计算机系统 程序的机器级表示:控制

湖南大学

《计算机系统》课程教学组

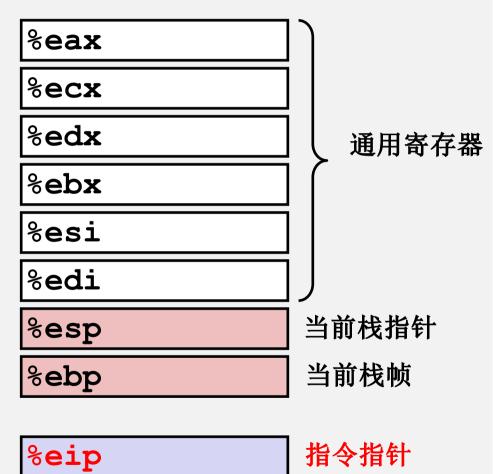




处理器状态

• 当前运行程序的相关信息

- 临时数据 (%eax, ...)
- 运行栈帧的地址 (%ebp,%esp)
- 即将要执行的指令地址 (%eip, ...)
- 标志位 (CF, ZF, SF, OF)



条件码

• 每个条件码占一个bit

- CF 最高位产生了进位,无符号操作数的溢出
- SF 符号标志,操作结果为负数
- ZF 零标志
- **OF** 溢出标志 (有符号数 signed)

```
例如: addl/addq Src, Dest \leftrightarrow t = a+b CF set , 如果t溢出
```

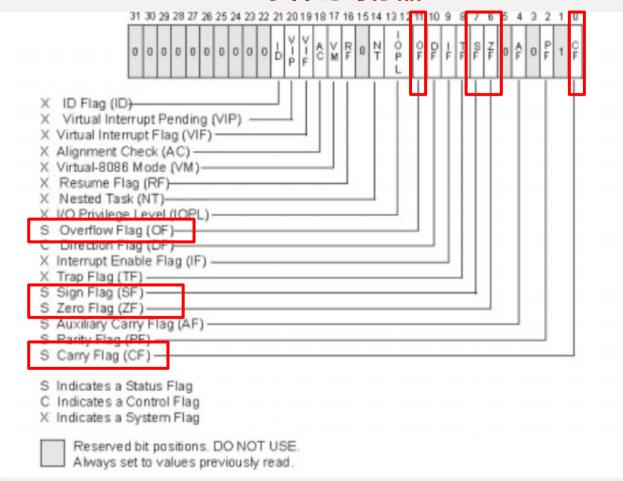
ZF set , 如果 t == 0

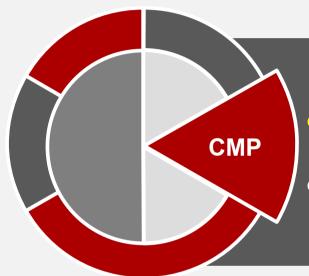
SF set , 如果 t < 0 (as signed)

OF set,如果有符号数溢出

lea/mov 指令不设置条件码







cmp1 Src, Dest; Dest - Src, 影响标志位

cmpl b,a 等价于计算 a-b,但不改变a与b的值。

CF set 无符号数运算时有进位

ZF set 如果 a == b

SF set 如果 有符号数 (a-b) < 0

OF set 如果有符号数运算溢出

以四个bit的数为例,有符号数能表示-8~7(1000~0111),无符号数能表示0~15(0000~1111)

$$6 > 3$$
 $6 - 3 = 3$.

$$SF = 0$$
, $OF = 0$, $SF \oplus OF = 0$

$$4 > -2$$
 4-(-2)=6,

$$SF=0$$
, $OF=0$

$$-2 < -4 -2 - (-4) = 2$$

$$6 < 7 \quad 6 - 7 = -1$$
.

$$SF=1$$
, $OF=0$, $SF\oplus OF=1$

$$-4 > -6 -4 - (-6) = 2$$

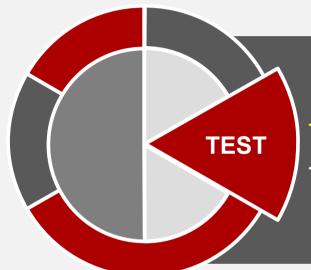
$$-4 < 7 - 4 - 7 = -11 = 10101 = 0101($$
IE $)$, SF=0, OF=1, SF \oplus OF=1

