Carnegie Mellon University

Database Systems

Two-Phase Locking

15-445/645 FALL 2024 » PROF. ANDY PAVLO

ADMINISTRIVIA

Project #3 is due Sunday Nov 17th @ 11:59pm

→ **Recitation: Monday Nov 4th @ 8:00pm (Zoom)**

UPCOMING DATABASE TALKS

Synnada (DB Seminar)
→ Monday Nov 4th @ 4:30pm
→ Zoom

InfluxDB (DB Seminar)
→ Monday Nov 11th @ 4:30pm
→ Zoom

GlareDB (DB Seminar)
→ Monday Nov 18th @ 4:30pm
→ Zoom

LAST CLASS

Conflict Serializable

- → Verify using either the “swapping” method or dependency graphs.
- → Any DBMS that says that they support “serializable” isolation does this.

View Serializable

- → No efficient way to verify.
- → No DBMS that supports this.

OBSERVATION

We need a way to guarantee that all execution schedules are correct (i.e., serializable) without knowing the entire schedule ahead of time.

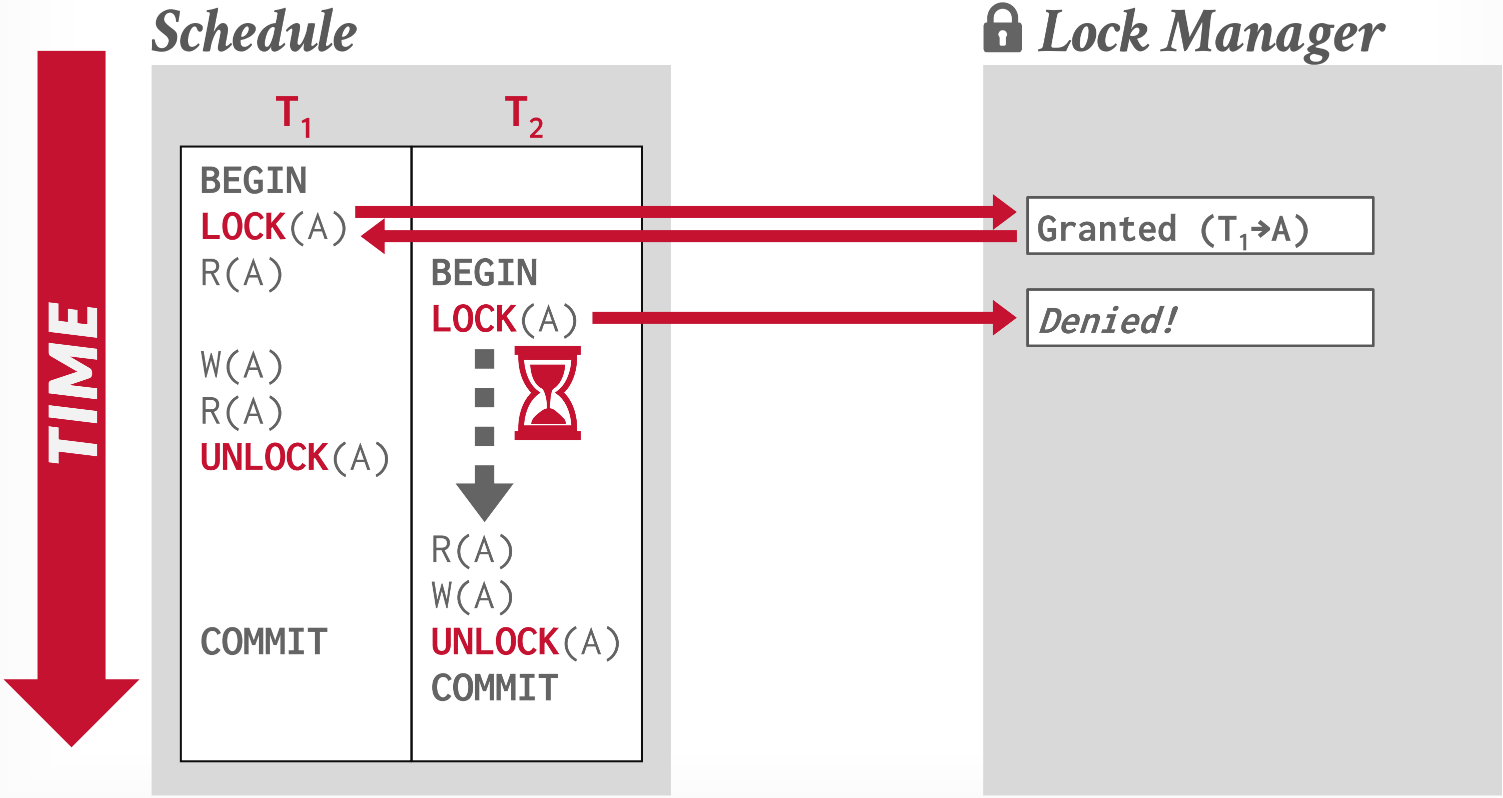
Solution: Use **locks** to protect database objects.

LOCKS VS. LATCHES

	Locks	*Latches*
Separate…	Transactions	Workers (threads, processes)
Protect…	Database Contents	In-Memory Data Structures
During…	Entire Transactions	Critical Sections
Modes…	Shared, Exclusive, Update, Intention	Read, Write
Deadlock	Detection & Resolution	Avoidance
…by…	Waits-for, Timeout, Aborts	Coding Discipline
Kept in…	Lock Manager 哈希表，跟踪锁的使用信息	Protected Data Structure

Source: Goetz Graefe

EXECUTING WITH LOCKS

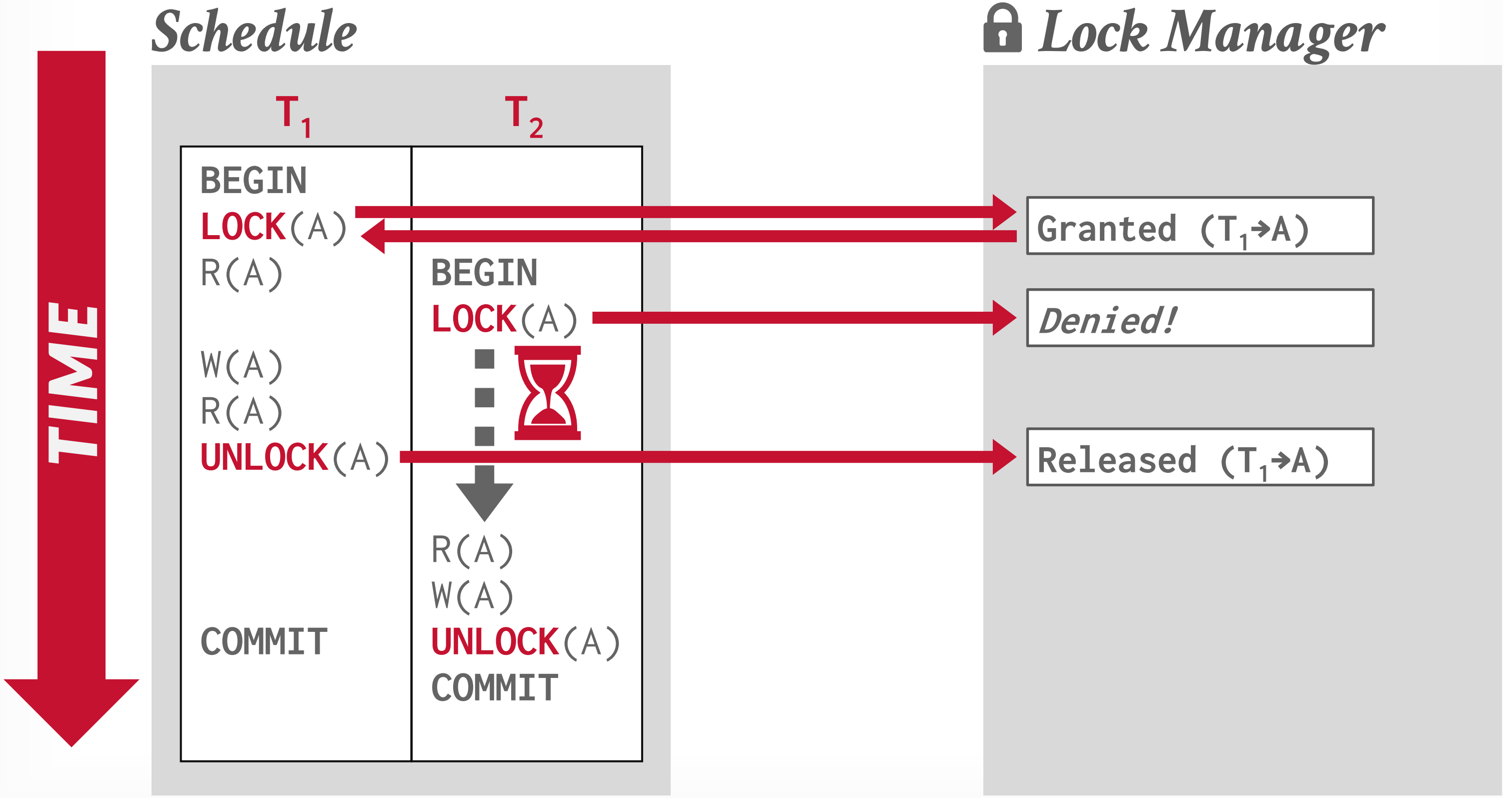


T_1	T_2
BEGIN	
LOCK(A)	
R(A)	BEGIN
	LOCK(A)
W(A)	
R(A)	
UNLOCK(A)	
	R(A)
	W(A)
COMMIT	UNLOCK(A)
	COMMIT

Granted (T_1→A)

Denied!

EXECUTING WITH LOCKS



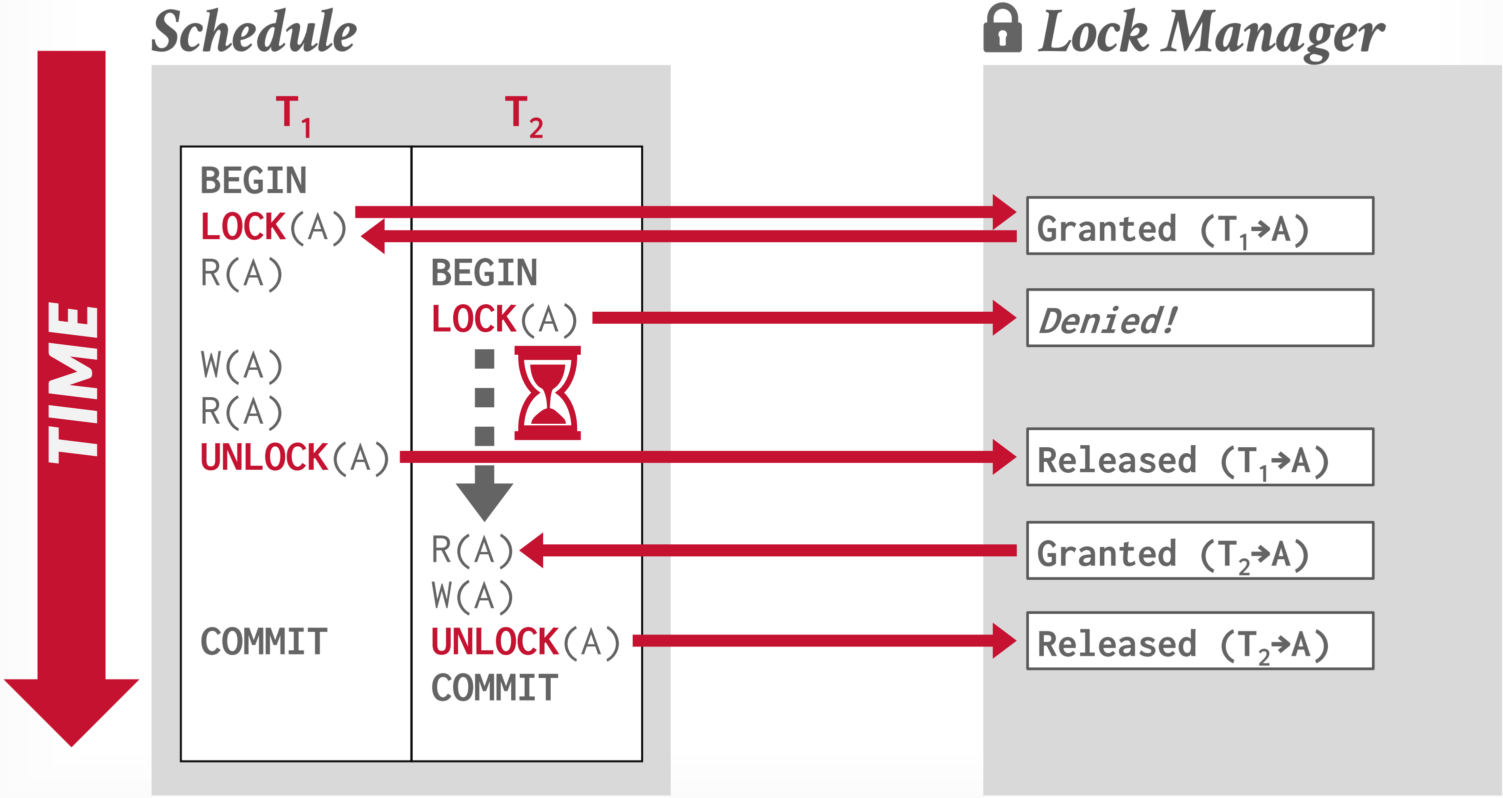
T_1	T_2
BEGIN	
LOCK(A)	
R(A)	BEGIN
	LOCK(A)
W(A)	
R(A)	
UNLOCK(A)	
	R(A)
	W(A)
COMMIT	UNLOCK(A)
	COMMIT

Granted (T_1→A)

Denied!

Released (T_1→A)

EXECUTING WITH LOCKS



TODAY'S AGENDA

Lock Types

Two-Phase Locking

Deadlock Detection + Prevention

Hierarchical Locking 通过锁定对象区域，而非单个对象来提高锁管理器的效率。

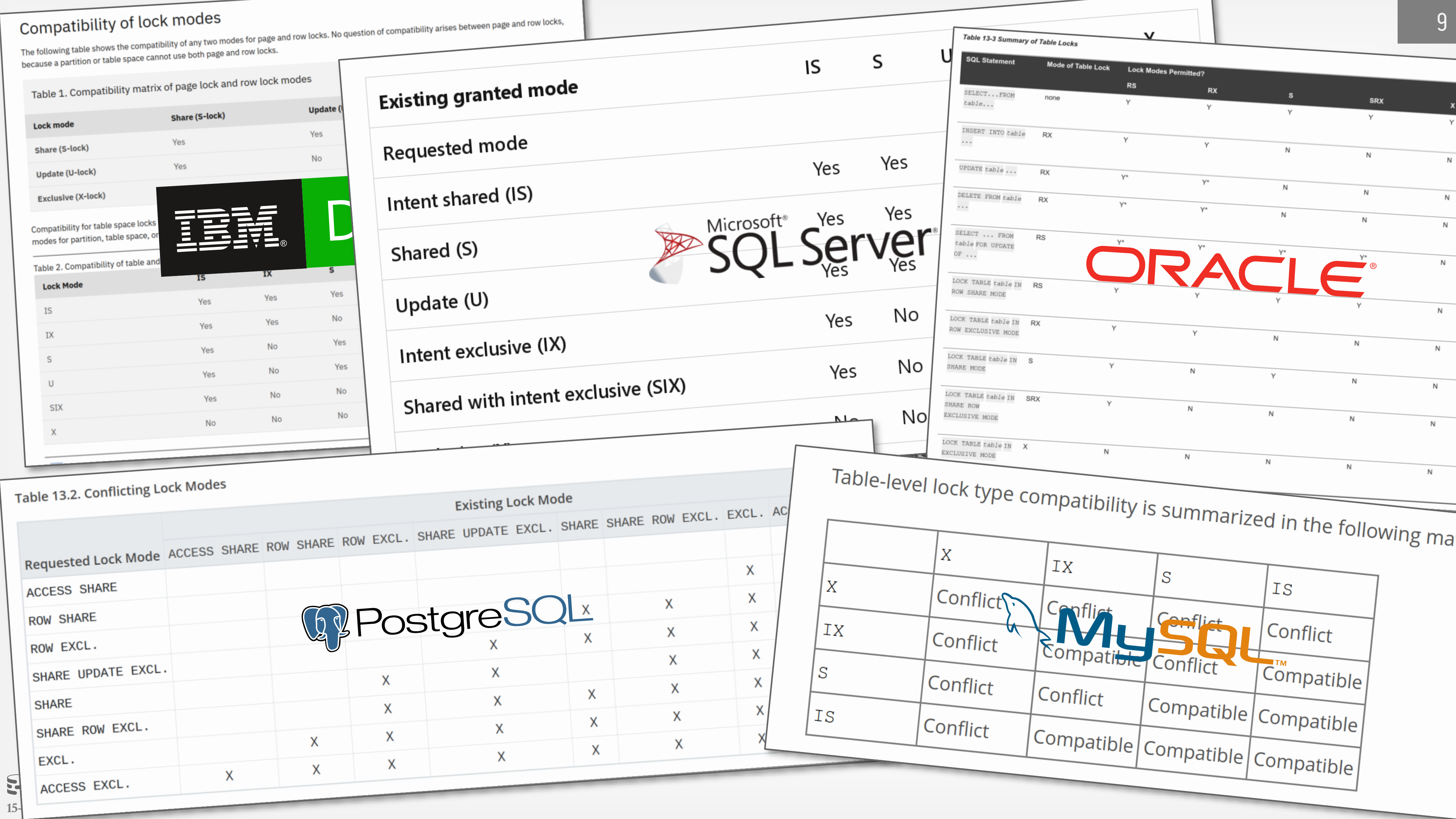
BASIC LOCK TYPES

S-LOCK: Shared locks for reads.

X-LOCK: Exclusive locks for writes.

Compatibility Matrix

	Shared S-LOCK	Exclusive X-LOCK
Shared S-LOCK	✓	✗
Exclusive X-LOCK	✗	✗



Compatibility of lock modes

The following table shows the compatibility of any two modes for page and row locks. No question of compatibility arises between page and row locks, because a partition or table space cannot use both page and row locks.

