

Chapter 7.1

并查集

The Disjoint Set ADT

Equivalence Class 等价类

1) Definition of Equivalence Class:

Suppose we have a set $U=\{1,2,\dots,n\}$ of n elements and a set $R=\{(i_1,j_1), (i_2,j_2), \dots, (i_r,j_r)\}$ of r relations. The relation R is an equivalence relation iff the following conditions are true (symbol ' \equiv ' represent the equivalence relation on sets, x,y,z are elements in set) :

- Reflexive $x \equiv x$. 自反
- Symmetric $x \equiv y, y \equiv x$ 对称
- Transitive $x \equiv y$ and $y \equiv z$, then $x \equiv z$ 传递

Equivalence Class

例如:

判别3个数a, b, c能否构成三角形的三条边?

能构成三角形的等价类:

$\{ (3,4,5), (4,5,6), (5,6,7), \dots \}$

不能构成三角形的等价类:

$\{ (1,2,3), (2,3,5), \dots \}$

Equivalence Class

2) Example:

set $s = \{0, 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11\}$

pairs of equivalence:

$(0\ 4), (3\ 1), (6\ 10), (8\ 9), (7\ 4), (6\ 8), (3\ 5), (2\ 11),$
 $(11\ 0)$

等价关系

Initial: $\{0\}, \{1\}, \{2\}, \{3\}, \{4\}, \{5\}, \{6\}, \{7\}, \{8\}, \{9\}, \{10\}, \{11\}$ \Rightarrow 合并等价类
构造并查集

$0 \equiv 4$ $\{0, 4\}, \{1\}, \{2\}, \{3\}, \{5\}, \{6\}, \{7\}, \{8\}, \{9\},$
 $\{10\}, \{11\}$ 等价类合并

$3 \equiv 1$ $\{0, 4\}, \{1, 3\}, \{2\}, \{5\}, \{6\}, \{7\}, \{8\}, \{9\},$
 $\{10\}, \{11\}$

Equivalence Class

$6 \equiv 10$ {0, 4}, {1, 3}, {2}, {5}, {6, 10}, {7}, {8},
{9}, {11}

$8 \equiv 9$ {0, 4}, {1, 3}, {2}, {5}, {6, 10}, {7},
{8, 9}, {11}

$7 \equiv 4$ ^{去重} {0, 4, 7}, {1, 3}, {2}, {5}, {6, 10}, {8, 9}, {11}

$6 \equiv 8$ {0, 4, 7}, {1, 3}, {2}, {5}, {6, 8, 9, 10}, {11}

$3 \equiv 5$ {0, 4, 7}, {1, 3, 5}, {2}, {6, 8, 9, 10}, {11}

$2 \equiv 11$ {0, 4, 7}, {1, 3, 5}, {2, 11}, {6, 8, 9, 10}

$11 \equiv 0$ {0, 4, 7, 2, 11}, {1, 3, 5}, {6, 8, 9, 10}

Equivalence Class

- 3) Online equivalence class operation 逻辑上允许的操作
- **Combine(a,b) : combine the equivalence classes that contains elements a and b into a single class**
 - **Find(e) : determine the class that currently contains element e.**



Combine(a,b) is equivalent to

i=Find(a); j=Find(b); if(i!=j) Union(i,j);

物理层

Equivalence Class

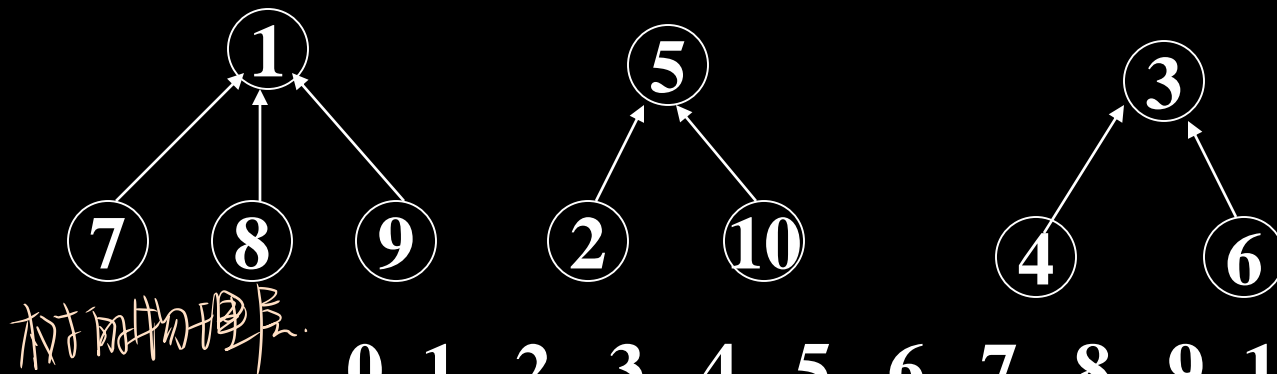
数组 \rightarrow 森林 \rightarrow 并查集

4) Tree Representation(Union-Find sets)

