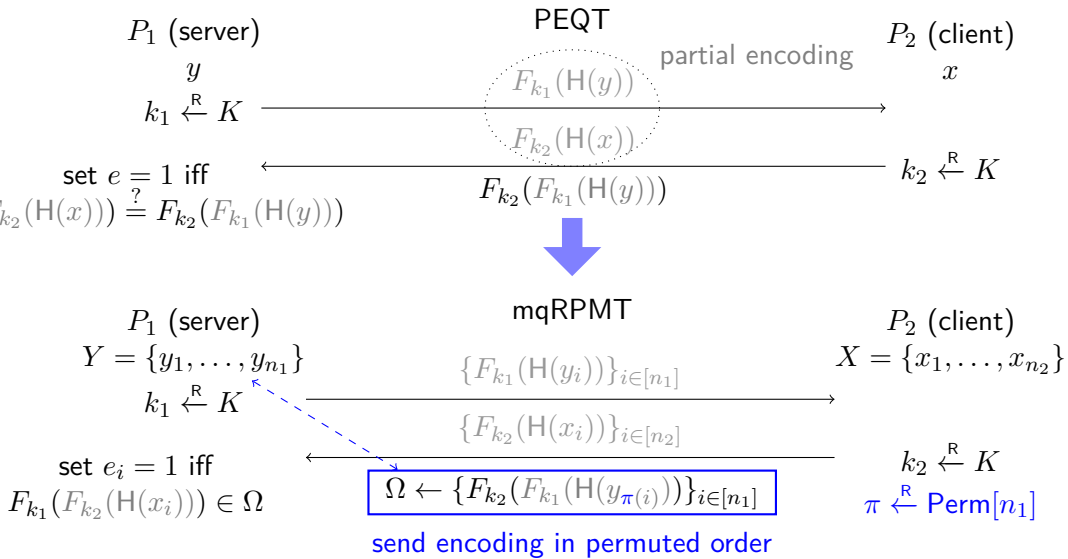

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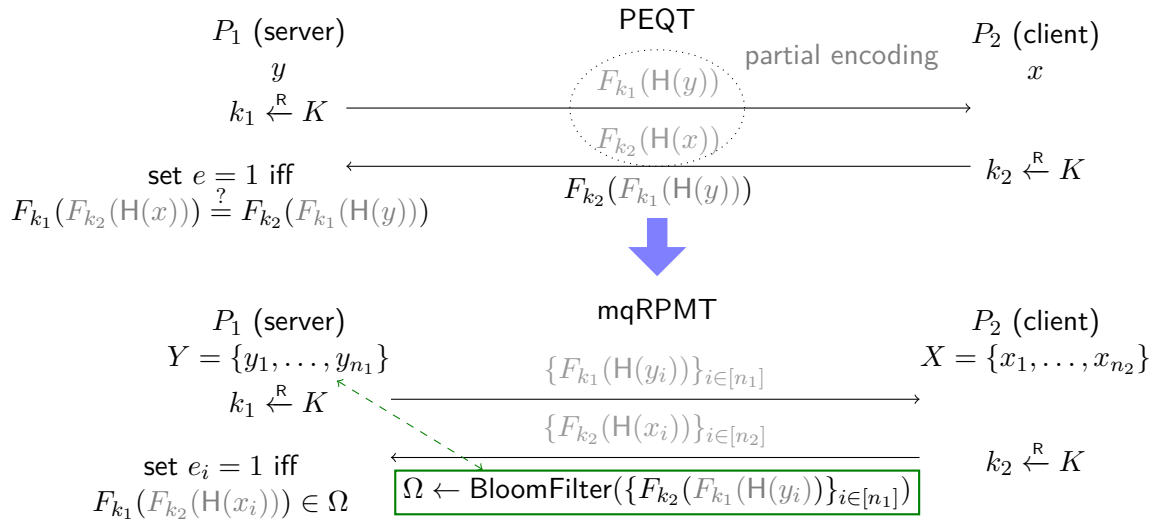


more comm. efficient: admit small false positive

Complexity Analysis

Consider the balanced setting: $n_1 = n_2 = n$

Table: Complexity of cwPRF-based mqRPMT.

