## ON THE CODEGREE DENSITY OF $PG_m(q)^*$

TAO ZHANG† AND GENNIAN GE‡

**Abstract.** For an r-graph G, the minimum (r-1)-degree  $\delta(G)$  is the largest integer t such that every (r-1)-subset of V(G) is contained in at least t edges of G. Given an r-graph F, the codegree density  $\gamma(F)$  is the largest  $\gamma>0$  such that there are F-free r-graphs G on n vertices with  $\delta(G)\geq (\gamma-o(1))n$ . In this paper, we consider the codegree density of projective geometries. Employing the moment identity of a subset of  $PG_m(q)$ , we prove (1)  $\gamma(PG_2(q))=\frac{1}{2}$  for prime power  $q\equiv 2\pmod 3$ ; and (2)  $\gamma(PG_3(q))=\frac{2}{3}$  for prime power  $q\equiv 1\pmod 2$  or  $q\equiv 2\pmod 3$ . Our results partially solve an open problem proposed by Keevash and Zhao  $[J.\ Combin.\ Theory\ Ser.\ B, 97\ (2007),\ pp.\ 919-928]$ . Previously, the codegree density problems for projective geometries were settled only for  $PG_2(2)$ ,  $PG_3(2)$ ,  $PG_3(3)$ , and  $PG_2(q)$  with odd prime power q.

Key words. codegree, hypergraph, projective geometry

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1. Introduction. In this paper, an r-graph is always an r-uniform hypergraph. Let H be an r-graph. An r-graph G is called H-free if G contains no copy of H as its subhypergraph. The  $Tur\'{a}n$  number ex(n,H) is the maximum number of edges in an n-vertex H-free r-graph. The determination of the Tur\'{a}n number ex(n,H) is a fundamental problem in extremal combinatorics. An easy averaging argument shows that the nonnegative sequence  $ex(n,H)/\binom{n}{r}$  is nonincreasing and hence converges to a limit as n goes to infinity. This limit is usually called  $Tur\'{a}n$  density of H and is denoted by  $\pi(H)$ .

For simple graphs (r=2), Turán [16] determined the exact value of  $\operatorname{ex}(n,K_t)$  for all complete graphs  $K_t$ . Later, Erdős and Stone [2] showed that  $\pi(H)=1-\frac{1}{\chi(H)-1}$ , where  $\chi(H)$  denotes the chromatic number of H. In contrast to the simple graphs, there are only a few results for the hypergraph Turán problems. For example, even the value of  $\pi(K_t^{(r)})$  is still unknown for any t>r>2, where  $K_t^{(r)}$  denotes the complete r-graph on t vertices. It is conjectured that  $\pi(K_4^{(3)})=\frac{5}{9}$ . Recently, there has been some new progress for the hypergraph Turán problems (for example, see [5, 6, 7, 10, 14, 17]). For more extremal results of hypergraphs, we refer the reader to the survey [9].

A natural variation of Turán problem is to consider how large the minimum degree can be in an H-free r-graph. For any r-graph G, let  $\delta(G)$  be the minimum d(S) of a set S of r-1 vertices, where d(S) is the number of edges of G containing S. The codegree extremal number co-ex(n, H) is the maximum possible value of  $\delta(G)$ , where

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G is an H-free r-graph on n vertices. In [15], Mubayi and Zhao showed that

$$\gamma(H) = \lim_{n \to \infty} \frac{\operatorname{co-ex}(n, H)}{n}$$

exists, which is called the *codegree density* of H.

It seems that for most hypergraphs, the codegree Turán problems are not easier than the original Turán problems. There are only a few r-graphs for which the codegree densities have been determined. In [13], Mubayi showed that  $\gamma(PG_2(2)) = \frac{1}{2}$ . Later, DeBiasio and Jiang [1] gave an alternative proof for codegree threshold for the Fano plane. Keevash and Zhao [11] studied the codegree density for other projective geometries and constructed a family of 3-graphs whose codegree densities are  $1 - \frac{1}{t}$  for all integers t > 1. In [3], Falgas-Ravry et al. determined  $\gamma(F_{3,2})$ , where  $F_{3,2}$  is the 3-graph on  $\{1,2,3,4,5\}$  with edges 123, 124, 125, 345. Falgas-Ravry et al. [4] proved that  $\gamma(K_4^{3-}) = \frac{1}{4}$ , where  $K_4^{3-}$  is the 3-graphs on four vertices with three edges. In [12], Lo and Zhao considered the codegree density of complete r-graphs.