7 > -5 7-(-5)=12=01100=1100(
$$\textcircled{\texttt{f}}$$
), SF=1, OF=1, SF⊕OF=0

当出现:

正+正=负、负+负=正、正-负=负、负-正=正

的情况,就是产生了溢出,OF=1



test1/testq Src, Dest; Dest & Src,影响标志位

testl b,a 等价于计算a&b(但不改变a或b的值)

ZF set

SF set

如果 a&b == 0

如果 a&b < 0

SetX 指令:根据条件码的组合将一个字节设置为0或1

	SetX	Condition	Description	
	sete/setz	ZF	Equal / Zero	
setne/setnz		~ZF	Not Equal / Not Zero	
	sets	SF	Negative	
	setns	~SF	Nonnegative	
\lceil	setg	~(SF^OF) &~ZF	Greater (Signed)	
	setge	~(SF^OF)	Greater or Equal (Signed)	
	setl	(SF^OF)	Less (Signed)	
	setle	(SF^OF) ZF	Less or Equal (Signed)	
٢	seta	~CF&~ZF	Above (unsigned)	
	setb	CF	Below (unsigned)	

有符号数

无符_ 号数

设置条件码

```
int gt (int x, int y)
{
  return x > y;
}
```

Body

```
movl 12(%ebp),%eax # eax = y
cmpl %eax,8(%ebp) # Compare x : y
setg %al # al = x > y
movzbl %al,%eax # Zero rest of %eax
```

%eax %ah %al



跳转指令

jX 指令:根据不同的条件跳转到某条指令处执行

jХ	Condition	Description
jmp	1	无条件跳转
je/jz	ZF	Equal / Zero
Jne/jnz	~ZF	Not Equal / Not Zero
js	SF	Negative
jns	~SF	Nonnegative
jg	~(SF^OF) &~ZF	Greater (Signed)
jge	~(SF^OF)	Greater or Equal (Signed)
jl	(SF^OF)	Less (Signed)
jle	(SF^OF) ZF	Less or Equal (Signed)
ja	~CF&~ZF	Above (unsigned)
jb	CF	Below (unsigned)

有符号数 符数

```
int absdiff(int x, int y)
{
    int result;
    if (x > y) {
        result = x-y;
    } else {
        result = y-x;
    }
    return result;
}
```

条件跳转

```
absdiff:
     pushl %ebp
                             Setup
     movl %esp, %ebp
%edx=xmovl 8(%ebp), %edx
%eax=ymovl 12(%ebp), %eax
                             Body1
     cmpl %eax, %edx
     jle .L6
     subl %eax, %edx
                             Body2a
            %edx, %eax
     movl
     jmp
           .L7
     .L6:
                             Body2b
           %edx, %eax
     subl
     .L7:
           %ebp
     popl
                             Finish
     ret
```

```
int goto_ad(int x, int y)
{
   int result;
   if (x <= y) goto Else;
   result = x-y;
   goto Exit;
Else:
   result = y-x;
Exit:
   return result;
}</pre>
```

条件跳转

```
absdiff:
     pushl %ebp
                             Setup
     movl %esp, %ebp
%edx=xmovl 8(%ebp), %edx
%eax=ymovl 12(%ebp), %eax
                             Body1
     cmpl %eax, %edx
     jle .L6
     subl
            %eax, %edx
                             Body2a
            %edx, %eax
     movl
     jmp
            .L7
     .L6:
                             Body2b
     subl
            %edx, %eax
     .L7:
     popl
            %ebp
     ret
```

- C 中可以采用goto语句进行跳转,与机器级的语言风格类似
- 但通常被认为是一种比较 "low" 的编程 风格

分支跳转

C Code

```
val = Test ? Then_Expr: Else_Expr;
```

```
val = x>y ? x-y : y-x;
```

Goto Version

```
nt = !Test;
if (nt) goto Else;
val = Then_Expr;
goto Done;
Else:
  val = Else_Expr;
Done:
    . . .
```

