

On k -Connectivity and Minimum Vertex Degree in Random s -Intersection Graphs

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Abstract

Random s -intersection graphs have recently received much interest in a wide range of application areas. Broadly speaking, a random s -intersection graph is constructed by first assigning each vertex a set of items in some *random* manner, and then putting an undirected edge between all pairs of vertices that share at least s items (the graph is called a *random intersection graph* when $s = 1$). A special case of particular interest is a *uniform random s -intersection graph*, where each vertex independently selects the same number of items uniformly at random from a common item pool. Another important case is a *binomial random s -intersection graph*, where each item from a pool is independently assigned to each vertex with the same probability. Both models have found numerous applications thus far including cryptanalysis, and modeling recommender systems, secure sensor networks, online social networks, trust networks and small-world networks (uniform random s -intersection graphs), as well as clustering analysis, classification, and the design of integrated circuits (binomial random s -intersection graphs).

In this paper, for binomial/uniform random s -intersection graphs, we present results related to k -connectivity and minimum vertex degree. Specifically, we derive the asymptotically exact probabilities and zero-one laws for the following three properties: (i) k -vertex-connectivity, (ii) k -edge-connectivity and (iii) the property of minimum vertex degree being at least k .

Keywords—Random intersection graph, random key graph, connectivity, secure sensor network.

1 Introduction

Random s -intersection graphs have received considerable attention recently [1–13]. In such a graph, each vertex is equipped with a set of items in some *random* manner, and two vertices establish an undirected edge in between if and only if they have at least s items in common. A large amount of work [17–36] study the case of s being 1, under which the graphs are simply referred to as *random intersection graphs*.

Random (s -)intersection graphs have been used to model secure wireless sensor networks [12–22], wireless frequency hopping [22], epidemics in human populations [1, 2], small-world networks [39], trust networks [40, 41], social networks [2–6] such as collaboration networks [2–4] and common-interest networks [5, 6]. Random intersection graphs also motivated Beer *et al.* [42, 43] to introduce a general concept of *vertex random graphs* that subsumes any graph model where *random* features are assigned to vertices, and edges are drawn based on deterministic relations between the features of the vertices.

Among different models of random s -intersection graphs, two widely studied models are the so-called *uniform random s -intersection graph* and *binomial random s -intersection graph* defined in detail below.

1.1 Graph models

Uniform random s -intersection graph. A *uniform random s -intersection graph*, denoted by $G_s(n, K_n, P_n)$, is defined on n vertices as follows. Each vertex *independently* selects K_n different items *uniformly at random* from a pool of P_n distinct items. Two vertices have an edge in between if and only if they share at least s items. The notion “uniform” means that all vertices have the same number of items (but likely different sets of items). Here K_n and P_n are both functions of n , while s does not scale with n . It holds that $1 \leq s \leq K_n \leq P_n$. Under $s = 1$, the graph is also known as a *random key graph* [37–39].

Binomial random s -intersection graph. A *binomial random s -intersection graph*, denoted by $H_s(n, t_n, P_n)$, is defined on n vertices as follows. Each item from a pool of P_n distinct items is assigned to each vertex *independently* with probability t_n . Two vertices establish an edge in between if and only if they have at least s items in common. The term “binomial” is used since the number of items assigned to each vertex follows a binomial distribution with parameters P_n (the number of trials) and t_n (the success probability in each trial). Here t_n and P_n are both functions of n , while s does not scale with n . Also it holds that $1 \leq s \leq P_n$.

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1.2 Problem Statement. Our goal in this paper is to investigate properties related to k -connectivity and *minimum vertex degree* of random s -intersection graphs (k -vertex-connectivity and k -edge-connectivity are called together as k -connectivity for convenience). In particular, we wish to answer the following question:

For a uniform random s -intersection graph $G_s(n, K_n, P_n)$ (resp., a binomial random s -intersection graph $H_s(n, t_n, P_n)$), with parameters K_n (resp., t_n) and P_n scaling with the number of vertices n , what is the asymptotic behavior of the probabilities for $G_s(n, K_n, P_n)$ (resp., $H_s(n, t_n, P_n)$) (i) being k -vertex-connected, (ii) being k -edge-connected, and (iii) having a minimum vertex degree at least k , respectively, as n grows large? A graph is said to be k -vertex-connected if the remaining graph is still connected despite the deletion of at most $(k-1)$ arbitrary vertices, and k -edge-connectivity is defined similarly for the deletion of edges [48]; with $k=1$, these definitions reduce to the standard notion of *graph connectivity* [44, 45]. The degree of a vertex is defined as the number of edges incident on it. The three graph properties considered here are related to each other in that k -vertex-connectivity implies k -edge-connectivity, which in turn implies that the minimum vertex degree is at least k [48].

1.3 Summary of Results. We summarize our results below, first for a uniform random s -intersection graph and then for a binomial random s -intersection graph. Throughout the paper, both s and k are positive integers and do not scale with n . Also, naturally we consider $1 \leq s \leq K_n \leq P_n$ for graph $G_s(n, K_n, P_n)$ and $1 \leq s \leq P_n$ for graph $H_s(n, t_n, P_n)$. We use the standard Landau asymptotic notation $\Omega(\cdot)$, $\omega(\cdot)$, $O(\cdot)$, $o(\cdot)$, $\Theta(\cdot)$. $\mathbb{P}[\mathcal{E}]$ denotes the probability that event \mathcal{E} happens.

k -Connectivity & minimum vertex degree in uniform random s -intersection graphs:

For a uniform random s -intersection graph $G_s(n, K_n, P_n)$ under $P_n = \Omega(n)$, with sequence α_n defined by

$$(1.1) \quad \frac{1}{s!} \cdot \frac{K_n^{2s}}{P_n^s} = \frac{\ln n + (k-1) \ln \ln n + \alpha_n}{n},$$

then as $n \rightarrow \infty$, if $\alpha_n \rightarrow \alpha^* \in [-\infty, \infty]$, the following convergence results hold:

$$\mathbb{P}[G_s(n, K_n, P_n) \text{ is } k\text{-vertex-connected.}] \rightarrow e^{-\frac{e^{-\alpha^*}}{(k-1)!}},$$

$$\mathbb{P}[G_s(n, K_n, P_n) \text{ is } k\text{-edge-connected.}] \rightarrow e^{-\frac{e^{-\alpha^*}}{(k-1)!}},$$

and

$$\mathbb{P}\left[G_s(n, K_n, P_n) \text{ has a minimum vertex degree at least } k.\right] \rightarrow e^{-\frac{e^{-\alpha^*}}{(k-1)!}}.$$

k -Connectivity & minimum vertex degree in binomial random s -intersection graphs:

For a binomial random s -intersection graph $H_s(n, t_n, P_n)$ under $P_n = \Omega(n)$ for $s \geq 2$ or $P_n = \Omega(n^c)$ for $s = 1$ with some constant $c > 1$, with sequence β_n defined by

$$\frac{1}{s!} \cdot t_n^{2s} P_n^s = \frac{\ln n + (k-1) \ln \ln n + \beta_n}{n},$$

then as $n \rightarrow \infty$, if $\beta_n \rightarrow \beta^* \in [-\infty, \infty]$, the following convergence results hold:

$$\mathbb{P}[H_s(n, t_n, P_n) \text{ is } k\text{-vertex-connected.}] \rightarrow e^{-\frac{e^{-\beta^*}}{(k-1)!}},$$

$$\mathbb{P}[H_s(n, t_n, P_n) \text{ is } k\text{-edge-connected.}] \rightarrow e^{-\frac{e^{-\beta^*}}{(k-1)!}},$$

and

$$\mathbb{P}\left[H_s(n, t_n, P_n) \text{ has a minimum vertex degree at least } k.\right] \rightarrow e^{-\frac{e^{-\beta^*}}{(k-1)!}}.$$

Since the probability $e^{-\frac{e^{-\alpha^*}}{(k-1)!}}$ (resp., $e^{-\frac{e^{-\beta^*}}{(k-1)!}}$) equals 1 if $\alpha^* = \infty$ (resp., $\beta^* = \infty$) and 0 if $\alpha^* = -\infty$ (resp., $\beta^* = -\infty$), the above results of asymptotically exact probabilities also imply the corresponding zero-one laws, where a zero-one law [37] means that the probability that the graph has certain property asymptotically converges to 0 under some conditions and converges to 1 under some other conditions.

1.4 Comparison with related work. Table 1 summarizes relevant work in the literature on uniform/binomial random s -intersection graphs in terms of k -vertex-connectivity, k -edge connectivity, and the property of minimum vertex degree being at least k .

Among the related work, Bloznelis and Rybarczyk [7] recently also derived the asymptotically exact probabilities of *uniform* random s -intersection graphs (but not of *binomial* random s -intersection graphs) for the three properties above (the easily implied results on k -edge-connectivity were not explicitly mentioned). Yet, when s is a constant or $O(1)$ as in many applications, their results require $K_n = O((\ln n)^{\frac{1}{5s}})$ for k -connectivity (k -vertex-connectivity and k -edge-connectivity), under a scaling the same as in Equation (1.1). In other words, the one-law part of k -connectivity is as follows: under certain conditions including $K_n = O((\ln n)^{\frac{1}{5s}})$, if $\frac{1}{s!} \cdot \frac{K_n^{2s}}{P_n^s} = \frac{\ln n + (k-1) \ln \ln n + \omega(1)}{n}$, then a uniform random s -intersection graph $G_s(n, K_n, P_n)$ is k -connected with a probability converging to 1 as $n \rightarrow \infty$. From $K_n = O((\ln n)^{\frac{1}{5s}})$ and $\frac{1}{s!} \cdot \frac{K_n^{2s}}{P_n^s} = \frac{\ln n + (k-1) \ln \ln n + \omega(1)}{n}$, it is straightforward to derive $P_n = O(n^{\frac{1}{s}} (\ln n)^{-\frac{3}{5s}})$; i.e., $P_n = \tilde{O}(n^{\frac{1}{s}})$ ignoring the $\ln n$ terms. However, in secure wireless sensor network applications where uniform random s -

intersection graphs are widely investigated, conditions $K_n = O((\ln n)^{\frac{1}{5s}})$ and $P_n = \tilde{O}(n^{\frac{1}{s}})$ are both likely impractical because K_n and P_n are often at least on the order of $\ln n$ and n , respectively, to ensure that the network has reasonable resiliency against sensor capture attacks [16–18]. The results reported in this paper cover the practical range where P_n is at least on the order of n .

1.5 Roadmap. We organize the rest of the paper as follows. We detail the main results in Section 2. Sections 3 and 4 detail the steps of establishing the theorems. We conclude the paper in Section 5. The Appendix provides additional arguments used in proving the theorems.

2 Main Results

Below we explain the main results of uniform random s -intersection graphs and binomial random s -intersection graphs, respectively.

2.1 Results of uniform random s -intersection graphs. The following theorem presents results on k -connectivity and minimum vertex degree in a uniform random s -intersection graph $G_s(n, K_n, P_n)$.