Computation	$4n\times$ $F_k(\cdot)$ + $2n \times$ H$(\cdot)$ hash-to-domain
Communication	$3n \times \lvert D\rvert$ or $2n \times \lvert D\rvert + c \cdot \lambda \ll \lvert D\rvert$

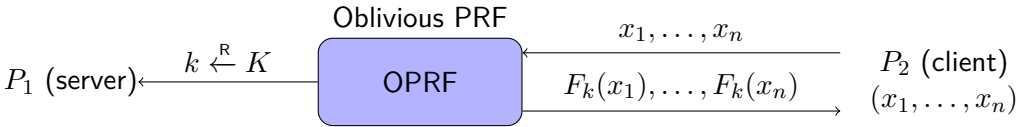
cwPRF-based mqRPMT is optimal in the sense that both computation and communication complexities are strictly linear in n

Instantiating the PSO framework with cwPRF-based mqRPMT, DDH assumption strikes back with the first strictly linear PSU protocol

incredibly simple and efficient

Outline

Starting Point: mqPMT/PSI from OPRF



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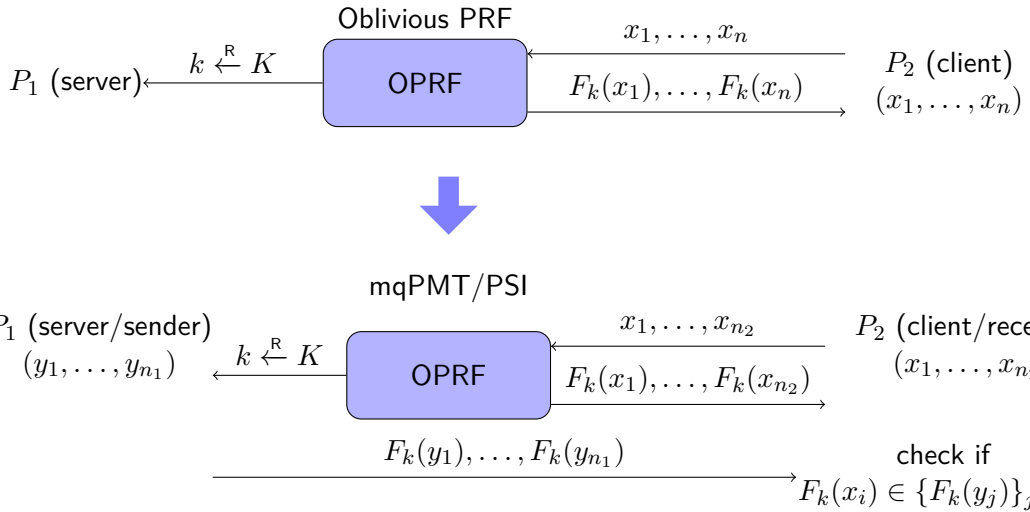

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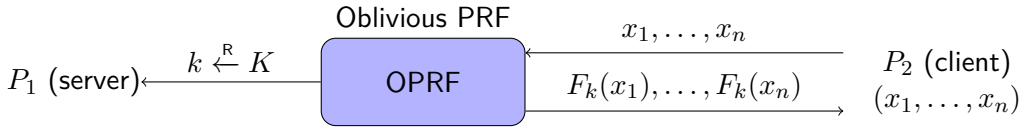