Test 是一个返回整数值的表达式

- = 0 逻辑假
- ≠0逻辑真
- 为每一个分支都产生一段代码
- 根据条件执行合适的代码段

条件传送指令——满足条件才传送 Instruction supports:

if (Test) Dest← Src

- 先计算一个条件操作的两种结果,然后根据条件 选择某一个
- 优势:能够更好的匹配现代处理器的特性
 - 流水线
 - 分支预测

C Code

```
tval = Then_Expr;
result = Else_Expr;
t = Test;
if (t) result = tval;
return result;
```

C Code

```
int comvdiff(int x, int y)
{
  int tval = y-x;
  int rval = x-y;
  int test = x < y;
  if (test) rval = tval;
  result rval;
}</pre>
```

条件传送

```
comvdiff:
movl 8(%ebp), %ecx
movl 12(%ebp), %edx
movl %edx, %ebx
subl %ecx, %ebx
movl %ecx, %eax
subl %edx, %eax
cmpl %edx, %ecx
cmovl %ebx, %eax
```

避免了跳转指令

- CPU无需做分支预测, 避免预测错误的代价
- 流水线效率更高

条件跳转

```
absdiff:
pushl
      %ebp
movl
      %esp, %ebp
movl 8(%ebp), %edx
movl
      12(%ebp), %eax
cmpl %eax, %edx
ile .L6
subl %eax, %edx
movl %edx, %eax
jmp .L7
.L6:
subl %edx, %eax
.L7:
popl %ebp
ret
```

C代码

```
int absdiff(int a, int b)
{
    return a>b ? a-b : b-a;
}
```

非优化编译:gcc-Stest.c:

```
8(%ebp), %eax
      movl
      cmpl
             12(%ebp), %eax
      jle
              .L2
             12(%ebp), %eax
      movl
      movl 8(%ebp), %edx
      movl
             %edx, %ecx
      subl
             %eax. %ecx
      movl
             %ecx, %eax
      jmp
              .L3
L2:
              8(%ebp), %eax
      movl
```

汇编代码出现了jle这样的跳转指令。 对于使用了流水线的CPU,这样的跳转是存在隐患的(P140),分支预测 失败就会刷新掉所有流水线中取到而 未执行的指令,影响运行性能。

优化: gcc -S O1 test.c

```
(%esp), %ecx
movl
        12(%esp), %edx
movl
movl
       %ecx, %eax
       %edx, %eax
subl
       %edx, %ebx
movl
subl
       %ecx, %ebx
       %edx, %ecx
cmpl
cmovle
       %ebx, %eax
popl
       %ebx
```

经过优化后的代码没有了 跳转指令,取而代之的是 一个条件传送指令— cmovle。使得控制流不依 赖于数据,流水线也更容 易保持满状态。

计算代价

- 两个计算过程 都需要运行
- 一般来说,只 有两个计算过 程都比较简单 的时候,才能 够发挥优势

val = Test(x)? Hard1(x) : Hard2(x);