THEOREM 2.1. *For a uniform random s -intersection graph $G_s(n, K_n, P_n)$ under*

$$(2.2) \quad P_n = \Omega(n),$$

with sequence α_n defined by

$$(2.3) \quad \frac{1}{s!} \cdot \frac{K_n^{2s}}{P_n^s} = \frac{\ln n + (k-1) \ln \ln n + \alpha_n}{n},$$

it holds that

$$\begin{aligned} & \lim_{n \rightarrow \infty} \mathbb{P}[G_s(n, K_n, P_n) \text{ is } k\text{-vertex-connected.}] \\ &= \lim_{n \rightarrow \infty} \mathbb{P}[G_s(n, K_n, P_n) \text{ is } k\text{-edge-connected.}] \\ &= \lim_{n \rightarrow \infty} \mathbb{P}\left[G_s(n, K_n, P_n) \text{ has a minimum}\right. \\ & \quad \left.\text{vertex degree at least } k.\right] \\ &= \begin{cases} e^{-\frac{e^{-\alpha^*}}{(k-1)!}}, & \text{if } \lim_{n \rightarrow \infty} \alpha_n = \alpha^* \in (-\infty, \infty), \\ 0, & \text{if } \lim_{n \rightarrow \infty} \alpha_n = -\infty, \\ 1, & \text{if } \lim_{n \rightarrow \infty} \alpha_n = \infty. \end{cases} \end{aligned}$$

For graph $G_s(n, K_n, P_n)$, Theorem 2.1 presents the asymptotically exact probabilities and zero-one laws for the following three properties: (i) k -vertex connectivity, (ii) k -edge-connectivity and (iii) the property of minimum vertex degree being at least k .

By Lemma 6.2 on Page 9, under (2.2) and (2.3) with constrained $|\alpha_n| = O(\ln \ln n)$, we can show that the left hand side of (2.3), i.e., $\frac{1}{s!} \cdot \frac{K_n^{2s}}{P_n^s}$, is

asymptotically equivalent to the edge probability of graph $G_s(n, K_n, P_n)$. As given in Lemma 13 of the full version [9], with q_n denoting the edge probability of graph $G_s(n, K_n, P_n)$, if condition (2.3) is replaced by $q_n = \frac{\ln n + (k-1) \ln \ln n + \alpha_n}{n}$, and condition (2.2) is kept unchanged, then all results in Theorem 2.1 still follow. Hence, the uniform random s -intersection graph model under condition (2.2) exhibits the same behavior as the well-known Erdős-Rényi graph model [45], in the sense that for each of (i) k -vertex-connectivity, (ii) k -edge-connectivity and (iii) the property of minimum vertex degree being at least k , a common point for the phase transition from a zero-law to a one-law occurs when the edge probability equals $\frac{\ln n + (k-1) \ln \ln n}{n}$.

2.2 Results for binomial random s -intersection graphs. The following theorem presents results on k -connectivity and minimum vertex degree in a binomial random s -intersection graph $H_s(n, t_n, P_n)$.

THEOREM 2.2. *For a binomial random s -intersection graph $H_s(n, t_n, P_n)$ under*

$$(2.4) \quad \begin{cases} P_n = \Omega(n), & \text{for } s \geq 2, \\ P_n = \Omega(n^c) \text{ for some constant } c > 1, & \text{for } s = 1, \end{cases}$$

with sequence β_n defined by

$$(2.5) \quad \frac{1}{s!} \cdot t_n^{2s} P_n^s = \frac{\ln n + (k-1) \ln \ln n + \beta_n}{n},$$

it holds that

$$\begin{aligned} & \lim_{n \rightarrow \infty} \mathbb{P}[H_s(n, t_n, P_n) \text{ is } k\text{-vertex-connected.}] \\ &= \lim_{n \rightarrow \infty} \mathbb{P}[H_s(n, t_n, P_n) \text{ is } k\text{-edge-connected.}] \\ &= \lim_{n \rightarrow \infty} \mathbb{P}\left[H_s(n, t_n, P_n) \text{ has a minimum}\right. \\ & \quad \left.\text{vertex degree at least } k.\right] \\ &= \begin{cases} e^{-\frac{e^{-\beta^*}}{(k-1)!}}, & \text{if } \lim_{n \rightarrow \infty} \beta_n = \beta^* \in (-\infty, \infty), \\ 0, & \text{if } \lim_{n \rightarrow \infty} \beta_n = -\infty, \\ 1, & \text{if } \lim_{n \rightarrow \infty} \beta_n = \infty. \end{cases} \end{aligned}$$

For graph $H_s(n, t_n, P_n)$, Theorem 2.2 presents the asymptotically exact probabilities and zero-one laws for the following three properties: (i) k -vertex connectivity, (ii) k -edge-connectivity and (iii) the property of minimum vertex degree being at least k .

By Lemma 12 of the full version [9], under (2.4) and (2.5) with constrained $|\beta_n| = O(\ln \ln n)$, we can show that the left hand side of (2.5), i.e., $\frac{1}{s!} \cdot t_n^{2s} P_n^s$, is asymptotically equivalent to the edge probability

Graph		Property	Results	Work
uniform random s -intersection graph	$G_s(n, K_n, P_n)$	k -connectivity & min. vertex degree $\geq k$	exact probabilities	this paper [7] (only for $K_n = O((\ln n)^{\frac{1}{5s}})$)
		connectivity & min. vertex degree ≥ 1	exact probabilities	[8] (only for $K_n = O((\ln n)^{\frac{1}{5s}})$)
	$G_1(n, K_n, P_n)$	k -connectivity & min. vertex degree $\geq k$	exact probabilities	[25]
		connectivity & min. vertex degree ≥ 1	zero-one laws	[19, 20]
			exact probabilities	[28]
			zero-one laws	[36, 37]
binomial random s -intersection graph	$H_s(n, t_n, P_n)$	k -connectivity & min. vertex degree $\geq k$	exact probabilities	this paper
		connectivity & min. vertex degree ≥ 1	zero-one laws	
			exact probabilities	
			zero-one laws	
	$H_1(n, t_n, P_n)$	k -connectivity & min. vertex degree $\geq k$	exact probabilities	[25]
		connectivity & min. vertex degree ≥ 1	zero-one laws	[20]
			exact probabilities	[21]
			zero-one laws	[29, 30]

Table 1: Comparison of our results with related work. k -vertex connectivity and k -edge connectivity are together written as k -connectivity. Note that results on k -connectivity and min. vertex degree $\geq k$ also imply the corresponding results on connectivity and min. vertex degree ≥ 1 by setting k as 1. Also, results of exact probabilities imply the corresponding results of zero-one laws by monotonicity arguments.

of graph $H_s(n, t_n, P_n)$. As given in Lemma 14 of the full version [9], with ρ_n denoting the edge probability of graph $H_s(n, t_n, P_n)$, if condition (2.5) is replaced by $\rho_n = \frac{\ln n + (k-1) \ln \ln n + \beta_n}{n}$, and condition (2.4) is kept unchanged, then all results in Theorem 2.2 still follow. Therefore, the binomial random s -intersection graph model under condition (2.4) exhibits the same behavior with Erdős-Rényi graph model, in the sense that for each of (i) k -vertex-connectivity, (ii) k -edge-connectivity, and (iii) the property of minimum vertex degree being at least k , a common point for the phase transition from a zero-law to a one-law occurs when the edge probability equals $\frac{\ln n + (k-1) \ln \ln n}{n}$.

The condition (2.4) has $P_n = \Omega(n^s)$ for $s \geq 2$, and requires a stronger one for $s = 1$: $P_n = \Omega(n^c)$ for some constant $c > 1$. The range $P_n = \Theta(n)$ is covered by $P_n = \Omega(n)$, but not by $P_n = \Omega(n^c)$ with $c > 1$. For $s = 1$ and $P_n = \Theta(n)$, results for k -vertex-connectivity, k -edge connectivity, and the property of minimum vertex degree being at least k use a scaling different from (2.5), as given by [20, Theorem 4].

3 Establishing Theorem 2.1

Theorem 2.1 in the special case of $s = 1$ is proved by us [25]. Below we explain the steps of establishing Theorem 2.1 for $s \geq 2$. In Section 3.1, we show that $|\alpha_n|$ can be confined as $O(\ln \ln n)$ in proving Theorem 2.1. In Section 3.2, we consider the relationships

between vertex connectivity, edge connectivity, and minimum vertex degree.

3.1 Confining $|\alpha_n|$. To confine $|\alpha_n|$ as $O(\ln \ln n)$ in proving Theorem 2.1, we will demonstrate

$$(3.6)$$

Theorem 2.1 under $|\alpha_n| = O(\ln \ln n) \Rightarrow$ Theorem 2.1.

Note that k -vertex-connectivity, k -edge-connectivity, and the property of minimum vertex degree being at least k are all monotone increasing¹. For any monotone increasing property \mathcal{I} , the probability that a spanning subgraph (resp., supergraph) of graph G has \mathcal{I} is at most (resp., at least) the probability of G having \mathcal{I} . Therefore, to show (3.6), it suffices to prove the following lemma.

LEMMA 3.1. (a) For graph $G_s(n, K_n, P_n)$ under $P_n = \Omega(n)$ and

$$(3.7) \quad \frac{1}{s!} \cdot \frac{K_n^{2s}}{P_n^s} = \frac{\ln n + (k-1) \ln \ln n + \alpha_n}{n}$$

with $\lim_{n \rightarrow \infty} \alpha_n = -\infty$, there exists graph $G_s(n, \widetilde{K}_n, \widetilde{P}_n)$ under $\widetilde{P}_n = \Omega(n)$ and

$$(3.8) \quad \frac{1}{s!} \cdot \frac{\widetilde{K}_n^{2s}}{\widetilde{P}_n^s} = \frac{\ln n + (k-1) \ln \ln n + \widetilde{\alpha}_n}{n}$$

¹A graph property is called monotone increasing if it holds under the addition of edges [48, 49].

with $\lim_{n \rightarrow \infty} \widetilde{\alpha}_n = -\infty$ and $\widetilde{\alpha}_n = -O(\ln \ln n)$, such that there exists a graph coupling² under which $G_s(n, K_n, P_n)$ is a spanning subgraph of $G_s(n, \widetilde{K}_n, \widetilde{P}_n)$.

(b) For graph $G_s(n, K_n, P_n)$ under $P_n = \Omega(n)$ and

$$(3.9) \quad \frac{1}{s!} \cdot \frac{K_n^{2s}}{P_n^s} = \frac{\ln n + (k-1) \ln \ln n + \alpha_n}{n}$$

with $\lim_{n \rightarrow \infty} \alpha_n = \infty$, there exists graph $G_s(n, \widehat{K}_n, \widehat{P}_n)$ under $\widehat{P}_n = \Omega(n)$ and

$$(3.10) \quad \frac{1}{s!} \cdot \frac{\widehat{K}_n^{2s}}{\widehat{P}_n^s} = \frac{\ln n + (k-1) \ln \ln n + \widehat{\alpha}_n}{n}$$

with $\lim_{n \rightarrow \infty} \widehat{\alpha}_n = \infty$ and $\widehat{\alpha}_n = O(\ln \ln n)$, such that there exists a graph coupling under which $G_s(n, K_n, P_n)$ is a spanning supergraph of $G_s(n, \widehat{K}_n, \widehat{P}_n)$.