In this paper, we focus on the codegree problem for projective geometries. The following result can be found in [11].

THEOREM 1.1 (see [11]). The codegree density of projective geometries satisfies  $\gamma(PG_m(q)) \leq 1 - \frac{1}{m}$ . Equality holds whenever m=2 and q is 2 or odd, and whenever m=3 and q is 2 or 3.

The behavior of  $\gamma(PG_2(4))$  seems different from others, and Keevash and Zhao [11] proved that

$$\gamma(PG_2(4)) \ge \frac{1}{3}.$$

Later, by the hypergraph regularity method, Keevash [8] showed the following upper bound:

$$\gamma(PG_2(4)) < \frac{1}{2} - \epsilon$$

for some  $\epsilon > 0$ .

Keevash and Zhao [11] also proposed the following open problem.

PROBLEM 1.2. Is  $\gamma(PG_m(q)) > 0$  for all m and q?

In this paper, we will continue this investigation and determine more values of  $\gamma(PG_m(q))$ . For m=2, we prove the following result.

THEOREM 1.3. Let q be a prime power with  $q \equiv 2 \pmod{3}$ ; then we have

$$\gamma(PG_2(q)) = \frac{1}{2}.$$

In particular, we have  $\gamma(PG_2(2^{2k+1})) = \frac{1}{2}$  for any integer  $k \geq 0$ . For m = 3, we have the following result.

Theorem 1.4. Let q be a prime power with  $q \equiv 1 \pmod 2$  or  $q \equiv 2 \pmod 3$ ; then we have

$$\gamma(PG_3(q)) = \frac{2}{3}.$$

The rest of this paper is organized as follows. In section 2, we give some basics of projective geometries. In section 3, we prove Theorem 1.3. Section 4 contains the proof of Theorem 1.4. Section 5 concludes this paper.

2. Preliminaries. Let  $\mathbb{F}_q$  denote the finite field with q elements. The projective geometry of dimension m over  $\mathbb{F}_q$ , denoted by  $PG_m(q)$ , is the following (q+1)-graph. Its vertex set is the set of all one-dimensional subspaces of  $\mathbb{F}_q^{m+1}$ . Its edges are the twodimensional subspaces of  $\mathbb{F}_q^{m+1}$ , in which for each two-dimensional subspace, the set of one-dimensional subspaces that it contains is an edge of the hypergraph  $PG_m(q)$ .

Let K be a subset of  $PG_m(q)$ ; a line  $\ell$  of  $PG_m(q)$  is an i-secant of K if  $|\ell \cap K| = i$ . Let  $\tau_i$  denote the total number of *i*-secants to K; then we have the following lemma.

LEMMA 2.1. Let K be a subset of  $PG_m(q)$  and let |K| = k; then the following equations hold:

$$\sum_{i=0}^{q+1} \tau_i = \frac{(q^m - 1)(q^{m+1} - 1)}{(q - 1)(q^2 - 1)},$$

$$\sum_{i=1}^{q+1} i\tau_i = k \sum_{i=0}^{m-1} q^i,$$

$$\sum_{i=2}^{q+1} \binom{i}{2} \tau_i = \binom{k}{2}.$$

*Proof.* The equations express in different ways the cardinality of the following sets:

- 1.  $\{\ell : \ell \text{ is a line of } PG_m(q)\};$
- 2.  $\{(P,\ell): P \in (K \cap \ell), \ell \text{ is a line of } PG_m(q)\};$
- 3.  $\{(\{P,Q\},\ell): P,Q \in (K \cap \ell), \ell \text{ is a line of } PG_m(q)\}.$
- **3.** Dimension 2: Proof of Theorem 1.3. Let  $q \equiv 2 \pmod{3}$ . By Theorem 1.1, it suffices to show that  $\gamma(PG_2(q)) \geq 1/2$ . Let V be a set of n vertices. Partition it as  $V = V_1 \cup V_2$  so that  $||V_i| - \frac{n}{2}| < 1$  for i = 1, 2. Let G be the (q+1)graph such that if e is an edge of G, then  $(|e \cap V_1|, |e \cap V_2|) \in \{(i, q+1-i) : i \not\equiv 0\}$  $\pmod{3}, 1 \leq i \leq q$ . Then it is easy to see that  $\delta(G) > (\frac{1}{2} - o(1))n$ .