Example:

$S_1=\{1,7,8,9\}$, $S_2=\{5,2,10\}$, $S_3=\{3,4,6\}$, they all belong to $S=\{1,2,3,\dots,10\}$

整个并查集是一个森林



树的物理层

parent

	0	1	2	3	4	5	6	7	8	9	10
parent	///	0	5	0	3	0	3	1	1	1	5

双亲表示法

父结点
(根是0)

Equivalence Class

simple tree solution to union-find problem

① void Initialize(int n)

```
{ parent=new int[n+1];  
  for(int e=1;e<=n;e++)  
    parent[e]=0;  
}
```

初始化时是n个树(每个树有根).

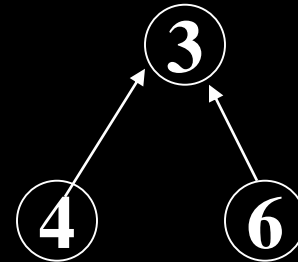
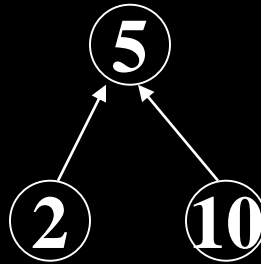
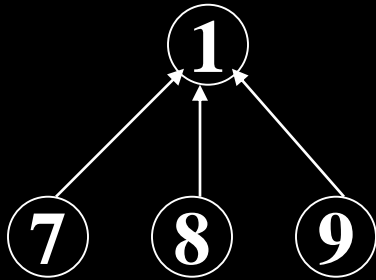
② int Find(int e) 返回树根标号

```
{ while(parent[e] 不是根)  
    e=parent[e];  
  return e;  
}
```

③ void Union(int i, int j)

```
{ parent[j]=i;  
}
```

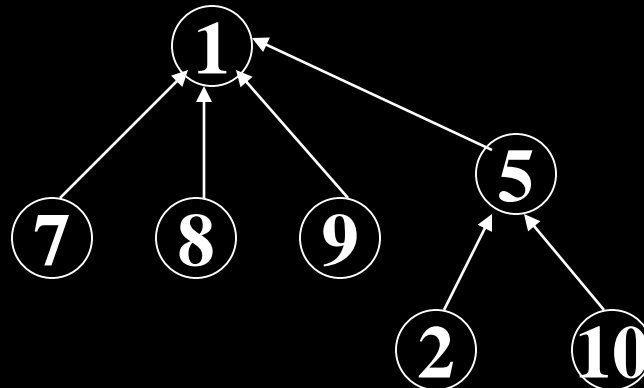

Equivalence Class



parent

	0	1	2	3	4	5	6	7	8	9	10
		0	5	0	3	10	3	1	1	1	5

Union (1,5)



Equivalence Class

Java

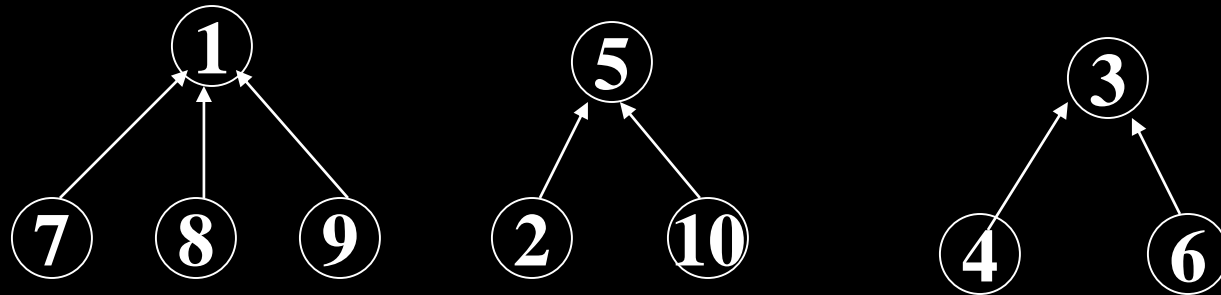
```
public class DisjSets
{ public DisjSets( int numElements )
    public void union( int root1, int root2 )
    public int find( int x )
    private int [ ] s;
}
```

```
public DisjSets( int numElements )
{ s = new int [ numElements ];
    for( int i = 0; i < s.length; i++ )
        s[ i ] = -1; //一个根结点
}
```