The proof of Lemma 3.1 is provided in Section 6.2 in the Appendix.

3.2 Relationships between vertex connectivity, edge connectivity, and minimum vertex degree.

Recall that the vertex connectivity of a graph is defined as the minimum number of vertices needing to be deleted to have the remaining graph disconnected, and the edge connectivity is defined similarly for the deletion of edges [48]. For graph $G_s(n, K_n, P_n)$, we use κ_v, κ_e and δ to denote the vertex connectivity, the edge connectivity, and the minimum vertex degree, respectively. Then k -vertex-connectivity, k -edge-connectivity, and the property of minimum vertex degree being at least k , are given by events $\kappa_v \geq k$, $\kappa_e \geq k$, and $\delta \geq k$, respectively. For any graph, the vertex connectivity is at most the edge connectivity, and the edge connectivity is at most the minimum vertex degree [46–48]. Therefore, $\kappa_v \leq \kappa_e \leq \delta$ holds. Then

$$(3.11) \quad \mathbb{P}[\kappa_v \geq k] \leq \mathbb{P}[\kappa_e \geq k] \leq \mathbb{P}[\delta \geq k],$$

and

$$(3.12) \quad \begin{aligned} \mathbb{P}[\kappa_v \geq k] &= \mathbb{P}[\delta \geq k] - \mathbb{P}[(\kappa_v < k) \cap (\delta \geq k)] \\ &\geq \mathbb{P}[\delta \geq k] - \sum_{\ell=0}^{k-1} \mathbb{P}[(\kappa_v = \ell) \cap (\delta > \ell)]. \end{aligned}$$

²As used by Rybarczyk [20, 21], a coupling of two random graphs G_1 and G_2 means a probability space on which random graphs G'_1 and G'_2 are defined such that G'_1 and G'_2 have the same distributions as G_1 and G_2 , respectively. If G'_1 is a spanning subgraph (resp., supergraph) of G'_2 , we say that under the coupling, G_1 is a spanning subgraph (resp., supergraph) of G_2 , which yields that for any monotone increasing property \mathcal{I} , the probability of G_1 having \mathcal{I} is at most (resp., at least) the probability of G_2 having \mathcal{I} .

Therefore, the proof is completed once we show Lemmas 3.2 and 3.3 below. Note that since k is a constant, condition (2.3) with $|\alpha_n| = O(\ln \ln n)$ in Theorem 2.1 implies condition (3.14) in Lemma 3.3.

LEMMA 3.2. ([15, THEOREM 2] (OUR WORK)) For uniform random s -intersection graph $G_s(n, K_n, P_n)$ under $P_n = \Omega(n)$, if there exists sequence α_n satisfying $|\alpha_n| = O(\ln \ln n)$ such that

$$\frac{1}{s!} \cdot \frac{K_n^{2s}}{P_n^s} = \frac{\ln n + (k-1) \ln \ln n + \alpha_n}{n},$$

then with δ denoting the minimum vertex degree, it holds that

$$(3.13)$$

$$\lim_{n \rightarrow \infty} \mathbb{P}[\delta \geq k] = e^{-\frac{e^{-\alpha^*}}{(k-1)!}}, \text{ if } \lim_{n \rightarrow \infty} \alpha_n = \alpha^* \in [-\infty, \infty].$$

LEMMA 3.3. For uniform random s -intersection graph $G_s(n, K_n, P_n)$ under $P_n = \Omega(n)$ and

$$(3.14) \quad \frac{1}{s!} \cdot \frac{K_n^{2s}}{P_n^s} = \frac{\ln n \pm O(\ln \ln n)}{n},$$

then with κ_v denoting the vertex connectivity and δ denoting the minimum vertex degree, it holds for constant integer ℓ that

$$(3.15) \quad \mathbb{P}[(\kappa_v = \ell) \cap (\delta > \ell)] = o(1).$$

We detail the proof of Lemma 3.3 below.

3.3 The proof of Lemma 3.3. For graph $G_s(n, K_n, P_n)$, let the set of vertices be $\mathcal{V}_n = \{v_1, v_2, \dots, v_n\}$. Also, for $i = 1, 2, \dots, n$, we let S_i denote the set of items on vertex v_i . We introduce event $\mathcal{E}(\mathbf{J})$ in the following manner:

$$(3.16) \quad \mathcal{E}(\mathbf{J}) = \bigcup_{\substack{T \subseteq \mathcal{V}_n, \\ |T| \geq 2}} [|\cup_{v_i \in T} S_i| \leq J_{|T|}],$$

where $\mathbf{J} = [J_2, J_3, \dots, J_n]$ is an $(n-1)$ -dimensional integer valued array, with J_i defined through

$$(3.17) \quad J_i = \begin{cases} \max\{\lfloor (1 + \varepsilon_1) K_n \rfloor, \lfloor \lambda_1 K_n i \rfloor\}, & i = 2, \dots, r_n, \\ \lfloor \mu_1 P_n \rfloor, & i = r_n + 1, \dots, n, \end{cases}$$

for an arbitrary constant $0 < \varepsilon_1 < 1$ and some positive constants λ_1, μ_1 in Lemma 3.4 below, where $r_n := \min(\lfloor \frac{P_n}{K_n} \rfloor, \lfloor \frac{n}{2} \rfloor)$.

By a crude bounding argument, we get

$$(3.18) \quad \begin{aligned} \mathbb{P}[(\kappa_v = \ell) \cap (\delta > \ell)] \\ \leq \mathbb{P}[\mathcal{E}(\mathbf{J})] + \mathbb{P}[(\kappa_v = \ell) \cap (\delta > \ell) \cap \overline{\mathcal{E}(\mathbf{J})}]. \end{aligned}$$

Hence, a proof of Lemma 3.3 consists of proving two lemmas below. Under (2.3) with $|\alpha_n| = O(\ln \ln n)$, we have $\frac{1}{s!} \cdot \frac{K_n^{2s}}{P_n^s} = \frac{\ln n \pm O(\ln \ln n)}{n} = o(1)$ and $\frac{K_n}{P_n} = o(1)$, enabling us to use Lemmas 3.4 and 3.5.

LEMMA 3.4. ([5, PROPOSITION 3] (OUR WORK))
If $P_n = \Omega(n)$ and $\frac{K_n}{P_n} = o(1)$, then for an arbitrary constant $0 < \varepsilon_1 < 1$ and some selected positive constants λ_1, μ_1 , it holds that

$$(3.19) \quad \mathbb{P}[\mathcal{E}(\mathbf{J})] = o(1).$$

LEMMA 3.5. For uniform random s -intersection graph $G_s(n, K_n, P_n)$ under $P_n = \Omega(n)$ and

$$(3.20) \quad \frac{1}{s!} \cdot \frac{K_n^{2s}}{P_n^s} = \frac{\ln n \pm O(\ln \ln n)}{n},$$

then

$$(3.21) \quad \mathbb{P}[(\kappa_v = \ell) \cap (\delta > \ell) \cap \overline{\mathcal{E}(\mathbf{J})}] = o(1).$$

The proof of Lemma 3.5 is given in Section 6.3 in the Appendix.

4 Establishing Theorem 2.2

Similar to the idea of confining $|\alpha_n|$ in Theorem 2.1, here we confine $|\beta_n|$ as $O(\ln \ln n)$ in Theorem 2.2. Specifically, we will demonstrate

$$(4.22)$$

Theorem 2.2 under $|\beta_n| = O(\ln \ln n) \Rightarrow$ Theorem 2.2.

Since k -vertex-connectivity, k -edge-connectivity, and the property of minimum vertex degree being at least k , are all monotone increasing, then to show (4.22), it suffices to prove the following lemma.

LEMMA 4.1. (a) For graph $H_s(n, t_n, P_n)$ under

$$(4.23) \quad \frac{1}{s!} \cdot t_n^{2s} P_n^s = \frac{\ln n + (k-1) \ln \ln n + \beta_n}{n}$$

with $\lim_{n \rightarrow \infty} \beta_n = -\infty$, there exists graph $H_s(n, \tilde{t}_n, \tilde{P}_n)$ under

$$(4.24) \quad \frac{1}{s!} \cdot \tilde{t}_n^{2s} \tilde{P}_n^s = \frac{\ln n + (k-1) \ln \ln n + \tilde{\beta}_n}{n}$$

with $\lim_{n \rightarrow \infty} \tilde{\beta}_n = -\infty$ and $\tilde{\beta}_n = -O(\ln \ln n)$ such that there exists a graph coupling under which $H_s(n, t_n, P_n)$ is a spanning subgraph of $H_s(n, \tilde{t}_n, \tilde{P}_n)$.

(b) For graph $H_s(n, t_n, P_n)$ under

$$(4.25) \quad \frac{1}{s!} \cdot t_n^{2s} P_n^s = \frac{\ln n + (k-1) \ln \ln n + \beta_n}{n}$$

with $\lim_{n \rightarrow \infty} \beta_n = \infty$, there exists graph $H_s(n, \hat{t}_n, \hat{P}_n)$ under

$$(4.26) \quad \frac{1}{s!} \cdot \hat{t}_n^{2s} \hat{P}_n^s = \frac{\ln n + (k-1) \ln \ln n + \hat{\beta}_n}{n}$$

with $\lim_{n \rightarrow \infty} \hat{\beta}_n = \infty$ and $\hat{\beta}_n = O(\ln \ln n)$ such that there exists a graph coupling under which $H_s(n, t_n, P_n)$ is a spanning supergraph of $H_s(n, \hat{t}_n, \hat{P}_n)$.

The proof of Lemma 4.1 is detailed in Section 6.4 in the Appendix.