Suppose G contains a copy of  $PG_2(q)$ . Let  $K = P\overline{G}_2(q) \cap V_1$  and let k = |K|. For  $1 \le i \le q$ , let  $\tau_i$  denote the total number of *i*-secants to K. Then by Lemma 2.1, we have

(1) 
$$\sum_{i \neq 0} \tau_i = q^2 + q + 1,$$

(1) 
$$\sum_{i \not\equiv 0 \pmod{3}} \tau_i = q^2 + q + 1,$$
(2) 
$$\sum_{i \not\equiv 0 \pmod{3}} i\tau_i = k(q+1),$$

(3) 
$$\sum_{i \not\equiv 0 \pmod{3}} {i \choose 2} \tau_i = {k \choose 2}.$$

Note that

$$\begin{pmatrix} s \\ 2 \end{pmatrix} \pmod{3} \equiv \begin{cases} 0 & \text{if } s \equiv 0 \text{ or } 1 \pmod{3}, \\ 1 & \text{if } s \equiv 2 \pmod{3}. \end{cases}$$

Taking both sides of (1), (2), and (3) modulo 3, we obtain

(4) 
$$\sum_{i \equiv 1 \pmod{3}} \tau_i + \sum_{i \equiv 2 \pmod{3}} \tau_i \equiv 1 \pmod{3},$$

(5) 
$$\sum_{i \equiv 1 \pmod{3}} \tau_i - \sum_{i \equiv 2 \pmod{3}} \tau_i \equiv 0 \pmod{3},$$

(6) 
$$\sum_{i \equiv 2 \pmod{3}} \tau_i \equiv \binom{k}{2} \pmod{3}.$$

From (4) and (5), we obtain

$$2 \sum_{i \equiv 2 \pmod{3}} \tau_i \equiv 1 \pmod{3}.$$

Combining with (6), we have  $\binom{k}{2} \equiv 2 \pmod{3}$ , which is a contradiction.

- **4. Dimension 3: Proof of Theorem 1.4.** In this section, we will determine the values of  $\gamma(PG_3(q))$  for prime power q with  $q \equiv 2 \pmod 3$  or  $q \equiv 1 \pmod 2$ . By Theorem 1.1, it suffices to show that  $\gamma(PG_3(q)) \ge 2/3$ . We divide our discussion into three cases.
- **4.1.**  $q \equiv 2 \pmod{3}$ . Let V be a set of n vertices. Partition it as  $V = V_1 \cup V_2 \cup V_3$  so that  $||V_i| \frac{n}{3}| < 1$  for i = 1, 2, 3. We say that a set A of vertices has type (a, b, c) if  $(|A \cap V_1| \pmod{3}, |A \cap V_2| \pmod{3}, |A \cap V_3| \pmod{3}) = (a, b, c)$ . We define a (q+1)-graph G on V as follows. The edges are all the sets of q+1 vertices of V having type in S, where

$$S = \{(0,1,2), (0,2,1), (1,2,0), (1,0,2), (2,1,0), (2,0,1)\}.$$

First, we verify the codegree property, i.e., all possible types (a,b,c) for a set of q vertices can be obtained from a type of G by decreasing one number by 1. Noting that  $q \equiv 2 \pmod{3}$ , it is easy to see that (a,b,c) is a permutation of (1,1,0),(2,0,0), or (1,2,2). Since S is closed under permutation, we will only consider one case of permutations of (a,b,c). If (a,b,c)=(1,1,0), then it can be obtained from (2,1,0) or (1,2,0). If (a,b,c)=(2,0,0), then it can be obtained from (2,1,0) or (2,0,1). If (a,b,c)=(1,2,2), then it can be obtained from (1,0,2) or (1,2,0). Hence  $\delta(G)>(\frac{2}{3}-o(1))n$ .