Equivalence Class

```
public void union( int root1, int root2 )  
{ s[ root2 ] = root1;  
}
```

```
public int find( int x )  
{ if( s[x] < 0 ) 是根  
    return x;  
    else  
        return find( s[ x ] );  
}
```



	0	1	2	3	4	5	6	7	8	9	10
S:		-1	5	-1	3	-1	3	1	1	1	5

Equivalence Class

5) Performance Evaluation

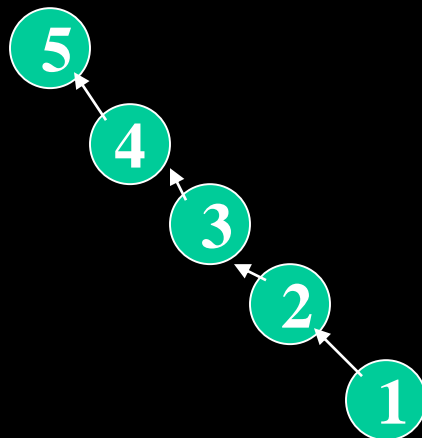
Time complexity: Find-- $O(h)$,

Union-- $\theta(1)$

Assume that u times unions and f times finds are to be performed, $f > u$,

in the worst case a tree with m elements can have a height of m :

Union(2,1), Union(3,2), Union(4,3), Union(5,4)...



Equivalence Class

★improve Union

two rules:

- **Weight rule:** ^{数重量} if the number of nodes in tree i is less than the number in tree j, then make j the parent of i; otherwise, make i the parent of j.
- **Height rule:** ^{高度} if the height of tree i is less than that of tree j, then make j the parent of i; otherwise, make i the parent of j.

目标: 树的高度尽量低

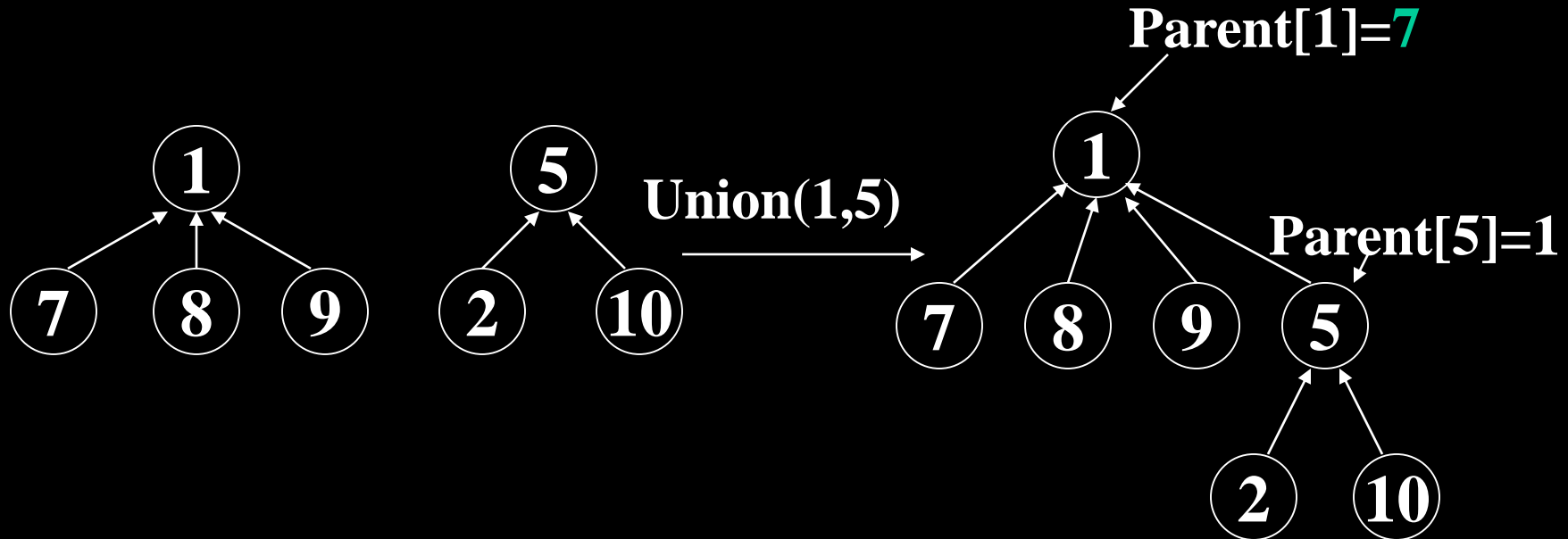
Equivalence Class

6) Performance Enhancement 2个优化.