Now we use Theorem 2.1 to prove Theorem 2.2 with confined $|\beta_n| = O(\ln \ln n)$. Here, the main idea is to exploit a coupling result between the uniform s -intersection graph and a binomial s -intersection graph. Let \mathcal{I}^* denote either one of the following graph properties: k -vertex-connectivity, k -edge-connectivity, and the property of minimum vertex degree being at least k . With K_n^- and K_n^+ defined by

$$(4.27) \quad K_n^\pm = t_n P_n \pm \sqrt{3 \ln n (\ln n + t_n P_n)},$$

we have from Lemma 6.5 that if $t_n P_n = \omega(\ln n)$, then

$$(4.28) \quad \begin{aligned} &\mathbb{P}[\text{Graph } G_s(n, K_n^-, P_n) \text{ has } \mathcal{I}^*] - o(1) \\ &\leq \mathbb{P}[\text{Graph } H_s(n, t_n, P_n) \text{ has } \mathcal{I}^*] \\ &\leq \mathbb{P}[\text{Graph } G_s(n, K_n^+, P_n) \text{ has } \mathcal{I}^*] + o(1). \end{aligned}$$

Under conditions (2.4), (2.5), and $|\beta_n| = O(\ln \ln n)$, we now show that $t_n P_n = \omega(\ln n)$. From (2.5) and $|\beta_n| = O(\ln \ln n)$, we first get

$$(4.29) \quad \frac{1}{s!} \cdot t_n^{2s} P_n^s = \frac{\ln n \pm O(\ln \ln n)}{n} = \frac{\ln n}{n} \cdot [1 \pm o(1)].$$

From (2.4) and (4.29), it follows that

$$(4.30) \quad \begin{aligned} t_n P_n &= \sqrt{t_n^2 P_n} \cdot \sqrt{P_n} \\ &= \begin{cases} \{s! n^{-1} \ln n \cdot [1 \pm o(1)]\}^{\frac{1}{2s}} \cdot \sqrt{\Omega(n)}, & \text{for } s \geq 2, \\ \{s! n^{-1} \ln n \cdot [1 \pm o(1)]\}^{\frac{1}{2s}} \cdot \sqrt{\Omega(n^c)}, & \text{for } s = 1, \end{cases} \\ &= \begin{cases} \Omega(n^{\frac{1}{2} - \frac{1}{2s}} (\ln n)^{\frac{1}{2s}}), & \text{for } s \geq 2, \\ \Omega(n^{\frac{c-1}{2}} (\ln n)^{\frac{1}{2}}), & \text{for } s = 1, \end{cases} \end{aligned}$$

yielding $t_n P_n = \omega(\ln n)$ in view of $c > 1$, so we can use (4.28). Using (4.27) and (4.30), we further obtain

$$(4.31) \quad \begin{aligned} \frac{(K_n^\pm)^{2s}}{P_n^s} &= \frac{[t_n P_n \pm \sqrt{3 \ln n (\ln n + t_n P_n)}]^{2s}}{P_n^s} \\ &= \frac{(t_n P_n)^{2s}}{P_n^s} \cdot \left[1 \pm \sqrt{\frac{3 \ln n}{t_n P_n} \left(\frac{\ln n}{t_n P_n} + 1 \right)} \right]^{2s} \\ &= t_n^{2s} P_n^s \cdot \left[1 \pm o\left(\frac{1}{\ln n}\right) \right], \end{aligned}$$

where in the last step we use $t_n P_n = \omega((\ln n)^3)$, which follows from (4.30) due to constant $c > 1$.

Applying (2.5) and (4.29) to (4.31), we have

$$(4.32) \quad \frac{1}{s!} \cdot \frac{(K_n^\pm)^{2s}}{P_n^s} = \frac{\ln n + (k-1) \ln \ln n + \beta_n \pm o(1)}{n}.$$

In view of (4.32) and $P_n = \Omega(n)$, we use Theorem 2.1 to obtain

$$(4.33) \quad \lim_{n \rightarrow \infty} \mathbb{P}[\text{Graph } G_s(n, K_n^\pm, P_n) \text{ has } \mathcal{I}^*] = e^{-\frac{e^{-\lim_{n \rightarrow \infty} [\beta_n \pm o(1)]}}{(k-1)!}} = e^{-\frac{e^{-\lim_{n \rightarrow \infty} \beta_n}}{(k-1)!}}.$$

The proof of Theorem 2.2 is completed by (4.28) and (4.33).

5 Conclusion

Random s -intersection graphs have been used in a wide range of applications. Two extensively studied models are a uniform random s -intersection graph and a binomial random s -intersection graph. In this paper, for a uniform/binomial random s -intersection graph, we derive exact asymptotic expressions for the probabilities of the following three properties: (i) k -vertex-connectivity, (ii) k -edge-connectivity and (iii) the property that each vertex has degree at least k .

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6 Appendix

We first present in Section 6.1 additional lemmas used in proving the theorems. Afterwards, we detail the proofs of the lemmas.

6.1 Additional lemmas. Some additional lemmas are given below. The relation " \sim " stands for an asymptotical equivalence; i.e., $f_n \sim g_n$ means $\lim_{n \rightarrow \infty} (f_n/g_n) = 1$.

LEMMA 6.1. If $\frac{1}{s!} \cdot \frac{K_n^{2s}}{P_n^s} = \frac{\ln n \pm O(\ln \ln n)}{n}$ and $P_n = \Omega(n^c)$ for constant c , then $K_n = \Omega(n^{\frac{c}{2} - \frac{1}{2s}} (\ln n)^{\frac{1}{2s}})$.

LEMMA 6.2. The following properties (a) and (b) hold, where q_n is the edge probability in uniform random s -intersection graph $G_s(n, K_n, P_n)$.

- (a) If $P_n = \Omega(n)$ and $\frac{1}{s!} \cdot \frac{K_n^{2s}}{P_n^s} = \frac{\ln n \pm O(\ln \ln n)}{n}$, then $q_n \sim \frac{1}{s!} \cdot \frac{K_n^{2s}}{P_n^s}$ and $|q_n - \frac{1}{s!} \cdot \frac{K_n^{2s}}{P_n^s}| = o(\frac{1}{n})$.
- (b) If $P_n = \Omega(n)$ and $q_n = \frac{\ln n \pm O(\ln \ln n)}{n}$, then $q_n \sim \frac{1}{s!} \cdot \frac{K_n^{2s}}{P_n^s}$ and $|q_n - \frac{1}{s!} \cdot \frac{K_n^{2s}}{P_n^s}| = o(\frac{1}{n})$.

LEMMA 6.3. For uniform random s -intersection graph $G_s(n, K_n, P_n)$ under $K_n = \omega(1)$, the following properties (a) (b) and (c) hold for $i = r+1, r+2, \dots, n$ (i.e., vertex $v_i \notin \{v_1, v_2, \dots, v_r\}$), where E_{ij} denotes the event that an edge exists between vertices v_i and v_j , S_i is the number of items on vertex v_i , and q_n is the edge probability.

- (a) If $|\bigcup_{j=1}^r S_j| \geq \lfloor (1 + \varepsilon_1)K_n \rfloor$ for some positive constant ε_1 , then for any positive constant $\varepsilon_2 < (1 + \varepsilon_1)^s - 1$, it holds for all n sufficiently large that

$$(6.34) \quad \mathbb{P} \left[\bigcap_{j=1}^r \overline{E_{ij}} \mid S_1, S_2, \dots, S_r \right] \leq e^{-q_n(1+\varepsilon_2)}.$$

- (b) If $|\bigcup_{j=1}^r S_j| \geq \lfloor \lambda_1 r K_n \rfloor$ for some positive constant λ_1 , then for any positive constant $\lambda_2 < \lambda_1^s$, it holds for all n sufficiently large that

$$(6.35) \quad \mathbb{P} \left[\bigcap_{j=1}^r \overline{E_{ij}} \mid S_1, S_2, \dots, S_r \right] \leq e^{-\lambda_2 r q_n}.$$

- (c) If $|\bigcup_{j=1}^r S_j| \geq \lfloor \mu_1 P_n \rfloor$ for some positive constant μ_1 , then for any positive constant $\mu_2 < (s!)^{-1} \mu_1^s$, it holds for all n sufficiently large that

$$(6.36) \quad \mathbb{P} \left[\bigcap_{j=1}^r \overline{E_{ij}} \mid S_1, S_2, \dots, S_r \right] \leq e^{-\mu_2 K_n}.$$

LEMMA 6.4. For uniform random s -intersection graph $G_s(n, K_n, P_n)$ under $P_n = \Omega(n)$, $K_n = \omega(1)$ and $r_n := \min(\lfloor \frac{P_n}{K_n} \rfloor, \lfloor \frac{n}{2} \rfloor) = \omega(1)$, the following properties (a) (b) and (c) hold for any constant integer $R \geq 2$, where ε_1 , λ_1 and μ_1 are specified in Lemma 3.4, and events $\mathcal{A}_{\ell, r}$ and $\mathcal{E}(\mathbf{J})$ are defined in Sections 6.3 and 3.3, respectively.

- (a) Let ε_3 be any positive constant with $\varepsilon_3 < (1 + \varepsilon_1)^s - 1$. For all n sufficiently large, it holds for $r = 2, 3, \dots, R$ that

$$\mathbb{P} [\mathcal{A}_{\ell, r} \cap \overline{\mathcal{E}(\mathbf{J})}] \leq r^{r-2} q_n^{r-1} (r q_n)^\ell e^{-q_n n(1+\varepsilon_3)}.$$

- (b) Let λ_2 be any positive constant with $\lambda_2 < \lambda_1^s$. For all n sufficiently large, it holds for $r = R+1, R+2, \dots, r_n$ that

$$\mathbb{P} [\mathcal{A}_{\ell, r} \cap \overline{\mathcal{E}(\mathbf{J})}] \leq r^{r-2} q_n^{r-1} e^{-\lambda_2 r q_n n/3}.$$

- (c) Let μ_2 be any positive constant with $\mu_2 < (s!)^{-1} \mu_1^s$. For all n sufficiently large, it holds for $r = r_n + 1, r_n + 2, \dots, \lfloor \frac{n-\ell}{2} \rfloor$ that

$$\mathbb{P} [\mathcal{A}_{\ell, r} \cap \overline{\mathcal{E}(\mathbf{J})}] \leq e^{-\mu_2 K_n n/3}.$$

LEMMA 6.5. ([12, LEMMA 4]) Let K_n^- and K_n^+ denote $t_n P_n - \sqrt{3 \ln n (\ln n + t_n P_n)}$ and $t_n P_n + \sqrt{3 \ln n (\ln n + t_n P_n)}$, respectively. If $t_n P_n = \omega(\ln n)$, then for any monotone increasing graph property \mathcal{I} , it holds that

$$\begin{aligned} & \mathbb{P} [\text{Graph } G_s(n, K_n^-, P_n) \text{ has } \mathcal{I}.] - o(1) \\ & \leq \mathbb{P} [\text{Graph } H_s(n, t_n, P_n) \text{ has } \mathcal{I}.] \\ & \leq \mathbb{P} [\text{Graph } G_s(n, K_n^+, P_n) \text{ has } \mathcal{I}.] + o(1). \end{aligned}$$

6.2 Proof of Lemma 3.1

Proving property (a). We define $\widetilde{\alpha}_n^*$ by

$$(6.37) \quad \widetilde{\alpha}_n^* = \max\{\alpha_n, -\ln \ln n\},$$

and define \widetilde{K}_n^* such that

$$(6.38) \quad \frac{1}{s!} \cdot \frac{(\widetilde{K}_n^*)^{2s}}{P_n^s} = \frac{\ln n + (k-1) \ln \ln n + \widetilde{\alpha}_n^*}{n}.$$

We set

$$(6.39) \quad \widetilde{K}_n := \lfloor \widetilde{K}_n^* \rfloor,$$

and

$$(6.40) \quad \widetilde{P}_n := P_n.$$

From (3.9) (6.37) and (6.38), it holds that

$$(6.41) \quad K_n \leq \widetilde{K}_n^*.$$

Then by (6.39) (6.41) and the fact that K_n and \widetilde{K}_n are both integers, it follows that

$$(6.42) \quad K_n \leq \widetilde{K}_n.$$

From (6.40) and (6.42), by [12, Lemma 3], there exists a graph coupling under which $G_s(n, K_n, P_n)$ is a spanning subgraph of $G_s(n, \widetilde{K}_n, \widetilde{P}_n)$. Therefore, the proof of property (a) is completed once we show $\widetilde{\alpha}_n$ defined in (3.8) satisfies

$$(6.43) \quad \lim_{n \rightarrow \infty} \widetilde{\alpha}_n = -\infty,$$

$$(6.44) \quad \widetilde{\alpha}_n = -O(\ln \ln n).$$

We first prove (6.43). From (3.8) (6.38) and (6.39), it holds that

$$(6.45) \quad \widetilde{\alpha}_n \leq \widetilde{\alpha}_n^*,$$

which together with (6.37) and $\lim_{n \rightarrow \infty} \alpha_n = -\infty$ yields (6.43).