Table 1. Compatibility matrix of page lock and row lock modes

Lock mode	Share (S-lock)	Update (
Share (S-lock)	Yes	Yes
Update (U-lock)	Yes	No
Exclusive (X-lock)		

Compatibility for table space locks
modes for partition, table space, or

Table 2. Compatibility of table and

Lock Mode	IS	IX	S
IS	Yes	Yes	Yes
IX	Yes	Yes	No
S	Yes	No	Yes
U	Yes	No	Yes
SIX	Yes	No	No
X	No	No	No

Existing granted mode	IS	S	U
Requested mode			
Intent shared (IS)	Yes	Yes	
Shared (S)	Yes	Yes	
Update (U)	Yes	Yes	
Intent exclusive (IX)	Yes	No	
Shared with intent exclusive (SIX)	Yes	No	
	No	No	

Table 13-3 Summary of Table Locks

SQL Statement	Mode of Table Lock	Lock Modes Permitted?				
		RS	RX	S	SRX	X
SELECT...FROM table...	none	Y	Y	Y	Y	Y
INSERT INTO table ...	RX	Y	Y	N	N	N
UPDATE table ...	RX	Y*	Y*	N	N	N
DELETE FROM table ...	RX	Y*	Y*	N	N	N
SELECT ... FROM table FOR UPDATE OF ...	RS	Y*	Y*	Y*	Y*	N
LOCK TABLE table IN ROW SHARE MODE	RS	Y	Y	Y	Y	N
LOCK TABLE table IN ROW EXCLUSIVE MODE	RX	Y	Y	N	N	N
LOCK TABLE table IN SHARE MODE	S	Y	N	Y	N	N
LOCK TABLE table IN SHARE ROW EXCLUSIVE MODE	SRX	Y	N	N	N	N
LOCK TABLE table IN EXCLUSIVE MODE	X	N	N	N	N	N

Table 13.2. Conflicting Lock Modes

Requested Lock Mode	Existing Lock Mode							
	ACCESS SHARE	ROW SHARE	ROW EXCL.	SHARE UPDATE EXCL.	SHARE	SHARE ROW EXCL.	EXCL.	AC
ACCESS SHARE								
ROW SHARE							X	
ROW EXCL.					X	X	X	
SHARE UPDATE EXCL.				X	X	X	X	
SHARE			X	X		X	X	
SHARE ROW EXCL.			X	X	X	X	X	
EXCL.		X	X	X	X	X	X	
ACCESS EXCL.	X	X	X	X	X	X	X	

Table-level lock type compatibility is summarized in the following ma

	X	IX	S	IS
X	Conflict	Conflict	Conflict	Conflict
IX	Conflict	Compatible	Conflict	Compatible
S	Conflict	Conflict	Compatible	Compatible
IS	Conflict	Compatible	Compatible	Compatible

EXECUTING WITH LOCKS

Transactions request locks (or upgrades). 锁可以升级，例如共享锁升级为排它锁。

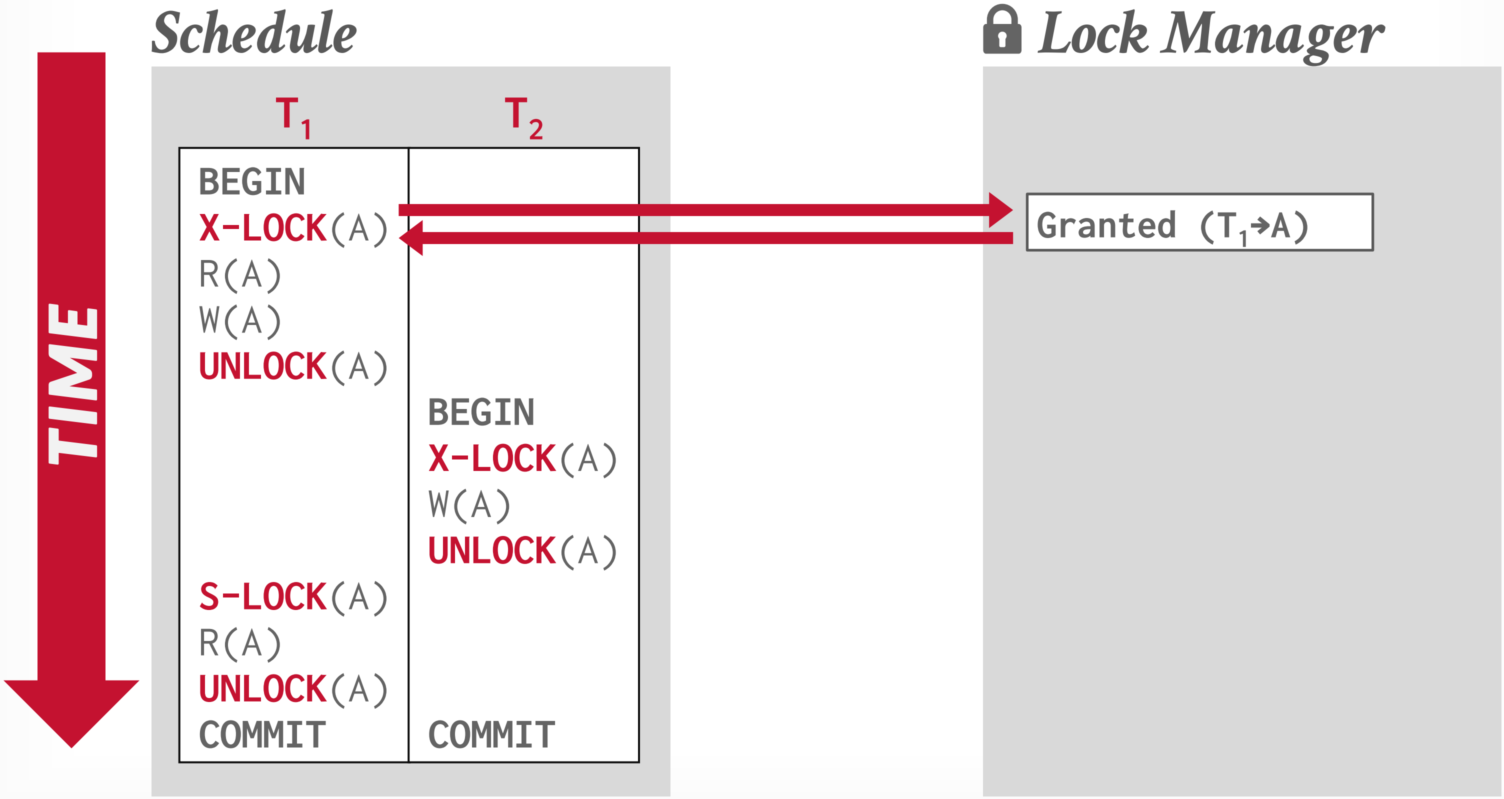
Lock manager grants or blocks requests.

Transactions release locks.

Lock manager updates its internal lock-table.

→ It keeps track of what transactions hold what locks and what transactions are waiting to acquire any locks.

EXECUTING WITH LOCKS



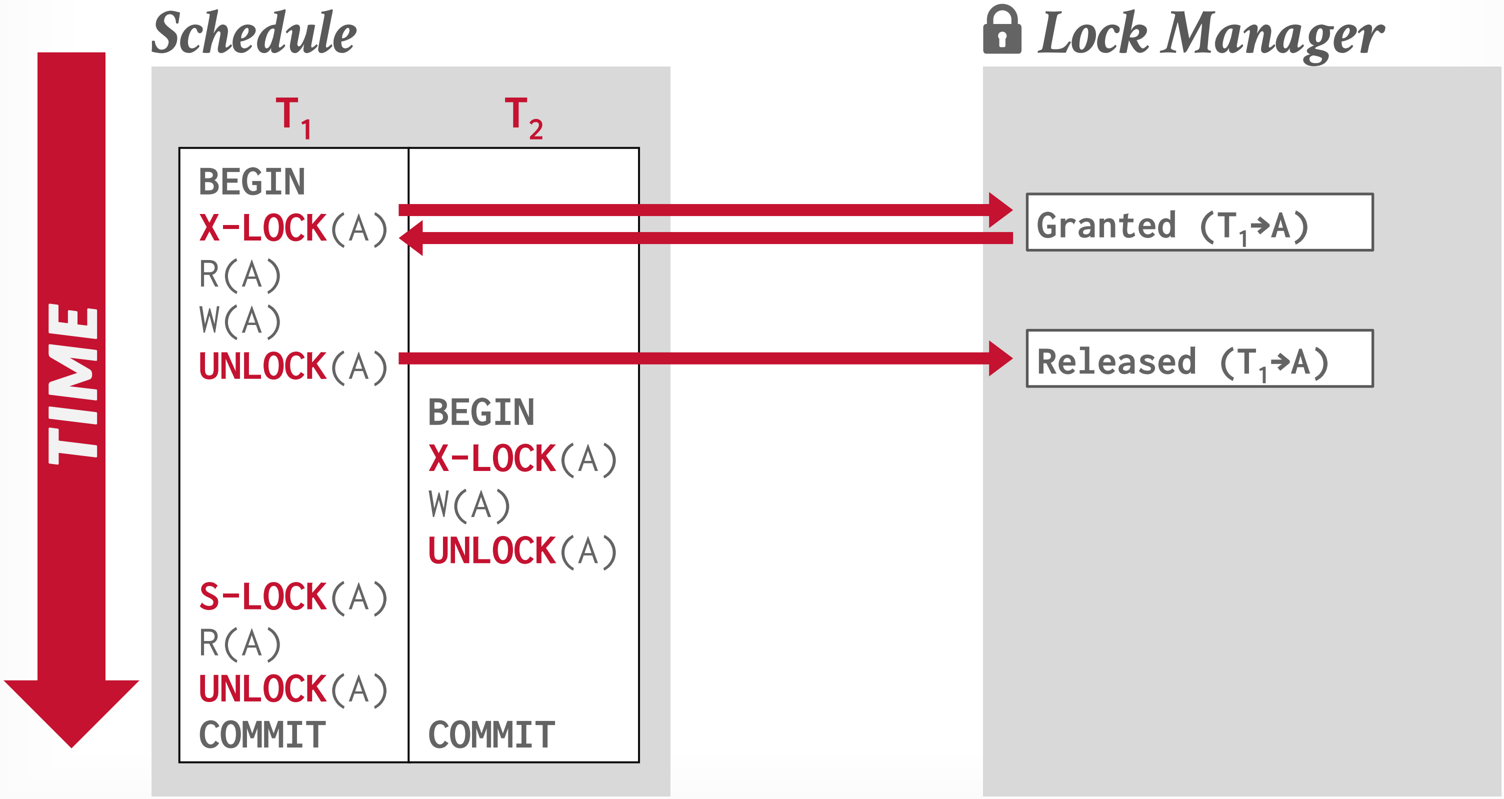
Schedule

T_1	T_2
BEGIN	
X-LOCK(A)	
R(A)	
W(A)	
UNLOCK(A)	
	BEGIN
	X-LOCK(A)
	W(A)
	UNLOCK(A)
S-LOCK(A)	
R(A)	
UNLOCK(A)	
COMMIT	COMMIT

Lock Manager

Granted (T_1→A)

EXECUTING WITH LOCKS



Schedule

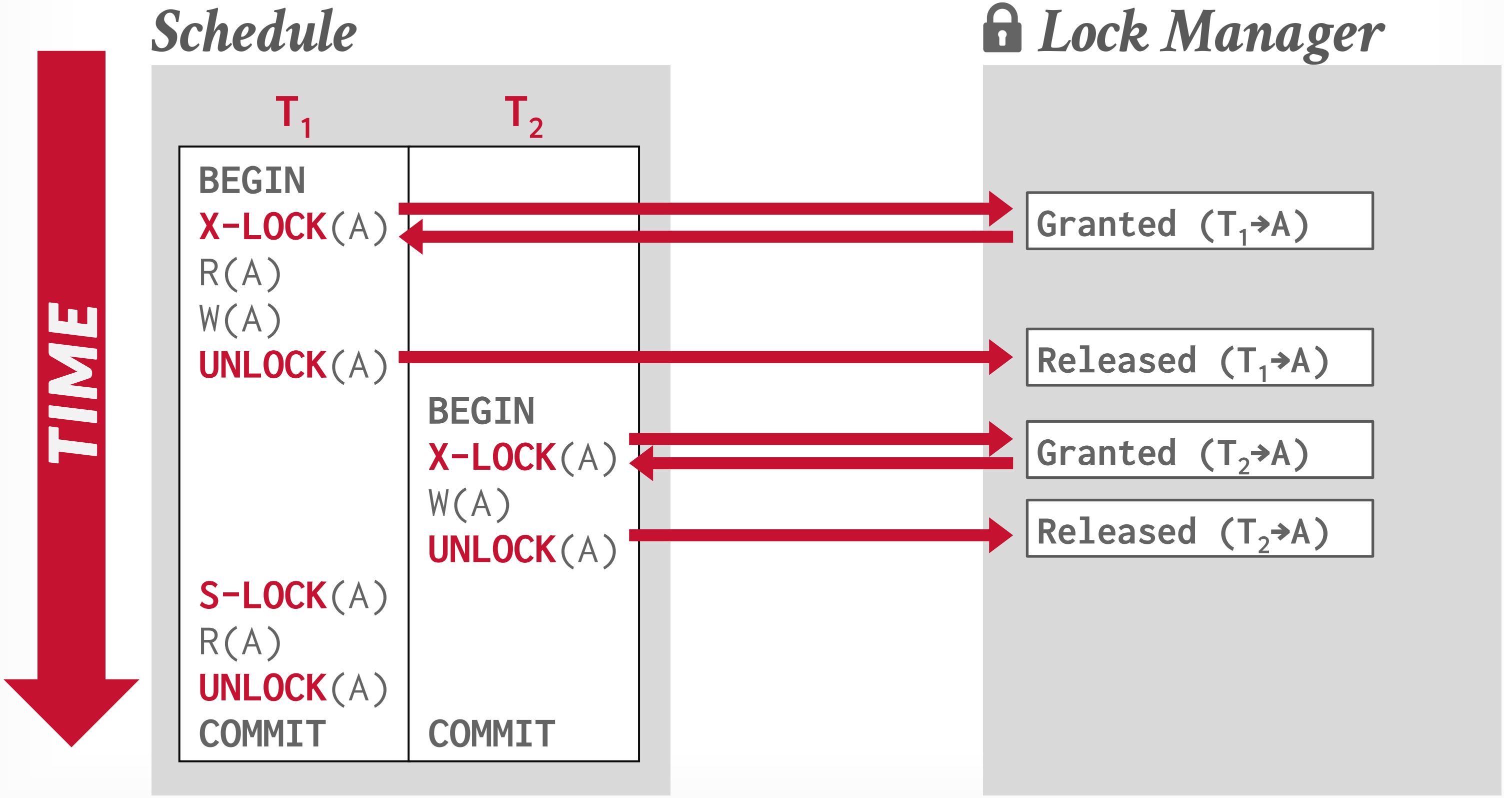
T_1	T_2
BEGIN	
X-LOCK(A)	
R(A)	
W(A)	
UNLOCK(A)	
	BEGIN
	X-LOCK(A)
	W(A)
	UNLOCK(A)
S-LOCK(A)	
R(A)	
UNLOCK(A)	
COMMIT	COMMIT

Lock Manager

Granted (T_1→A)

Released (T_1→A)

EXECUTING WITH LOCKS



Schedule

T_1	T_2
BEGIN	
X-LOCK(A)	
R(A)	
W(A)	
UNLOCK(A)	
	BEGIN
	X-LOCK(A)
	W(A)
	UNLOCK(A)
S-LOCK(A)	
R(A)	
UNLOCK(A)	
COMMIT	COMMIT

Lock Manager

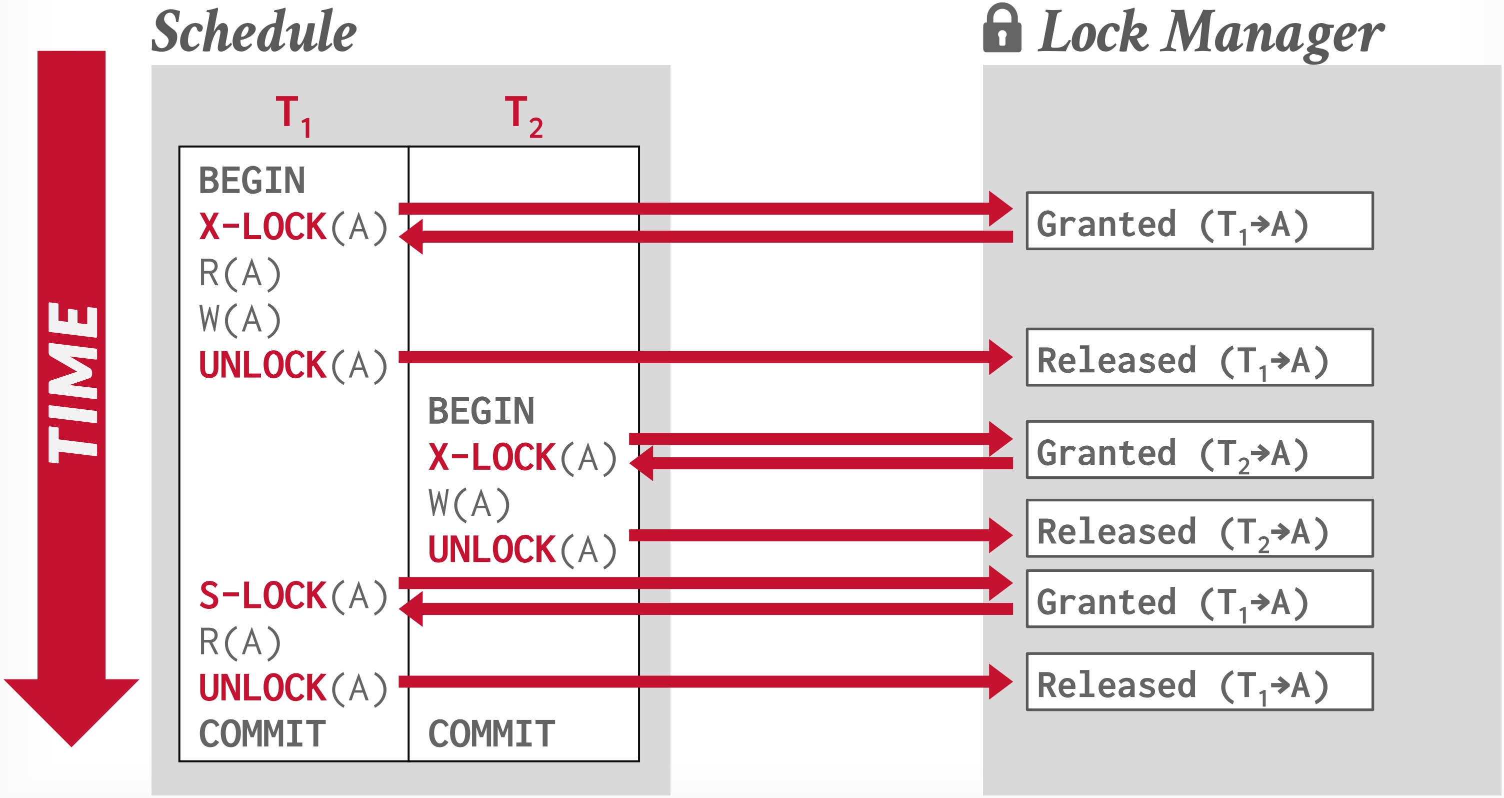
Granted (T_1→A)

Released (T_1→A)

Granted (T_2→A)

Released (T_2→A)

EXECUTING WITH LOCKS



Schedule

T_1	T_2
BEGIN	
X-LOCK(A)	
R(A)	
W(A)	
UNLOCK(A)	
	BEGIN
	X-LOCK(A)
	W(A)
	UNLOCK(A)
S-LOCK(A)	
R(A)	
UNLOCK(A)	
COMMIT	COMMIT

Lock Manager

Granted (T_1→A)

Released (T_1→A)

Granted (T_2→A)

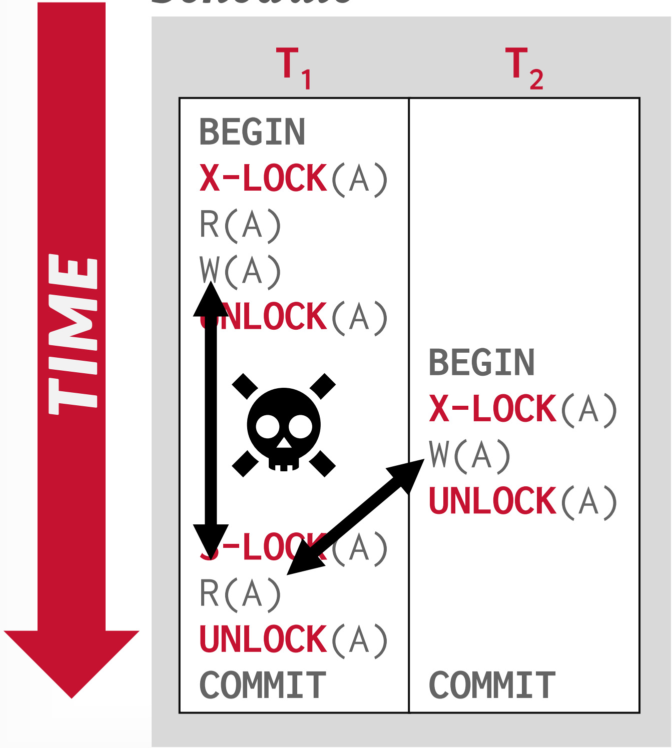
Released (T_2→A)

Granted (T_1→A)

Released (T_1→A)

EXECUTING WITH LOCKS

Schedule



T_1	T_2
BEGIN	
X-LOCK(A)	
R(A)	
W(A)	
UNLOCK(A)	
	BEGIN
	X-LOCK(A)
	W(A)
	UNLOCK(A)
S-LOCK(A)	
R(A)	
UNLOCK(A)	
COMMIT	COMMIT

T1 不可重复读

Lock Manager

Granted (T_1→A)

Released (T_1→A)

Granted (T_2→A)

Released (T_2→A)

Granted (T_1→A)

Released (T_1→A)

T1 释放锁之后又重新获取锁

CONCURRENCY CONTROL PROTOCOL

Two-phase locking (2PL) is a concurrency control protocol that determines whether a txn can access an object in the database at runtime. 并发控制协议。

The protocol does not need to know all the queries that a txn will execute ahead of time.