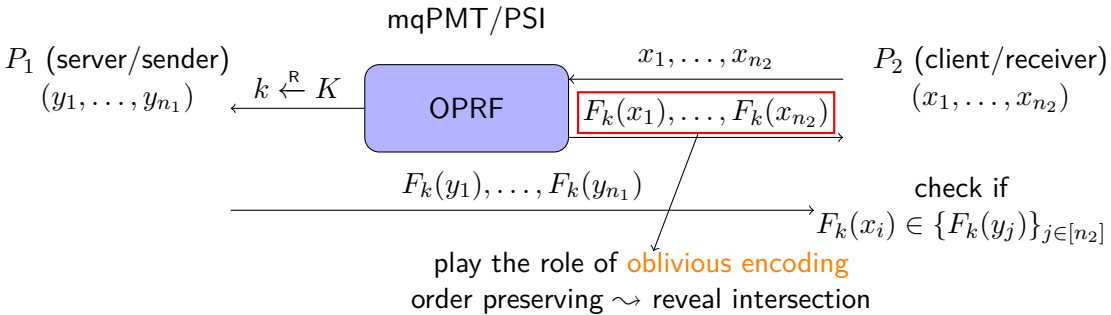

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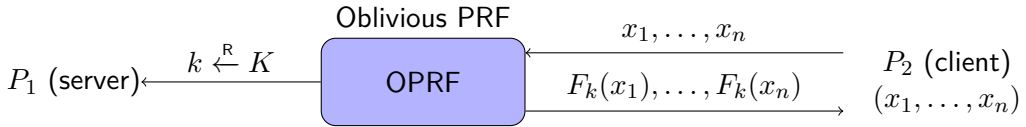
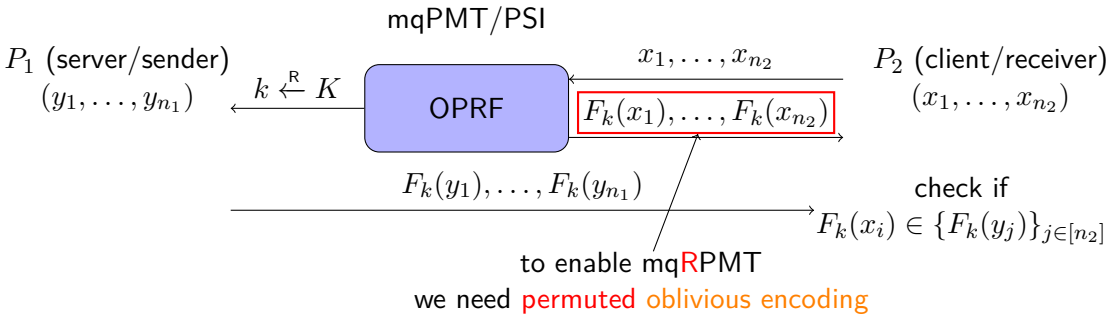
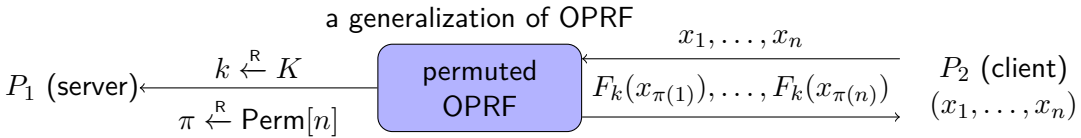
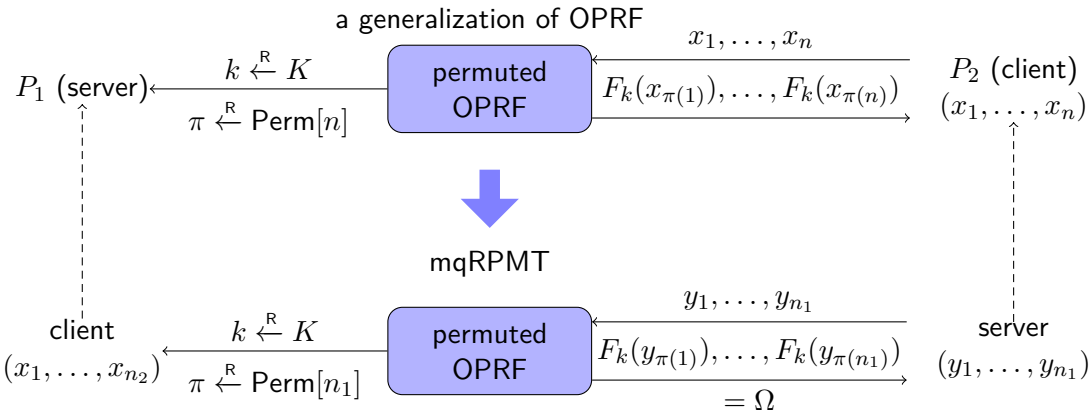




