

Lecture 16: Rank-Sparsity Matrix Decomposition

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Parts of the notes are based on [Chandrasekaran et al. \[2009, 2011\]](#).

Setting: Let $C = A^* + B^*$ with $A^* \in \mathbb{R}^{n \times n}$ being a sparse matrix and $B^* \in \mathbb{R}^{n \times n}$ a low-rank matrix, where both A^* and B^* are unknown. In this notes, we restrict ourselves to square matrices in $\mathbb{R}^{n \times n}$, but the analysis can be extended to rectangular matrices $\mathbb{R}^{n_1 \times n_2}$ if we simply replace n by $\max\{n_1, n_2\}$.

Goal: Given C , we want to recover A^* and B^* without any prior information about the sparsity pattern of A^* or the rank/singular vectors of B^* .

Solution: Consider the following optimization problem:

$$\begin{aligned} \arg \min_{A, B} [\gamma \|A\|_1 + \|B\|_*] \\ \text{subject to } A + B = C. \end{aligned} \quad (1)$$

Here, $\|A\|_1 = \sum_{i,j} |A_{ij}|$ is the elementwise L_1 -norm of a matrix A , $\|B\|_* = \sum_k \sigma_k(B)$ is the nuclear norm, which is the sum of the singular values of B , and γ is a tuning parameter that provides a trade-off between the low-rank and sparse components.

Remark 1. This optimization problem (1) is convex and can be written as a semi-definite program (SDP; [Vandenberghe and Boyd 1996](#)), for which there exist polynomial-time general-purpose solvers; see Appendix A in [Chandrasekaran et al. \[2011\]](#). Under a mild tightening of the conditions for fundamental identifiability, the minimizer of (1) is unique and recover A^*, B^* . Essentially, these conditions require that the sparse matrix does not have support concentrated within a single row/column, while the low-rank matrix does not have row/column spaces closely aligned with the coordinate axes [[Chandrasekaran et al., 2009](#)].

Notations: We begin by introducing several algebraic varieties¹. The set of rank-constrained matrices is defined as:

$$\mathcal{P}(k) = \{M \in \mathbb{R}^{n \times n} : \text{rank}(M) \leq k\}.$$

This is an algebraic variety with dimension $k(2n - k) = n^2 - (n - k)^2$, since it can be defined through the vanishing of all $(k + 1) \times (k + 1)$ minors of the matrix M . Let $M = UDV^T \in \mathbb{R}^{n \times n}$ be the singular value decomposition of M with $U, V \in \mathbb{R}^{n \times k}$ and $\text{rank}(M) = k$. The tangent space at M is defined as:

$$T(M) = \{UX^T + YV^T : X, Y \in \mathbb{R}^{n \times n}\},$$

which consists of the span of all matrices with either the same row space as M or the same column space as M . We also define

$$\Omega(M) = \{N \in \mathbb{R}^{n \times n} : \text{support}(N) \subseteq \text{support}(M)\},$$

which is the tangent space of $\{M \in \mathbb{R}^{n \times n} : |\text{support}(M)| \leq m\}$. Consider the following two quantities:

$$\xi(M) = \max_{N \in T(M), \|N\|_2 \leq 1} \|N\|_\infty$$

¹Recall that an algebraic variety is defined as the zero set of a system of polynomial equations [[Hartshorne, 2013](#)].

which will be small when (appropriately scaled) elements of the tangent space $T(M)$ are “diffuse” (i.e., these elements are not too sparse), and

$$\mu(M) = \max_{N \in \Omega(M), \|N\|_\infty \leq 1} \|N\|_2$$

which will be small when the spectrum of any matrix in $\Omega(M)$ is “diffuse” (i.e., the singular values of these elements are not too large). Here, $\|\cdot\|_\infty$ denotes the largest entry in magnitude and $\|\cdot\|_2$ is the spectral norm (i.e., the largest singular value).