Suppose G contains a copy of  $PG_3(q)$ . Let  $K_i = PG_3(q) \cap V_i$  and let  $k_i = |K_i|$ . For  $0 \le j \le q$ , let  $\tau_{ij}$  denote the total number of j-secants to  $K_i$ . Then by Lemma 2.1, we have

$$\sum_{j=2}^{q} {j \choose 2} \tau_{ij} = {k_i \choose 2}.$$

Let

$$L_i = \{e : e \text{ is a line in } PG_3(q), |e \cap K_i| \equiv 2 \pmod{3} \}.$$

Then it is easy to see that  $L_1, L_2, L_3$  form a partition of the lines of  $PG_3(q)$ . Hence, we have

$$|L_1| + |L_2| + |L_3| = (q^2 + 1)(q^2 + q + 1) \equiv 2 \pmod{3}.$$

On the other hand,  $|L_i| = \sum_{j \equiv 2 \pmod{3}} \tau_{ij} \equiv \sum_{j=2}^q {j \choose 2} \tau_{ij} \equiv {k_i \choose 2} \pmod{3}$ . Then we have

(7) 
$$2 \equiv |L_1| + |L_2| + |L_3| \equiv \binom{k_1}{2} + \binom{k_2}{2} + \binom{k_3}{2} \pmod{3}.$$

Noting that  $K_1, K_2, K_3$  partition the vertices of  $PG_3(q)$ , then we have

(8) 
$$k_1 + k_2 + k_3 = q^3 + q^2 + q + 1 \equiv 0 \pmod{3}.$$

Equation (8) implies  $k_1 \equiv k_2 \equiv k_3 \pmod{3}$  or  $\{k_1, k_2, k_3\} \pmod{3} = \{0, 1, 2\}$ , which contradicts (7).

**4.2.**  $q \equiv 3 \pmod{4}$ . Let V be a set of n vertices. Partition it as  $V = V_1 \cup V_2 \cup V_3$  so that  $||V_i| - \frac{n}{3}| < 1$  for i = 1, 2, 3. We say that a set A of vertices has type (a, b, c) if  $(|A \cap V_1| \pmod{4}, |A \cap V_2| \pmod{4}, |A \cap V_3| \pmod{4}) = (a, b, c)$ . We define a (q + 1)-graph G on V as follows. The edges are all the sets of q + 1 vertices of V having type in S, where

$$S = \{(3,1,0), (3,0,1), (1,3,0), (1,0,3), (0,1,3), (0,3,1), (2,1,1), (1,2,1), (1,1,2), (3,3,2), (3,2,3), (2,3,3)\}.$$

First, we verify the codegree property, i.e., all possible types (a,b,c) for a set of q vertices can be obtained from a type of G by decreasing one number by 1. Noting that  $q \equiv 3 \pmod{4}$ , it is easy to see that (a,b,c) is a permutation of (1,1,1), (2,1,0), (3,0,0), (3,2,2), or (3,3,1). Since S is closed under permutation, we will only consider one case of permutations of (a,b,c). If (a,b,c)=(1,1,1), it can be obtained from (2,1,1), (1,2,1), or (1,1,2). If (a,b,c)=(2,1,0), then it can be obtained from (3,1,0) or (3,1,0) or (2,1,1). If (a,b,c)=(3,0,0), then it can be obtained from (3,3,2) or (3,2,3). If (a,b,c)=(3,3,1), then it can be obtained from (3,0,1), (0,3,1), or (3,3,2). Hence  $\delta(G)>(\frac{2}{3}-o(1))n.$ 

Suppose G contains a copy of  $PG_3(q)$ . Let H be the set of vertices of any hyperplane of  $PG_3(q)$  and  $H_i = H \cap V_i$  for i = 1, 2, 3. We claim that each  $|H_i|$  is odd. Since  $H_1, H_2, H_3$  partition the vertices of H, then we have

