Moments & Deviations

3.1 Markov's Inequality

Slide 1

Markov's Inequality

Some useful bounds for tail probabilities, which are useful, for example, in analyzing algorithms.

We start with an easy to prove, but also very weak bound:

Theorem 3.1: Markov's Inequality. Let $X \ge 0$ be a r.v. Then for Slide 4 any a > 0

$$\Pr(X \ge a) \le \frac{E(X)}{a}$$

Proof: Let $I = \mathbb{1}(X \ge a)$. Then by construction (i) $I \le X/a$, and (ii) $E(I) = \Pr(I = 1) = \Pr(X \ge a)$.

$$\implies E(I) = \Pr(X \ge a) \le E(X/a) = E(X)/a.$$

Example: X = # heads in n fair coin flips, $X \sim \text{Bin}(n, \frac{1}{2})$.

$$\implies \Pr(X \ge \frac{3}{4}n) \le \frac{E(X)}{\frac{3}{4}n} = \frac{\frac{n}{2}}{\frac{3}{4}n} = \frac{2}{3}$$

Proof: Var(X + Y) =

$$E\{[(X + Y) - E(X + Y)]^2\}$$

$$= E\left\{ [X + Y - EX - EY)]^2 \right\}$$

$$= E\{[(X - EX) + (Y - EY)]^2\}$$

$$= E\{(X - EX)^2\} + E\{(Y - EY)^2\} + 2E\{(X - EX)(Y - EY)\}$$

$$= Var(X) + Var(Y) + 2 Cov(X, Y)$$

More on (Co-)Variance

Theorem 3.3: If X, Y are independent r.v's (we write " $X \perp Y$ "), then

$$E(X \cdot Y) = E(X) \cdot E(Y)$$

but, the same is not true for dependent r.v's.

Proof:
$$E(XY) = \sum_{i} \sum_{j} (i \cdot j) \underbrace{\Pr(X = i, Y = j)}_{=p_X(i) \cdot p_Y(j)} = \underbrace{\left(\sum_{i} ip_X(i)\right) \cdot \left(\sum_{i} jp_Y(j)\right)}_{=E(X) \cdot E(Y)}$$

Corollary 3.4: $X \perp Y \implies \text{Cov}(X, Y) = 0$ and therefore

Var(X + Y) = Var(X) + Var(Y).

Independent, joint probability function = product of the marginal probability functions

Proof: Cov(X, Y) =

$$=E\left\{(X-EX)\cdot(Y-EY)\right\}=E(XY)-EX\cdot EY-EX\cdot EY+EX\cdot EY=0$$

using
$$E(XY) = EX \cdot EY$$
.

Variance & Moments 3.2

The first moment is the expected value Slide 2

Variance & Moments

We can get better bounds by using $E(X^k)$:

Definition 3.1: $E(X^k) = k$ -th moment of a r.v. X.

Definition 3.2:
$$Var(X) = E\{(X - EX)^2\} = \dots = E(X^2) - (EX)^2$$
.

And standard deviation $\sigma(X) = \sqrt{\operatorname{Var}(X)}$,

or just σ (if the r.v. X is understood)

Proof: of the identity for Var(X): $E(X) = \sum x^* Px(x)$

$$E\{(X - EX)^2\} = E\{X^2 - 2X \cdot EX + (EX)^2\} = Slide$$

$$= E(X^2) - 2EX \cdot E(X) + (EX)^2 = E(X^2) - (EX)^2$$

Slide 5

Example: Binomial Variance

Recall binomial r.v's. Let $X \sim Bin(n, p)$, with

Example: Binomial Variance

$$p_X(j) = \binom{n}{j} p^{J} (1-p)^{(n-j)}.$$

We find Var(X). Some observations

3.3 Covariance

Positive correlation, Cov(X,Y)>0 Slide 3 Negative correlation, Cov(X,Y)<0

Covariance

While E(X + Y) = E(X) + E(Y), the same is not true for Var(X).

We need:

 $Cov(X, Y) = \sum (X - EX)(Y - EY)P(X,Y)$

Definition 3.3: Covariance, $Cov(X, Y) = E\{(X, Y)\}$

1

Theorem 3.2. Var(X + Y) = Var(X) + Var(Y) + 2Cov(X, Y)

Binomial theorem:

$$\sum_{\ell=0}^{k} {k \choose \ell} p^{\ell} (1-p)^{(k-\ell)} = 1.$$

• We will find $Var(X) = E(X^2) - (EX)^2$. In the expansion we will twice use the binomial theorem, with k = n - 2 and k = n - 1, respectively.