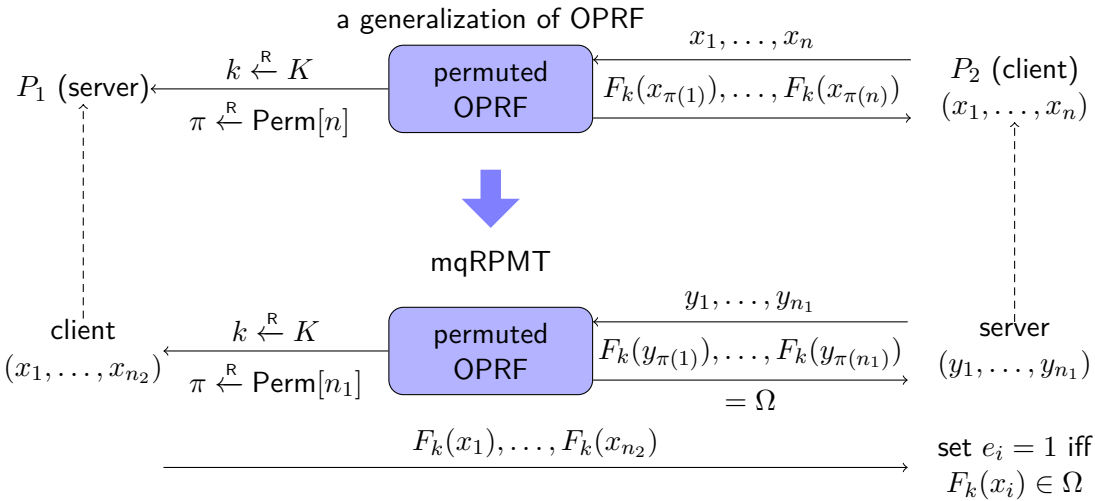
mqRPMT from Permuted OPRF



mqRPMT from Permuted OPRF

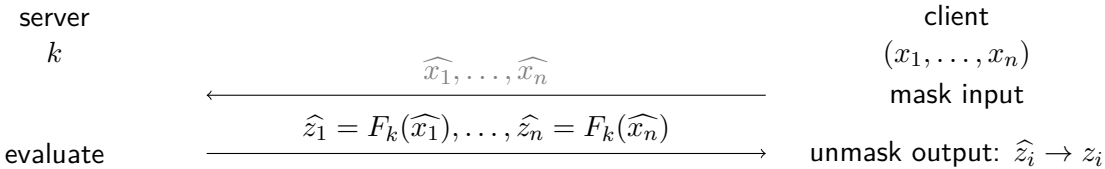


mqRPMT from Permuted OPRF



Build Permuted OPRF from cwPRP

A common approach to build OPRF is "mask-then-unmask" via homomorphism



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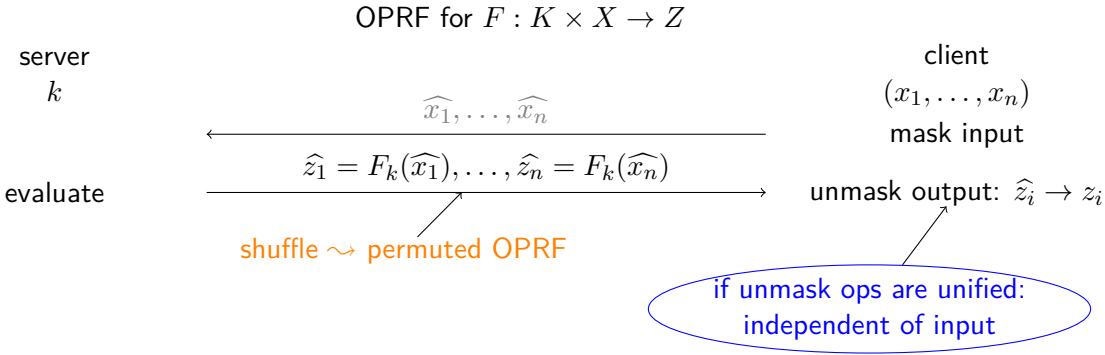

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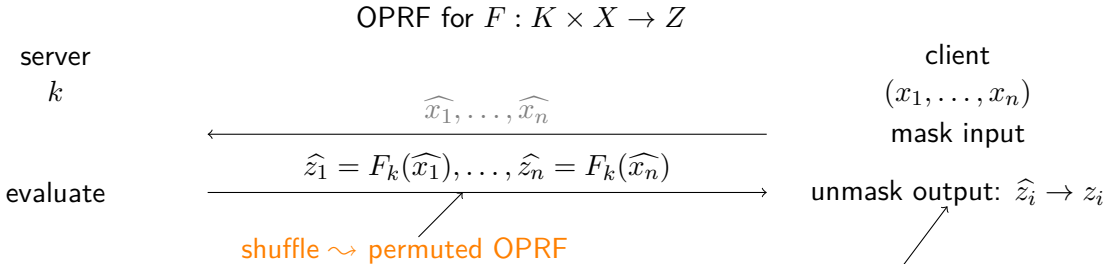