Now we establish (6.44). From (6.39), we have $\widetilde{K}_n > \widetilde{K}_n^* - 1$. Then from (3.8) and (6.40), it holds that

$$(6.46) \quad \begin{aligned} \widetilde{\alpha}_n &= n \cdot \frac{1}{s!} \cdot \frac{\widetilde{K}_n^{2s}}{P_n^s} - [\ln n + (k-1) \ln \ln n] \\ &> n \cdot \frac{1}{s!} \cdot \frac{(\widetilde{K}_n^* - 1)^{2s}}{P_n^s} - [\ln n + (k-1) \ln \ln n]. \end{aligned}$$

By $\lim_{n \rightarrow \infty} \alpha_n = -\infty$, it holds that $\alpha_n \leq 0$ for all n sufficiently large. Then from (6.37), it follows that

$$(6.47) \quad \widetilde{\alpha}_n^* = -O(\ln \ln n),$$

which along with Lemma 6.1, equation (6.38) and condition $P_n = \Omega(n)$ induces

$$(6.48) \quad \widetilde{K}_n^* = \Omega((\ln n)^{\frac{1}{2s}}).$$

Hence, we have $\lim_{n \rightarrow \infty} \widetilde{K}_n^* = \infty$ and it further holds for all n sufficient large that

$$(6.49) \quad (\widetilde{K}_n^* - 1)^{2s} > (\widetilde{K}_n^*)^{2s} - 3s(\widetilde{K}_n^*)^{2s-1}.$$

Applying (6.49) to (6.46) and then using (6.38), Lemma 6.1 and $P_n = \Omega(n)$, it follows that

$$(6.50) \quad \begin{aligned} \widetilde{\alpha}_n &> \frac{n}{s!} \cdot \frac{(\widetilde{K}_n^*)^{2s} - 3s(\widetilde{K}_n^*)^{2s-1}}{P_n^s} - [\ln n + (k-1) \ln \ln n] \\ &= \widetilde{\alpha}_n^* - \frac{3s}{s!} \cdot n \cdot \Theta(P_n^{-\frac{1}{2}} n^{-\frac{2s-1}{2s}} (\ln n)^{\frac{2s-1}{2s}}) \\ &= \widetilde{\alpha}_n^* - O(n^{-\frac{1}{2} + \frac{1}{2s}} (\ln n)^{1 - \frac{1}{2s}}). \end{aligned}$$

As noted at the beginning of Section 3, our proof is for $s \geq 2$ since the case of $s = 1$ already is proved by us [25]. Using $s \geq 2$ in (6.50), it holds that $\widetilde{\alpha}_n > \widetilde{\alpha}_n^* + o(1)$, which along with (6.45) and (6.47) yields (6.44).

Proving property (b). We define $\widehat{\alpha}_n^*$ by

$$(6.51) \quad \widehat{\alpha}_n^* = \min\{\alpha_n, \ln \ln n\},$$

and define \widehat{K}_n^* such that

$$(6.52) \quad \frac{1}{s!} \cdot \frac{(\widehat{K}_n^*)^{2s}}{P_n^s} = \frac{\ln n + (k-1) \ln \ln n + \widehat{\alpha}_n^*}{n}.$$

We set

$$(6.53) \quad \widehat{K}_n := \lceil \widehat{K}_n^* \rceil,$$

and

$$(6.54) \quad \widehat{P}_n := P_n.$$

From (3.9) (6.51) and (6.52), it holds that

$$(6.55) \quad K_n \geq \widehat{K}_n^*.$$

Then by (6.53) (6.55) and the fact that K_n and \widehat{K}_n are both integers, it follows that

$$(6.56) \quad K_n \geq \widehat{K}_n.$$

From (6.54) and (6.56), by [12, Lemma 3], there exists a graph coupling under which $G_s(n, K_n, P_n)$ is a spanning supergraph of $G_s(n, \widehat{K}_n, \widehat{P}_n)$. Therefore, the proof of property (b) is completed once we show $\widehat{\alpha}_n$ defined in (3.10) satisfies

$$(6.57) \quad \lim_{n \rightarrow \infty} \widehat{\alpha}_n = \infty,$$

$$(6.58) \quad \widehat{\alpha}_n = O(\ln \ln n).$$

We first prove (6.57). From (3.10) (6.52) and (6.53), it holds that

$$(6.59) \quad \widehat{\alpha}_n \geq \widehat{\alpha}_n^*,$$

which together with (6.51) and $\lim_{n \rightarrow \infty} \alpha_n = \infty$ yields (6.57).

Now we establish (6.58). From (6.53), we have $\widehat{K}_n < \widehat{K}_n^* + 1$. Then from (3.10) and (6.54), it holds that

$$(6.60) \quad \begin{aligned} \widehat{\alpha}_n &= n \cdot \frac{1}{s!} \cdot \frac{\widehat{K}_n^{2s}}{P_n^s} - [\ln n + (k-1) \ln \ln n] \\ &< n \cdot \frac{1}{s!} \cdot \frac{(\widehat{K}_n^* + 1)^{2s}}{P_n^s} - [\ln n + (k-1) \ln \ln n]. \end{aligned}$$

By $\lim_{n \rightarrow \infty} \alpha_n = \infty$, it holds that $\alpha_n \geq 0$ for all n sufficiently large. Then from (6.51), it follows that

$$(6.61) \quad \widehat{\alpha}_n^* = O(\ln \ln n),$$

which along with Lemma 6.1, equation (6.52) and condition $P_n = \Omega(n)$ induces

$$(6.62) \quad \widehat{K}_n^* = \Omega((\ln n)^{\frac{1}{2s}}).$$

Hence, we have $\lim_{n \rightarrow \infty} \widehat{K}_n^* = \infty$ and it further holds for all n sufficient large that

$$(6.63) \quad (\widehat{K}_n^* + 1)^{2s} < (\widehat{K}_n^*)^{2s} + 3s(\widehat{K}_n^*)^{2s-1}.$$

Applying (6.63) to (6.60) and then using (6.52), Lemma 6.1 and $P_n = \Omega(n)$, it follows that

$$\begin{aligned}
 (6.64) \quad & \widehat{\alpha}_n \\
 & < \frac{n}{s!} \cdot \frac{(\widehat{K}_n^*)^{2s} + 3s(\widehat{K}_n^*)^{2s-1}}{P_n^s} - [\ln n + (k-1) \ln \ln n] \\
 & = \widehat{\alpha}_n^* + \frac{3s}{s!} \cdot n \cdot \Theta(P_n^{-\frac{1}{2}} n^{-\frac{2s-1}{2s}} (\ln n)^{\frac{2s-1}{2s}}) \\
 & = \widehat{\alpha}_n^* + O(n^{-\frac{1}{2} + \frac{1}{2s}} (\ln n)^{1 - \frac{1}{2s}}).
 \end{aligned}$$

As noted at the beginning of Section 3, our proof is for $s \geq 2$ since the case of $s = 1$ already is proved by Rybarczyk [28]. Using $s \geq 2$ in (6.64), it holds that $\widehat{\alpha}_n < \widehat{\alpha}_n^* + o(1)$, which along with (6.59) and (6.61) yields (6.58).

6.3 The proof of Lemma 3.5. By the analysis in [5, Section IV], we obtain [5, Equation (148)]. Namely, with some events defined as follows:

- \mathcal{C}_r : event that the induced subgraph of $G_s(n, K_n, P_n)$ defined on vertex set $\{v_1, v_2, \dots, v_r\}$ is connected,
- $\mathcal{B}_{\ell,r}$: event that any vertex in $\{v_{r+1}, v_{r+2}, \dots, v_{r+\ell}\}$ has an edge with at least one vertex in $\{v_1, v_2, \dots, v_r\}$,
- $\mathcal{D}_{\ell,r}$: event that any vertex in $\{v_{r+\ell+1}, v_{r+\ell+2}, \dots, v_n\}$ and any vertex in $\{v_1, v_2, \dots, v_r\}$ has no edge in between, and
- $\mathcal{A}_{\ell,r}$: event that events \mathcal{C}_r , $\mathcal{B}_{\ell,r}$ and $\mathcal{D}_{\ell,r}$ all happen,

it holds that

$$\begin{aligned}
 (6.65) \quad & \mathbb{P}[(\kappa = \ell) \cap (\delta > \ell) \cap \overline{\mathcal{E}(\mathcal{J})}] \\
 & \leq \sum_{r=2}^{\lfloor \frac{n-\ell}{2} \rfloor} \binom{n}{\ell} \binom{n-\ell}{r} \mathbb{P}[\mathcal{A}_{\ell,r} \cap \overline{\mathcal{E}(\mathcal{J})}].
 \end{aligned}$$

The proof of Lemma 3.5 is completed once we show the following three results:

$$(6.66) \quad \sum_{r=2}^R \binom{n}{\ell} \binom{n-\ell}{r} \mathbb{P}[\mathcal{A}_{\ell,r} \cap \overline{\mathcal{E}(\mathcal{J})}] = o(1),$$

$$(6.67) \quad \sum_{r=R+1}^{r_n} \binom{n}{\ell} \binom{n-\ell}{r} \mathbb{P}[\mathcal{A}_{\ell,r} \cap \overline{\mathcal{E}(\mathcal{J})}] = o(1),$$

and

$$(6.68) \quad \sum_{r=r_n+1}^{\lfloor \frac{n-\ell}{2} \rfloor} \binom{n}{\ell} \binom{n-\ell}{r} \mathbb{P}[\mathcal{A}_{\ell,r} \cap \overline{\mathcal{E}(\mathcal{J})}] = o(1),$$

where $r_n = \min(\lfloor \frac{P_n}{K_n} \rfloor, \lfloor \frac{n}{2} \rfloor)$.