Slide 6

Slide 8

$$E(X^{2}) = \sum_{j} p_{X}(j) \cdot j^{2} = \sum_{j=0}^{n} {n \choose j} p^{j} (1-p)^{n-j} \cdot ((j^{2}-j)+j)$$

$$= \sum_{j=0}^{n} \frac{n!(j^{2}-j)}{(n-j)! \ j!} p^{j} (1-p)^{n-j} + \sum_{j=0}^{n} \frac{n!}{(n-j)! \ j!} p^{j} (1-p)^{n-j} \cdot j$$

$$= n(n-1)p^{2} \sum_{j=2}^{n} \frac{n! -2, \text{k: n-j, n-k: (n-2)-(n-j)=j-2}}{(n-j)! (j-2)!} p^{j-2} (1-p)^{n-j} + proof$$

$$= n(n-1)p^{2} \sum_{j=1}^{n} \frac{(n-1)!}{(n-j)! (j-1)!} p^{j-1} (1-p)^{n-j} + proof$$

$$= n \sum_{j=1}^{n} \frac{(n-1)!}{(n-j)! (j-1)!} p^{j-1} (1-p)^{n-j}$$

$$= 1 \text{ (bin. thm., } k = n-1, \ell = j-1)$$

Chebyshev's Inequality

Recall Markov's inequality, using only E(X). Using also the 2nd moment, Var(X) we get a better bound

Theorem 3.6: Chebyshev's Inequality. If X is a r.v. with Var(X), then for any a > 0,

$$\Pr(|X - EX| \ge a) \le \frac{\operatorname{Var}(X)}{a^2} \quad \operatorname{Var}(X) = \operatorname{E}\{(X - \operatorname{EX})^2\}$$

Proof: Markov inequality with $D = (X - EX)^2$.

Corollary 3.7. For any t > 1

$$\Pr(|X - EX| \ge t\sigma(X)) \le \frac{1}{t^2}$$

and

$$\Pr(|X - EX| \ge tEX) \le \frac{\operatorname{Var}(X)}{t^2(EX)^2}$$

and therefore $Var(X) = E(X^2) - (EX)^2 = \dots = np(1-p)$, E(X) = np

n choose k = n!/k!(n-k)!

 $= n(n-1)p^2 + np$

3.7 Median & Mean

Slide 9

3.5 Examples

Slide 7

Median and Mean

Example

Suppose that it is known that the number of items X produced in a factory during a week is a random variable with mean $\mu = 50$.

(a) What can be said about $Pr(X \ge 75)$?

Solution: Markov's inequality gives $Pr(X \ge 75) \le \mu/75 = 2/3$.

(b) If $Var(X) = \sigma^2 = 25$, then what can be said about Pr(40 < X < 60)?

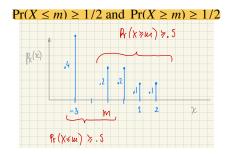
Solution: Consider
$$D = (X - \mu)^2$$
.
Note $Var(X) = E[(X - \mu)^2] = E(D)$.

By Markov inequality (for $D \dots$)

$$\begin{aligned} & \Pr(40 < X < 60) = \Pr((X - \mu) < 10) \\ & 1 - \Pr(40 < X < 60) = \Pr(D > 10^2) \le \frac{E(D)}{100} = \frac{\sigma^2}{100} = \frac{1}{4} \end{aligned}$$

² (Note: see Chebyshev's inequality).

Median: the *median* of a r.v. X, Md(X), is any value m with



Example: $X \sim \text{Unif}(\{x_1, \dots, x_{2k+1}\})$, uniform over an odd # of (sorted) values. Then $Md(X) = x_{k+1}$.

If $Y \sim \text{Unif}(\{x_1, \dots, x_{2k}\})$, then Md(Y) = m for any value $x_k < m < x_{k+1}$.

Slide 10 $E(X) = \sum x^* Px(x)$

Theorem 3.9: For any r.v. X with finite E(X) and Md(X)

- (a) $E(X) = \mu$ is the value c that minimizes $E[(X c)^2]$, and
- (b) Md(X) = m is a value c that minimizes E[|X c|] (need not be unique).

¹In the lecture video we used X > 75. Considering the integer nature of X (as a count) we could actually get a sharper bound by $Pr(X > 75) = Pr(X ≥ 76) \le \mu/76$.

²In the lecture the previous line was missing $1 - \dots$!

Proof: Part (a) follows from taking the derivative in

$$E[(X-c)^2] = E(X^2) - 2cEX + c^2$$
, giving

$$-2 EX + 2c = 0 \implies c = EX$$
.

when derivative = 0, get minimizes

Slide 11

For (b), will show $E(|X-c|-|X-m|) \ge 0$ for any $c \ne m$. First consider a value c with c > m and therefore $Pr(X \ge c) < 1$ (and then a similar argument for c < m),

Write the expectation, breaking it down into three terms:

$$|x-c|-|x-m| = \begin{cases} -(c-m) & \text{for } x \ge c \\ c+m-2x & \text{for } m < x < c \\ c-m & \text{for } x \le m, \end{cases}$$

implying

$$E(|X - c| - |X - m|) = -(c - m) \underbrace{\Pr(X \ge c)}_{+} + \sum_{m < x < c} (c + m - 2x) p_X(x) + (c - m) \Pr(X \le m)$$

Case Pr(m < X < c) = 0: We are done:

 $|x-c| - |x-m| = -(c-m)\Pr(X \ge c) + (c-m)\Pr(X < c) > 0$ For those following along in the book, we are now skipping §3.5.

