### Ve 280

#### Programming and Elementary Data Structures

#### Linear List; Stack

#### **Learning Objectives:**

Understand what is a linear list and what is a stack

Know how they can be implemented

Discover some applications of the stack data structure

## Outline

- Linear List
- Stack
  - Implementation
  - Application

### Linear List ADT

- Recall the IntSet ADT
  - A collection of zero or more integers, with **no duplicates**.
  - It supports insertion and removal, but by value.
- A related ADT: linear list
  - A collection of zero or more integers; duplicates possible.
    - $L = (e_0, e_1, ..., e_{N-1})$
  - It supports insertion and removal by position.

#### Linear List ADT

#### Insertion

```
void insert(int i, int v) // if 0 <= i <= N</pre>
// (N is the size of the list), insert v at
// position i; otherwise, throws BoundsError
// exception.
Example: L1 = (1, 2, 3)
L1.insert(0, 5) = (5, 1, 2, 3);
L1.insert(1, 4) = (1, 4, 2, 3);
L1.insert(3, 6) = (1, 2, 3, 6);
L1.insert(4, 0) throws BoundsError
```

#### Linear List ADT

#### Removal

```
void remove(int i) // if 0 <= i < N (N is</pre>
 // the size of the list), remove the i-th
 // element; otherwise, throws BoundsError
 // exception.
Example: L2 = (1, 2, 3)
L2.remove(0) = (2, 3);
L2.remove(1) = (1, 3);
L2.remove(2) = (1, 2);
L2.remove(3) throws BoundsError
```



### Which Answers Are Correct?

Suppose we want to implement a linear list ADT so that it can grow as large as the user wants. It can be implemented by:

**\( \big|** 

- **A.** a static array
- **B.** a dynamic array
- C. a singly-linked list
- **D.** a doubly-linked list



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### Stack

- A "pile" of objects where new object is put on **top** of the pile and the top object is removed first.
  - LIFO access: last in, first out.
  - Restricted form of a **linear list**: insert and remove only at the end of the list.



### Methods of Stack

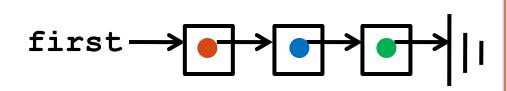
- size(): number of elements in the stack.
- isEmpty (): checks if stack has no elements.
- push (Object o): add object o to the top of stack.
- **pop()**: remove the top object if stack is not empty; otherwise, throw **stackEmpty**.
- Object &top(): return a reference to the top element.

# Stacks Using Arrays

```
Array[MAXSIZE]: 2314
```

- Maintain an integer **size** to record the size of the stack.
- size():return size;
- isEmpty():return (size == 0);
- push (Object o): add object o at index size of the array and increment size. Allocate more space if necessary.
- pop(): If isEmpty(), throw stackEmpty; otherwise, decrement size.
- Object &top(): return a reference to the top element Array[size-1]

## Stacks Using Linked Lists

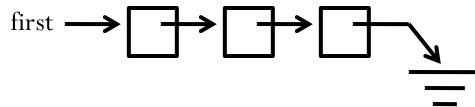


For single-ended linked list, which end is preferred to be the top? Why?

- size():LinkedList::size();
- isEmpty():LinkedList::isEmpty();
- push (Object o): insert object at the beginning LinkedList::insertFirst(Object o);
- pop(): remove the first node
  LinkedList::removeFirst();
- **Object &top()**: return a reference to the object stored in the first node.

## LinkedList::size()

• How to get the size of a linked list?



```
int LinkedList::size() {
  int count = 0;
  node *current = first;
  while(current) {
    count++;
    current = current->next;
  }
  return count;
}
```

A fast version: add a size data member

#### Linked List with a size data member

- Question: do we need to change any other parts of the code?
  - We need to increment/decrement **size** when nodes are inserted/removed.

### Array vs. Linked List: Which is Better?

- Linked list with a size data member is better
- Array
  - not memory-efficient: need to allocate a big enough array
- Linked list with a size data member
  - memory-efficient: a new item just needs extra constant amount of memory
  - All operations are of constant runtime, same as array

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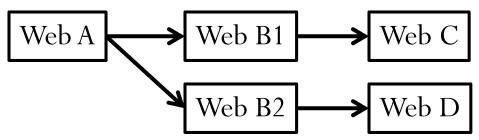
# Application of Stacks

• Function calls in C++

• Web browser's "back" feature

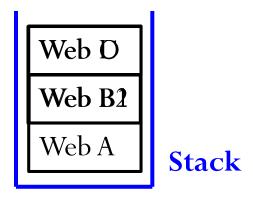
Parentheses Matching

### Web Browser's "back" Feature



#### Visiting order

- Web A
- Web B1
- Web C
- Back (to Web B1)
- Back (to Web A)
- Web B2
- Web D



## Parentheses Matching

• Output pairs (u,v) such that the left parenthesis at position u is matched with the right parenthesis at v.

```
( (a + b) * c + d - e) / (f + g)
0 1 2 3 4 5 6 7 8 9 10 12 14 16 18
• Output is: (1, 5); (0, 12); (14, 18);
```

```
(a+b)) * ((c+d)
0 1 2 3 4 5 6 7 8 9 10 12
```

Output is

(0,4);

Right parenthesis at 5 has no matching left parenthesis; (8, 12);

Left parenthesis at 7 has no matching right parenthesis

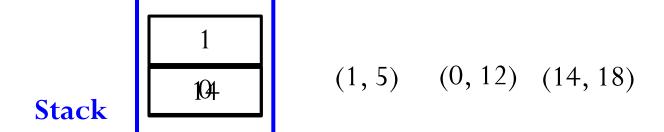
## How to Realize Parentheses Matching?

```
( ( a + b ) * c + d - e ) / (f + g )
0 1 2 3 4 5 6 7 8 9 10 12 14 16 18
```

- Scan expression from left to right.
- When a **left** parenthesis is encountered, push its position to the stack.
- When a **right** parenthesis is encountered, pop the top position from the stack, which is the position of the **matching left** parenthesis.
  - If the stack is empty, the **right** parenthesis is not matched.
- If string is scanned over but the stack is not empty, there are not-matched **left** parentheses.

# Parentheses Matching

```
((a+b) * c + d - e) / (f + g)
0 1 2 3 4 5 6 7 8 9 10 12 14 16 18
```





# A Stack Can Be Used:

Select all the correct answers.

• A. to manage any arithmetic expression

- **B.** to undo operations (such as a text editor)
- **C.** to reverse an array
- **D.** to implement an output buffer



### Reference

- **Problem Solving with C++ (8<sup>th</sup> Edition)**, by *Walter Savitch*, Addison Wesley Publishing (2011)
  - Chapter 13.2 Stack