cwPRP enables simplest unified mask-then-unmask

mask: $\hat{x} \leftarrow F_s(x)$

evaluate: $\hat{z} \leftarrow F_k(\hat{x})$

unmask: $z \leftarrow F_s^{-1}(F_k(\hat{x}) = F_k(F_s^{-1}(\hat{x})) = F_k(x)$

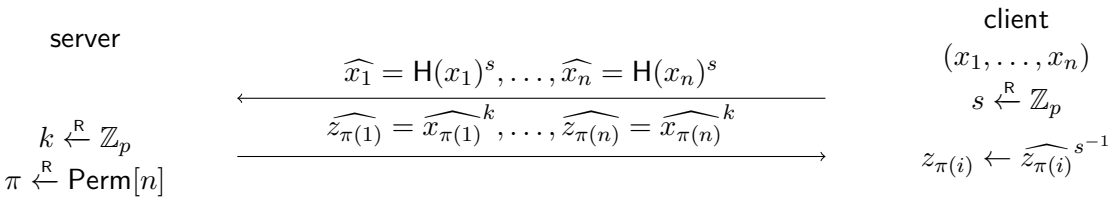
Permuted OPRF from DDH-based cwPRP

Observe that the DDH-based cwPRF is acturally a cwPRP $F : \mathbb{Z}_p \times \mathbb{G} \to \mathbb{G}$.

- combine $\mathsf{H} : \{0,1\}^* \to \mathbb{G} \Rightarrow$ permuted OPRF protocol for $G : \mathbb{Z}_p \times \{0,1\}^* \to \mathbb{G}$ defined as $G_k(x) = F_k(\mathsf{H}(x))$.

---

pOPRF for $G_k(x) = F_k(\mathsf{H}(x))$

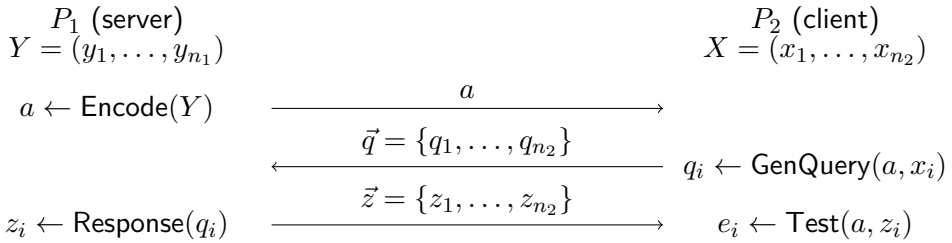


Outline

Sigma-mqPMT

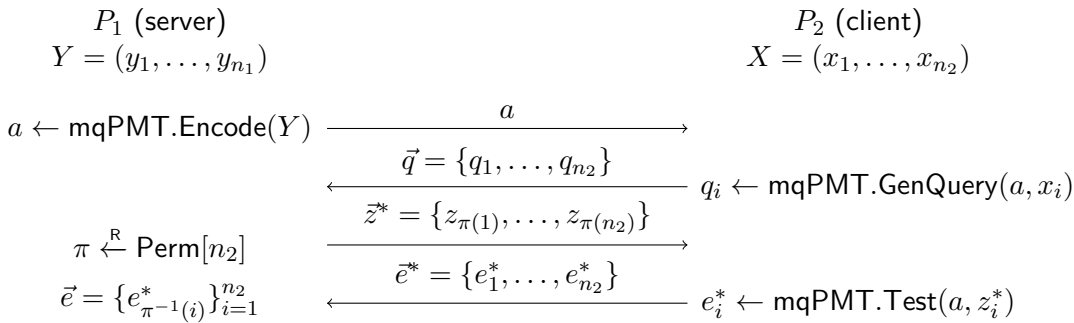
Given the efficiency gap between PSI and other PSO protocols, it is intriguing to study the connection between mqPMT and mqRPMT.

- Towards this goal, we first abstract a category of mqPMT called Sigma-mqPMT.



- **Reusable:** a (best interpreted as encoding of Y) can be safely reused.
- **Context-independent:** q_i is only related to a, x_i under test and P_2's randomness.
- **Stateless test:** Test algorithm can work without knowing (x_i, q_i).