(9) 
$$|H_1| + |H_2| + |H_3| = q^2 + q + 1 \equiv 1 \pmod{4}.$$

Let

$$L_i = \{e : e \text{ is a line in } H, |e \cap H_i| \pmod{4} \in \{2,3\}\},\$$
  
 $L_{233} = \{e : e \text{ is a line in } H, \text{ and has type } (2,3,3), (3,2,3), \text{ or } (3,3,2)\}.$ 

Then  $L_{233} \in L_i$  for i = 1, 2, 3 and the sets  $L_1 \setminus L_{233}$ ,  $L_2 \setminus L_{233}$ ,  $L_3 \setminus L_{233}$ , and  $L_{233}$  form a partition of the lines of H. Hence we have

$$q^{2} + q + 1 = |L_{1} \setminus L_{233}| + |L_{2} \setminus L_{233}| + |L_{3} \setminus L_{233}| + |L_{233}|$$
$$= |L_{1}| + |L_{2}| + |L_{3}| - 2|L_{233}|.$$

Therefore,

$$|L_1| + |L_2| + |L_3| \equiv 1 \pmod{2}$$
.

Let  $\ell_j(H_i)$  denote the number of lines with exactly j points in  $H_i$ . Then  $|L_i| = \sum_{j \equiv 2.3 \pmod{4}} \ell_j(H_i)$ . On the other hand, by Lemma 2.1,

Note that

$$\begin{pmatrix} s \\ 2 \end{pmatrix} \pmod{2} \equiv \begin{cases} 0 & \text{if } s \equiv 0 \text{ or } 1 \pmod{4}, \\ 1 & \text{if } s \equiv 2 \text{ or } 3 \pmod{4}. \end{cases}$$

Taking both sides of (10) modulo 2, we obtain

Hence, we obtain the following equations:

$$\begin{aligned} |H_1| + |H_2| + |H_3| &= q^2 + q + 1 \equiv 1 \pmod{4}, \\ \binom{|H_1|}{2} + \binom{|H_2|}{2} + \binom{|H_3|}{2} &\equiv |L_1| + |L_2| + |L_3| \equiv 1 \pmod{2}. \end{aligned}$$

Considering Table 1, where  $|H_1| + |H_2| + |H_3| \equiv 1 \pmod{4}$ , then we have

(11) 
$$|H_1| \equiv |H_2| \equiv |H_3| \equiv 1 \pmod{2}.$$

Table 1

$\{ H_1 ,  H_2 ,  H_3 \} \pmod{4}$	$\binom{ H_1 }{2} + \binom{ H_2 }{2} + \binom{ H_3 }{2} \equiv 1 \pmod{2}$
$\{0, 0, 1\}$	0
$\{0, 2, 3\}$	0
$\{1, 1, 3\}$	1
$\{1, 2, 2\}$	0

Let  $U_i = V(PG_3(q)) \cap V_i$  for i = 1, 2, 3, where  $V(PG_3(q))$  is the vertex set of graph  $PG_3(q)$ . Fixing a point  $x \in U_1$ , we count the size of the following set:

$$\{(y,H): y \in U_2, x,y \in H, H \text{ is a hyperplane of } PG_3(q)\}.$$

For each y, there are q + 1 hyperplanes containing x and y, so the number of pairs (y, H) is  $(q + 1)|U_2|$ , which is even. On the other hand, there are  $q^2 + q + 1$  planes containing x, and each contains an odd number of points from  $U_2$  by the above claim (see (11)), which implies that the number of pairs (y, H) is odd, a contradiction.

**4.3.**  $q \equiv 1 \pmod{4}$ . Let V be a set of n vertices. Partition it as  $V = V_1 \cup V_2 \cup V_3$  so that  $||V_i| - \frac{n}{3}| < 1$  for i = 1, 2, 3. We say that a set A of vertices has type (a, b, c) if  $(|A \cap V_1| \pmod{4}, |A \cap V_2| \pmod{4}, |A \cap V_3| \pmod{4}) = (a, b, c)$ . We define a (q+1)-graph G on V as follows. The edges are all the sets of q+1 vertices of V having type in S, where