TWO-PHASE LOCKING

Phase #1: Growing 增长阶段

→ Each txn requests the locks that it needs from the DBMS's lock manager.

→ The lock manager grants/denies lock requests.

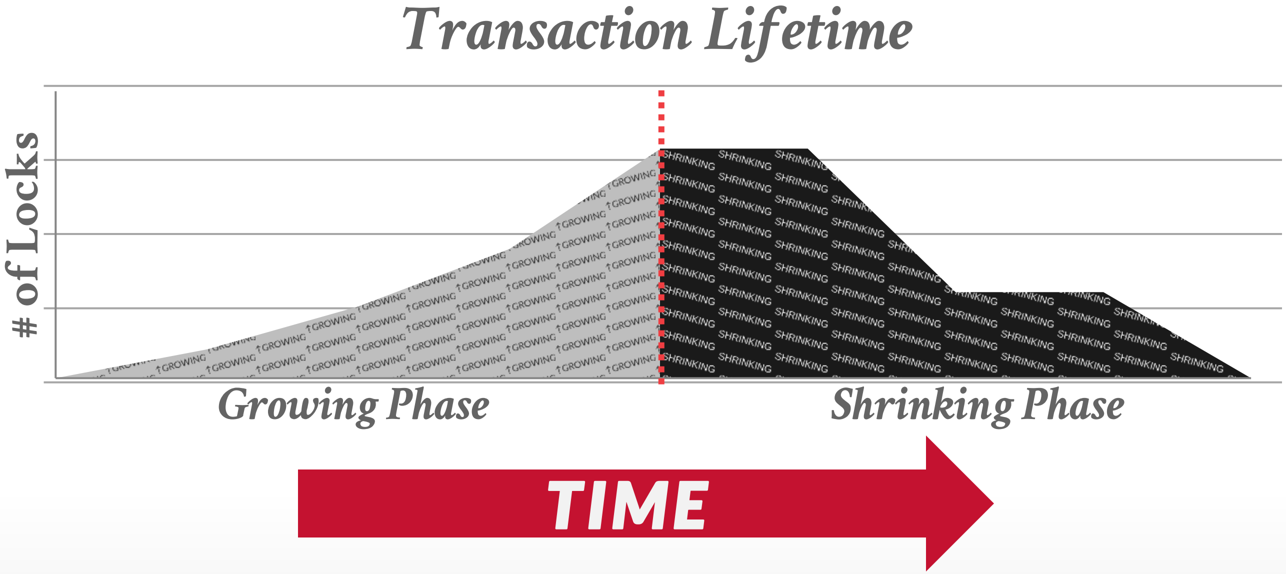
Phase #2: Shrinking 收缩阶段

→ The txn is allowed to only release/downgrade locks that it previously acquired. It cannot acquire new locks.

一旦事务开始释放锁，就自动进入收缩阶段，不可重新获取锁。

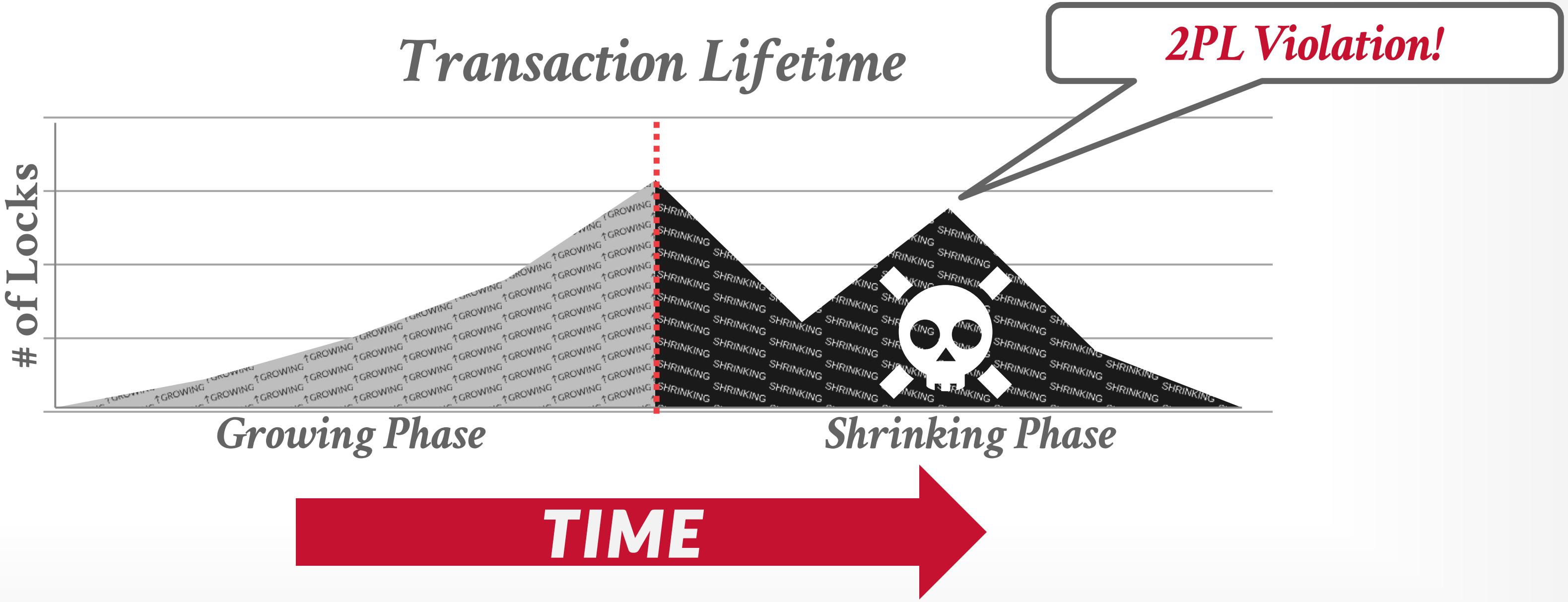
TWO-PHASE LOCKING

The txn is not allowed to acquire/upgrade locks after the growing phase finishes.

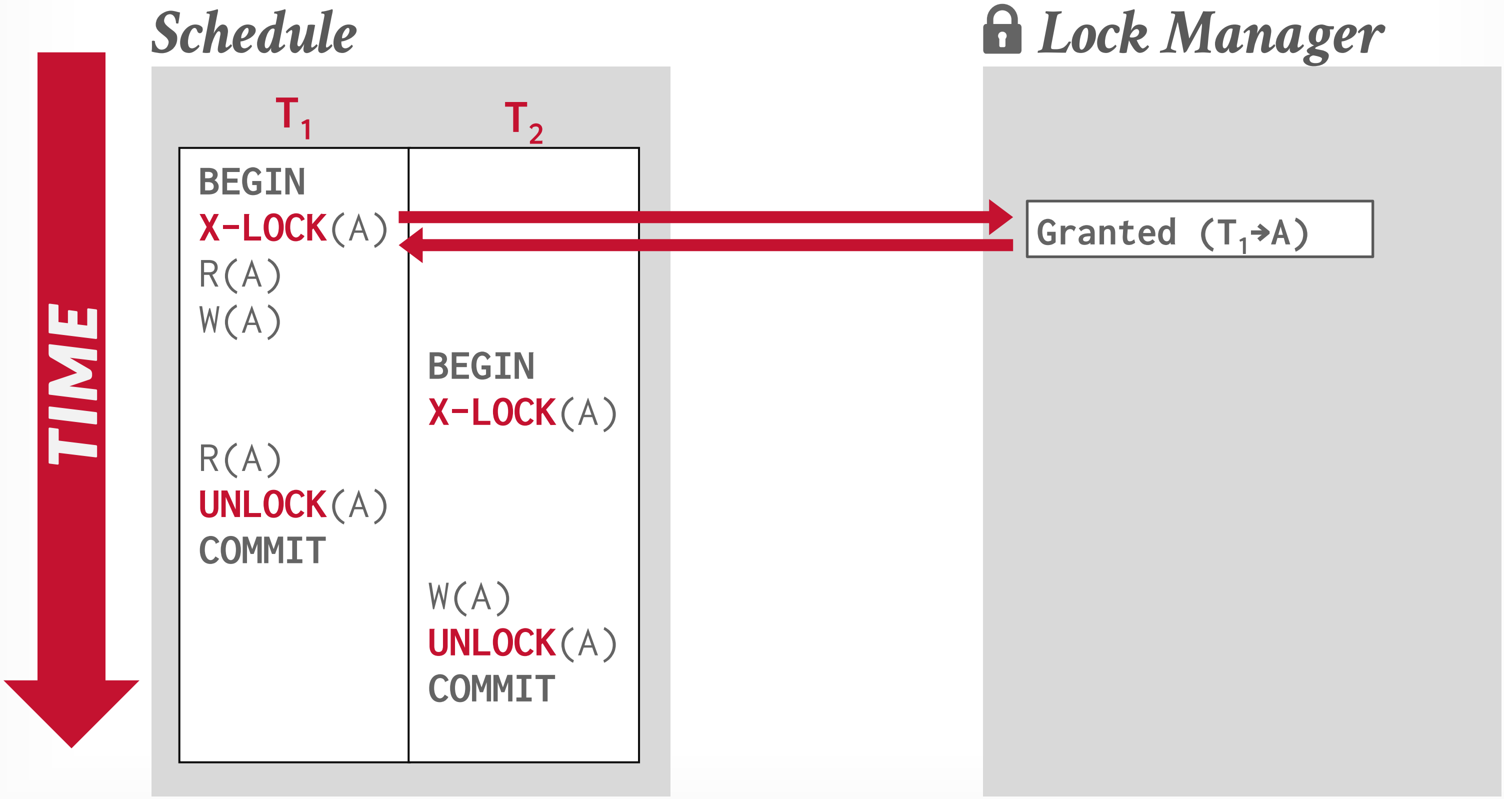


TWO-PHASE LOCKING

The txn is not allowed to acquire/upgrade locks after the growing phase finishes.



EXECUTING WITH 2PL



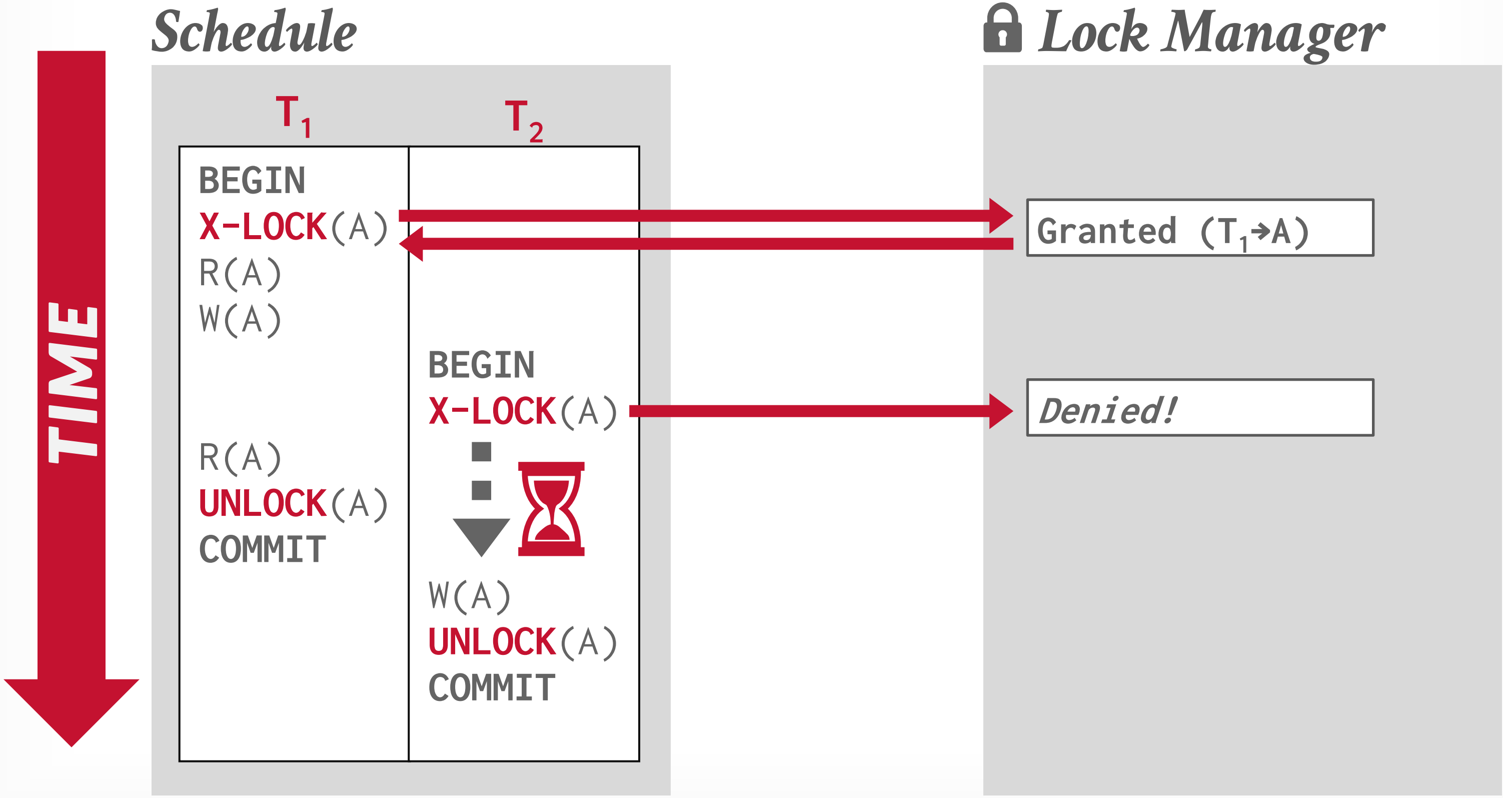
Schedule

T_1	T_2
BEGIN	
X-LOCK(A)	
R(A)	
W(A)	
	BEGIN
	X-LOCK(A)
R(A)	
UNLOCK(A)	
COMMIT	
	W(A)
	UNLOCK(A)
	COMMIT

Lock Manager

Granted (T_1→A)

EXECUTING WITH 2PL

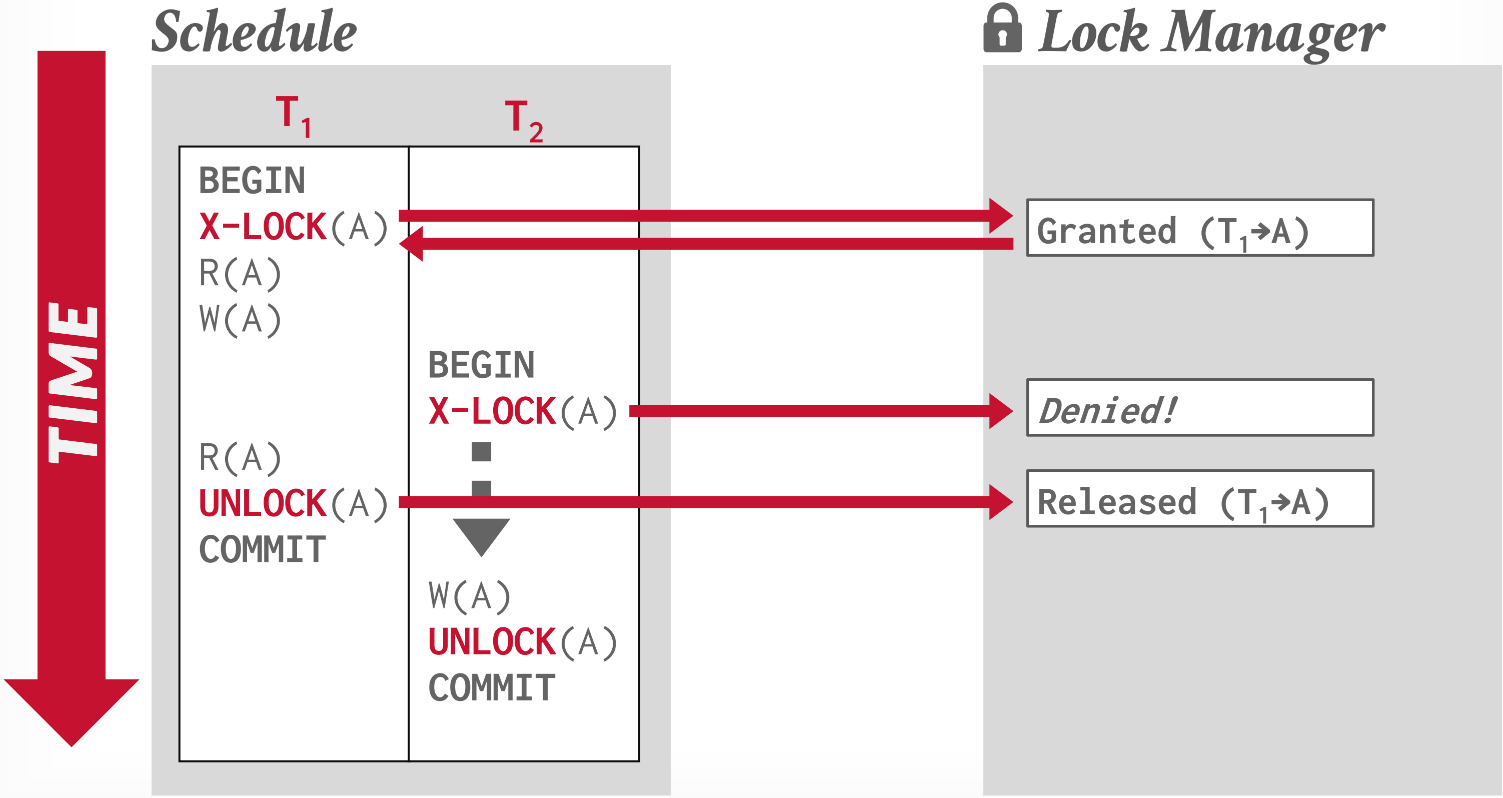


T_1	T_2
BEGIN	
X-LOCK(A)	
R(A)	
W(A)	
	BEGIN
	X-LOCK(A)
R(A)	
UNLOCK(A)	
COMMIT	
	W(A)
	UNLOCK(A)
	COMMIT

Granted (T_1→A)