非法操作

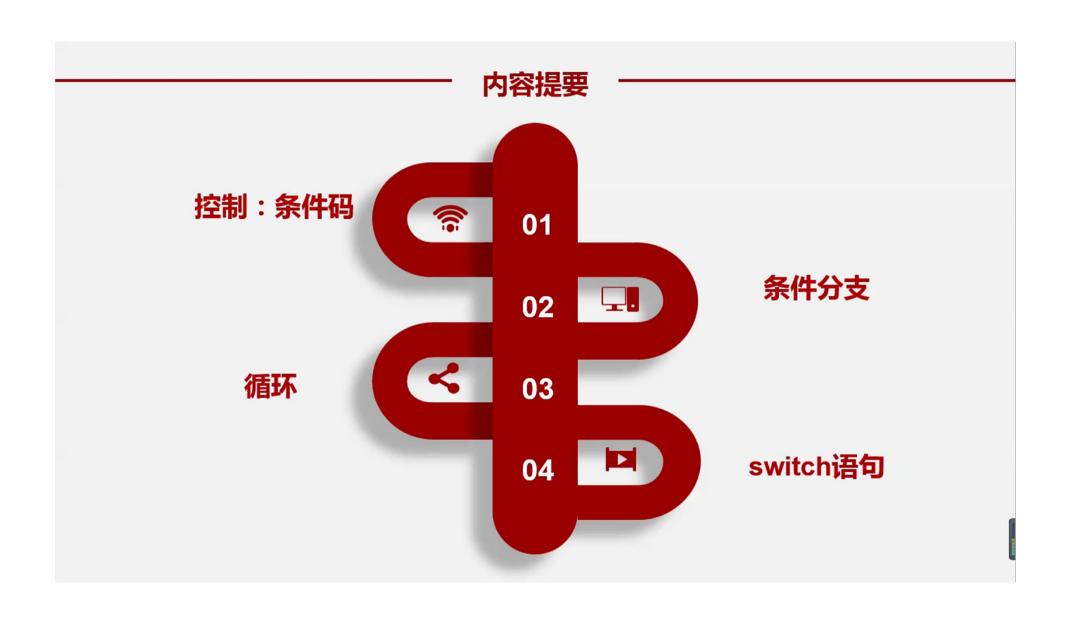
在p为0的时候,仍 然会去引用*p, 从而产生非法操作

val = p? *p : 0;

副作用

- 两个表达式都 进行了计算
- 产生了意料之 外的赋值过程

val = x > 0? x*=7 : x+=3;



C Code

```
int pcount_do(unsigned x) {
  int result = 0;
  do {
    result += x & 0x1;
    x >>= 1;
  } while (x);
  return result;
}
```

do-while 循环

- ◆ 计算x中有多少个1("popcount")
- 使用条件跳转指令来进行条件判断

```
寄存器
%edx x
%ecx result
```

while 循环

C Code for while loop

```
int pcount_while(unsigned x) {
  int result = 0;
  while (x) {
    result += x & 0x1;
    x >>= 1;
  }
  return result;
}
```

C code for do loop

```
int pcount_do(unsigned x)
{
  int result = 0;
  do {
    result += x & 0x1;
    x >>= 1;
  }
  while (x)
  return result;
}
```

While 和 do-while二者的代码是否完全一致?

- 都是条件测试失败退出循环
- do-while循环至少执行一次循环体

for 循环

C Code

```
#define WSIZE 8*sizeof(int)
int pcount_for(unsigned x) {
  int i;
  int result = 0;
  for (i = 0; i< WSIZE; i++) {
    unsigned mask = 1 <<i;
    result += (x & mask) != 0;
  }
  return result;
}</pre>
```

各循环对比

For Version

```
for (Init; Test; Update)

Body
```



While Version

```
Init;
while (Test) {

Body

Update;
}
```

Do-while Version

```
Init;

if (!Test)

goto done;

do

Body

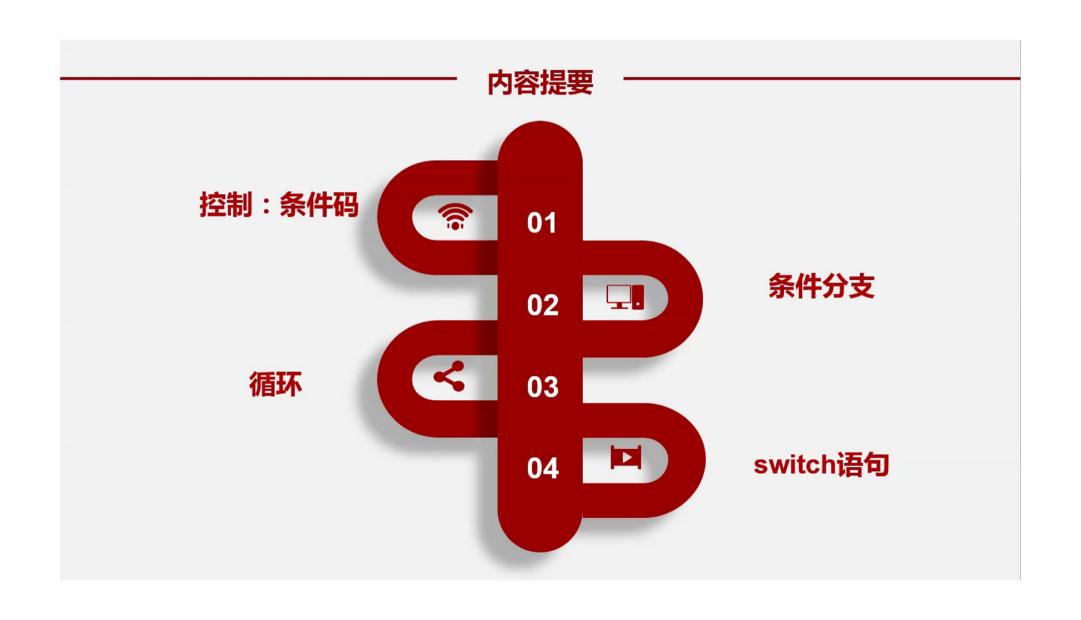
Update

while(Test);

done:
```