Remark 2. One can show that

$$\deg_{\min}(M) \leq \mu(M) \leq \deg_{\max}(M),$$

where $\deg_{\max}(M)$ is the maximum number of nonzero entries per row/column and $\deg_{\min}(M)$ is the minimum number of nonzero entries per row/column; see Proposition 3 in [Chandrasekaran et al. \[2011\]](#). Analogously, we can bound $\xi(M)$ as:

$$\text{inc}(M) \leq \xi(M) \leq 2 \cdot \text{inc}(M),$$

where $\text{inc}(M) = \max\{\beta(\text{row-space}(M)), \beta(\text{column-space}(M))\}$ is the incoherence of the row/column spaces of a matrix $M \in \mathbb{R}^{n \times n}$ with $\beta(S) = \max_i \|P_S e_i\|_2$ as the incoherence of a subspace $S \subset \mathbb{R}^n$. Here, $\{e_1, \dots, e_n\}$ is the standard basis of \mathbb{R}^n , P_S denotes the projection onto the subspace S , and $\|\cdot\|_2$ is the vector ℓ_2 -norm.

1 Basic Properties

Proposition 1. If $\mu(A^*)\xi(B^*) < 1$ for two matrices $A^*, B^* \in \mathbb{R}^{n \times n}$, then $\Omega(A^*) \cap T(B^*) = \{0\}$.

We may choose γ properly to have $\mu(A^*)\xi(B^*) < 1/6$, which guarantees the recoveries of A^* and B^* . To establish Proposition 1, we leverage the following lemma.

Lemma 2. $\max_{N \in T(B^*), \|N\|_2 \leq 1} \|P_{\Omega(A^*)}(N)\|_2 \leq \mu(A^*) \cdot \xi(B^*)$, where $P_{\Omega(A^*)}(N)$ is the projection of N on the space $\Omega(A^*)$.

Proof of Lemma 2. We have the following sequence of inequalities:

$$\begin{aligned} \max_{N \in T(B^*), \|N\|_2 \leq 1} \|P_{\Omega(A^*)}(N)\|_2 &\leq \max_{N \in T(B^*), \|N\|_2 \leq 1} \mu(A^*) \|P_{\Omega(A^*)}(N)\|_\infty \\ &\leq \max_{N \in T(B^*), \|N\|_2 \leq 1} \mu(A^*) \|N\|_\infty \\ &= \mu(A^*) \cdot \xi(B^*), \end{aligned}$$

where the first inequality follows from the definition of $\mu(A^*)$ as $P_{\Omega(A^*)}(N) \in \Omega(A^*)$ and the second inequality is due to $\|P_{\Omega(A^*)}(N)\|_\infty \leq \|N\|_\infty$. \square

Proof of Proposition 1. Suppose that there exists $\tilde{N} \neq 0$ and $\tilde{N} \in \Omega(A^*) \cap T(B^*)$. Given that $\tilde{N} \in T(B^*)$, we can scale \tilde{N} so that $\|\tilde{N}\|_2 = 1$. Thus, by Lemma 2,

$$\mu(A^*)\xi(B^*) \geq \max_{N \in T(M), \|N\|_2 \leq 1} \|P_{\Omega(A^*)}(N)\|_2 \geq \|P_{\Omega(A^*)}(\tilde{N})\|_2 = 1$$

contradicting to $\mu(A^*)\xi(B^*) < 1$. The result follows. \square

One important consequence of Proposition 1 is the following rank-sparsity uncertainty principle.

Theorem 3 (Rank-Sparsity Uncertainty Principle). *For a matrix $M \neq 0$, we have that*

$$\xi(M) \cdot \mu(M) \geq 1.$$

Proof. Notice that $M \in \Omega(M) \cap T(M)$. By Proposition 1, we know that $\xi(M) \cdot \mu(M) < 1$, leading to a contradiction. \square

2 Optimality Condition

Consider the Lagrangian function of (1) as:

$$\mathcal{L}(A, B, Q) = \gamma \|A\|_1 + \|B\|_* + \langle Q, C - A - B \rangle.$$

From the optimality conditions of a convex program, (A^*, B^*) is a minimizer of (1) if and only if the dual matrix $Q \in \mathbb{R}^{n \times n}$ satisfies

$$Q \in \gamma \partial \|A^*\|_1 \quad \text{and} \quad Q \in \partial \|B^*\|_*. \quad (2)$$