- ★improve Union in order to decrease the time each find take, so that the height of tree will not increase linearly .
- ★.Improvement of Find –path compression

Equivalence Class

Let's discuss the weight rule(C++) :



Besides the *parent* field, each node has a boolean field *root*. The *root* field is true iff the node is presently a root node. The *parent* field of each root node is used to keep a count of the total number of nodes in the tree.

Equivalence Class

Union with the weight rule

```
void Initialize(int n)
{ root=new bool[n+1]; 是否是根节点
  parent=new int[n+1]; 记录父节点.
  for(int e=1;e<=n;e++)
  { parent[e]=1;
    root[e]=true;
  }
}

int Find(int e)
{ while(!root[e])
  { e=parent[e];
  }
  return e;
}
```


Equivalence Class

```
void Union(int i, int j)
{ if(parent[i]<parent[j]) // i becomes subtree of j
  { parent[j]=parent[j]+parent[i];
    root[i]=false;
    parent[i]=j;  让父节点是j
  }
  else {  让i是父节点
    parent[i]=parent[i]+parent[j];
    root[j]=false;
    parent[j]=i;
  }
}
```

Equivalence Class

Java(高度规则) 不用 root 标志位

用一个数组来实现，根结点中放负数，而且是代表高度。绝对值

不是根，则记录父节点下标。

```
public void union( int root1, int root2 )
```

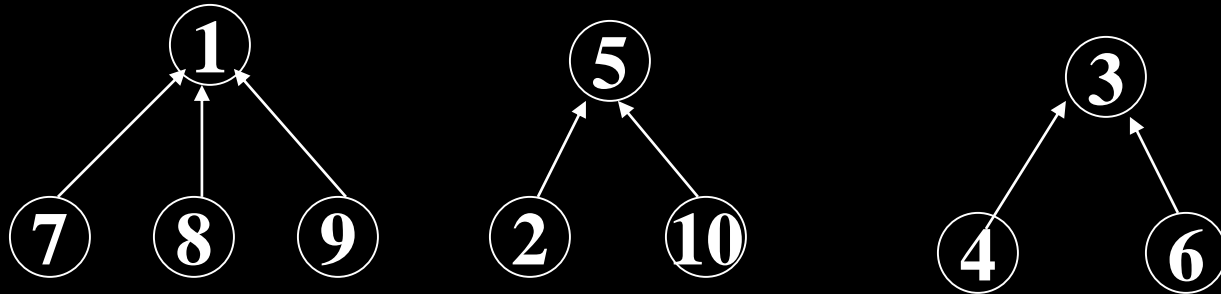
```
{ if( s[ root2 ] < s[ root1 ] )
```

```
    s[ root1 ] = root2; 且 s[root2] 不变
```

```
    else { if( s[ root1 ] == s[ root2 ] )
```

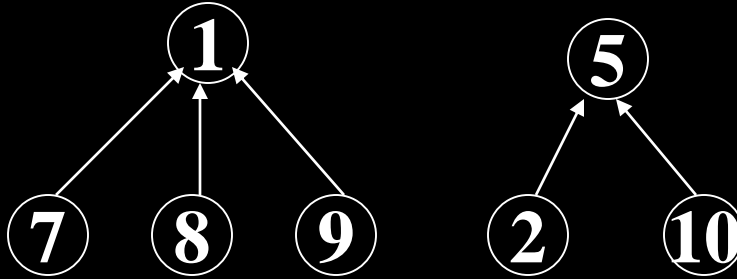
```
        s[ root1 ]--;
```

```
        s[ root2 ] = root1; } }
```