Goto Version

```
Init;
if (!Test)
goto done;
loop:
Body
Update
if (Test)
goto loop;
done:
```



switch语句

```
long switch eg
   (long x, long y, long z)
    long w = 1;
    switch(x) {
    case 1:
       w = y*z;
       break;
    case 2:
       w = y/z;
       /* Fall Through */
    case 3:
       w += z;
       break;
    case 5:
    case 6:
       w -= z;
       break;
    default:
       w = 2;
    return w;
```

- 多重分支
- 使用跳转表
- x = 2时无break

跳转表

Switch Form

```
switch(x) {
  case val_0:
    Block 0
  case val_1:
    Block 1
    • • •
  case val_n-1:
    Blockn-1
}
```

Jump Table

```
Targ_0
Targ_1
Targ_2

Targ_n-1
```

Jump Targets

Targ_0: Code Block 0

Targ_1: Code Block 1

Targ_2: Code Block 2

Approximate Translation

```
target = JTab[x];
goto *target;
```

Targ_*n*-1:

Code Block *n*–1

```
long switch eg(long x, long y,
long z)
    long w = 1;
    switch(x) {
    return w;
setup:
```

switch语句

```
Jump Table
jtab:
       Tarq0
       Tarq1
       Targ2
      Targn-1
```

Jump table

```
.section .rodata
 .align 4
.L4:
         .L8 \# x = 0
 .long
        .L3 \# x = 1
 .long
         .L5 \# x = 2
 .long
         .L9 \# x = 3
 .long
        .L8 \# x = 4
 .long
          .L7 \# x = 5
 .long
 .long
         .L7 \# x = 6
```

switch eg: 注意: 此处W没有进行初始化

```
pushl
         %ebp
                        # Setup
movl
        %esp,%ebp
                        # Setup
       8(\%ebp),\%eax # \%eax = x
movl
```

cmpl \$6,%eax # Compare x:6

jа .L8 # If unsigned '>' goto default

*.L4(,%eax,4) # Goto *JTab[x] jmp

跳转表结构

- ●每个跳转地址需要4字节
- ●基址在 .L4

跳转

- Direct : jmp .L2
- ●直接跳转,地址为.L2
- •Indirect : jmp *.L4(,%eax,4)
- ●跳转表基地址:.L4
- ●4是因为每个地址占4个字节
- ●取到有效地址.L4+ eax*4
 - $0 \le x \le 6$

Jump Table

```
.section .rodata
  .align 4
.L4:
  .long .L8  # x = 0
  .long .L3  # x = 1
  .long .L5  # x = 2
  .long .L9  # x = 3
  .long .L8  # x = 4
  .long .L7  # x = 5
  .long .L7  # x = 6
```

跳转表

Jump Table

```
.section .rodata
  .align 4
.L4:
  .long .L8  # x = 0
  .long .L3  # x = 1
  .long .L5  # x = 2
  .long .L9  # x = 3
  .long .L8  # x = 4
  .long .L7  # x = 5
  .long .L7  # x = 6
```

```
switch(x) {
case 1: // .L3
  w = y*z;
  break;
case 2: // .L5
  w = y/z;
  /* Fall Through */
case 3: // .L9
  w += z;
  break;
case 5:
case 6: // .L7
 w -= z;
  break;
default: // .L8
  w = 2;
```

代码块1

```
.L3: # x == 1
  movl 12(%ebp),%eax # y
  imull 16(%ebp),%eax # w = y*z
  jmp .L2 # Goto done
```