$$S = \{(0,1,1), (1,0,1), (1,1,0), (0,3,3), (3,0,3), (3,3,0), (1,2,3), (1,3,2), (2,1,3), (2,3,1), (3,1,2), (3,2,1)\}.$$

First, we verify the codegree property, i.e., all possible types (a,b,c) for a set of q vertices can be obtained from a type of G by decreasing one number by 1. Noting that  $q \equiv 1 \pmod{4}$ , it is easy to see that (a,b,c) is a permutation of (1,0,0),(1,1,3),(1,2,2),(2,0,3), or (3,3,3). Since S is closed under permutation, we will only consider one case of permutations of (a,b,c). If (a,b,c)=(1,0,0), it can be obtained from (1,0,1) or (1,1,0). If (a,b,c)=(1,1,3), then it can be obtained from (1,1,0),(1,2,3), or (2,1,3). If (a,b,c)=(1,2,2), then it can be obtained from (1,2,3) or (1,3,2). If (a,b,c)=(2,0,3), then it can be obtained from (2,1,3) or (3,0,3). If (a,b,c)=(3,3,3), then it can be obtained from (0,3,3),(3,0,3), or (3,3,0). Hence  $\delta(G)>(\frac{2}{3}-o(1))n$ .

Suppose G contains a copy of  $PG_3(q)$ . Let H be the set of vertices of any hyperplane of  $PG_3(q)$  and  $H_i = H \cap V_i$  for i = 1, 2, 3. We claim that each  $|H_i|$  is odd. Since  $H_1, H_2, H_3$  partition the vertices of H, we then have

(12) 
$$|H_1| + |H_2| + |H_3| = q^2 + q + 1 \equiv 3 \pmod{4}.$$

Let

$$L_i = \{e : e \text{ is a line in } H, |e \cap H_i| \pmod{4} \in \{0, 1\}\},$$
  
 $L_{011} = \{e : e \text{ is a line in } H, \text{ and has type } (0, 1, 1), (1, 0, 1), \text{ or } (1, 1, 0)\}.$ 

Then  $L_{011} \in L_i$  for i = 1, 2, 3 and the sets  $L_1 \setminus L_{011}$ ,  $L_2 \setminus L_{011}$ ,  $L_3 \setminus L_{011}$ , and  $L_{011}$  form a partition of the lines of H. Therefore, we have

$$q^{2} + q + 1 = |L_{1} \setminus L_{011}| + |L_{2} \setminus L_{011}| + |L_{3} \setminus L_{011}| + |L_{011}|$$
$$= |L_{1}| + |L_{2}| + |L_{3}| - 2|L_{011}|.$$

Hence

$$|L_1| + |L_2| + |L_3| \equiv 1 \pmod{2}$$
.

Let  $\ell_j(H_i)$  denote the number of lines with exactly j points in  $H_i$ . Then  $|L_i| = \sum_{j\equiv 0,1\pmod{4}}\ell_j(H_i)$ . On the other hand, by Lemma 2.1, we have

(13) 
$$\sum_{j=0}^{q+1} \ell_j(H_i) = q^2 + q + 1,$$

(14) 
$$\sum_{j\geq 2} \binom{j}{2} \ell_j(H_i) = \binom{|H_i|}{2}.$$

Note that

$$\binom{s}{2} \pmod{2} \equiv \begin{cases} 0 & \text{if } s \equiv 0 \text{ or } 1 \pmod{4}, \\ 1 & \text{if } s \equiv 2 \text{ or } 3 \pmod{4}. \end{cases}$$

Taking both sides of (14) modulo 2, we obtain

$$\binom{|H_i|}{2} \equiv \sum_{j \equiv 2, 3 \pmod{4}} \ell_j(H_i) \pmod{2}.$$

Then we have

$$|L_i| = \sum_{j \equiv 0, 1 \pmod{4}} \ell_j(H_i)$$

$$= \sum_{j=0}^{q+1} \ell_j(H_i) - \sum_{j \equiv 2, 3 \pmod{4}} \ell_j(H_i)$$

$$\stackrel{\text{(13)}}{=} q^2 + q + 1 - \sum_{j \equiv 2, 3 \pmod{4}} \ell_j(H_i)$$

$$\equiv 1 - \binom{|H_i|}{2} \pmod{2}.$$

Hence, we obtain the following equations:

$$|H_1| + |H_2| + |H_3| = q^2 + q + 1 \equiv 3 \pmod{4},$$
  
 $\binom{|H_1|}{2} + \binom{|H_2|}{2} + \binom{|H_3|}{2} \equiv 3 - (|L_1| + |L_2| + |L_3|) \equiv 0 \pmod{2}.$ 