Based on the subdifferentials of $\|\cdot\|_1$ and $\|\cdot\|_*$, we know that (2) is equivalent to

$$P_{\Omega(A^*)}(Q) = \gamma \text{sign}(A^*), \|P_{\Omega(A^*)}(Q)\|_\infty \leq \gamma \quad \text{and} \quad P_{T(B^*)}(Q) = UV^T, \|P_{T(B^*)^\perp}(Q)\|_2 \leq 1, \quad (3)$$

where $U, V \in \mathbb{R}^{n \times k}$ comes from $B^* = U\Sigma V^T$. (Recall that $\partial \|B^*\|_* = \{UV^T + W : U^T W = W V^T = 0\}$.) Notice that (3) are necessary and sufficient conditions for (A^*, B^*) be a minimizer of (1). To ensure the uniqueness for the solution to (1), we need to tighten the conditions in (2) and (3) as the following proposition.

Proposition 4 (Uniqueness of the Optimal Solution). *Suppose that $C = A^* + B^*$. Then, $(\hat{A}, \hat{B}) = (A^*, B^*)$ is the unique minimizer of (1) if the following conditions are satisfied:*

1. $\Omega(A^*) \cap T(B^*) = \{0\}$.
2. *There exists a dual matrix $Q \in \mathbb{R}^{n \times n}$ such that*
 - (a) $P_{T(B^*)}(Q) = UV^T$;
 - (b) $P_{\Omega(A^*)}(Q) = \gamma \cdot \text{sign}(A^*)$;
 - (c) $\|P_{T(B^*)^\perp}(Q)\|_2 < 1$;
 - (d) $\|P_{\Omega(A^*)^c}(Q)\|_\infty < \gamma$.

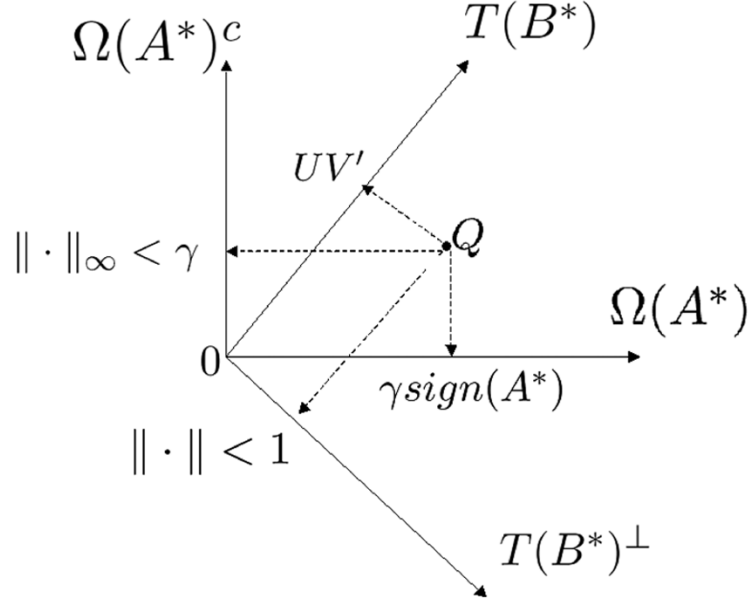
Proof of Proposition 4. Notice that (A^*, B^*) is an optimum by the condition 2 in Proposition 4. To avoid cluttered notation, we let $\Omega = \Omega(A^*)$, $T = T(B^*)$, $\Omega^c = \Omega(A^*)^c$, and $T_\perp(B^*) = T^\perp$.

Suppose that there is another feasible solution $(A^* + N_A, B^* + N_B)$ that also minimizes (1). Since $A^* + B^* = C = (A^* + N_A) + (B^* + N_B)$, we must have $N_A + N_B = 0$. For any subgradient (Q_A, Q_B) of the function $\gamma \|A\|_1 + \|B\|_*$ at (A^*, B^*) , we have that

$$\gamma \|A^* + N_A\|_1 + \|B^* + N_B\|_* \geq \gamma \|A^*\|_1 + \|B^*\|_* + \langle Q_A, N_A \rangle + \langle Q_B, N_B \rangle. \quad (4)$$

Since (Q_A, Q_B) is a subgradient of the function $\gamma \|A\|_1 + \|B\|_*$ at (A^*, B^*) , we must have from (3) that

- $Q_A = \gamma \cdot \text{sign}(A^*) + P_{\Omega^c}(Q_A)$ with $\|P_{\Omega^c}(Q_A)\|_\infty \leq \gamma$;
- $Q_B = UV^T + P_{T^\perp}(Q_B)$ with $\|P_{T^\perp}(Q_B)\|_2 \leq 1$.