-Fall Through

```
long w = 1;
...
switch(x) {
...
case 2:
    w = y/z;
    y'* Fall Through */
case 3:
    w += z;
    break;
...
}
case 2:
    w = y/z;
    goto merge;

case 3:
    w = 1;
    merge:
    w += z;
```

代码块 2&3

代码块

5&6&default

```
switch(x) {
    . . .
    case 5: // .L7
    case 6: // .L7
    w -= z;
    break;
    default:// .L8
    w = 2;
}
```

·代码块:结束·

return w;

```
.L2: # done:
popl %ebp
ret
```

使用跳转表是一种非常有效的实现多重分支的方法

目标代码

准备阶段

- Label . L8 becomes address 0x80484b8
- Label . L4 becomes address 0x8048680

Assembly Code

```
switch_eg:
    . . .
    ja    .L8  # If unsigned >goto default
    jmp *.L4(,%eax,4) # Goto *JTab[x]
```

Disassembled Object Code

目标代码:跳转表:

跳转表

- ●在反汇编代码中无法直接看到,可以通过 GDB来观察
- •gdb switch
- •(gdb) x/7xw 0x8048680
 - Examine 7 hexadecimal format "words" (4-bytes each)—x/7xw

0x8048680: 0x080484b8 0x08048492 0x0804849b 0x080484a4

0x8048690: 0x080484b8 0x080484ae 0x080484ae

跳转表解释

0x8048680: 0x080484b8 0x08048492 0x8048690: 0x080484b8 0x080484ae

Address	Value	X
0x8048680	0x80484b8	0
0x8048684	0x8048492	1
0x8048688	0x804849b	2
0x804868c	0x80484a4	3
0x8048690	0x80484b8	4
0x8048694	0x80484ae	5
0x8048698	0x80484ae	6

0x0804849b 0x080484a4 0x080484ae

.section	.rodata
.align 4	
.L4:	
.long	.L8 # x = 0
.long	.L3 # x = 1
.long	.L5 # x = 2
.long	.L9 # x = 3
.long	.L8 # x = 4
.long	.L7 # x = 5
.long	.L7 # x = 6

反汇编

```
long switch eg(long x, long y, long z)
        8048492: 8b 45 0c
                 Of af 45 10
        8048495:
                                         long w = 1;
                                         switch(x) {
        8048499: eb 22
                                         case 1:
        804849b: 8b 45 0c
                                            w = y*z;
                                            break;
        804849e: 99
                                         case 2:
        804849f: f7 7d 10
                                            w = y/z;
                                           /* Fall Through */
        80484a2: eb 05
                                         case 3:
        80484a4: b8 01 00 00 00
                                            w += z;
                                            break;
        80484a9: 03 45 10
                                         case 5:
        80484ac: eb 0f
                                         case 6:
                                            w -= z;
        80484ae: b8 01 00 00 00
                                            break:
        80484b3: 2b 45 10
                                         default:
                                            w = 2;
        80484b6: eb 05
default
        80484b8: b8 02 00 00 00
                                         return w;}
```

反汇编

	8048492	8b 45	0c	mov	0xc(%ebp),%eax
	8048495	Of af	45 10	imul	0x10(%ebp),%eax
Value	8048499	eb 22		jmp	80484bd
	804849b	8b 45	0c	mov	0xc(%ebp),%eax
0x80484b8	804849e	99		cltd	
0x8048492	804849f	f7 7d	10	idiv]	l 0x10(%ebp)
0x804849b	80484a2	eb 05		jmp	80484a9
0x80484a4	80484a4	b8 01	00 00 0	0 mov	\$0x1,%eax
0110010101	80484a9	03 45	10	add	0x10(%ebp),%eax
0x80484b8	80484ac	eb 0f		jmp	80484bd
0x80484ae	80484ae	b8 01	00 00 0	0 mov	\$0x1,%eax
0x80484ae	80484b3	2b 45	10	sub	0x10(%ebp),%eax
	\\\80484b6	eb 05		jmp	80484bd
	80484b8	b8 02	00 00 0	0 mov	\$0x2,%eax

下一节:过程

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