	0	1	2	3	4	5	6	7	8	9	10
S:		-2	5	-2	3	-2	3	1	1	1	5

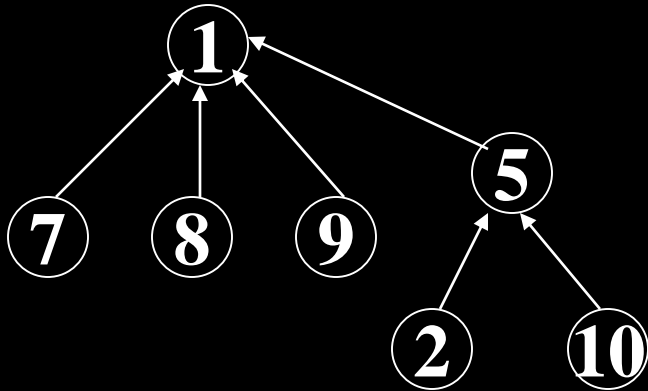
Equivalence Class



union (1, 5)

s[root1]--

s[root2] = root1



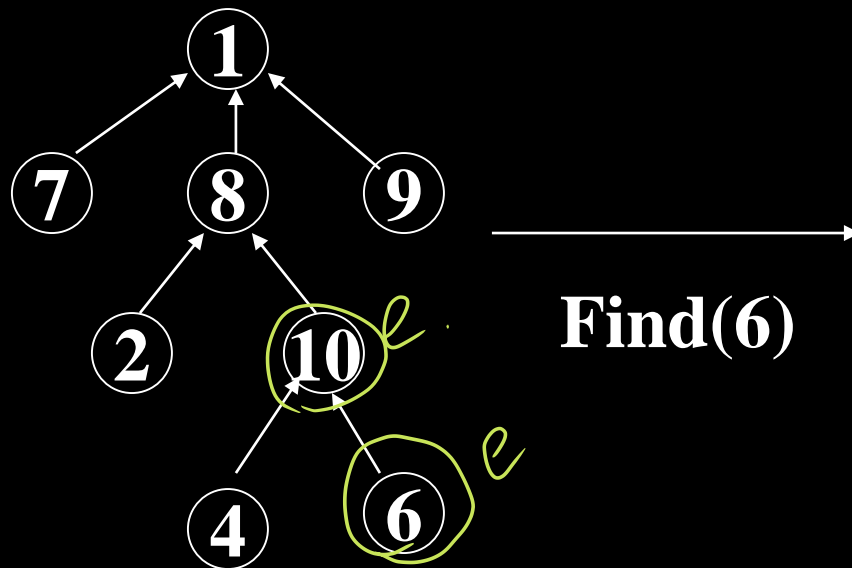
Equivalence Class

★.Improvement of *Find* –path compression

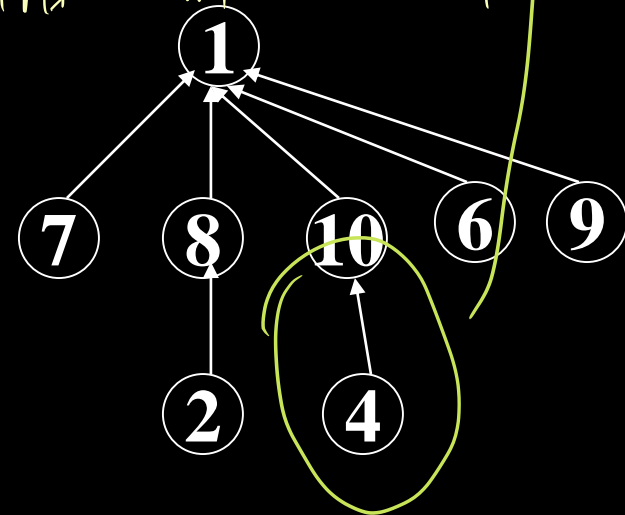
When processing a equivalence pair, we need to operate *Find* twice, *WeightUnion* once.

Example of improvement:

访问过的节点挂到根上
不增加复杂度且树高度降低



Find(6)



Equivalence Class

走两遍

```
int Find( int e) { /* C++ */ }
```

```
{ int j=e;
```

找根

```
while(!root[j]) j=parent[j];
```

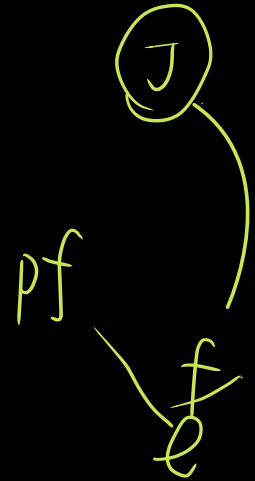
```
int f=e;
```

```
while(f!=j)
```

```
{ int pf=parent[f]; parent[f]=j; f=pf; }
```

```
}
```

挂到根上



Equivalence Class

Java

```
public int find( int x )
```

```
{   if( s[ x ] < 0 )
```

```
    return x;
```

```
    else
```

```
        return s[ x ] = find( s[ x ] );
```

```
}
```

返回根 并把根赋给 s[x]

递归