Figure 1: Geometric interpretation of optimality conditions: the existence of a dual matrix Q .

Thus, we calculate that

$$\begin{aligned}
 \langle Q_A, N_A \rangle &= \langle \gamma \cdot \text{sign}(A^*) + P_{\Omega^c}(Q_A), N_A \rangle \\
 &= \langle P_{\Omega}(Q) + P_{\Omega^c}(Q_A), N_A \rangle \quad \text{using (b) in Condition 2} \\
 &= \langle P_{\Omega^c}(Q_A) - P_{\Omega^c}(Q), N_A \rangle + \langle Q, N_A \rangle \quad \text{by } P_{\Omega}(Q) = Q - P_{\Omega^c}(Q).
 \end{aligned}$$

Similarly, we have that

$$\begin{aligned}
 \langle Q_B, N_B \rangle &= \langle UV^T + P_{T^\perp}(Q_B), N_B \rangle \\
 &= \langle P_T(Q) + P_{T^\perp}(Q_B), N_B \rangle \quad \text{using (a) in Condition 2} \\
 &= \langle P_{T^\perp}(Q_B) - P_{T^\perp}(Q), N_B \rangle + \langle Q, N_B \rangle \quad \text{by } P_T(Q) = Q - P_{T^\perp}(Q).
 \end{aligned}$$

Adding the above two equalities together gives us that

$$\begin{aligned}
 \langle Q_A, N_A \rangle + \langle Q_B, N_B \rangle &= \langle P_{\Omega^c}(Q_A) - P_{\Omega^c}(Q), N_A \rangle + \langle Q, N_A \rangle + \langle P_{T^\perp}(Q_B) - P_{T^\perp}(Q), N_B \rangle + \langle Q, N_B \rangle \\
 &= \langle P_{\Omega^c}(Q_A) - P_{\Omega^c}(Q), P_{\Omega^c}(N_A) \rangle + \langle P_{T^\perp}(Q_B) - P_{T^\perp}(Q), P_{T^\perp}(N_B) \rangle,
 \end{aligned} \tag{5}$$

where we use the fact that $N_A + N_B = 0$ and the projection matrices $P_{\Omega^c}, P_{T^\perp}$ are idempotent.

Given that any subgradient (Q_A, Q_B) of the function $\gamma \|A\|_1 + \|B\|_\star$ at (A^*, B^*) will satisfy the above equality, we can choose (Q_A, Q_B) as follows:

- Take Q_A so that $P_{\Omega^c}(Q_A) = \gamma \cdot \text{sign}(P_{\Omega^c}(N_A))$ with $\|P_{\Omega^c}(Q_A)\|_\infty \leq \gamma$ and $\langle P_{\Omega^c}(Q_A), P_{\Omega^c}(N_A) \rangle = \gamma \|P_{\Omega^c}(N_A)\|_1$.
- Given the singular value decomposition of $P_{T^\perp}(N_B) = \tilde{U} \tilde{\Sigma} \tilde{V}^T$, we choose Q_B so that $P_{T^\perp}(Q_B) = \tilde{U} \tilde{V}^T$ with $\|P_{T^\perp}(Q_B)\|_2 = 1$ and $\langle P_{T^\perp}(Q_B), P_{T^\perp}(N_B) \rangle = \|P_{T^\perp}(N_B)\|_\star$.