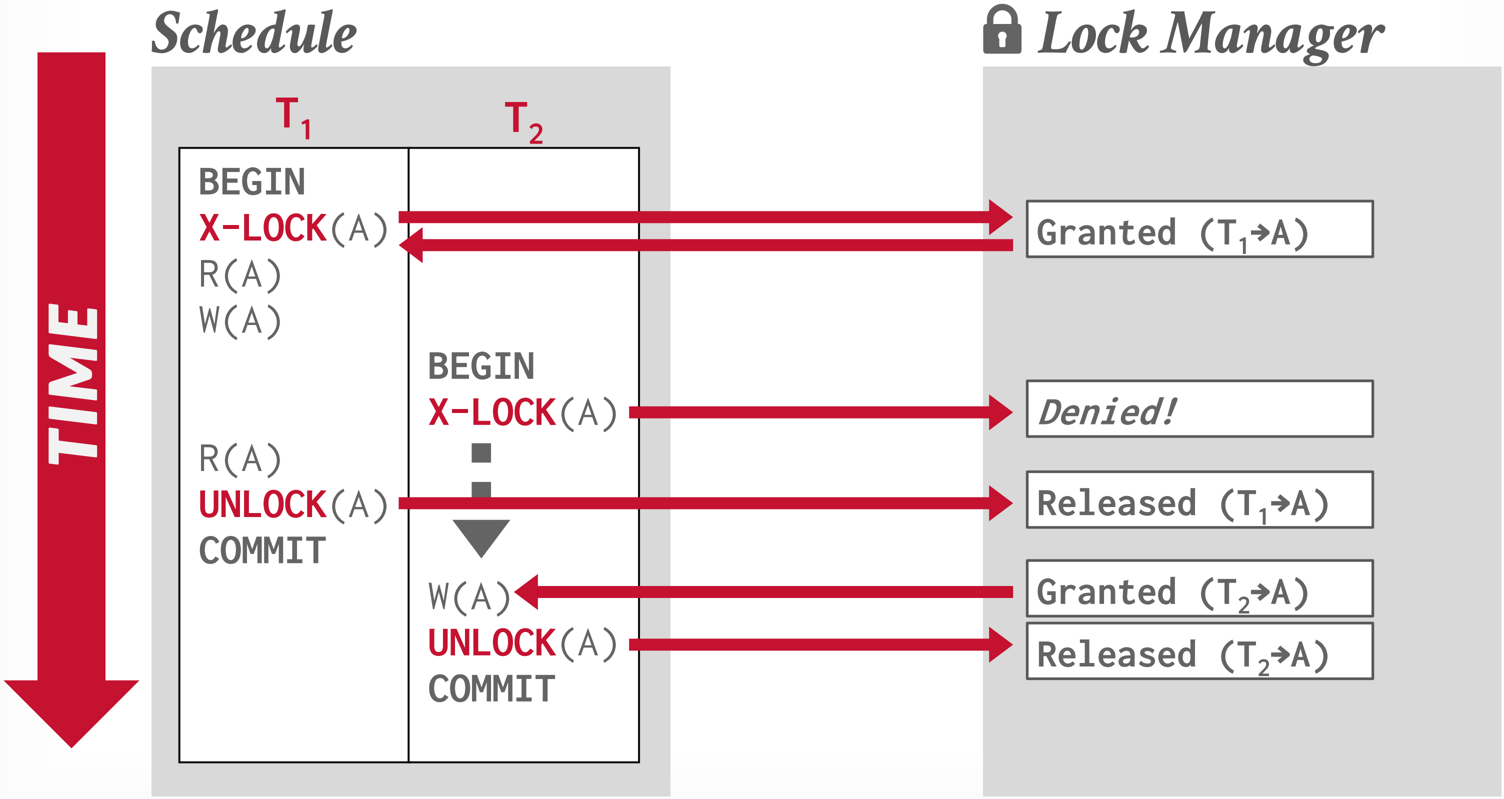
Denied!

EXECUTING WITH 2PL



T_1	T_2
BEGIN	
X-LOCK(A)	
R(A)	
W(A)	
	BEGIN
	X-LOCK(A)
R(A)	
UNLOCK(A)	
COMMIT	
	W(A)
	UNLOCK(A)
	COMMIT

EXECUTING WITH 2PL



Schedule

T_1	T_2
BEGIN	
X-LOCK(A)	
R(A)	
W(A)	
	BEGIN
	X-LOCK(A)
R(A)	
UNLOCK(A)	
COMMIT	
	W(A)
	UNLOCK(A)
	COMMIT

Lock Manager

Granted (T_1→A)

Denied!

Released (T_1→A)

Granted (T_2→A)

Released (T_2→A)

TWO-PHASE LOCKING

2PL on its own is sufficient to guarantee conflict serializability because it generates schedules whose precedence graph is acyclic.

But it is subject to **cascading aborts**. 级联中断，这是一个性能问题。

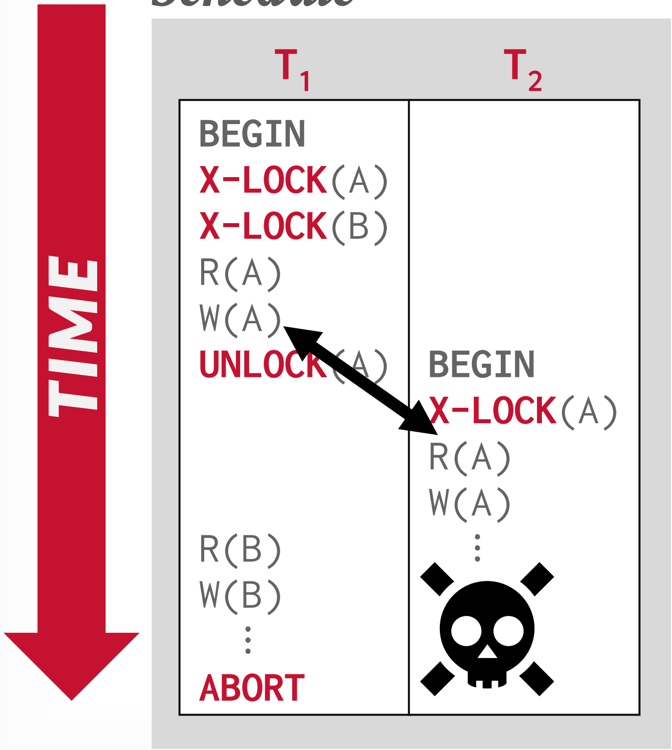
2PL: CASCADING ABORTS

Schedule

T_1	T_2
BEGIN	
X-LOCK(A)	
X-LOCK(B)	
R(A)	
W(A)	
UNLOCK(A)	BEGIN
	X-LOCK(A)
	R(A)
	W(A)
R(B)	⋮
W(B)	
⋮	
ABORT	

2PL: CASCADING ABORTS

Schedule



T_1	T_2
BEGIN	
X-LOCK(A)	
X-LOCK(B)	
R(A)	
W(A)	
UNLOCK(A)	BEGIN
	X-LOCK(A)
	R(A)
	W(A)
	⋮
R(B)	
W(B)	
⋮	
ABORT	

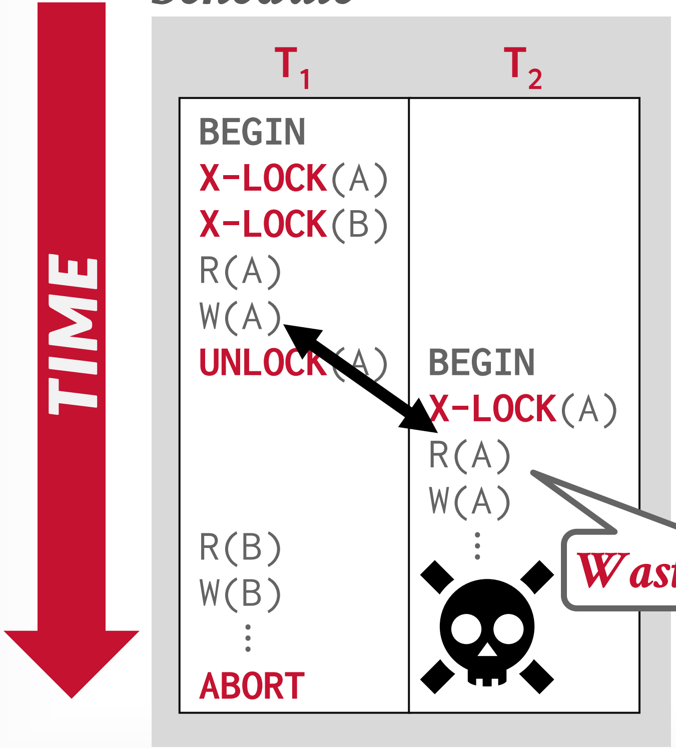
This is a permissible schedule in 2PL, but the DBMS has to also abort T_2 when T_1 aborts.

Any information about T_1 cannot be "leaked" to the outside world.

Any computation performed must be rolled back.

2PL: CASCADING ABORTS

Schedule



T_1	T_2
BEGIN	
X-LOCK(A)	
X-LOCK(B)	
R(A)	
W(A)	
UNLOCK(A)	BEGIN
	X-LOCK(A)
	R(A)
	W(A)
R(B)	⋮
W(B)	
⋮	
ABORT	

Wasted work!

This is a permissible schedule in 2PL, but the DBMS has to also abort T_2 when T_1 aborts.

T1 中断后，T2 也需要回滚相应的操作。

Any information about T_1 cannot be “leaked” to the outside world.

Any computation performed must lled back.

2PL OBSERVATIONS

There are potential schedules that are serializable but would not be allowed by 2PL because locking limits concurrency.

→ Most DBMSs prefer correctness before performance.

May still have "dirty reads".

→ Solution: **Strong Strict 2PL (aka Rigorous 2PL)**

May lead to deadlocks.

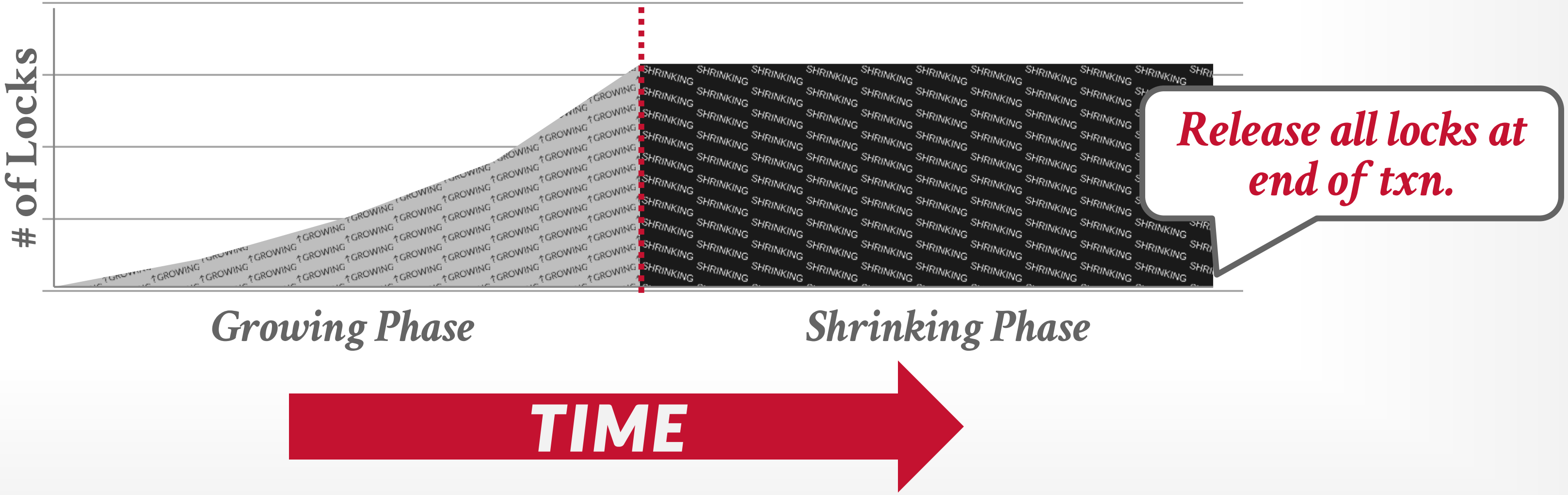
→ Solution: **Detection** or **Prevention**

STRONG STRICT TWO-PHASE LOCKING

The txn is only allowed to release locks after it has ended (i.e., committed or aborted). 所有锁在事务结束时全部释放。

还有一种严格两阶段锁，支持事务在收缩阶段释放共享锁。

Allows only conflict serializable schedules, but it is often stronger than needed for some apps.



Release all locks at end of txn.

STRONG STRICT TWO-PHASE LOCKING

A schedule is **strict** if a value written by a txn is not read or overwritten by other txns until that txn finishes.

Advantages:

→ Does not incur cascading aborts.
→ Aborted txns can be undone by just restoring original values of modified tuples.

EXAMPLES

T_1 – Move $100 from Andy's account (**A**) to his bookie's account (**B**).

T_2 – Compute the total amount in all accounts and return it to the application.

T_1

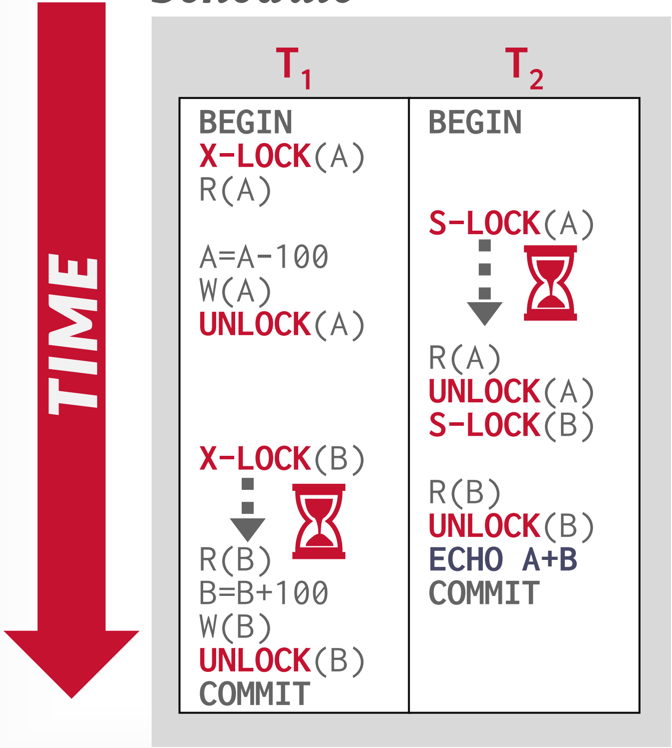
```
BEGIN
A=A-100
B=B+100
COMMIT
```

T_2

```
BEGIN
ECHO A+B
COMMIT
```

NON-2PL EXAMPLE

Schedule



T_1	T_2
BEGIN	BEGIN
X-LOCK(A)	
R(A)	
	S-LOCK(A)
A=A-100	
W(A)	
UNLOCK(A)	
	R(A)
	UNLOCK(A)
	S-LOCK(B)
X-LOCK(B)	
	R(B)
	UNLOCK(B)
R(B)	ECHO A+B
B=B+100	COMMIT
W(B)	
UNLOCK(B)	
COMMIT	

Initial Database State

A=1000, B=1000

T_2 Output

A+B=1900

2PL EXAMPLE

Schedule

TIME

T1	T2
BEGIN	BEGIN
X-LOCK(A)	
R(A)	S-LOCK(A)
A=A-100	⌛
W(A)	
X-LOCK(B)	
UNLOCK(A)	R(A)
	S-LOCK(B)
R(B)	⌛
B=B+100	
W(B)	
UNLOCK(B)	R(B)
COMMIT	UNLOCK(A)
	UNLOCK(B)
	ECHO A+B
	COMMIT

Initial Database State

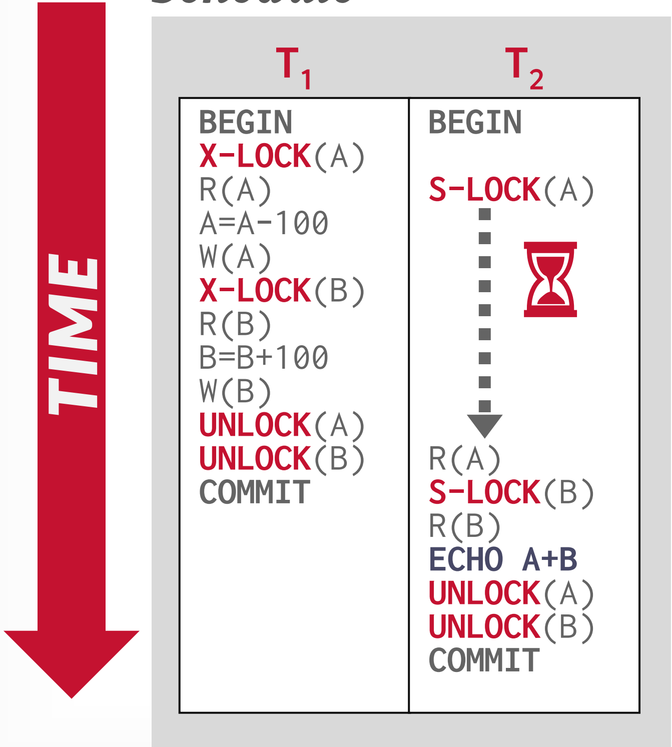
A=1000, B=1000

T2 Output

A+B=2000

STRONG STRICT 2PL EXAMPLE

Schedule



T_1	T_2
BEGIN	BEGIN
X-LOCK(A)	
R(A)	S-LOCK(A)
A=A-100	
W(A)	
X-LOCK(B)	
R(B)	
B=B+100	
W(B)	
UNLOCK(A)	
UNLOCK(B)	R(A)
COMMIT	S-LOCK(B)
	R(B)
	ECHO A+B
	UNLOCK(A)
	UNLOCK(B)
	COMMIT

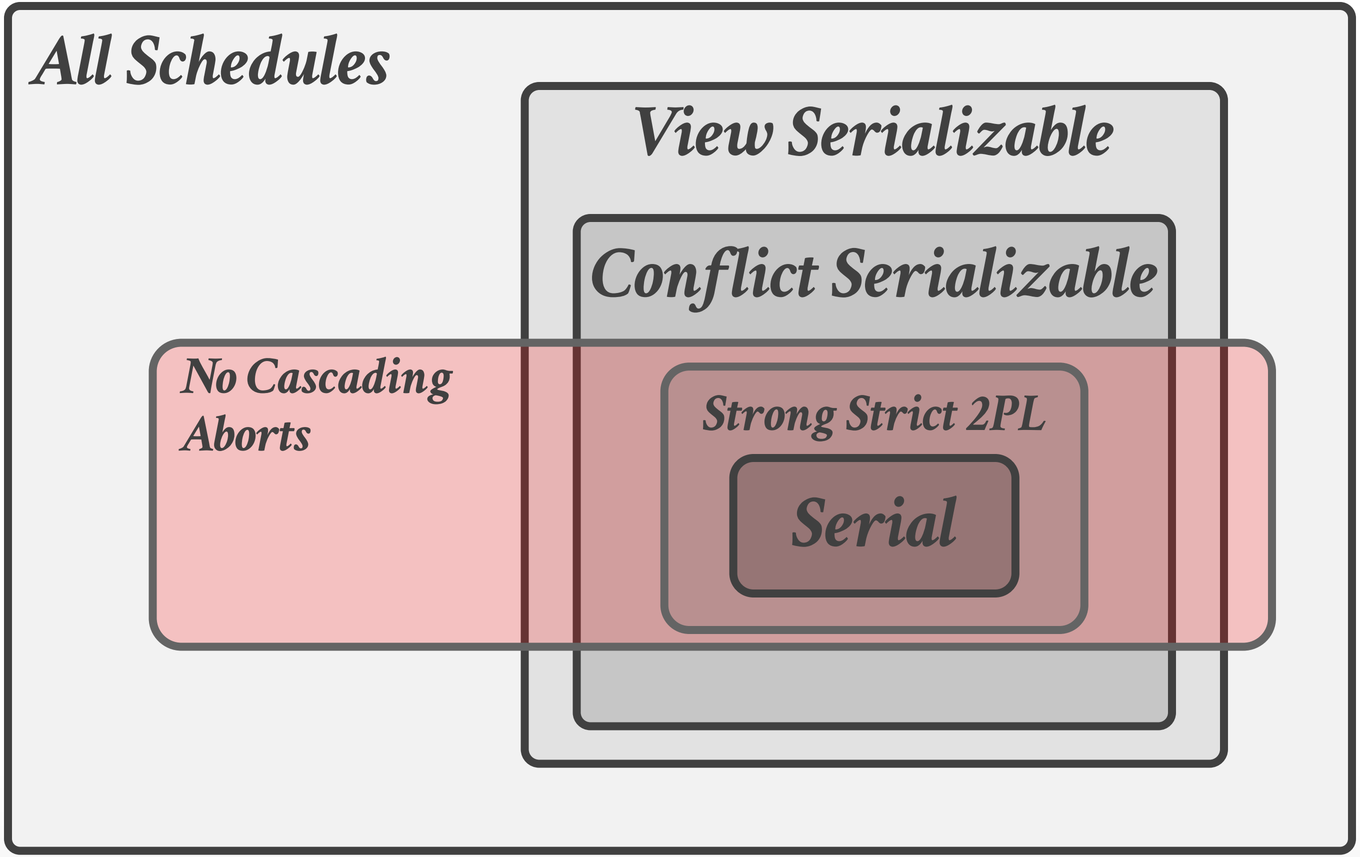
Initial Database State

A=1000, B=1000

T_2 Output

A+B=2000

UNIVERSE OF SCHEDULES



2PL OBSERVATIONS

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→ Most DBMSs prefer correctness before performance.