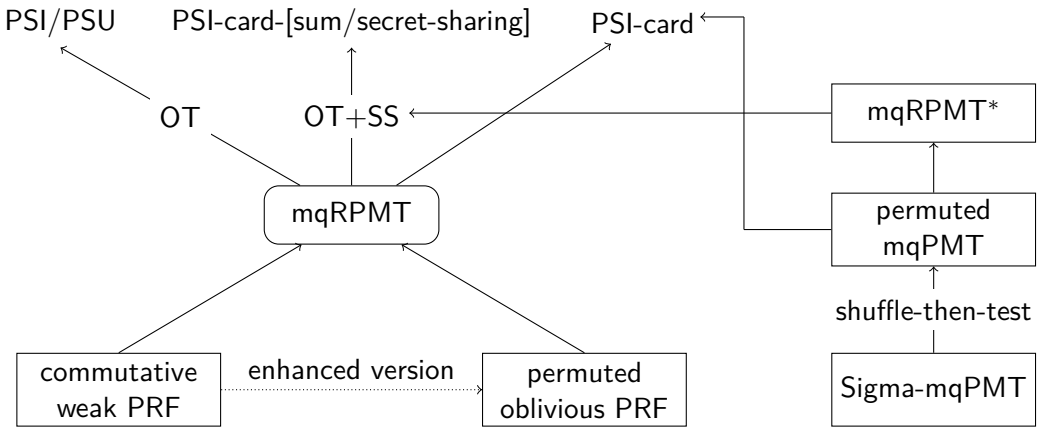
mqRPMT* from Sigma-mqPMT



Via the "permute-then-test" approach, we can tweak Sigma-mqPMT to mqRPMT* (additionally reveal intersection size to client).

- translate a category of PSI protocols (such as [Mea86, FIPR05, CLR17]) to other PSO protocols (allowing both parties learn the intersection size).
- make the initial step towards establishing the connection between mqRPMT and mqPMT.

Summary of Main Results



Outline

Cryptographic Engineering Matters

We implement our PSO framework via the following vein

$$\text{EC groups} \quad \text{DDH-based cwPRF} \rightsquigarrow \text{mqRPMT} \rightsquigarrow \text{PSO framework}$$

Cryptographic Engineering Matters

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$$\text{EC groups} \quad \text{DDH-based cwPRF} \rightsquigarrow \text{mqRPMT} \rightsquigarrow \text{PSO framework}$$

1. `NIST P-256` ⧫ ▼ (also known as `secp256r1` and `prime256v1`)
   - hash-to-point operation is expensive $\approx$ non-fixed Exp
   - point compression halving communication at cost
     $\rightsquigarrow$ point decompression is expensive $\approx$ non-fixed Exp

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   - hash-to-point operation is expensive $\approx$ non-fixed Exp
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2. `Curve25519` ★ (*de facto* alternative of `NIST P-256`)
   - numerous merits: no backdoor, fast Exp, immunity against side-channel attacks
   - allow fast Exp in compressed form $\rightsquigarrow$ halving comm. without decompression
   - any 32-byte bit string (interpreting as X-coordinate) can be ambiguously identified as a valid point $\rightsquigarrow$ hash-to-point operation is almost free

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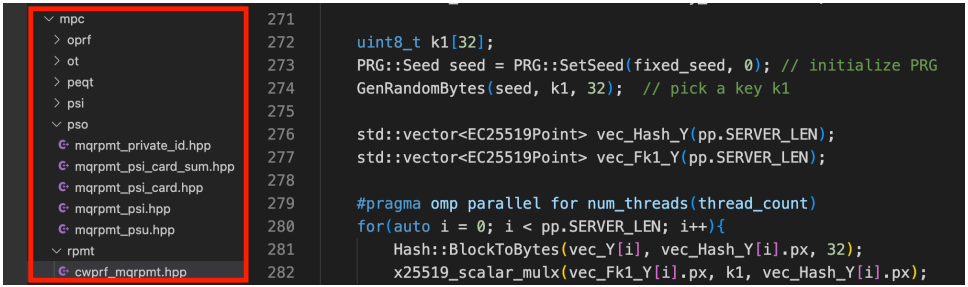
For the first time, `Curve25519` fully unleashes its power in PSO area.
Correct the prejudice that "public-key operations are expensive":

- By leveraging optimized implementation, their performances are comparable with symmetric-key operations

Implementation Features

- Modular design: admit flexible combination to support various scenarios
- Minimum dependency: only require OpenSSL and OpenMP
- Multi-platforms: run smoothly on Linux and MacOS
- Rich functionality: support all PSO operations
- Highly parallelizable: scalable $\rightsquigarrow$ support large-scale applications



Implementation Details

Dev/Test environment	Other Parameters
CPU = Intel i7 2.50 GHZ	$\kappa = 128$, $\lambda = 40$
Physical core = 8	item length= 128 bits
RAM = 8GB	set sizes= $\{2^{12}, 2^{16}, 2^{20}\}$
OS = Ubuntu 20.04	LAN = 10Gbps, WAN = 50Mbps, RTT= 80ms

Protocols:

- mqRPMT, PSI, PSI-card, PSI-card-sum, PSU, Private-ID

Test items:

- Functionality
- Computation cost: total running time
- Communication cost: sum of two parties

Core protocol: mqRPMT

Protocol	T	Running time (s)						Commu. (MB)		
		LAN			WAN			total		
		2^{12}	2^{16}	2^{20}	2^{12}	2^{16}	2^{20}	2^{12}	2^{16}	2^{20}
mqRPMT$^{\blacklozenge}$	1	0.50	7.20	114.16	1.39	9.68	136.27	0.52	8.35	133.6
	2	0.31	3.89	62.09	1.14	6.54	86.60			
	4	0.22	2.37	40.41	1.11	5.08	62.77			
Speedup		1.6-2.3 ×	1.9-3.0 ×	1.8-2.8 ×	1.2-1.3 ×	1.5-1.9 ×	1.6-2.2 ×	–	–	–
mqRPMT$^{\blacktriangledown}$	1	0.50	8.00	128.00	1.35	10.15	141.52	0.27	4.35	69.6
	2	0.32	5.05	80.69	1.18	7.11	94.19			
	4	0.23	3.54	58.40	1.08	5.54	71.26			
Speedup		1.6-2.2 ×	1.6-2.3 ×	1.6-2.2 ×	1.1-1.3×	1.4-1.8 ×	1.5-2 ×	–	–	–
mqRPMT$^{\bigstar}$	1	0.26	3.51	54.85	0.81	5.41	68.68	0.26	4.23	67.66
	2	0.15	1.79	28.24	0.75	3.83	41.38			
	4	0.10	1.07	15.32	0.72	3.09	28.31			
Speedup		1.7-2.6 ×	2.0-3.3 ×	1.9-3.6 ×	1.1-1.1 ×	1.4-1.8 ×	1.7-2.4 ×	–	–	–

strict linear complexity & high parallelism

2^{20} scale: #time < 15s using 4 threads on laptop, #communication < 70M

PSI: Performance and Comparison

PSI	Running time (s)						Comm. (MB)		
	LAN			WAN			total		
	2^{12}	2^{16}	2^{20}	2^{12}	2^{16}	2^{20}	2^{12}	2^{16}	2^{20}
[PRTY19]*	5.51	88.64	1418.20	5.82	90.79	1498.67	0.30	4.74	76.60
Our PSI$^\blacklozenge$	**0.50**	**7.24**	**114.66**	**1.71**	**10.50**	**142.45**	**0.67**	**10.38**	**165.77**
Our PSI$^\blacktriangledown$	**0.55**	**8.04**	**128.18**	**1.73**	**11.02**	**148.18**	**0.41**	**6.38**	**101.63**
Our PSI*	**0.29**	**3.56**	**55.11**	**1.19**	**6.38**	**75.56**	**0.40**	**6.25**	**99.71**
DH-PSI*	**0.22**	**3.39**	**54.79**	**0.92**	**5.57**	**69.31**	**0.28**	**4.57**	**74.1**

compared to existing DH-PSI implementation: # time speeds up 4.9-25.7×

PSI	Running time (ms)						Comm. (KB)		
	LAN			WAN			total		
	2^8	2^9	2^{10}	2^8	2^9	2^{10}	2^8	2^9	2^{10}
[RT21]*	50.0	71.0	147.3	224.1	260.2	457.9	17.9	34.1	66.3
Our PSI*	**41.9**	**69.5**	**99.3**	**577.0**	**582.9**	**646.1**	**38.6**	**63.5**	**113.3**
DH-PSI*	**16.49**	**31.80**	**56.91**	**210.42**	**227.33**	**252.32**	**18.48**	**36.68**	**72.8**

achieve the fastest speed in small set setting ($< 2^{10}$)

PSI-card: Performance and Comparison

PSI-card	Running time (s)						Comm. (MB)		
	LAN			WAN			total		
	2^{12}	2^{16}	2^{20}	2^{12}	2^{16}	2^{20}	2^{12}	2^{16}	2^{20}
[GMR+21]	1.00	8.41	126.01	8.60	27.46	323.52	2.93	55.49	1030
Our PSI-card$^{\blacklozenge}$	**0.49**	**7.20**	**114.31**	**1.30**	**9.68**	**136.06**	**0.52**	**8.36**	**133.71**
Our PSI-card$^{\blacktriangledown}$	**0.53**	**8.00**	**128.00**	**1.35**	**10.16**	**141.31**	**0.27**	**4.35**	**69.6**
Our PSI-card$^{\bigstar}$	**0.27**	**3.51**	**54.89**	**0.82**	**5.42**	**68.31**	**0.26**	**4.23**	**67.70**