Under this choice of (Q_A, Q_B) , we simplify (5) as:

$$\langle Q_A, N_A \rangle + \langle Q_B, N_B \rangle = \langle P_{\Omega^c}(Q_A) - P_{\Omega^c}(Q), P_{\Omega^c}(N_A) \rangle + \langle P_{T^\perp}(Q_B) - P_{T^\perp}(Q), P_{T^\perp}(N_B) \rangle$$

$$\begin{aligned} &\geq (\gamma - \|P_{\Omega^c}(Q)\|_\infty) \|P_{\Omega^c}(N_A)\|_1 + (1 - \|P_{T^\perp}(Q)\|_2) \|P_{T^\perp}(N_B)\|_* \\ &> 0 \end{aligned}$$

unless $P_{\Omega^c}(N_A) = P_{T^\perp}(N_B) = 0$, where we obtain the last positivity based on (c) and (d) in Condition 2. However, if $P_{\Omega^c}(N_A) \neq 0$ or $P_{T^\perp}(N_B) \neq 0$, we know from (4) that

$$\gamma \|A^* + N_A\|_1 + \|B^* + N_B\|_* > \gamma \|A^*\|_1 + \|B^*\|_*,$$

which violates the optimality of $(A^* + N_A, B^* + N_B)$. Now, when $P_{\Omega^c}(N_A) = P_{T^\perp}(N_B) = 0$, $P_\Omega(N_A) + P_T(N_B) = 0$ as well because of $N_A + N_B = 0$. In other words,

$$P_\Omega(N_A) = -P_T(N_B).$$

This is only possible if $P_\Omega(N_A) = P_T(N_B) = 0$ because $\Omega \cap T = \{0\}$ by Condition 1, which in turn implies that $N_A = N_B = 0$. The proof of uniqueness is completed. \square

While Proposition 4 sheds light on the sufficient conditions for uniquely recovering (A^*, B^*) , we now discuss the existence of an appropriate dual matrix Q entailed by Proposition 4. From Proposition 1, we already know that Condition 1 in Proposition 4 ($\Omega(A^*) \cap T(B^*) = \{0\}$) is valid when $\mu(A^*)\xi(B^*) < 1$. If we slightly strengthen the condition as $\mu(A^*)\xi(B^*) < \frac{1}{6}$, there will be a dual matrix Q satisfying the requirements in Condition 2 of Proposition 4 as well.

Theorem 5. *Given $C = A^* + B^*$ with $\mu(A^*)\xi(B^*) < \frac{1}{6}$, the unique minimizer (\hat{A}, \hat{B}) of (1) will be (A^*, B^*) for the following range of γ :*

$$\gamma \in \left(\frac{\xi(B^*)}{1 - 4\mu(A^*)\xi(B^*)}, \frac{1 - 3\mu(A^*)\xi(B^*)}{\mu(A^*)} \right).$$

Specifically, $\gamma = \frac{[3\xi(B^)]^p}{[2\mu(A^*)]^{1-p}}$ for any choice of $p \in [0, 1]$ is always inside the above range and thus guarantees exact recovery of (A^*, B^*) .*

The detailed proof of Theorem 5 can be found in Theorem 2 of [chandrasekaran2011rank]. The high-level idea is that we consider candidates for the dual matrix Q in the direct sum $\Omega(A^*) \oplus T(B^*)$ of the tangent spaces. Since $\mu(A^*)\xi(B^*) < \frac{1}{6}$, $\Omega(A^*) \cap T(B^*) = \{0\}$ by Proposition 1 and there exists a *unique* element $\hat{Q} \in \Omega(A^*) \oplus T(B^*)$ satisfying $P_{T(B^*)}(\hat{Q}) = UV^T$ and $P_{\Omega(A^*)}(\hat{Q}) = \gamma \cdot \text{sign}(A^*)$. The proof proceeds by showing that if $\mu(A^*)\xi(B^*) < \frac{1}{6}$, then the projections of \hat{Q} onto the orthogonal spaces $\Omega(A^*)^\perp$ and $T(B^*)^\perp$ are small, and Condition 2 of Proposition 4 is thus satisfied.

Other further reading for the course:

- Chandrasekaran, V., Recht, B., Parrilo, P. A., & Willsky, A. S. (2012). The convex geometry of linear inverse problems. *Foundations of Computational Mathematics*, 12, 805-849.

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