May still have "dirty reads".

→ Solution: **Strong Strict 2PL (aka Rigorous 2PL)**

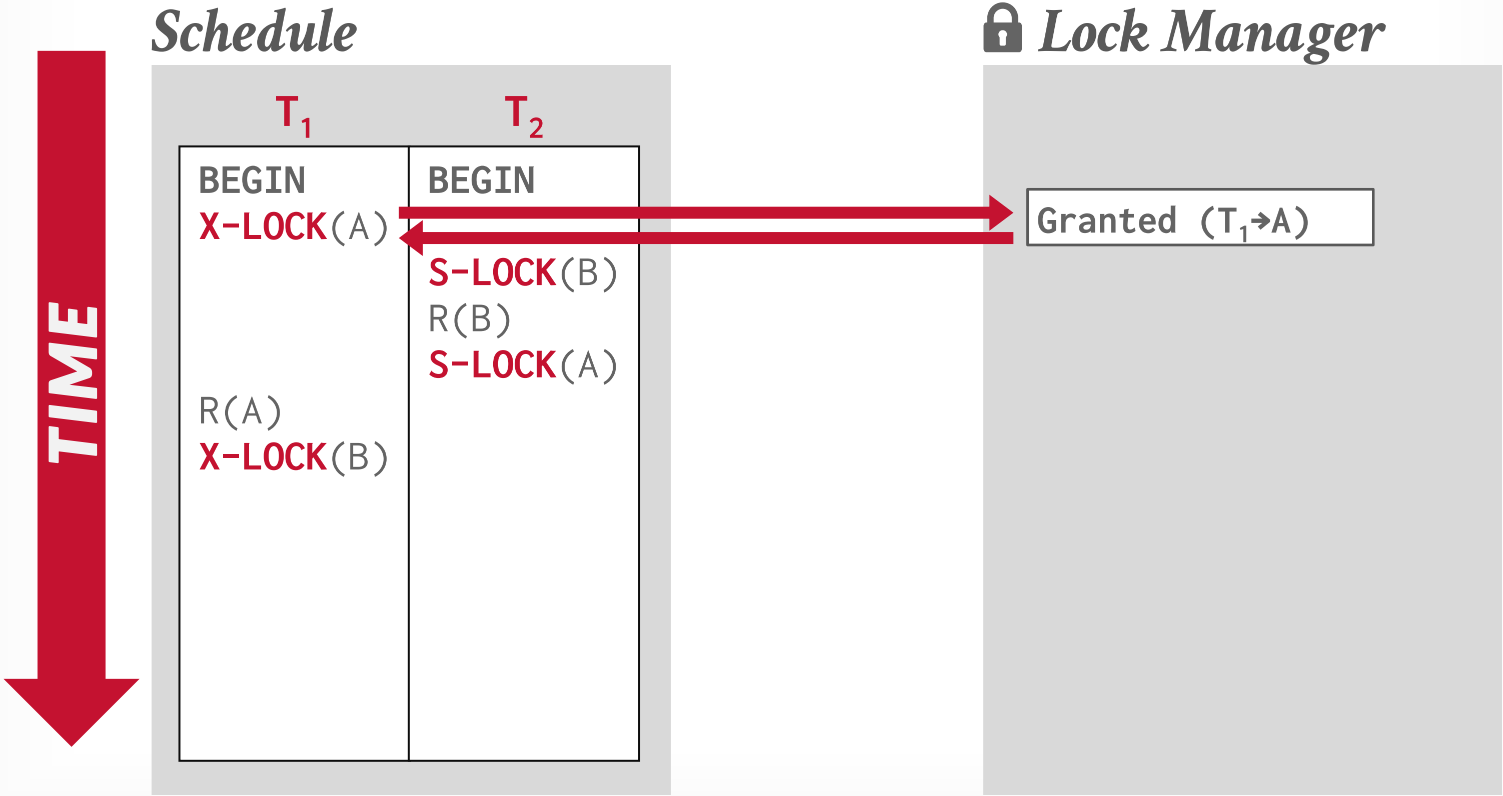
May lead to deadlocks.

→ Solution: **Detection** or **Prevention**

死锁检测：后台线程或进程检测死锁何时发生，并尝试解除它。

死锁预防：以某种方式对事务获取锁的顺序进行安排，确保不可能发生死锁。

IT JUST GOT REAL



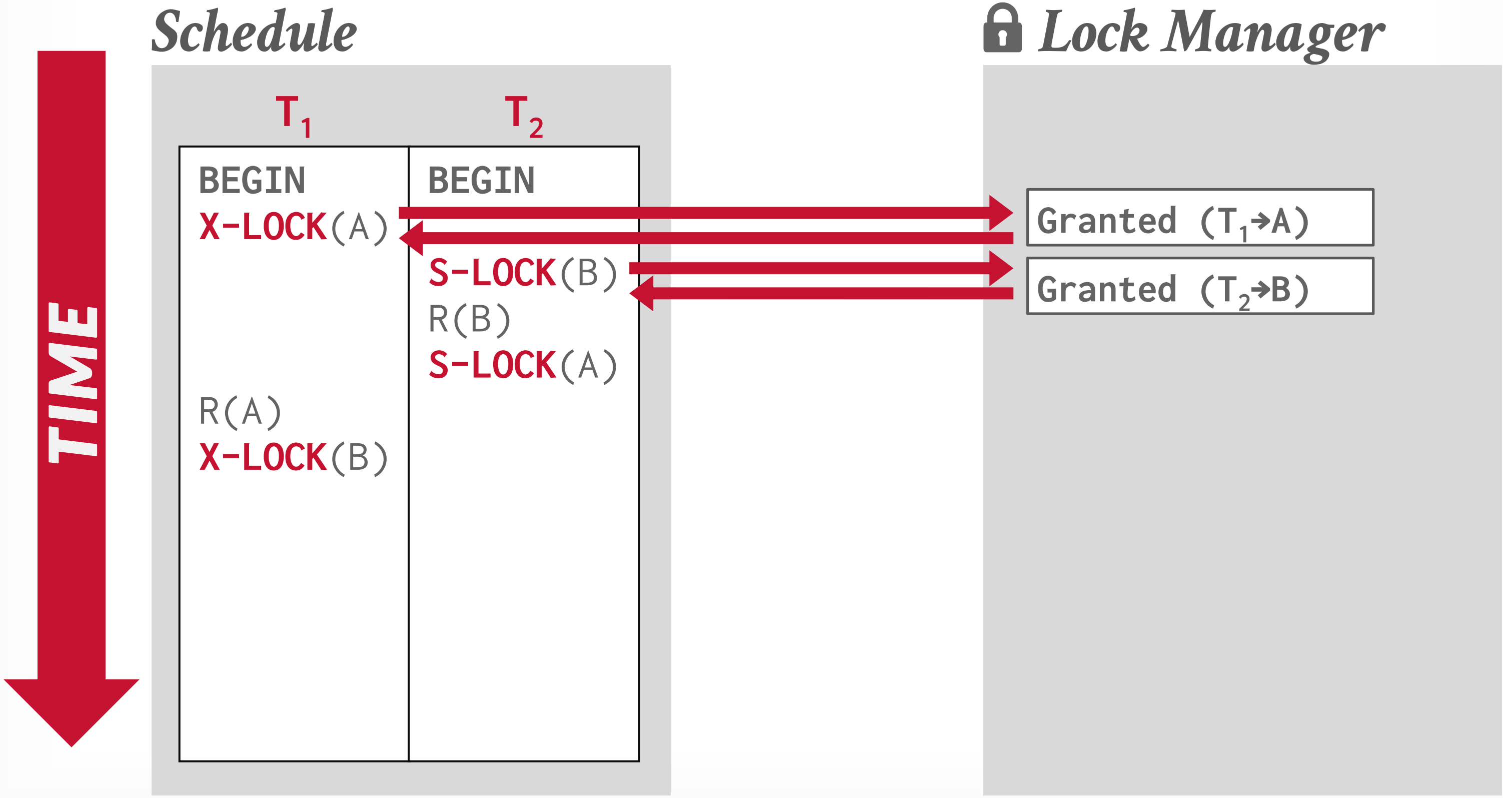
Schedule

T_1	T_2
BEGIN	BEGIN
X-LOCK(A)	
	S-LOCK(B)
	R(B)
	S-LOCK(A)
R(A)	
X-LOCK(B)	

Lock Manager

Granted (T_1→A)

IT JUST GOT REAL



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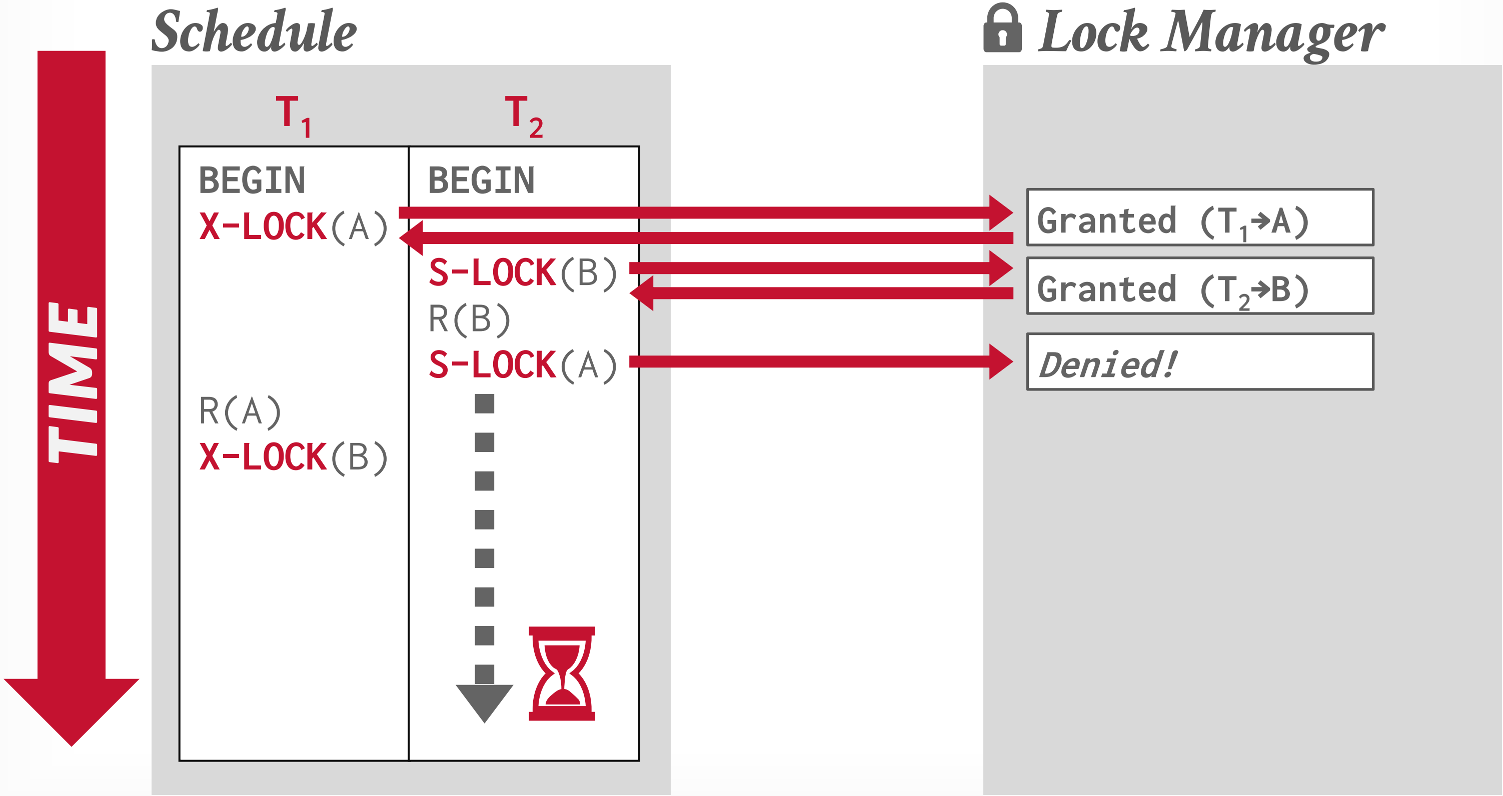


Schedule

T_1	T_2
BEGIN	BEGIN
X-LOCK(A)	
	S-LOCK(B)
	R(B)
	S-LOCK(A)
R(A)	
X-LOCK(B)	

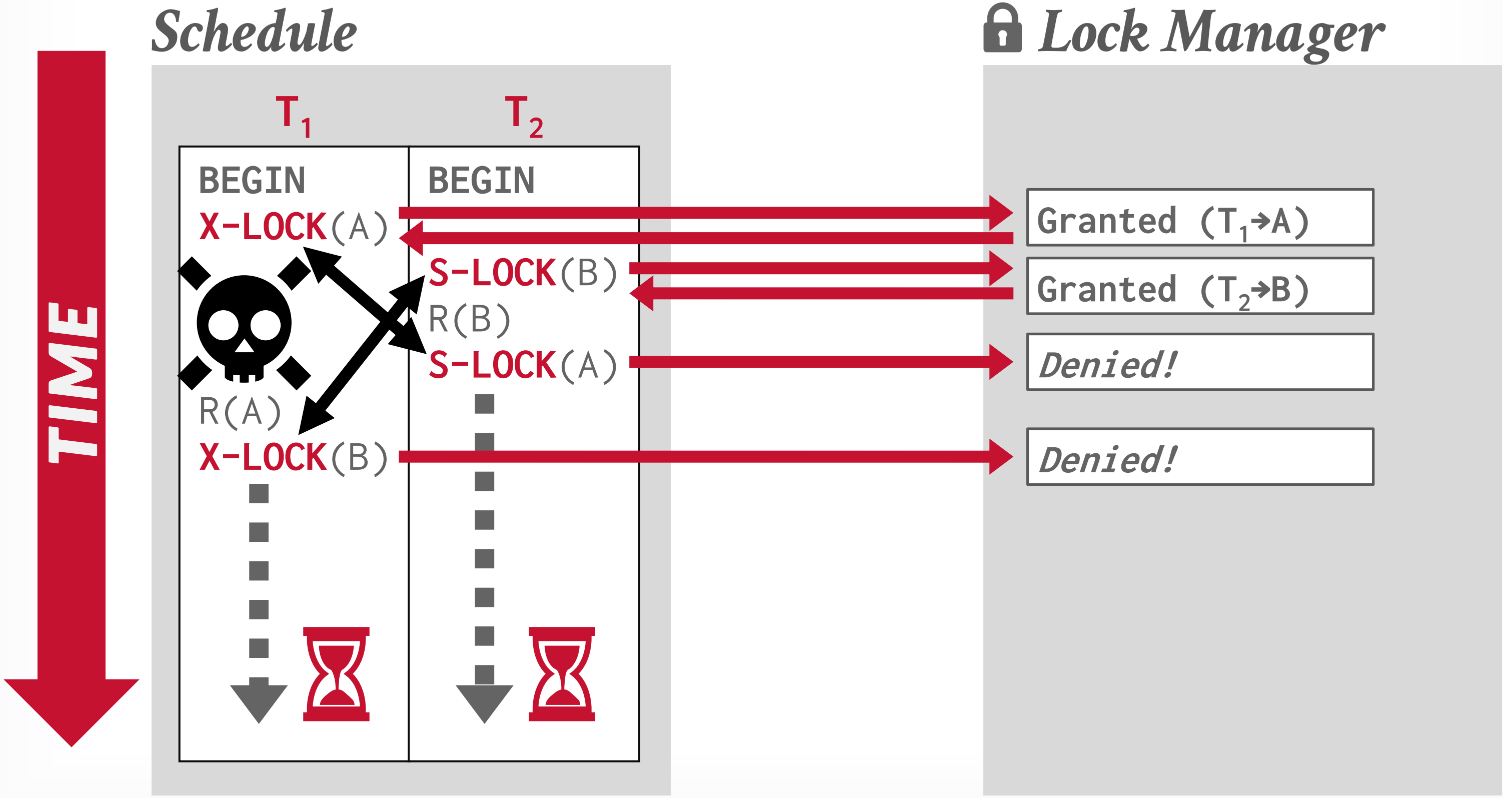
Lock Manager

Granted (T_1→A)

Granted (T_2→B)

Denied!

IT JUST GOT REAL



2PL DEADLOCKS

A **deadlock** is a cycle of transactions waiting for locks to be released by each other.

Two ways of dealing with deadlocks:
→ **Approach #1: Deadlock Detection**
→ **Approach #2: Deadlock Prevention**

DEADLOCK DETECTION

The DBMS creates a **waits-for** graph to keep track of what locks each txn is waiting to acquire:

→ Nodes are transactions

→ Edge from T_i to T_j if T_i is waiting for T_j to release a lock.

The system periodically checks for cycles in ***waits-for*** graph and then decides how to break it.