compared to the SOTA

time speeds up 2.3-10.5×, # communication reduces 11.3-15.2×

PSI-card-sum: Performance and Comparison

PSI-card-sum	Running time (s)						Comm. (MB)		
	LAN			WAN			total		
	2^{12}	2^{16}	2^{20}	2^{12}	2^{16}	2^{20}	2^{12}	2^{16}	2^{20}
[IKN+20]▼ (deployed)	23.64	176.34	–	30.10	186.29	–	2.72	43.24	–
Our PSI-card-sum♦	**0.51**	**7.22**	**113.66**	**1.46**	**9.68**	**136.27**	**0.64**	**9.89**	**157.80**
Our PSI-card-sum▼	**0.57**	**8.12**	**129.66**	**1.94**	**11.83**	**157.66**	**0.38**	**5.87**	**93.74**
Our PSI-card-sum★	**0.31**	**3.73**	**57.44**	**1.36**	**6.53**	**76.16**	**0.37**	**5.75**	**91.70**

Google google / private-join-and-compute Public

compared to the SOTA

time speeds up 22.1-76.3×, # communication reduces 7.4-7.5×

PSU: Performance and Comparison

PSU	Running time (s)						Comm. (MB)		
	LAN			WAN			total		
	2^{12}	2^{16}	2^{20}	2^{12}	2^{16}	2^{20}	2^{12}	2^{16}	2^{20}
[GMR+21]	1.16	10.06	151.34	10.34	38.52	349.43	3.85	67.38	1155
[ZCL+23]$^\blacklozenge$	4.87	12.19	141.38	5.78	15.75	182.88	1.35	21.41	342.38
[ZCL+23]$^\blacktriangledown$	5.10	15.13	187.29	5.82	17.37	210.06	0.77	12.20	195.17
[JSZ+22]	2.29	8.50	516.04	5.33	27.00	736.30	3.59	70.37	1341.55
Our PSU$^\blacklozenge$	**0.52**	**7.27**	**114.44**	**1.70**	**10.56**	**143.29**	**0.68**	**10.38**	**165.77**
Our PSU$^\blacktriangledown$	**0.57**	**8.04**	**128.20**	**1.76**	**10.92**	**148.15**	**0.41**	**6.38**	**101.63**
Our PSU$^\bigstar$	**0.30**	**3.55**	**55.48**	**1.19**	**6.38**	**74.96**	**0.40**	**6.25**	**99.71**

compared to the SOTA: first achieves strict linear complexity

\# time speeds up 2.4-17×, # communication reduces 2×

Private-ID: Performance and Comparison

Private-ID	Running time (s)						Comm. (MB)		
	LAN			WAN			total		
	2^{12}	2^{16}	2^{20}	2^{12}	2^{16}	2^{20}	2^{12}	2^{16}	2^{20}
[GMR+21]	1.65	11.023	158.76	13.82	43.00	385.12	4.43	76.57	1293
[BKM+20]*	2.21	37.56	671.75	7.98	46.97	710.94	1.00	15.97	226.70
Our Private-ID$^\blacklozenge$	**0.77**	**8.40**	**114.45**	**2.91**	**13.62**	**148.48**	**1.46**	**20.52**	**330.40**
Our Private-ID$^\blacktriangledown$	**0.89**	**9.57**	**146.73**	**3.07**	**14.53**	**186.98**	**1.13**	**16.43**	**266.31**
Our Private-ID*	**0.61**	**5.11**	**71.57**	**2.83**	**10.06**	**114.66**	**1.13**	**16.31**	**265.15**

compared to the SOTA

time speeds up 2.7-4.9×, # communication is slightly larger

Outline

Summary of This Work

Unified PSO framework from mqRPMT

- show mqRPMT is complete for all PSO protocols
- greatly reduce the deployment and maintaining costs of PSO

Generic construction of mqRPMT

- cwPRF: DDH assumption is truly a golden goose
- permuted OPRF: make the concept of OPRF more useful
- mqRPMT* from Sigma-mqPMT: a initial step towards the connection to mqPMT

Efficient implementation

- identify expensive ECC operations in cheap disguise
- find the perfect match: `Curve25519`

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Likewise, we do not take shame in the simplicity of our construction :-)

Thanks for Your Attention!

Any Questions?

http://eprint.iacr.org/2022/652
https://yuchen1024.github.io

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