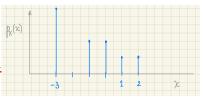
recall that we are in the case c > m and $Pr(X \ge c) < 1/2$.

3.8 Examples

Slide 14

Examples

Consider a discrete r.v. X with



$$p_X(x) = \begin{cases} 0.4 & \text{for } x = -3\\ 0.2 & \text{for } x = -1\\ 0.2 & \text{for } x = 0\\ 0.1 & \text{for } x = 1\\ 0.1 & \text{for } x = 2 \end{cases}$$

(a) Find a to minimize the expected sq distance $E((a-X)^2)$.

Solution: By Theorem 3.9a, a = E(X) = -1.1.

(b) Find b to minimize the expected absolute distance E(|b-X|).

Solution: By Theorem 3.9b, b = Md(X) = -1.

We are now ready to discuss statistical inference, using, for example, binomial r.v.'s and the rules for working with probabilities.

Slide 12

Case Pr(m < X < c) > 0:

E(|X-c|-|X-m|) =

$$|X - c| - |X - m| =$$

$$= -(c - m) \operatorname{Pr}(X \ge c)$$

$$+ \sum_{m < x < c} \underbrace{((c - x) - (x - m))}_{\ge -(c - m)} p_X(x) + (c - m) \operatorname{Pr}(X \le m)$$

>
$$-(c-m) \underbrace{\Pr(X > m)}_{<0.5}$$
 $+(c-m) \underbrace{\Pr(X \le m)}_{\ge 0.5}$
> $-(c-m) 0.5 + (c-m) 0.5 = 0.$

$$> -(c-m)0.5 + (c-m)0.5 = 0$$

In the 2nd line we used (c-x)-(x-m) > -(c-m): the difference of the two subintervals > - length of the entire interval.

Slide 13

Theorem 3.10: If X is a r.v. with finite $E(X) = \mu$, $Var(X) = \sigma^2$ and Md(X) = m, then $|\mu - m| < \sigma$. $Var(X) = E\{(X - EX)2\}$

Proof:

$$\begin{split} |EX-m| &= |E(X-m)| \leq E(|X-m|) \leq \\ &\leq E(|X-\mu|) = E[\sqrt{(X-\mu)^2}] \leq \sqrt{E[(X-\mu)^2]} = \sigma \\ &\underset{\text{using Jensen's inequality (1st and 3rd inequality) and}}{\text{Md}(X) = m \text{ is a value c that minimizes } E[|X-c|] \\ &= \sigma \\ &= \sigma$$

Theorem 3.9(b) (2nd inequality).

The Role of Probability in Statistical Inference

Slide 16

The Role of Probability in Statistical Inference

Probability: In the discussion so far, did you notice that we cheated? We always assumed that probabilities and distributions of r.v.'s were known.

For example "Let X be the number of cells with [...]. Assume $X \sim \text{Bin}(n, p)$, with p = 0.6. Find $\text{Pr}(X \ge 100)$."

In a real experiment, e.g., carried out in a lab, you never *know p*! Or the rate λ of the service time for a customer in the post office etc. This is where statistics comes in.

Statistics: We ask "Using an observed value x for $X \sim \text{Bin}(n, p)$, can we guess what p could be?", or "Could p be greater 0.5?"

That is, we change perspective. Assuming we have a good description of the experimental data as a r.v.'s, we try to report inference on the parameters, like p etc.

Slide 17

The change of perspective between probability and statistics naturally gives rise to some different vocabulary in statistics.

- *Sample*: a set of r.v.'s $X = \{X_i, i = 1, ..., n\}, X_i \sim F$, i.i.d. that are actually observed. We also just refer to X as the *data*. set of random variables
- *Statistic*: any function of the observed data, $S = f(X_1, ..., X_n)$. In probability we would have simply said S is another r.v. Of course, it is both!
- Parameters: This is tricky. In some pars of statistics
 (Bayesian statistics), parameters are r.v.'s that are not observed.
 In other contexts they index (describe) the distribution of the data. We usually use greek letters, like λ, μ, well, or p.
- *Hypothesis:* a hypothesis is simply an event for the parameters, like $A = \{\mu > 0\}$ (and we might use different names like H_0 etc.)

In summary, probability describes uncertainty, whereas statistics is about decisions in the face of such uncertainty.