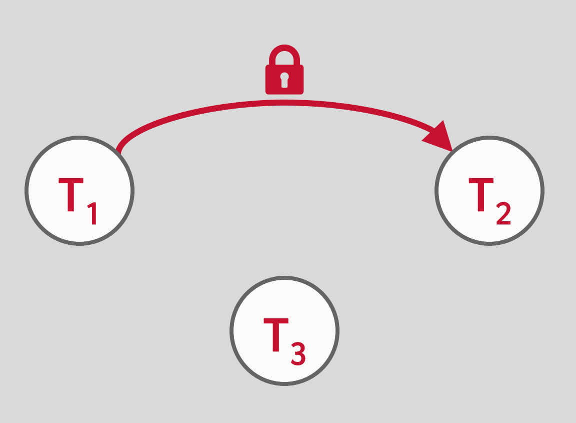
DEADLOCK DETECTION

Schedule

TIME

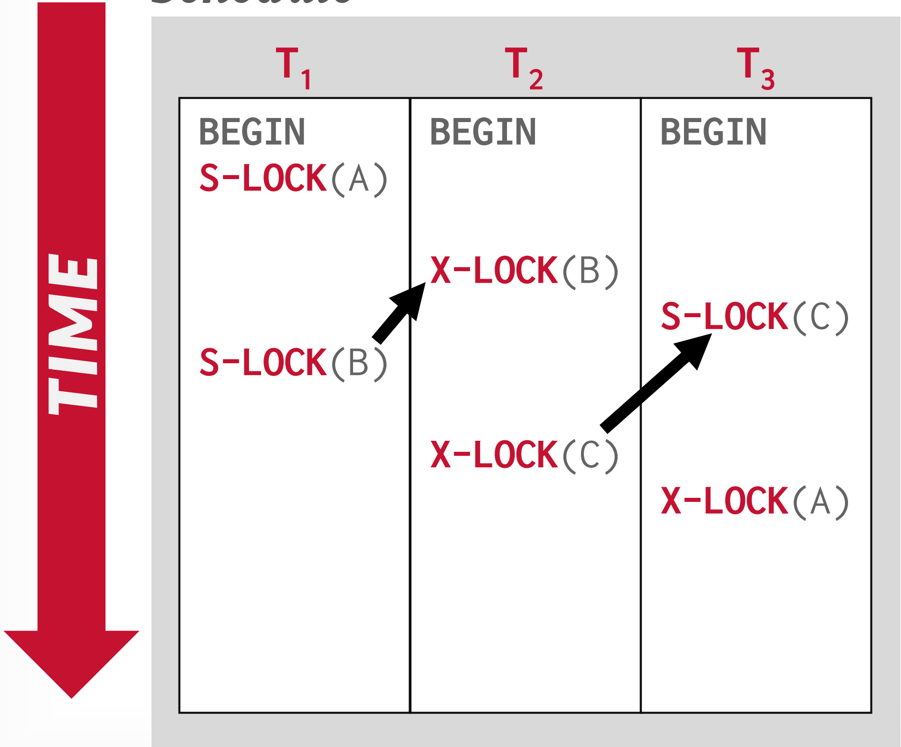
T_1	T_2	T_3
BEGIN	BEGIN	BEGIN
S-LOCK(A)		
	X-LOCK(B)	
		S-LOCK(C)
S-LOCK(B)		
	X-LOCK(C)	
		X-LOCK(A)

Waits-For Graph



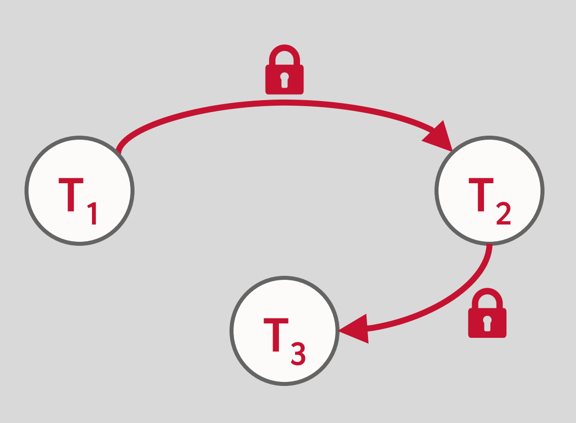
DEADLOCK DETECTION

Schedule



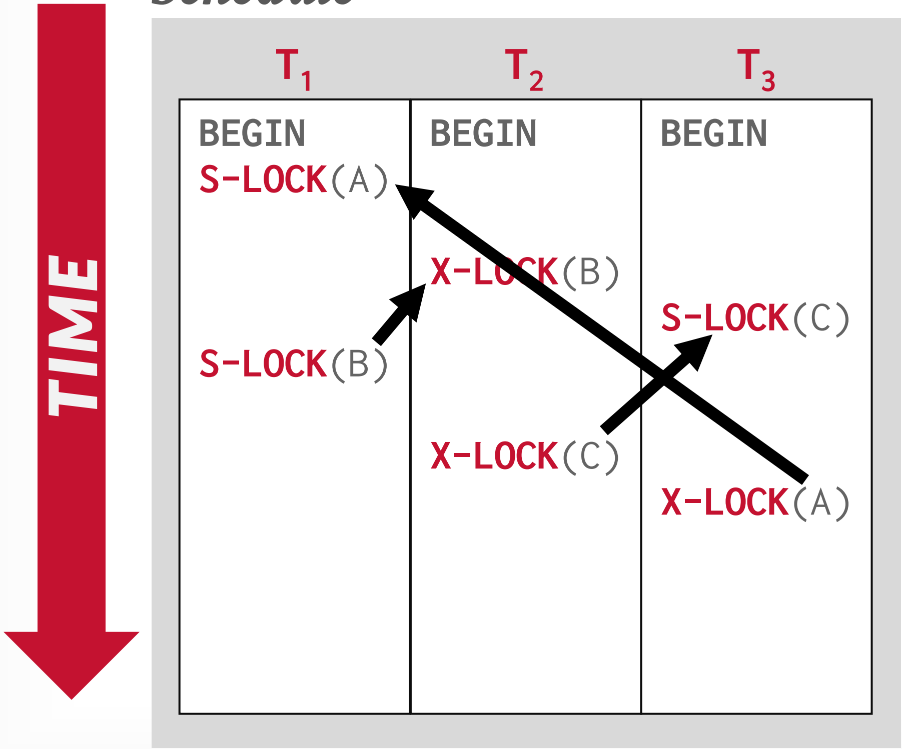
T_1	T_2	T_3
BEGIN	BEGIN	BEGIN
S-LOCK(A)		
	X-LOCK(B)	
		S-LOCK(C)
S-LOCK(B)		
	X-LOCK(C)	
		X-LOCK(A)

Waits-For Graph



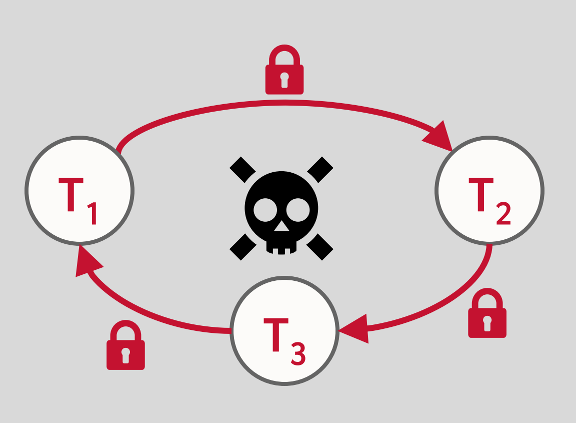
DEADLOCK DETECTION

Schedule



T_1	T_2	T_3
BEGIN	BEGIN	BEGIN
S-LOCK(A)		
	X-LOCK(B)	
		S-LOCK(C)
S-LOCK(B)		
	X-LOCK(C)	
		X-LOCK(A)

Waits-For Graph



DEADLOCK HANDLING

When the DBMS detects a deadlock, it will select a "victim" txn to rollback to break the cycle.

后台工作进程根据锁表和锁管理器运行死锁检测算法。

The victim txn will either restart or abort (more common) depending on how it was invoked.

There is a trade-off between the frequency of checking for deadlocks and how long txns wait before deadlocks are broken.

DEADLOCK HANDLING: VICTIM SELECTION

Selecting the proper victim depends on a lot of different variables….

→ By age (lowest timestamp) 新/老事务
→ By progress (least/most queries executed) 事务执行的查询数量
→ By the # of items already locked 已持有的锁数量
→ By the # of txns that we have to rollback with it 级联回滚的事务数量

We also should consider the # of times a txn has been restarted in the past to prevent starvation.

DEADLOCK HANDLING: ROLLBACK LENGTH

After selecting a victim txn to abort, the DBMS can also decide on how far to rollback the txn's changes.

Approach #1: Completely

→ Rollback entire txn and tell the application it was aborted.

Approach #2: Partial (Savepoints)

→ DBMS rolls back a portion of a txn (to break deadlock) and then attempts to re-execute the undone queries.

DEADLOCK PREVENTION

When a txn tries to acquire a lock that is held by another txn, the DBMS kills one of them to prevent a deadlock. 请求锁的事务等待或终止已持有锁的事务。

This approach does not require a ***waits-for*** graph or detection algorithm.

DEADLOCK PREVENTION

Assign priorities based on timestamps:

→ Older Timestamp = Higher Priority (e.g., $T_1 > T_2$)

Wait-Die ("Old Waits for Young")

→ If *requesting txn* has higher priority than *holding txn*, then *requesting txn* waits for *holding txn.*

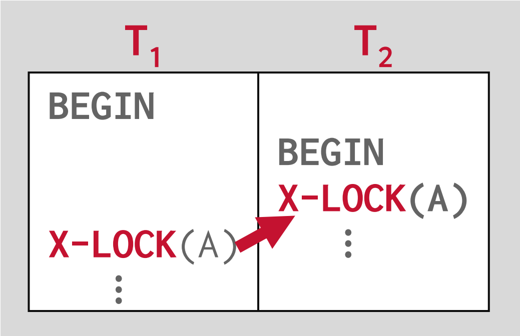
→ Otherwise *requesting txn* aborts.

Wound-Wait ("Young Waits for Old")

→ If *requesting txn* has higher priority than *holding txn*, then *holding txn* aborts and releases lock.

→ Otherwise *requesting txn* waits.

DEADLOCK PREVENTION



T_1	T_2
BEGIN	
	BEGIN
	X-LOCK(A)
X-LOCK(A)	⋮
⋮	

➡

Wait-Die

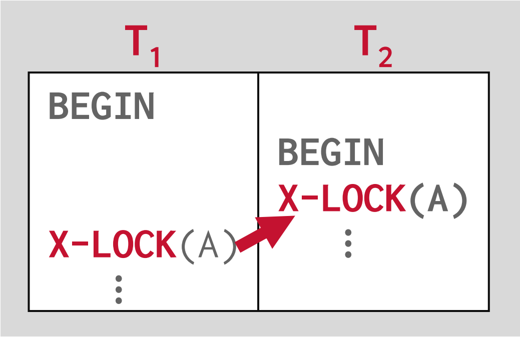
T_1 waits

Wound-Wait

T_2 aborts

T_1	T_2
BEGIN	
X-LOCK(A)	
⋮	
	BEGIN
	X-LOCK(A)
	⋮

DEADLOCK PREVENTION



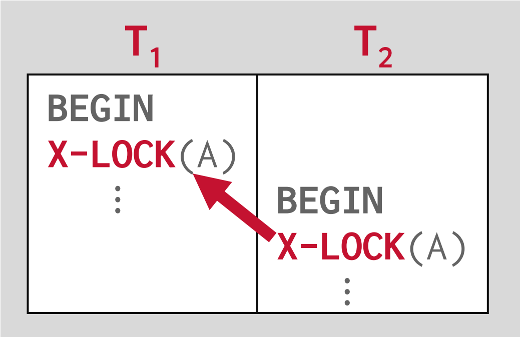
T_1	T_2
BEGIN	
	BEGIN
	X-LOCK(A)
	⋮
X-LOCK(A)	
⋮	

Wait-Die

T_1 waits

Wound-Wait

T_2 aborts



T_1	T_2
BEGIN	
X-LOCK(A)	
⋮	
	BEGIN
	X-LOCK(A)
	⋮

Wait-Die

T_2 aborts

Wound-Wait

T_2 waits

DEADLOCK PREVENTION

Why do these schemes guarantee no deadlocks?

Only one "type" of direction allowed when waiting for a lock. 获取锁的操作都朝着一个方向进行。

When a txn restarts, what is its (new) priority?

Its original timestamp to prevent it from getting starved for resources like an old man at a corrupt senior center.

OBSERVATION

All these examples have a one-to-one mapping from database objects to locks.

If a txn wants to update one billion tuples, then it must acquire one billion locks.

Acquiring locks is a more expensive operation than acquiring a latch even if that lock is available.

LOCK GRANULARITIES

When a txn wants to acquire a "lock", the DBMS can decide the granularity (i.e., scope) of that lock.
锁的粒度

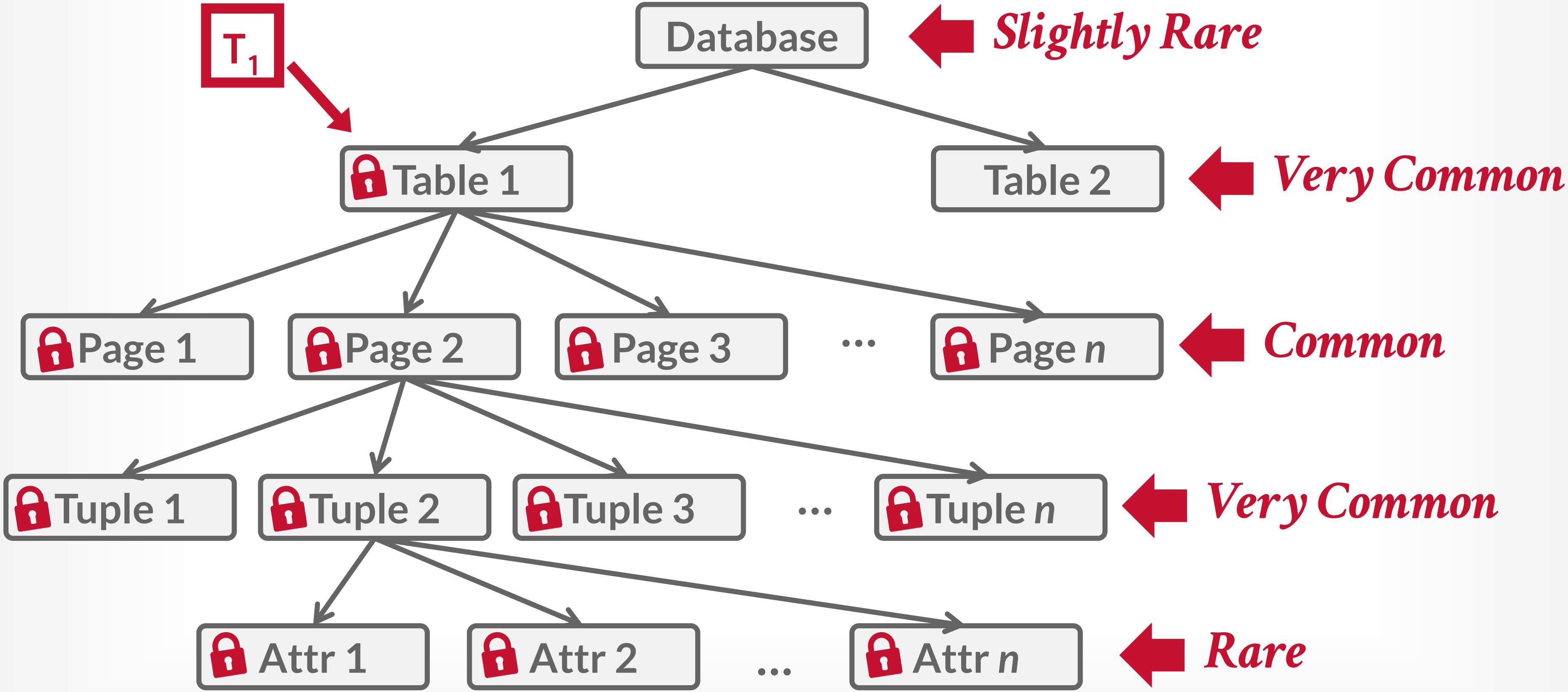
- → Attribute? Tuple? Page? Table?

The DBMS should ideally obtain fewest number of locks that a txn needs.

Trade-off between parallelism versus overhead.

- → Fewer Locks, Larger Granularity vs. More Locks, Smaller Granularity.

DATABASE LOCK HIERARCHY



INTENTION LOCKS

意向锁旨在解决多粒度锁系统中的效率问题。

An **intention lock** allows a higher-level node to be locked in **shared** or **exclusive** mode without having to check all descendent nodes.

意向锁允许高层节点以共享锁或排它锁被锁定，而不需要检查所有的低层节点的锁情况。

If a node is locked in an intention mode, then some txn is doing explicit locking at a lower level in the tree.

如果一个节点被加了意向锁，则表明某个事务在该节点的更低层级持有显式锁。

INTENTION LOCKS

Intention-Shared (IS)

- → Indicates explicit locking at lower level with S locks.
- → Intent to get S lock(s) at finer granularity.

Intention-Exclusive (IX)

- → Indicates explicit locking at lower level with X locks.
- → Intent to get X lock(s) at finer granularity.

Shared+Intention-Exclusive (SIX)

普通共享锁+意向排它锁

- → The subtree rooted by that node is locked explicitly in S mode and explicit locking is being done at a lower level with X locks.

COMPATIBILITY MATRIX

T_2 Wants

T_1 Holds

	IS	IX	S	SIX	X
IS	✓	✓	✓	✓	✗
IX	✓	✓	✗	✗	✗
S	✓	✗	✓	✗	✗
SIX	✓	✗	✗	✗	✗
X	✗	✗	✗	✗	✗

LOCKING PROTOCOL

Each txn obtains appropriate lock at highest level of the database hierarchy.

To get **S** or **IS** lock on a node, the txn must hold at least **IS** on parent node.

To get **X**, **IX**, or **SIX** on a node, must hold at least **IX** on parent node.

EXAMPLE

T_1 – Get the balance of Andy's off-shore bank account.

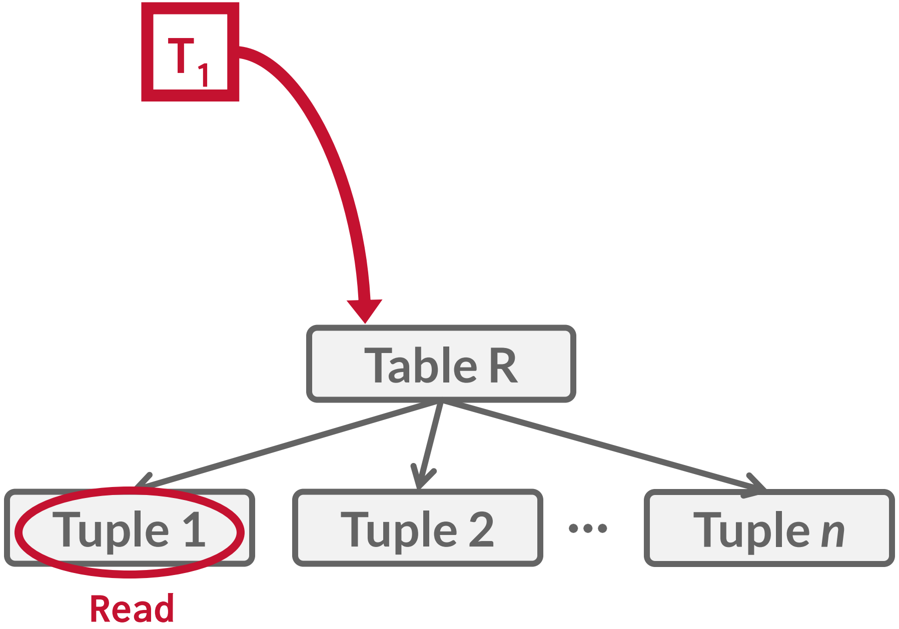
T_2 – Increase bookie's account balance by 1%.

What locks should these txns obtain?

→ **Exclusive** + **Shared** for leaf nodes of lock tree.
→ Special **Intention** locks for higher levels.