From condition (3.20), it follows that $\frac{K_n}{P_n} = o(1)$, yielding $r_n = \omega(1)$. From conditions (3.20) and $P_n = \Omega(n)$, we use Lemma 6.1 to derive $K_n = \omega(1)$. Therefore, we have $P_n = \Omega(n)$, $K_n = \omega(1)$ and $r_n = \omega(1)$, enabling us to use Lemma 6.4.

In addition, given conditions (3.20) and $P_n = \Omega(n)$, we use Lemma 6.2 to obtain

$$(6.69) \quad q_n = \frac{\ln n \pm O(\ln \ln n)}{n}.$$

Hence, it holds that

$$(6.70) \quad q_n \leq \frac{2 \ln n}{n}, \text{ for all } n \text{ sufficiently large,}$$

and there exists constant c_0 such that

$$(6.71) \quad q_n \geq \frac{\ln n - c_0 \ln \ln n}{n}, \text{ for all } n \text{ sufficiently large.}$$

6.3.1 Establishing (6.66). From $\binom{n}{\ell} \leq n^\ell$, $\binom{n-\ell}{r} \leq n^r$ and property (a) of Lemma 6.4, it follows that

$$\begin{aligned}
 (6.72) \quad & \binom{n}{\ell} \binom{n-\ell}{r} \mathbb{P}[\mathcal{A}_{\ell,r} \cap \overline{\mathcal{E}(\mathcal{J})}] \\
 & \leq n^\ell \cdot n^r \cdot r^{r-2} q_n^{r-1} (r q_n)^\ell \cdot e^{-q_n n(1+\varepsilon_3)} \\
 & = r^\ell r^{r-2} \cdot n^{\ell+r} q_n^{\ell+r-1} \cdot e^{-q_n n(1+\varepsilon_3)}.
 \end{aligned}$$

Applying (6.69) and (6.71) to (6.72), we get

$$\begin{aligned}
 & \binom{n}{\ell} \binom{n-\ell}{r} \mathbb{P}[\mathcal{A}_{\ell,r} \cap \overline{\mathcal{E}(\mathcal{J})}] \\
 & \leq r^\ell r^{r-2} n^{\ell+r} \left(\frac{2 \ln n}{n} \right)^{\ell+r-1} e^{-(1+\varepsilon_3)(\ln n - c_0 \ln \ln n)} \\
 & \leq 2^{\ell+r-1} r^{\ell+r-2} n^{-\varepsilon_3} (\ln n)^{\ell+r-1+c_0(1+\varepsilon_3)} \\
 & = o(1).
 \end{aligned}$$

Since R is a constant, (6.66) clearly follows.

6.3.2 Establishing (6.67). From $\binom{n}{\ell} \leq n^\ell$, $\binom{n-\ell}{r} \leq \left(\frac{e(n-\ell)}{r} \right)^r \leq \left(\frac{en}{r} \right)^r$ and property (b) of Lemma 6.4, we have

$$\begin{aligned}
 (6.73) \quad & \binom{n}{\ell} \binom{n-\ell}{r} \mathbb{P}[\mathcal{A}_{\ell,r} \cap \overline{\mathcal{E}(\mathcal{J})}] \\
 & \leq n^\ell \cdot \left(\frac{e(n-\ell)}{r} \right)^r \cdot r^{r-2} q_n^{r-1} e^{-\lambda_2 r q_n n/3} \\
 & \leq n^{\ell+r} e^r q_n^{r-1} e^{-\lambda_2 r q_n n/3}.
 \end{aligned}$$

Applying (6.69) and (6.71) to (6.73), we get

$$\begin{aligned}
 (6.74) \quad & \binom{n}{\ell} \binom{n-\ell}{r} \mathbb{P}[\mathcal{A}_{\ell,r} \cap \overline{\mathcal{E}(\mathbf{J})}] \\
 & \leq n^{\ell+r} e^r \cdot \left(\frac{2 \ln n}{n} \right)^{r-1} \cdot e^{-\lambda_2 r (\ln n - c_0 \ln \ln n)/3} \\
 & \leq n^{\ell+1} \cdot (2en^{-\lambda_2/3} (\ln n)^{c_0 \lambda_2/3+1})^r.
 \end{aligned}$$

Given $2en^{-\lambda_2/3} (\ln n)^{c_0 \lambda_2/3+1} = o(1)$ and (6.74), we obtain

$$\begin{aligned}
 & \sum_{r=R+1}^{r_n} \binom{n}{\ell} \binom{n-\ell}{r} \mathbb{P}[\mathcal{A}_{\ell,r} \cap \overline{\mathcal{E}(\mathbf{J})}] \\
 & \leq \sum_{r=R+1}^{\infty} n^{\ell+1} \cdot (2en^{-\lambda_2/3} (\ln n)^{c_0 \lambda_2/3+1})^r \\
 & = n^{\ell+1} \cdot \frac{(2en^{-\lambda_2/3} (\ln n)^{c_0 \lambda_2/3+1})^{R+1}}{1 - 2en^{-\lambda_2/3} (\ln n)^{c_0 \lambda_2/3+1}} \\
 (6.75) \quad & \sim n^{\ell+1-\lambda_2(R+1)/3} (2e(\ln n)^{c_0 \lambda_2/3+1})^{R+1}.
 \end{aligned}$$

We pick constant $R \geq \frac{3(\ell+1)}{\lambda_2}$ so that $\ell+1 - \lambda_2(R+1)/3 \leq -\frac{\lambda_2}{3}$. As a result, we obtain

$$\text{R.H.S. of (6.75)} = o(1)$$

and thus establish (6.67).

6.3.3 Establishing (6.68). From $\binom{n}{\ell} \leq n^\ell$ and property (c) of Lemma 6.4, it holds that

$$\begin{aligned}
 (6.76) \quad & \sum_{r=r_n+1}^{\lfloor \frac{n-\ell}{2} \rfloor} \binom{n}{\ell} \binom{n-\ell}{r} \mathbb{P}[\mathcal{A}_{\ell,r} \cap \overline{\mathcal{E}(\mathbf{J})}] \\
 & \leq n^\ell \cdot e^{-\mu_2 K_n n/3} \cdot \sum_{r=r_n+1}^{\lfloor \frac{n-\ell}{2} \rfloor} \binom{n-\ell}{r}.
 \end{aligned}$$

Given conditions $P_n = \Omega(n)$ and (3.20), we use Lemma 6.1 to derive

$$K_n = \Omega(n^{\frac{1}{2} - \frac{1}{2s}} (\ln n)^{\frac{1}{2s}}) = \omega(1),$$

which yields

$$(6.77) \quad \mu_2 K_n/3 \geq 2 \ln 2, \text{ for all } n \text{ sufficiently large.}$$

We have

$$\begin{aligned}
 (6.78) \quad & \sum_{r=r_n+1}^{\lfloor \frac{n-\ell}{2} \rfloor} \binom{n-\ell}{r} \leq \sum_{r=r_n+1}^{\lfloor \frac{n-\ell}{2} \rfloor} \binom{n}{r} \leq \sum_{r=0}^n \binom{n}{r} = 2^n.
 \end{aligned}$$

Applying (6.77) and (6.78) to (6.76), we finally obtain

$$\begin{aligned}
 & \sum_{r=r_n+1}^{\lfloor \frac{n-\ell}{2} \rfloor} \binom{n}{\ell} \binom{n-\ell}{r} \mathbb{P}[\mathcal{A}_{\ell,r} \cap \overline{\mathcal{E}(\mathbf{J})}] \\
 & \leq n^\ell \cdot 2^n \cdot e^{-\mu_2 K_n n/3} \\
 & = e^{\ell \ln n + n \ln 2 - \mu_2 K_n n/3} \\
 & \leq e^{\ell \ln n - n \ln 2}, \text{ for all } n \text{ sufficiently large.}
 \end{aligned}$$

The result (6.68) clearly follows with $n \rightarrow \infty$.

6.4 Proof of Lemma 4.1

(a)

$$(6.79) \quad \widetilde{P}_n = P_n,$$

and

$$(6.80) \quad \widetilde{\beta}_n = \max\{\beta_n, -\ln \ln n\}.$$

Given (6.80) and $\lim_{n \rightarrow \infty} \widetilde{\beta}_n = -\infty$, we clearly obtain $\lim_{n \rightarrow \infty} \widetilde{\beta}_n = -\infty$ and $\widetilde{\beta}_n = -O(\ln \ln n)$.

It holds from (6.80) that $\widetilde{\beta}_n \geq \beta_n$, which along with (4.23) (4.24) and (6.79) yields $t_n \leq \widetilde{t}_n$. Under $t_n \leq \widetilde{t}_n$ and $\widetilde{P}_n = P_n$, by [20, Section 3], there exists a graph coupling under which $H_s(n, t_n, P_n)$ is a spanning subgraph of $H_s(n, \widetilde{t}_n, \widetilde{P}_n)$.

(b) We set

$$(6.81) \quad \widehat{P}_n = P_n,$$

and

$$(6.82) \quad \widehat{\beta}_n = \min\{\beta_n, \ln \ln n\}.$$

Given (6.82) and $\lim_{n \rightarrow \infty} \beta_n = \infty$, we clearly obtain $\lim_{n \rightarrow \infty} \widehat{\beta}_n = \infty$ and $\widehat{\beta}_n = O(\ln \ln n)$.

It holds from (6.82) that $\widehat{\beta}_n \leq \beta_n$, which along with (4.25) (4.26) and (6.81) yields $t_n \geq \widehat{t}_n$. Under $t_n \geq \widehat{t}_n$ and $\widehat{P}_n = P_n$, by [20, Section 3], there exists a graph coupling under which $H_s(n, t_n, P_n)$ is a spanning supergraph of $H_s(n, \widehat{t}_n, \widehat{P}_n)$.

6.5 Proof of Lemma 6.1. From condition

$$(6.83) \quad \frac{1}{s!} \cdot \frac{K_n^{2s}}{P_n^s} = \frac{\ln n \pm O(\ln \ln n)}{n} \sim \frac{\ln n}{n},$$

it holds that

$$(6.84) \quad \frac{K_n^2}{P_n} = \Theta(n^{-\frac{1}{s}} (\ln n)^{\frac{1}{s}}),$$

which along with condition $P_n = \Omega(n^c)$ yields

$$\begin{aligned}
 (6.85) \quad & K_n = \sqrt{P_n \cdot \Theta(n^{-\frac{1}{s}} (\ln n)^{\frac{1}{s}})} = \Omega(n^{\frac{c}{2} - \frac{1}{2s}} (\ln n)^{\frac{1}{2s}}).
 \end{aligned}$$

6.6 Proof of Lemma 6.2.

(a) We still have (6.83) and (6.84) here. Then setting c as 1 in (6.85), it holds that

$$(6.86) \quad K_n = \Omega\left(n^{\frac{1}{2} - \frac{1}{2s}} (\ln n)^{\frac{1}{2s}}\right).$$

Given (6.84) and (6.86), we use [15, Lemma 1] and [5, Lemma 8] to have

$$(6.87) \quad q_n = \begin{cases} \frac{1}{s!} \left(\frac{K_n^2}{P_n}\right)^s [1 \pm O\left(\frac{K_n^2}{P_n^2}\right) \pm O\left(\frac{1}{K_n}\right)], & \text{for } s \geq 2, \\ \frac{K_n^2}{P_n} [1 \pm O\left(\frac{K_n^2}{P_n^2}\right)], & \text{for } s = 1. \end{cases}$$

Now we use (6.84) (6.86) and (6.87) to derive $q_n \sim \frac{1}{s!} \cdot \frac{K_n^{2s}}{P_n^s}$ and $|q_n - \frac{1}{s!} \cdot \frac{K_n^{2s}}{P_n^s}| = o\left(\frac{1}{n}\right)$.

