Data bases Lab1

Part 1

6 different superkeys : {EmpID}, {SSN}, {Email}, {Phone}, {EmpID, Name}, {SSN, Department} .

Candidate keys: EmpID, SSN, Email, Phone

I would choose EmpID as primary key, because it's more convenient to use, and it's shorter.

I guess no , but some of them can have one number in case if they use unity work number.

Task 1.2

Student.AdvisorID → Professor.ProfID

Student.Major → Department.DeptCode

Professor.Department → Department.DeptCode

Course.DepartmentCode → Department.DeptCode

Department.ChairID → Professor.ProfID

Enrollment.StudentID → Student.StudentID

Enrollment.CourseID → Course.CourseID

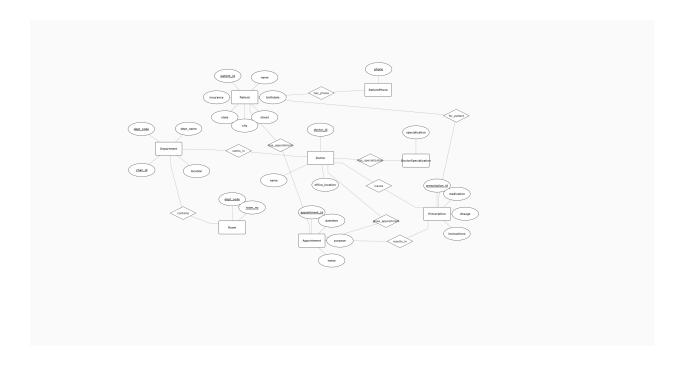
Part 2

2.1 Hospital Management System

Strong entities: Patient, Doctor, Department, Room, Appointment, Prescription

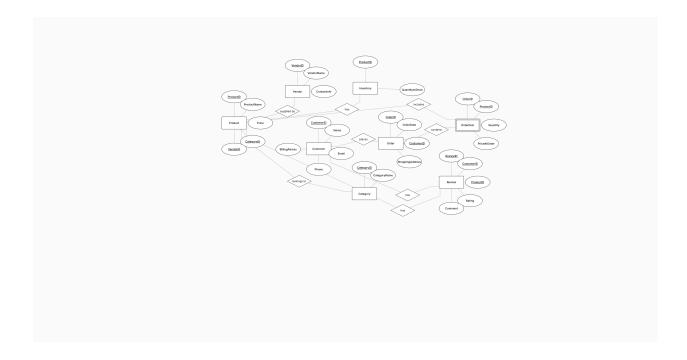
Weak entities: PatientPhone, DoctorSpecialization

Data bases Lab1



2.2 E-commerce Platform

1.



- 2. The weak entity is OrderItem.
- It cannot exist without both an Order and a Product.

- Its primary key is a composite key (OrderID + ProductID).
- It depends on the strong entities Order and Product for its identification.
- 3. The many-to-many relationship is between Order and Product, which is represented by the associative entity OrderItem.
- This relationship requires attributes such as Quantity and PriceAtOrder.
- Without these attributes, the system cannot track how many units of each product were ordered and at what price at the time of the order.

Part 4

4.1

1)

- StudentID → StudentName, StudentMajor (A student has one name and one major)
- 2. ProjectID \rightarrow ProjectTitle, ProjectType (a project has one title and one type)
- 3. SupervisorID → SupervisorName, SupervisorDept(a supervisor has one name and one department)
- 4. StudentID , ProjectID \rightarrow Role, HourdsWorked, startdate, enddate (A student's roles, hours, and dates are specific to a project)

2)

Redundancy: Student details (StudentName, StudentMajor) are repeated for every project a student is on. Project details are repeated for every student on that project.

Update Anomaly: Changing a student's major requires updating multiple rows.

Insert Anomaly: A new project or supervisor cannot be added unless a student is assigned to that project.

Delete Anomaly: Deleting the last row for a project removes all project information.

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3)

There are no 1NF violations. All attributes are atomic.

4)

The primary key is {StudentID, ProjectID}. The table is not in 2NF due to partial dependencies.