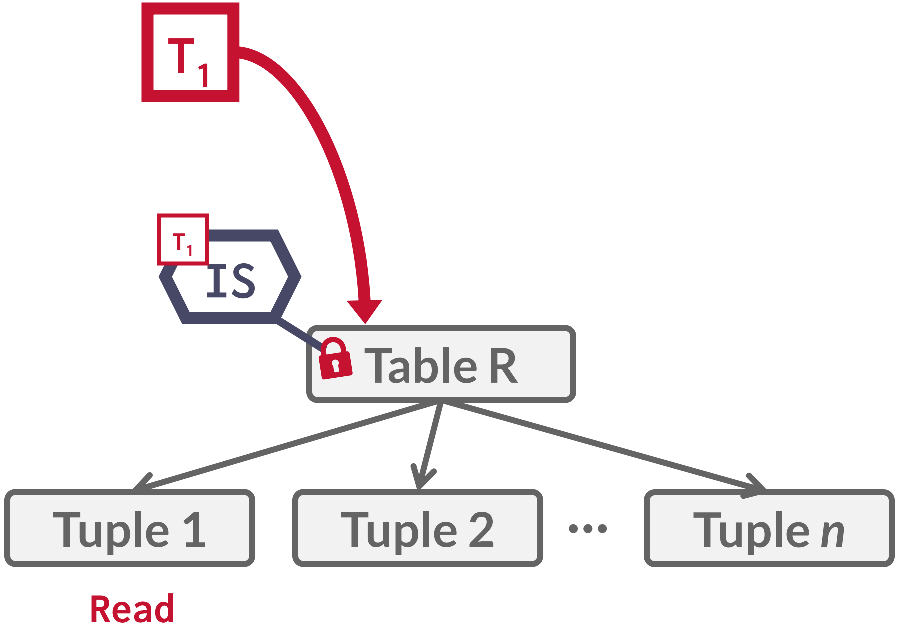
EXAMPLE – TWO-LEVEL HIERARCHY

Read Andy's record in **R**.



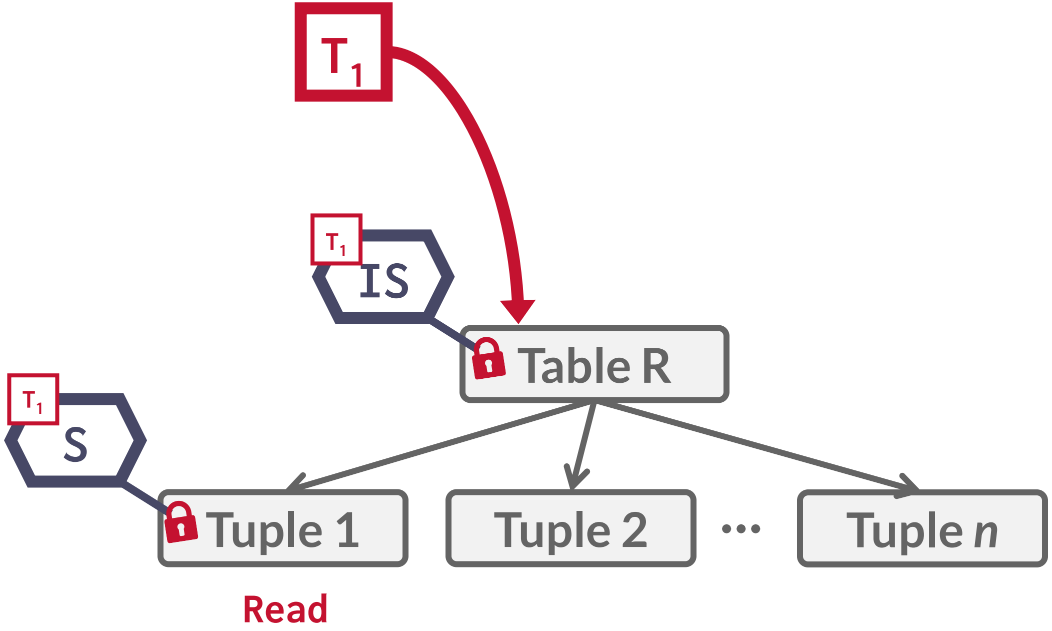
EXAMPLE – TWO-LEVEL HIERARCHY

Read Andy's record in **R**.



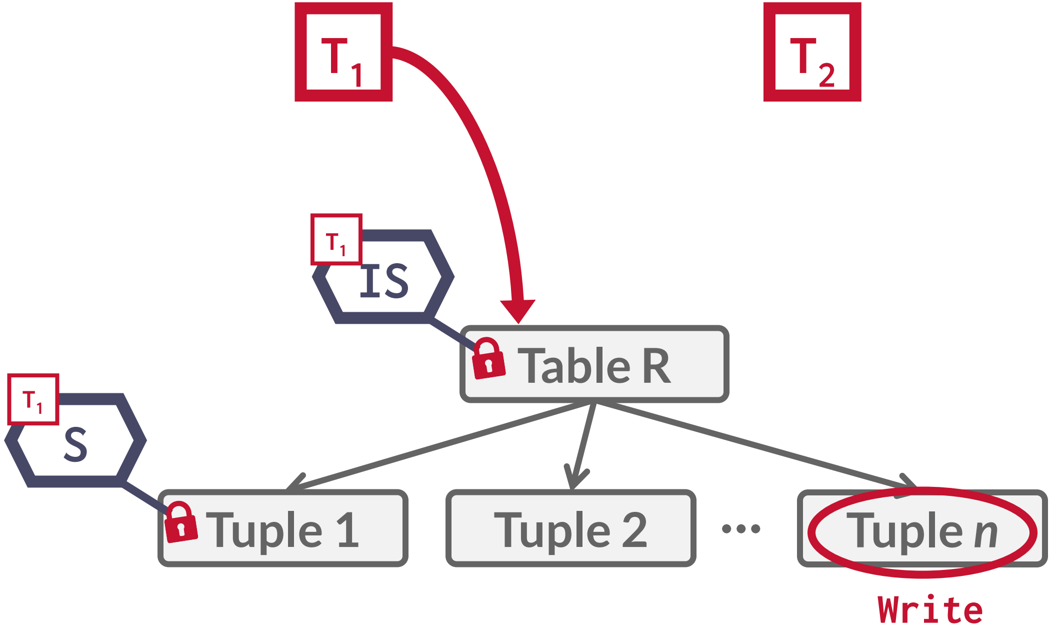
EXAMPLE – TWO-LEVEL HIERARCHY

Read Andy's record in **R**.



EXAMPLE – TWO-LEVEL HIERARCHY

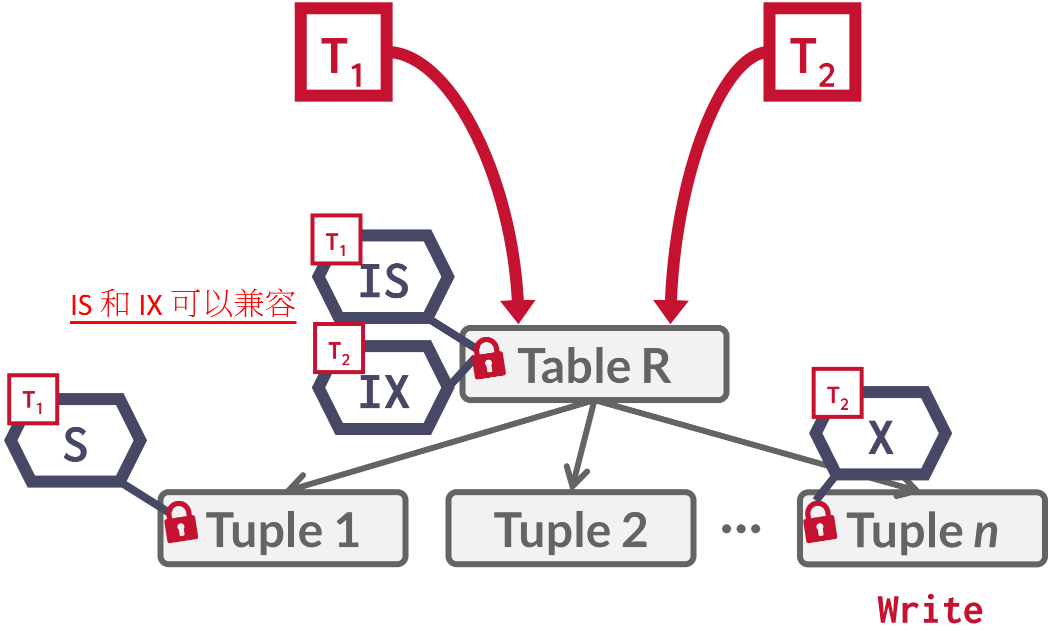
Update bookie's record in **R**.



	IS	IX	S	SIX	X
IS	✓	✓	✓	✓	×
IX	✓	✓	×	×	×
S	✓	×	✓	×	×
SIX	✓	×	×	×	×
X	×	×	×	×	×

EXAMPLE – TWO-LEVEL HIERARCHY

Update bookie's record in **R**.



IS 和 IX 可以兼容

	IS	IX	S	SIX	X
IS	✓	✓	✓	✓	×
IX	✓	✓	×	×	×
S	✓	×	✓	×	×
SIX	✓	×	×	×	×
X	×	×	×	×	×

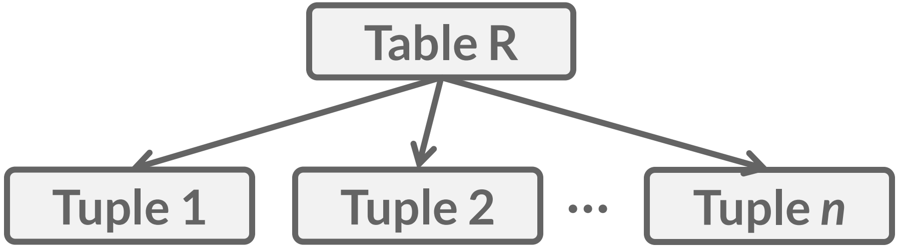
EXAMPLE – THREE TXNS

Assume three txns execute at same time:

→ **T_1** – Scan all tuples in **R** and update one tuple.

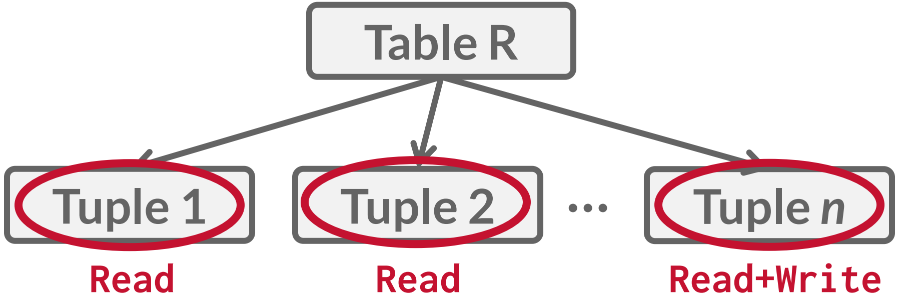
→ **T_2** – Read a single tuple in **R**.

→ **T_3** – Scan all tuples in **R**.



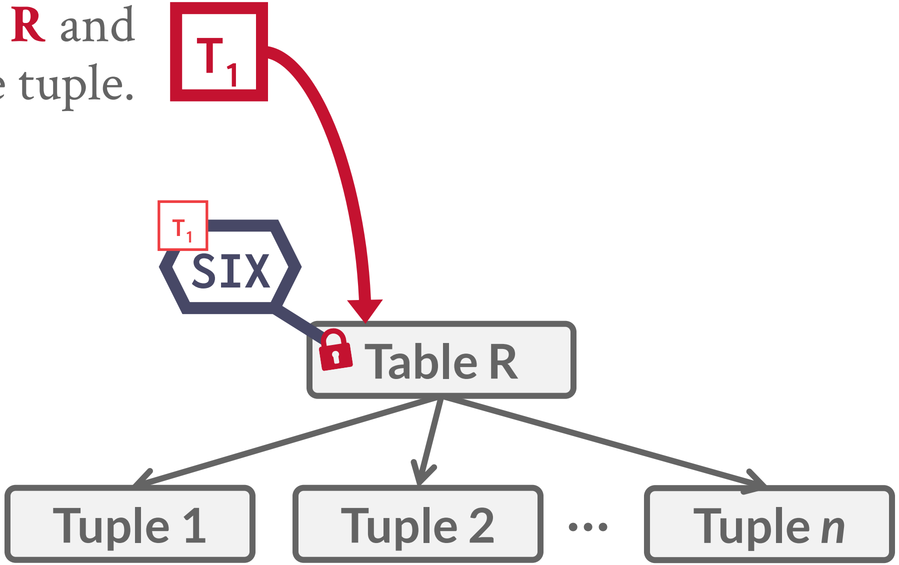
EXAMPLE – THREE TXNS

Scan all tuples in **R** and update one tuple. T_1



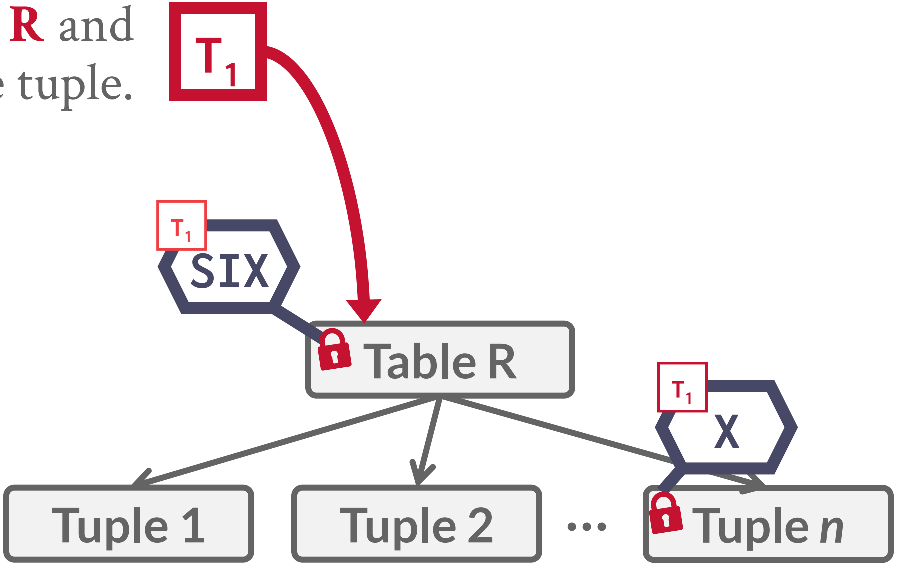
EXAMPLE – THREE TXNS

Scan all tuples in **R** and update one tuple.

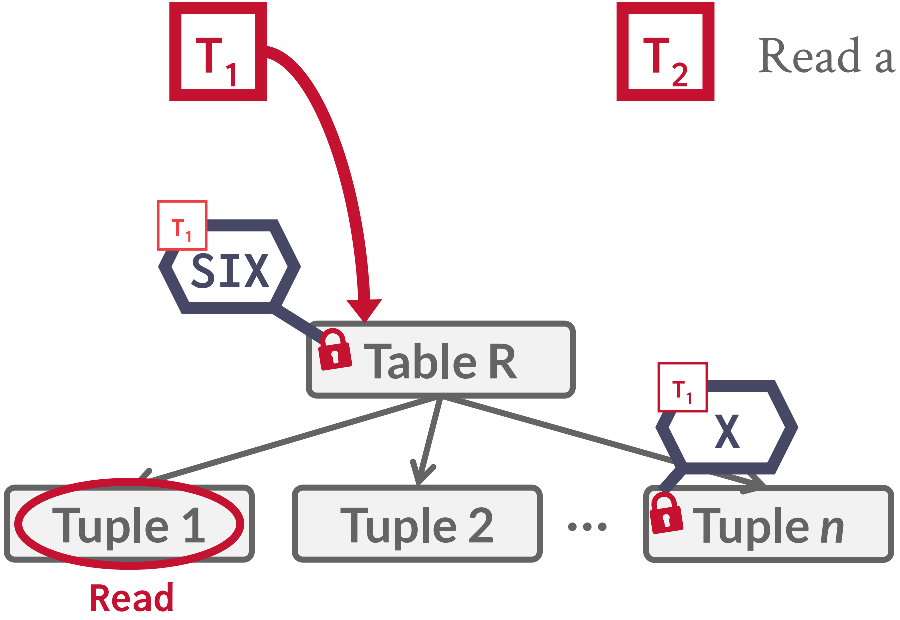


EXAMPLE – THREE TXNS

Scan all tuples in **R** and update one tuple.

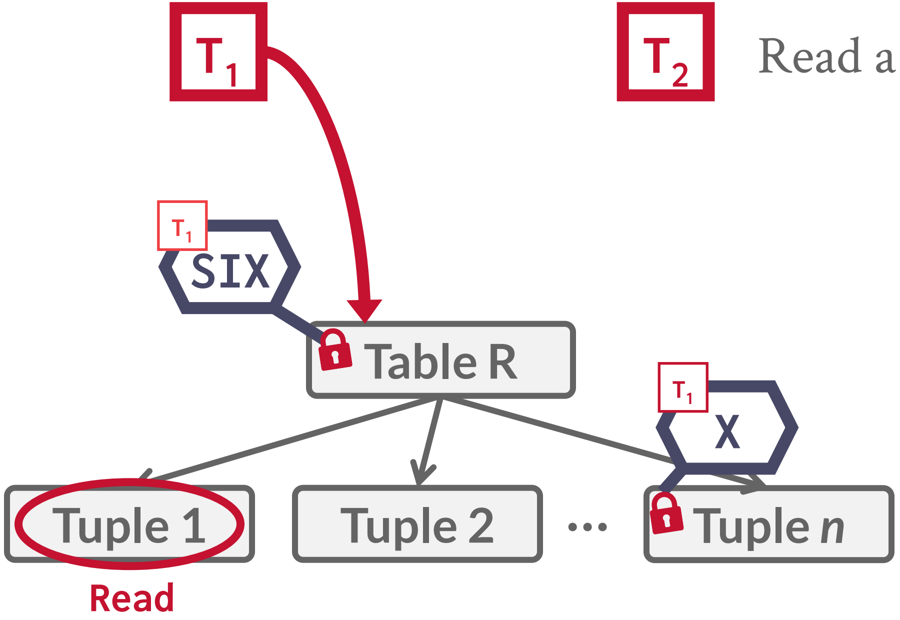


EXAMPLE – THREE TXNS



T_2 Read a single tuple in **R**.

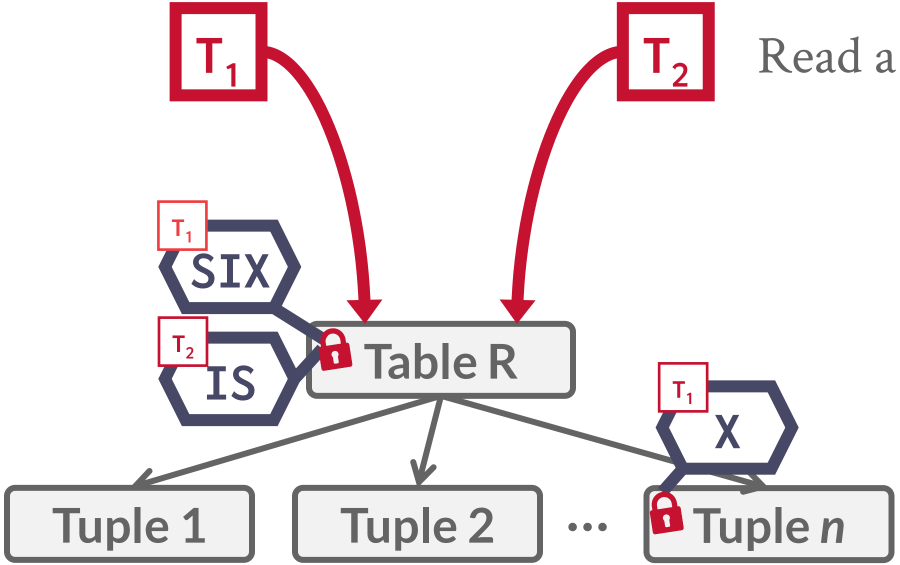
EXAMPLE – THREE TXNS



T_2 Read a single tuple in **R**.