First, (6.84) and (6.86) imply $\frac{K_n^2}{P_n} = o(1)$ and $K_n = \omega(1)$, respectively, which are used in (6.87) to derive $q_n \sim \frac{1}{s!} \cdot \frac{K_n^{2s}}{P_n^s}$. Second, applying (6.84) and (6.86) directly to (6.87), we obtain the following two cases:

(i) For $s \geq 2$, it holds that

$$\begin{aligned} & \left| q_n - \frac{1}{s!} \left(\frac{K_n^2}{P_n}\right)^s \right| \\ &= \Theta\left(\frac{\ln n}{n}\right) \left[\pm O\left(n^{-\frac{1}{s}} (\ln n)^{\frac{1}{s}}\right) \pm O\left(n^{\frac{1}{2s} - \frac{1}{2}} (\ln n)^{-\frac{1}{2s}}\right) \right] \\ &= \pm o\left(\frac{1}{n}\right). \end{aligned}$$

(ii) For $s = 1$, it holds that

$$\left| q_n - \frac{K_n^2}{P_n} \right| = \pm O\left(\left(\frac{\ln n}{n}\right)^2\right) = \pm o\left(\frac{1}{n}\right).$$

Summarizing cases (i) and (ii) above, we have proved property (a) of Lemma 6.2.

(b) By [3, Lemma 6], the edge probability q_n satisfies

$$(6.88) \quad q_n \leq \frac{\left[\binom{K_n}{s}\right]^2}{\binom{P_n}{s}}.$$

From (6.88) and condition $q_n = \frac{\ln n \pm O(\ln \ln n)}{n}$, it holds that

$$\frac{\left[\binom{K_n}{s}\right]^2}{\binom{P_n}{s}} \geq \frac{\ln n}{n} \cdot [1 - o(1)],$$

which along with $\frac{\left[\binom{K_n}{s}\right]^2}{\binom{P_n}{s}} \leq \frac{1}{s!} \cdot \frac{K_n^{2s}}{(P_n - s + 1)^s}$ and $P_n = \Omega(n)$ leads to

$$K_n^{2s} \geq s! (P_n - s + 1)^s \cdot \frac{\ln n}{n} [1 - o(1)] = \Omega(n^{s-1} \ln n).$$

Therefore, it follows that

$$(6.89) \quad K_n = \Omega\left(n^{\frac{1}{2} - \frac{1}{2s}} (\ln n)^{\frac{1}{2s}}\right).$$

Note that for some n , if $P_n < 2K_n - s$, then two vertices share at least s items with probability 1, resulting in $q_n = 1$. Therefore, given condition $q_n = \frac{\ln n \pm O(\ln \ln n)}{n}$, we know that for all n sufficiently large, $P_n \geq 2K_n - s$ holds, so the probability that two vertices share exactly s items is expressed by $\binom{K_n}{s} \binom{P_n - K_n}{K_n - s} / \binom{P_n}{K_n}$. Then

$$\begin{aligned} (6.90) \quad q_n &= \mathbb{P}[\text{Two vertices share exactly } s \text{ items.}] \\ &= \binom{K_n}{s} \binom{P_n - K_n}{K_n - s} / \binom{P_n}{K_n} \\ &= \frac{1}{s!} \cdot \left[\prod_{i=0}^{s-1} (K_n - i) \right]^2 \cdot \frac{\prod_{i=0}^{K_n - s - 1} (P_n - K_n)}{\prod_{i=0}^{K_n - 1} (P_n - i)} \\ &\geq \frac{1}{s!} \cdot \frac{(K_n - s + 1)^{2s}}{P_n^s}, \end{aligned}$$

which together with condition $q_n = \frac{\ln n \pm O(\ln \ln n)}{n}$ implies

$$(6.91) \quad \frac{(K_n - s + 1)^2}{P_n} = O\left(n^{-\frac{1}{s}} (\ln n)^{\frac{1}{s}}\right).$$

From (6.89) and the fact that s is a constant, it holds that $K_n - s + 1 \sim K_n$, which with (6.91) yields

$$(6.92) \quad \frac{K_n^2}{P_n} = O\left(n^{-\frac{1}{s}} (\ln n)^{\frac{1}{s}}\right).$$

Now we use (6.92) (6.89) and (6.87) to derive $q_n \sim \frac{1}{s!} \cdot \frac{K_n^{2s}}{P_n^s}$ and $|q_n - \frac{1}{s!} \cdot \frac{K_n^{2s}}{P_n^s}| = o\left(\frac{1}{n}\right)$, in a way similar to proving property (a) above.

First, (6.92) and (6.89) imply $K_n = \omega(1)$ and $\frac{K_n^2}{P_n} = o(1)$, respectively, which are used in (6.87) to derive $q_n \sim \frac{1}{s!} \cdot \frac{K_n^{2s}}{P_n^s}$. Second, applying (6.92) and (6.89) directly to (6.87), we still have cases (i) and (ii) in the proof of property (a) above. Hence, finally we also obtain $|q_n - \frac{1}{s!} \cdot \frac{K_n^{2s}}{P_n^s}| = o\left(\frac{1}{n}\right)$. Then property (b) is proved.

6.7 Proof of Lemma 6.3. Recall that E_{ij} denotes the event that an edge exists between vertices v_i and v_j , and S_i is the number of items on vertex v_i . Event E_{ij} occurs if and only if $|S_i \cap S_j| \geq s$. Therefore, event $\bigcap_{j=1}^r \overline{E_{ij}}$ is equivalent to $\bigcap_{j=1}^r (|S_i \cap S_j| < s)$, which clearly is implied by event

$|S_i \cap (\bigcup_{j=1}^r S_j)| < s$. Then

(6.93)

$$\begin{aligned} & \mathbb{P} \left[\bigcap_{j=1}^r \overline{E_{ij}} \mid S_1, S_2, \dots, S_r \right] \\ & \leq \mathbb{P} \left[\left| S_i \cap \left(\bigcup_{j=1}^r S_j \right) \right| < s \mid S_1, S_2, \dots, S_r \right] \\ & \leq 1 - \mathbb{P} \left[\left| S_i \cap \left(\bigcup_{j=1}^r S_j \right) \right| = s \mid S_1, S_2, \dots, S_r \right] \\ & = 1 - \frac{\binom{|\bigcup_{j=1}^r S_j|}{s} \binom{P_n - s}{K_n - s}}{\binom{P_n}{K_n}} \\ & = 1 - \frac{\binom{|\bigcup_{j=1}^r S_j|}{s} \binom{K_n}{s}}{\binom{P_n}{s}} \\ & \leq e^{-\frac{\binom{|\bigcup_{j=1}^r S_j|}{s} \binom{K_n}{s}}{\binom{P_n}{s}}}. \end{aligned}$$

First, we have (6.88) by [3, Lemma 6]. Applying (6.88) to (6.93), we obtain

(6.94)

$$\mathbb{P} \left[\bigcap_{j=1}^r \overline{E_{ij}} \mid S_1, S_2, \dots, S_r \right] \leq e^{-\frac{\binom{|\bigcup_{j=1}^r S_j|}{s} \binom{K_n}{s}}{\binom{P_n}{s}} q_n}.$$

Now we prove properties (a), (b), and (c) of Lemma 6.3, respectively.

(a) Given condition $K_n = \omega(1)$, it follows that $\lfloor (1 + \varepsilon_1)K_n \rfloor > s$ for all n sufficiently large. For property (a), we have condition $|\bigcup_{j=1}^r S_j| \geq \lfloor (1 + \varepsilon_1)K_n \rfloor$, which is used in (6.94) to derive

(6.95)

$$\mathbb{P} \left[\bigcap_{j=1}^r \overline{E_{ij}} \mid S_1, S_2, \dots, S_r \right] \leq e^{-\frac{\binom{\lfloor (1 + \varepsilon_1)K_n \rfloor}{s} \binom{K_n}{s}}{\binom{P_n}{s}} q_n}.$$

We have

$$\begin{aligned} (6.96) \quad \frac{\binom{\lfloor (1 + \varepsilon_1)K_n \rfloor}{s}}{\binom{K_n}{s}} &= \frac{\prod_{i=0}^{s-1} \{ \lfloor (1 + \varepsilon_1)K_n \rfloor - i \}}{\prod_{i=0}^{s-1} (K_n - i)} \\ &\geq \left[\frac{(1 + \varepsilon_1)K_n - 1 - s}{K_n} \right]^s. \end{aligned}$$

Given conditions $\varepsilon_2 < (1 + \varepsilon_1)^s - 1$ and $K_n = \omega(1)$, it follows that $K_n \geq \frac{s+1}{1 + \varepsilon_1 - \sqrt[s]{1 + \varepsilon_2}}$ for all n sufficiently large, yielding

$$\begin{aligned} (6.97) \quad \frac{(1 + \varepsilon_1)K_n - 1 - s}{K_n} &\geq (1 + \varepsilon_1) - (s + 1) \cdot \frac{1 + \varepsilon_1 - \sqrt[s]{1 + \varepsilon_2}}{s + 1} = \sqrt[s]{1 + \varepsilon_2}. \end{aligned}$$

Applying (6.97) to (6.96), we obtain

$$\frac{\binom{\lfloor (1 + \varepsilon_1)K_n \rfloor}{s}}{\binom{K_n}{s}} \geq (\sqrt[s]{1 + \varepsilon_2})^s = 1 + \varepsilon_2,$$

which is substituted into (6.95) to induce

$$\mathbb{P} \left[\bigcap_{j=1}^r \overline{E_{ij}} \mid S_1, S_2, \dots, S_r \right] \leq e^{-q_n(1 + \varepsilon_2)}.$$

(b) Given condition $K_n = \omega(1)$, it follows that $\lfloor \lambda_1 r K_n \rfloor > s$ for all n sufficiently large. For property (b), we have condition $|\bigcup_{j=1}^r S_j| \geq \lfloor \lambda_1 r K_n \rfloor$, which is used in (6.94) to derive

$$(6.98) \quad \mathbb{P} \left[\bigcap_{j=1}^r \overline{E_{ij}} \mid S_1, S_2, \dots, S_r \right] \leq e^{-\frac{\binom{\lfloor \lambda_1 r K_n \rfloor}{s} \binom{K_n}{s}}{\binom{P_n}{s}} q_n}.$$