Considering Table 2, where  $|H_1| + |H_2| + |H_3| \equiv 3 \pmod{4}$ , we then have

(15) 
$$|H_1| \equiv |H_2| \equiv |H_3| \equiv 1 \pmod{2}.$$

Table 2

$\{ H_1 ,  H_2 ,  H_3 \} \pmod{4}$	$\binom{ H_1 }{2} + \binom{ H_2 }{2} + \binom{ H_3 }{2} \equiv 1 \pmod{2}$
$\{0,0,3\}$	1
$\{0, 1, 2\}$	1
$\{1, 1, 1\}$	0
$\{1, 3, 3\}$	0
$\{2, 2, 3\}$	1

Let  $U_i = V(PG_3(q)) \cap V_i$  for i = 1, 2, 3, where  $V(PG_3(q))$  is the vertex set of graph  $PG_3(q)$ . Fixing a point  $x \in U_1$ , we count the size of the following set:

$$\{(y,H): y \in U_2, x,y \in H, H \text{ is a hyperplane of } PG_3(q)\}.$$

For each y, there are q + 1 hyperplanes containing x and y, so the number of pairs (y, H) is  $(q + 1)|U_2|$ , which is even. On the other hand, there are  $q^2 + q + 1$  planes containing x, and each contains an odd number of points from  $U_2$  by the above claim (see (15)), which implies that the number of pairs (y, H) is odd, a contradiction.

**5. Conclusion.** If  $a \leq b$ , then  $PG_a(q) \subseteq PG_b(q)$ , hence  $\gamma(PG_a(q)) \leq \gamma(PG_b(q))$ . Table 3 summarizes the known results of the codegree densities of projective geometries. It seems that the determination of  $\gamma(PG_m(4^k))$  is quite hard. In [8, 11], the authors proved that  $\frac{1}{3} \leq \gamma(PG_2(4)) < \frac{1}{2}$ , and the lower bound is based on the classification of the blocking sets of  $PG_2(4)$ . In general, we do not know the classification of the blocking sets of  $PG_m(q)$ .

In this paper, we determine some new families of codegree densities of projective geometries. Our main technique is employing the moment identity of a subset of  $PG_m(q)$  (Lemma 2.1). For m=4, we can get similar equations, such as (7) and (8), but there are many solutions since there are more variables. It would be interesting to determine  $\gamma(PG_m(q))$  for  $m \geq 4$ .

Table 3
Known results of  $\gamma(PG_m(q))$ .

q $m$	2	4	Odd q	$2^{2k+1}, k \ge 1$	$2^{2k},k\geq 2$
2	$\frac{1}{2}$	$\left[\frac{1}{3},\frac{1}{2}\right)$	$\frac{1}{2}$	$\frac{1}{2}$	$[0, \frac{1}{2}]$
3	$\frac{2}{3}$	$[\frac{1}{2}, \frac{2}{3}]$	$\frac{2}{3}$	$\frac{2}{3}$	$[0,\frac{2}{3}]$
4	$[\frac{2}{3}, \frac{3}{4}]$	$[\frac{1}{2}, \frac{3}{4}]$	$[\frac{2}{3}, \frac{3}{4}]$	$[\frac{2}{3}, \frac{3}{4}]$	$[0, \frac{3}{4}]$
$m \ge 5$	$\left[\frac{3}{4}, 1 - \frac{1}{m}\right]$	$\left[\frac{1}{2}, 1 - \frac{1}{m}\right]$	$\left[\frac{2}{3}, 1 - \frac{1}{m}\right]$	$\left[\frac{2}{3}, 1 - \frac{1}{m}\right]$	$[0,1-\frac{1}{m}]$

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