	IS	IX	S	SIX	X
IS	✓	✓	✓	✓	×
IX	✓	✓	×	×	×
S	✓	×	✓	×	×
SIX	✓	×	×	×	×
X	×	×	×	×	×

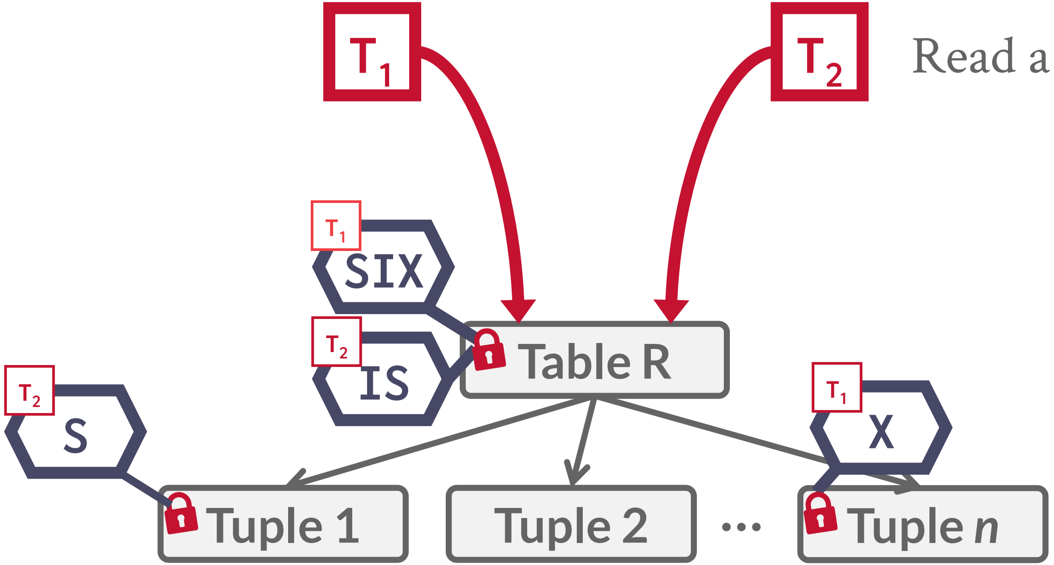
EXAMPLE – THREE TXNS



Read a single tuple in **R**.

	IS	IX	S	SIX	X
IS	✓	✓	✓	✓	×
IX	✓	✓	×	×	×
S	✓	×	✓	×	×
SIX	✓	×	×	×	×
X	×	×	×	×	×

EXAMPLE – THREE TXNS

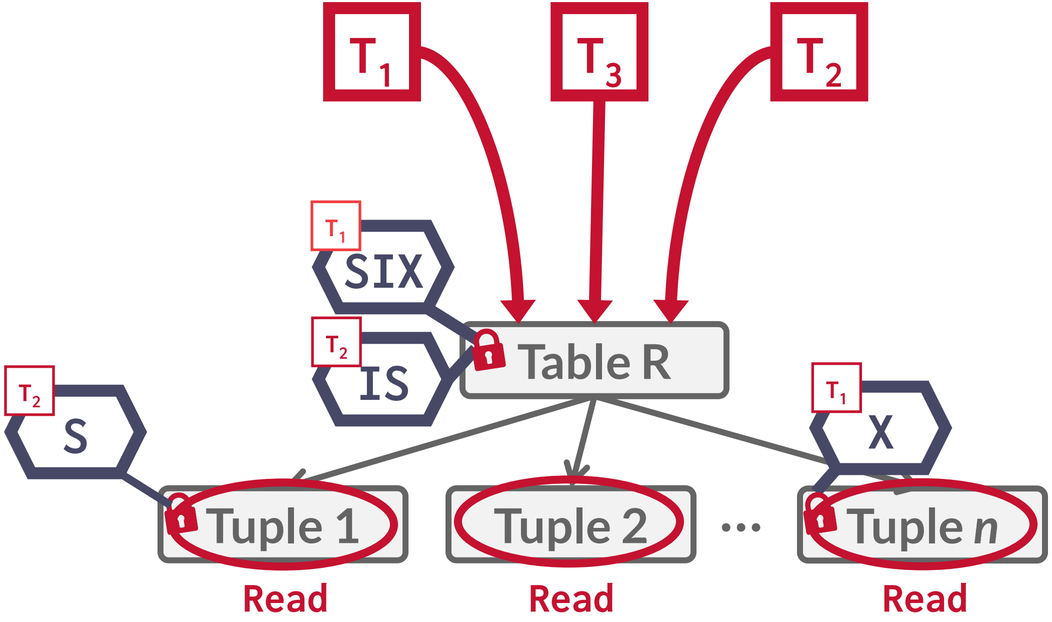


Read a single tuple in **R**.

	IS	IX	S	SIX	X
IS	✓	✓	✓	✓	×
IX	✓	✓	×	×	×
S	✓	×	✓	×	×
SIX	✓	×	×	×	×
X	×	×	×	×	×

EXAMPLE – THREE TXNS

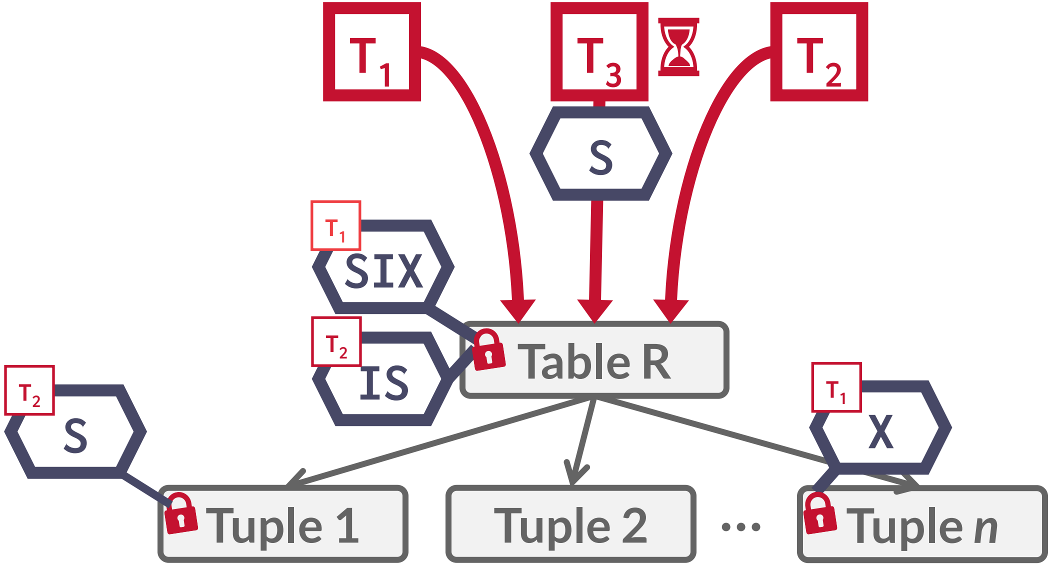
Scan all tuples in **R**.



	IS	IX	S	SIX	X
IS	✓	✓	✓	✓	×
IX	✓	✓	×	×	×
S	✓	×	✓	×	×
SIX	✓	×	×	×	×
X	×	×	×	×	×

EXAMPLE – THREE TXNS

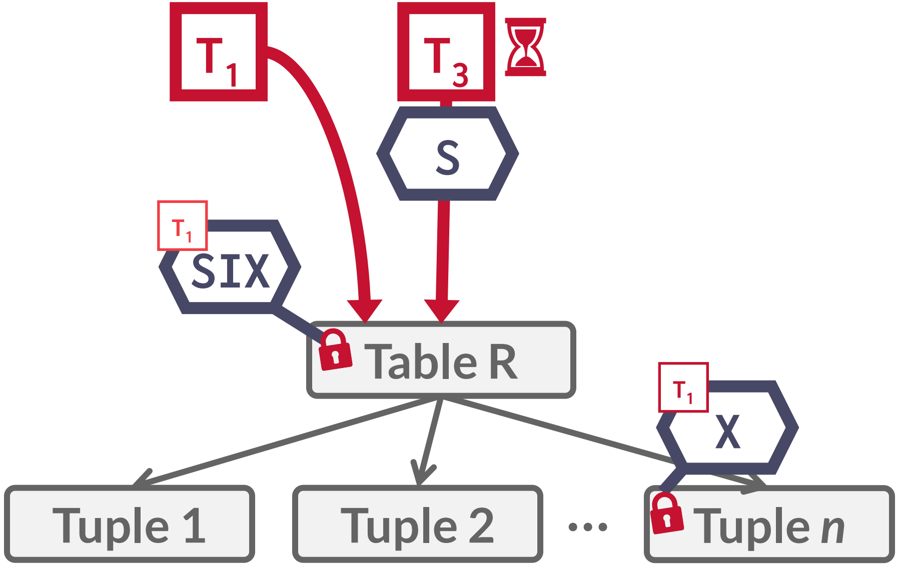
Scan all tuples in **R**.



	IS	IX	S	SIX	X
IS	✓	✓	✓	✓	×
IX	✓	✓	×	×	×
S	✓	×	✓	×	×
SIX	✓	×	×	×	×
X	×	×	×	×	×

EXAMPLE – THREE TXNS

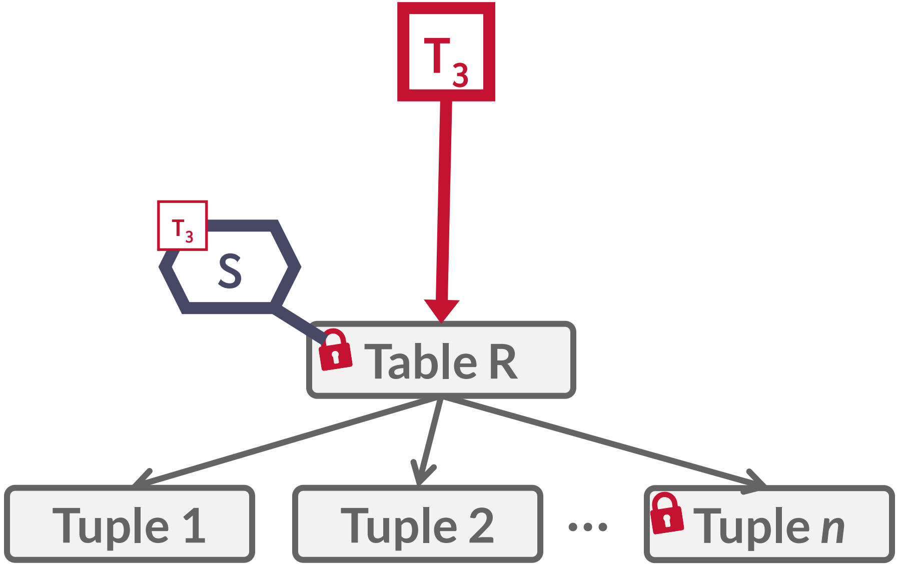
Scan all tuples in **R**.



	IS	IX	S	SIX	X
IS	✓	✓	✓	✓	×
IX	✓	✓	×	×	×
S	✓	×	✓	×	×
SIX	✓	×	×	×	×
X	×	×	×	×	×

EXAMPLE – THREE TXNS

Scan all tuples in **R**.



	IS	IX	S	SIX	X
IS	✓	✓	✓	✓	×
IX	✓	✓	×	×	×
S	✓	×	✓	×	×
SIX	✓	×	×	×	×
X	×	×	×	×	×

LOCK ESCALATION

锁升级：将多个细粒度锁升级为粗粒度锁。

The DBMS can automatically switch to coarser-grained locks when a txn acquires too many low-level locks.

This reduces the number of requests that the lock manager must process.

LOCKING IN PRACTICE

Applications typically do not acquire a txn's locks manually (i.e., explicit SQL commands).

Sometimes you need to provide the DBMS with hints to help it to improve concurrency.

→ Update a tuple after reading it.
→ Skip any tuple that is locked.

Explicit locks are also useful when doing major changes to the database.

SELECT...FOR UPDATE

Perform a **SELECT** and then sets an exclusive lock on the matching tuples.

Can also set shared locks:

→ Postgres: **FOR SHARE**

→ MySQL: **LOCK IN SHARE MODE**

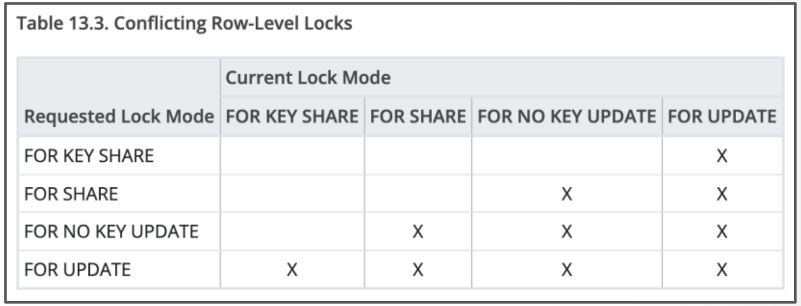


Table 13.3. Conflicting Row-Level Locks

	Current Lock Mode			
Requested Lock Mode	FOR KEY SHARE	FOR SHARE	FOR NO KEY UPDATE	FOR UPDATE
FOR KEY SHARE				X
FOR SHARE			X	X
FOR NO KEY UPDATE		X	X	X
FOR UPDATE	X	X	X	X

```
SELECT * FROM <table>
 WHERE <qualification> FOR UPDATE;
```

普通的 SELECT 语句将获取 S 锁，但增加 FOR UPDATE 表示申请 X 锁，后续将修改数据。

SELECT...SKIP LOCKED

Perform a **SELECT** and automatically ignore any tuples that are already locked in an incompatible mode.

→ Useful for maintaining queues inside of a DBMS.

```
SELECT * FROM <table>
 WHERE <qualification> SKIP LOCKED;
```

CONCLUSION

2PL is used in almost every DBMS.

Automatically generates correct interleaving:
→ Locks + protocol (2PL, SS2PL ...)
→ Deadlock detection + handling
→ Deadlock prevention

Many more things not discussed…
→ Nested Transactions
→ Savepoints

PROJECT #3 – QUERY EXECUTION

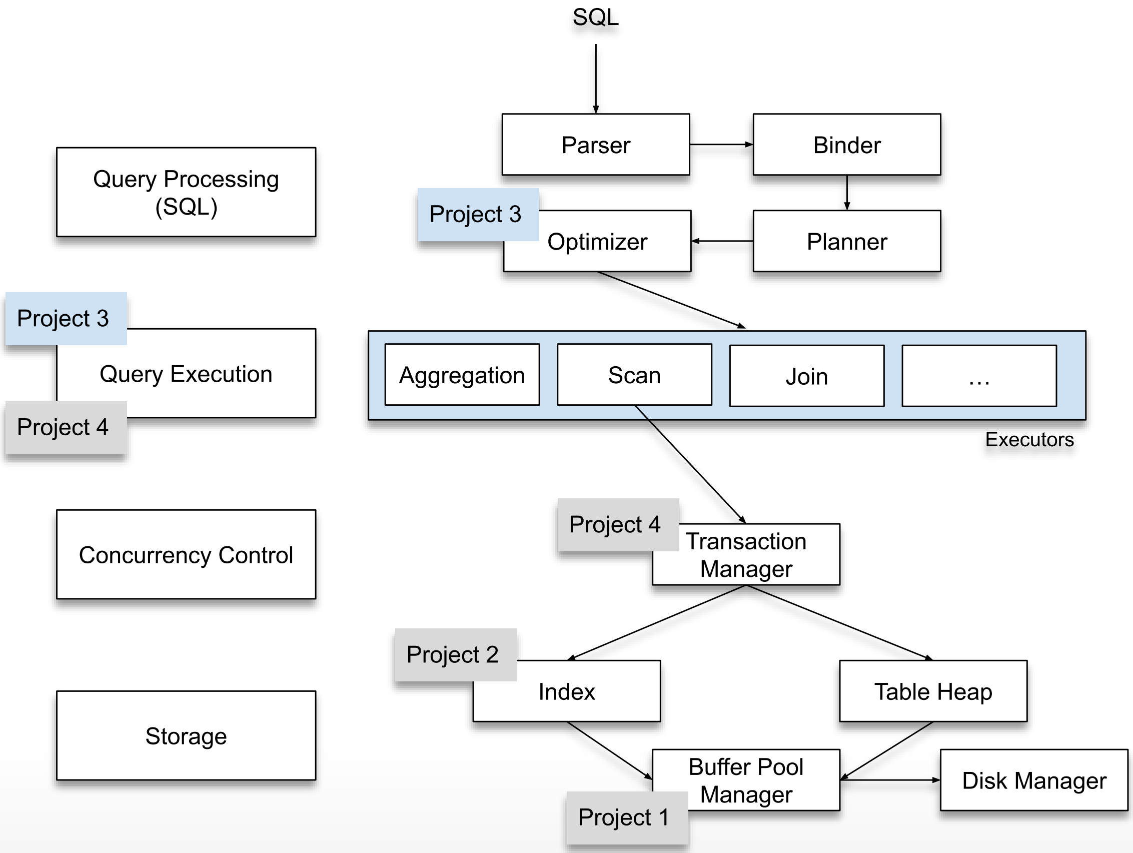
You will add support for optimizing and executing queries in BusTub.

BusTub supports (basic) SQL with a rule-based optimizer for converting AST into physical plans.

***Prompt**: Red and white bus driving on a road through mountains. The bus has a bubbly bath as its roof with a showerhead. The bus has crazy eyes and a mouth in the front. There are bold red letters that say "CMU" on the side of the bus*

https://15445.courses.cs.cmu.edu/fall2024/project3

PROJECT #3 – QUERY EXECUTION



PROJECT #3 – TASKS

Plan Node Executors

- → Access Methods: Sequential Scan, Index Scan
- → Modifications: Insert, Delete, Update
- → Joins: Nested-Loop, Index Nested-Loop Hash Join
- → Miscellaneous: External Merge-Sort, Limit

Optimizer Rule:

- → Convert Nested Loops to Hash Join
- → Convert Sequential Scan to Index Scan

PROJECT #3 - LEADERBOARD

The leaderboard requires you to add additional rules to the optimizer to generate query plans.

→ It will be impossible to get a top ranking by just having the fastest implementations in Project #1 + Project #2.

Tasks:

→ Column Pruning
→ More Aggressive Predicate Pushdown
→ Bloom Filter for Hash Join

DEVELOPMENT HINTS

Implement the **Insert** and **Sequential Scan** executors first so that you can populate tables and read from it.

You do not need to worry about transactions.

The aggregation hash table does not need to be backed by your buffer pool (i.e., use STL)

Gradescope is for meant for grading, not debugging. Write your own local tests.

THINGS TO NOTE

Do **not** change any file other than the ones that you submit to Gradescope.

Make sure you pull in the latest changes from the BusTub main branch.

Post your questions on Piazza or come to TA office hours.

Compare against our solution in your browser!

PLAGIARISM WARNING

The homework and projects must be your own original work. They are **not** group assignments.

You may **not** copy source code from other people or the web.

Plagiarism is **not** tolerated. You will get lit up.

→ Please ask me if you are unsure.

See CMU's Policy on Academic Integrity for additional information.

NEXT CLASS

Timestamp Ordering Concurrency Control

Isolation Levels