We have

$$(6.99) \quad \frac{\binom{\lfloor \lambda_1 r K_n \rfloor}{s}}{\binom{K_n}{s}} = \frac{\prod_{i=0}^{s-1} (\lfloor \lambda_1 r K_n \rfloor - i)}{\prod_{i=0}^{s-1} (K_n - i)} \geq \left(\frac{\lambda_1 r K_n - 1 - s}{K_n} \right)^s.$$

Given conditions $\lambda_2 < \lambda_1^s$ and $K_n = \omega(1)$, it follows that $K_n \geq \frac{s+1}{\lambda_1 - \sqrt[s]{\lambda_2}} \geq \frac{s+1}{r(\lambda_1 - \sqrt[s]{\lambda_2})}$ for all n sufficiently large, inducing

(6.100)

$$\frac{\lambda_1 r K_n - 1 - s}{K_n} \geq \lambda_1 r - (s + 1) \frac{r(\lambda_1 - \sqrt[s]{\lambda_2})}{s + 1} = \sqrt[s]{\lambda_2} r.$$

Applying (6.100) to (6.99), we obtain $\binom{\lfloor \lambda_1 r K_n \rfloor}{s} / \binom{K_n}{s} \geq (\sqrt[s]{\lambda_2} r)^s = \lambda_2 r^s \geq \lambda_2 r$, which is substituted into (6.98) to induce

$$\mathbb{P} \left[\bigcap_{j=1}^r \overline{E_{ij}} \mid S_1, S_2, \dots, S_r \right] \leq e^{-\lambda_2 r q_n}.$$

(c) From $P_n \geq K_n = \omega(1)$, it follows that $P_n = \omega(1)$. Then $\lfloor \mu_1 P_n \rfloor > s$ for all n sufficiently large. For property (c), we have condition $|\bigcup_{j=1}^r S_j| \geq \lfloor \mu_1 P_n \rfloor$, which is used in (6.93) to derive

(6.101)

$$\mathbb{P} \left[\bigcap_{j=1}^r \overline{E_{ij}} \mid S_1, S_2, \dots, S_r \right] \leq e^{-\frac{\binom{\lfloor \mu_1 P_n \rfloor}{s} \binom{K_n}{s}}{\binom{P_n}{s}}}.$$

We have

$$\begin{aligned} (6.102) \quad \frac{\binom{\lfloor \mu_1 P_n \rfloor}{s} \binom{K_n}{s}}{\binom{P_n}{s}} &\geq \frac{(s!)^{-1} (\lfloor \mu_1 P_n \rfloor - s)^s \cdot (s!)^{-1} (K_n - s)^s}{(s!)^{-1} (P_n)^s} \\ &\geq \frac{1}{s!} \cdot \left(\frac{\mu_1 P_n - 1 - s}{P_n} \right)^s \cdot (K_n - s)^s. \end{aligned}$$

Given $0 < \mu_2 < (s!)^{-1}\mu_1^s$ and $P_n \geq K_n = \omega(1)$, it follows that $P_n \geq \frac{s+1}{\mu_1 - \sqrt{\mu_1}^{2s/s!}\mu_2}$ and $K_n \geq \frac{s+1}{1 - \sqrt{\mu_1}^{s!}\mu_2^s}$ for all n sufficiently large, inducing

$$(6.103) \quad \frac{\mu_1 P_n - 1 - s}{P_n} \geq \mu_1 - (s+1) \cdot \frac{\mu_1 - \sqrt{\mu_1}^{2s/s!}\mu_2}{s+1} = \sqrt{\mu_1}^{2s/s!}\mu_2,$$

and

$$(6.104) \quad (K_n - s)^s \geq K_n - s \geq K_n - K_n \left(1 - \sqrt{\frac{s!\mu_2}{\mu_1^s}}\right) = \sqrt{\frac{s!\mu_2}{\mu_1^s}} K_n.$$

Applying (6.103) and (6.104) to (6.102), we obtain

$$\frac{\binom{\lfloor \mu_1 P_n \rfloor}{s} \binom{K_n}{s}}{\binom{P_n}{s}} \geq \frac{1}{s!} (\sqrt{\mu_1}^{2s/s!}\mu_2)^s \sqrt{\frac{s!\mu_2}{\mu_1^s}} K_n = \mu_2 K_n,$$

which is substituted into (6.101) to induce

$$\mathbb{P} \left[\bigcap_{j=1}^r \overline{E_{ij}} \mid S_1, S_2, \dots, S_r \right] \leq e^{-\mu_2 K_n}.$$

6.8 Proof of Lemma 6.4. We consider events $\mathcal{B}_{\ell,r}$, $\mathcal{D}_{\ell,r}$ and $\mathcal{A}_{\ell,r}$ defined in Section 6.3. By definitions, we have

$$\mathcal{B}_{\ell,r} := \bigcap_{i=r+1}^{r+\ell} \bigcup_{j=1}^r E_{ij},$$

$$\mathcal{D}_{\ell,r} := \bigcap_{i=r+\ell+1}^n \bigcap_{j=1}^r \overline{E_{ij}},$$

and

$$\mathcal{A}_{\ell,r} := \mathcal{B}_{\ell,r} \cap \mathcal{C}_r \cap \mathcal{D}_{\ell,r}.$$

Then considering that given S_1, S_2, \dots, S_r , events $\mathcal{B}_{\ell,r}$ and $\mathcal{D}_{\ell,r} \cap \overline{\mathcal{E}(\mathcal{J})}$ are conditionally independent, we obtain

$$(6.105) \quad \mathbb{P} [\mathcal{A}_{\ell,r} \cap \overline{\mathcal{E}(\mathcal{J})}] = \mathbb{P} [\mathcal{C}_r \cap \mathcal{B}_{\ell,r} \cap \mathcal{D}_{\ell,r} \cap \overline{\mathcal{E}(\mathcal{J})}] = \sum_{\substack{S_1, S_2, \dots, S_r: \\ \mathcal{C}_r \text{ happens.}}} \left\{ \mathbb{P}[S_1, S_2, \dots, S_r] \mathbb{P}[\mathcal{B}_{\ell,r} \mid S_1, S_2, \dots, S_r] \mathbb{P}[\mathcal{D}_{\ell,r} \cap \overline{\mathcal{E}(\mathcal{J})} \mid S_1, S_2, \dots, S_r] \right\}.$$

We have

$$\mathbb{P}[\mathcal{B}_{\ell,r} \mid S_1, S_2, \dots, S_r] = \left\{ \mathbb{P} \left[\bigcup_{j=1}^r E_{ij} \mid S_1, S_2, \dots, S_r \right] \right\}^{\ell}.$$

By the union bound,

$$\mathbb{P} \left[\bigcup_{j=1}^r E_{ij} \mid S_1, S_2, \dots, S_r \right] \leq \sum_{j=1}^r \mathbb{P}[E_{ij} \mid S_1, S_2, \dots, S_r] = \sum_{j=1}^r \mathbb{P}[E_{ij}] = r q_n.$$

Then

$$(6.106) \quad \mathbb{P}[\mathcal{B}_{\ell,r} \mid S_1, S_2, \dots, S_r] \leq \min\{(r q_n)^{\ell}, 1\}.$$

We have

$$\mathbb{P}[\mathcal{D}_{\ell,r} \cap \overline{\mathcal{E}(\mathcal{J})} \mid S_1, S_2, \dots, S_r] = \left\{ \mathbb{P} \left[\left(\bigcap_{j=1}^r \overline{E_{ij}} \right) \cap \overline{\mathcal{E}(\mathcal{J})} \mid S_1, S_2, \dots, S_r \right] \right\}^{n-\ell-r}.$$

By Lemma 6.3, for all n sufficiently large,

(a) for $r = 2, 3, \dots, R$, it holds that

$$\mathbb{P}[\mathcal{D}_{\ell,r} \cap \overline{\mathcal{E}(\mathcal{J})} \mid S_1, S_2, \dots, S_r] \leq e^{-q_n(1+\varepsilon_2)(n-\ell-r)} \leq e^{-q_n n(1+\varepsilon_3)}.$$

To see this, pick any $\varepsilon_3 < (1 + \varepsilon_1)^s - 1$, and use Lemma 6.3 with $\varepsilon_3 < \varepsilon_2 < (1 + \varepsilon_1)^s - 1$.

(b) for $r = 2, 3, \dots, r_n$, it holds that

$$\mathbb{P}[\mathcal{D}_{\ell,r} \cap \overline{\mathcal{E}(\mathcal{J})} \mid S_1, S_2, \dots, S_r] \leq e^{-\lambda_2 r q_n (n-\ell-r)} \leq e^{-\lambda_2 r q_n n/3}.$$

(c) for $r = r_n + 1, r_n + 2, \dots, \lfloor \frac{n-\ell}{2} \rfloor$, it holds that

$$\mathbb{P}[\mathcal{D}_{\ell,r} \cap \overline{\mathcal{E}(\mathcal{J})} \mid S_1, S_2, \dots, S_r] \leq e^{-\mu_2 K_n (n-\ell-r)} \leq e^{-\mu_2 K_n n/3}.$$

For simplicity, we use Λ to summarize the upper bounds on $\mathbb{P}[\mathcal{D}_{\ell,r} \cap \overline{\mathcal{E}(\mathcal{J})} \mid S_1, S_2, \dots, S_r]$ in cases (a) (b) and (c) above; i.e., $\Lambda = e^{-q_n n(1+\varepsilon_3)}$ for $r = 2, 3, \dots, R$, and $e^{-\lambda_2 r q_n n/3}$ for $r = R+1, R+2, \dots, r_n$, and $e^{-\mu_2 K_n n/3}$ for $r = r_n + 1, r_n + 2, \dots, \lfloor \frac{n-\ell}{2} \rfloor$. In view of Λ , (6.106) and $\mathbb{P}[\mathcal{C}_r] \leq \min\{r^{r-2} q_n^{r-1}, 1\}$ by [5, Lemma 11], we obtain from (6.105) that

$$(6.107) \quad \mathbb{P} [\mathcal{A}_{\ell,r} \cap \overline{\mathcal{E}(\mathcal{J})}] \leq \sum_{\substack{S_1, S_2, \dots, S_r: \\ \mathcal{C}_r \text{ happens.}}} \left\{ \mathbb{P}[S_1, S_2, \dots, S_r] \cdot \min\{(r q_n)^{\ell}, 1\} \cdot \Lambda \right\} = \mathbb{P}[\mathcal{C}_r] \cdot \min\{(r q_n)^{\ell}, 1\} \cdot \Lambda \leq \min\{r^{r-2} q_n^{r-1}, 1\} \cdot \min\{(r q_n)^{\ell}, 1\} \cdot \Lambda,$$

which clearly completes the proof of Lemma 6.4.