

# CAP Off: Local Reads and Linearizable Asynchronous Replication

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Today’s cloud services rely on high-performance distributed datastores for storing and accessing their data. These services are often characterized by read-dominant accesses and numerous concurrent requests [1, 5]. Thus, datastores must provide high throughput to meet the performance demands of online services while offering high availability, despite being deployed on failure-prone commodity hardware [6].

Data replication is a fundamental feature of high performance and resilient datastores. Data must be replicated across multiple servers to increase throughput because a single server often cannot keep up with the request load [4]. Replication is also necessary to guarantee that a crash of a server does not render the dataset inaccessible.

**Motivation.** Keeping the replicas consistent, to ensure that the services running on the datastore operate correctly, is a challenge, especially in the presence of crashes. A crash-tolerant replication protocol is responsible for keeping the replicas of a datastore consistent – even when crashes occur – by determining the necessary actions to execute reads and writes. Several crash-tolerant protocols favor performance by relaxing consistency (*RC protocols*). As such, their reads may return stale values leading to nasty surprises for both clients and developers [15, 17]. There exist however, protocols that offer strongly consistent (i.e., linearizable) reads, which are more desirable for correctness and programmability.

Crash-tolerant protocols can be synchronous or asynchronous based on the model they rely on to ensure consistency. Synchronous protocols, which depend on known bounded processing and communication delays, are easier to design. However, in the real world, distributed datastores are deployed over complex software stacks and virtualization layers [2, 14]. Consequently, the network and compute nodes of a distributed datastore experience asynchrony and other timing anomalies, which may lead to timing violations and compromise the safety of synchronous protocols. To tolerate such timing violations, safer protocols adopt the *asynchronous* model where there are no timing assumptions, implying that processing and communication delays can be arbitrary.

Existing crash-tolerant replication protocols that afford linearizable reads fall into two categories; local reads under synchrony (*LS protocols*), and remote reads under asynchrony (*RA protocols*). LS protocols offer cheap linearizable reads that are served locally on a single replica, without inter-replica communication, as in Hermes [12]. However, these protocols assume a synchronous model [7] (e.g., to exploit leases). In contrast, RA protocols, such as Raft [16] and Paxos [13], are safe under asynchrony, but for each read, they mandate

costly inter-replica coordination and the involvement of remote replicas. As Schwarzmann and Hadjistasi advocate [10], it is important to study the feasibility of crash-tolerant protocols that support linearizable local reads under asynchrony.

There exist fundamental theoretical results related to asynchronous replication, including the FLP result [8] and the CAP theorem [3, 9], but neither suffice to answer the above as both fall short in examining the performance of reads.

**Theory.** To address the existing gap in understanding, this work sheds light on a fundamental three-way tradeoff of crash-tolerant protocols, revealing a tension between consistency, performance, and the timing assumptions of a system. Briefly, we present the L<sup>2</sup>AW theorem, an impossibility result which asserts that *in any Linearizable Asynchronous read/write register implementation that tolerates even a single crash (Without blocking reads or writes), no reads are Local*.

We observe that the performance aspect of this tradeoff affects the latency but not necessarily the throughput of reads. Thus, asynchronous linearizable reads need not be as costly as in existing RA protocols, where each read incurs network and computation costs to remote replicas.

**Practice.** Capitalizing on the above insight, we introduce almost-local reads (ALRs), a scheme that affords low-cost reads in a linearizable and crash-tolerant manner under asynchrony. ALRs exploit the high volume of concurrent reads in online services by leveraging batching, but with a twist. Unlike traditional batching, all reads in an ALR-batch are executed against the local replica of a server, and only a lightweight *sync* operation per batch involves remote replicas. The sync incurs only a small network and computation cost regardless of the batch size. Moreover, the sync can sometimes be elided as existing writes can act as implicit syncs. As a result, ALRs incur little or no extra network and processing costs to remote replicas, thus achieving the performance of local reads while offering linearizability under asynchrony.

When applied to the protocols in all three corners of the design space (i.e., RA, RC, LS), ALRs add the missing piece: ❶ they improve the throughput for RA protocols (e.g., Raft) ❷ ensure linearizability for RC protocols (e.g., ZAB [11]) and ❸ they allow LS protocols (e.g., Hermes) to operate under asynchrony. Our experiments show that ALR-enhanced variants of ZAB and Hermes protocols on 95% reads come within 2% and 5% of their original throughput, respectively, while also ensuring linearizability under asynchrony. Finally, ALRs yield over 2.5× higher throughput on Raft in read-intensive workloads without sacrificing consistency or asynchrony.

<sup>\*</sup>This work occurred when the authors were at the University of Edinburgh.

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