- StudentID → StudentName, StudentMajor (depends on a subset of the key)
- ProjectID → ProjectTitle, ProjectType, SupervisorID (depends on a subset of the key)

2NF Decomposition:

- Student(StudentID, StudentName, StudentMajor)
- Project(ProjectID, ProjectTitle, ProjectType, SupervisorID)
- StudentProject(StudentID, ProjectID, Role, HoursWorked, StartDate, EndDate)

5)

The Project table from the 2NF decomposition has a transitive dependency.

ProjectID \rightarrow SupervisorID and SupervisorID \rightarrow SupervisorName, SupervisorDept.

3NF Decomposition (Final Schemas):

Student(StudentID, StudentName, StudentMajor)

Project(ProjectID, ProjectTitle, ProjectType, SupervisorID)

StudentProject(StudentID, ProjectID, Role, HoursWorked, StartDate, EndDate)

Supervisor(SupervisorID, SupervisorName, SupervisorDept)

4.2

1)

The primary key is {StudentID, CourseID}.

2)

StudentID → StudentMajor

CourseID → CourseName

InstructorID → InstructorName

{TimeSlot, Room} → Building

{CourseID, TimeSlot, Room} → InstructorID

 $\{StudentID, CourseID\} \rightarrow all other attributes.$

3)

The table is not in BCNF because there are several FDs where the determinant is not a superkey.

StudentID → StudentMajor

CourseID → CourseName

InstructorID → InstructorName

{TimeSlot, Room} → Building

{CourseID, TimeSlot, Room} → InstructorID

4) BCNF Decomposition:

Student Information: Student(StudentID, StudentMajor)

Course Information: Course(CourseID, CourseName)

Instructor Information: Instructor(InstructorID, InstructorName)

Room Information: Room(Room, Building)

Course Section Information: CourseSection(CourseID, TimeSlot, Room,

InstructorID)

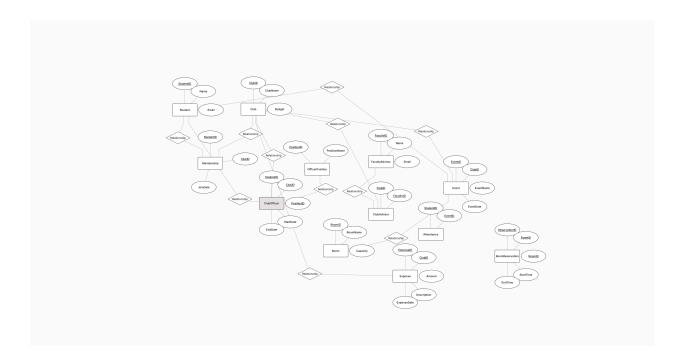
Student Enrollment: StudentEnrollment(StudentID, CourseID, TimeSlot, Room)

5)

The decomposition is lossless-join. No information is lost, but the original single table view is gone. To retrieve the complete course schedule for a student, you must perform joins between the decomposed tables. This is the goal of normalization, as it removes redundancy and prevents anomalies

Part 5

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2)

Student(StudentID PK, Name, Email)

Club(ClubID PK, ClubName, Budget)

Membership(StudentID PK, ClubID PK, JoinDate, FK → Student, FK → Club)

OfficerPosition(PositionID PK, PositionName)

ClubOfficer(StudentID PK, ClubID PK, PositionID PK, StartDate, EndDate, FK \rightarrow Student, FK \rightarrow Club, FK \rightarrow OfficerPosition)

FacultyAdvisor(FacultyID PK, Name, Email)

ClubAdvisor(ClubID PK, FacultyID PK, FK \rightarrow Club, FK \rightarrow FacultyAdvisor)

Event(EventID PK, ClubID FK, EventName, EventDate)

Attendance(StudentID PK, EventID PK, FK \rightarrow Student, FK \rightarrow Event)

Room(RoomID PK, RoomName, Capacity)

RoomReservation(ReservationID PK, EventID FK Unique, RoomID FK, StartTime, EndTime)

Expense(ExpenseID PK, ClubID FK, Amount, Description, ExpenseDate)

3)

One design decision was how to model club officers. I chose to create a separate entity OfficerPosition and a relationship table ClubOfficer, instead of storing officer role inside Membership. This allows students to hold different officer roles over time and keeps the schema normalized.

4)

- Find all students who are officers in the Computer Science Club.
- List all events scheduled for next week with their reserved rooms and times.
- Show the total expenses of each club in the past semester.