

# TREE OF ORIENTATIONS ON A NESTED COLLECTION OF CUTS

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## 1. INTRODUCTION

Let  $\mathcal{C} \subseteq 2^X$  be a collection of non-empty subsets of a set  $X$ . With the definitions in Section 2, we prove the following

**Theorem 1.1.** *If  $\mathcal{C}$  is nested, then the graph  $\mathcal{T}_{\mathcal{C}}$ , whose:*

- *Vertices are finitely-based orientations on  $\mathcal{C}$ ; and whose*
- *Neighbors of  $\mathcal{U} \in V(\mathcal{T}_{\mathcal{C}})$  are  $\mathcal{U} \triangle \{A, A^c\}$  for every minimal  $A \in \mathcal{U}$  with  $A^c \in \mathcal{C}$ ;*

*is acyclic. Furthermore,  $\mathcal{T}_{\mathcal{C}}$  is a tree iff  $\mathcal{C}$  is closed under complements.*

In particular, this applies to when  $(X, G)$  is a graph and  $\mathcal{C}$  is a nested collection of cuts on  $X$ .

## 2. PRELIMINARIES

**2.1. Orientations.** Since we do not assume that  $\mathcal{C}$  is closed under complements, we slightly modify the definition of orientations, as follows.

**Definition 2.1.** An *orientation* on  $\mathcal{C}$  is a subset  $\mathcal{U} \subseteq \mathcal{C}$  such that

1. (*Upward-closure*). If  $A \in \mathcal{U}$  and  $B \in \mathcal{C}$  contains  $A$ , then  $B \in \mathcal{U}$ .
2. (*Ultra*). If  $A, A^c \in \mathcal{C}$ , then either  $A \in \mathcal{U}$  or  $A^c \in \mathcal{U}$ , but not both.

**Remark 2.2.** This coincides with the standard definition when  $\mathcal{C}$  is a subposet of  $2^X$ .

**Lemma 2.3.** *If  $\mathcal{U} \subseteq \mathcal{C}$  is an orientation, then for any  $\subseteq$ -minimal  $A \in \mathcal{U}$  with  $A^c \in \mathcal{C}$ , so is  $\mathcal{U} \triangle \{A, A^c\}$ .*

*Proof.* That  $\mathcal{V} := \mathcal{U} \triangle \{A, A^c\} = \mathcal{U} \cup \{A^c\} \setminus \{A\}$  is upward-closed follows from  $\subseteq$ -minimality of  $A$ . Now, if  $B, B^c \in \mathcal{C}$  and  $B^c \notin \mathcal{V}$ , we need to show that  $B \in \mathcal{V}$ .

To this end, note that  $B^c \notin \mathcal{V}$  implies  $A \neq B$  and either  $B = A^c$  or  $B^c \notin \mathcal{U}$ . The former case follows from  $A^c \in \mathcal{V}$ , and for the latter, we have  $B \in \mathcal{U} \setminus \{A\}$  since  $\mathcal{U}$  is ultra. ■

**Remark 2.4.** In the above notations, clearly  $\mathcal{U} \neq \mathcal{U} \triangle \{A, A^c\}$ . Furthermore, for any other such orientation  $\mathcal{U}'$  and  $A' \in \mathcal{U}'$ , that  $\mathcal{U} = \mathcal{U}'$  and  $\mathcal{U} \triangle \{A, A^c\} = \mathcal{U}' \triangle \{A', A'^c\}$  together imply  $A = A'$ .

**Definition 2.5.** A *base* for an orientation  $\mathcal{U} \subseteq \mathcal{C}$  is a  $\subseteq$ -minimal subset  $\mathcal{B} \subseteq \mathcal{U}$  such that  $\mathcal{U} = \uparrow \mathcal{B}$ , where

$$\uparrow \mathcal{B} := \bigcup_{B \in \mathcal{B}} \uparrow B := \bigcup_{B \in \mathcal{B}} \{A \in \mathcal{C} : A \supseteq B\}.$$

**Definition 2.6.** A collection  $\mathcal{C}$  is said to be *nested* if every  $C_1, C_2 \in \mathcal{C}$  has an empty corner, i.e.,  $C_1^i \cap C_2^j = \emptyset$  for some  $i, j \in \{1, -1\}$ , where  $C^i := C$  if  $i = 1$  and  $C^i := C^c$  if  $i = -1$ .

**Remark 2.7.** If  $\mathcal{C}$  is nested, then every  $\subseteq$ -minimal  $B \in \mathcal{C}$  induces an orientation  $\uparrow B := \{A \in \mathcal{C} : A \supseteq B\}$ , called the *principal* orientation. Indeed,  $\uparrow B$  is clearly upward-closed, and if  $A, A^c \in \mathcal{C}$ , then, by  $\subseteq$ -minimality of  $B$  and nestedness of  $\mathcal{C}$ , either  $A \supseteq B$  or  $A^c \supseteq B$  (but clearly not both).

This construction generalizes to any collection  $\mathcal{B} \subseteq \mathcal{C}$  with each  $B \in \mathcal{B}$  being  $\subseteq$ -minimal, in that  $\uparrow \mathcal{B}$  is an orientation on  $\mathcal{C}$ .

**Definition 2.8.** An orientation  $\mathcal{U} \subseteq \mathcal{C}$  is said to be *finitely-based* if it admits a finite base.

**Remark 2.9.** If  $\mathcal{U} = \uparrow \{B_1, \dots, B_n\}$  is finitely-based, then for any  $\subseteq$ -minimal  $A \in \mathcal{U}$  with  $A^c \in \mathcal{C}$ , so is the orientation  $\mathcal{V} := \mathcal{U} \triangle \{A, A^c\}$ . Indeed,  $A = B_i$  for some  $1 \leq i \leq n$ , and  $\mathcal{V} = \uparrow(\{A\} \cup \{B_j\}_{j \neq i})$ .

3. THE GRAPH  $\mathcal{T}_{\mathcal{C}}$ 

Fix a nested collection of non-empty subsets of a set  $X$ . Using Lemma 2.3 and Remarks 2.4 and 2.9, we construct a graph  $\mathcal{T}_{\mathcal{C}}$  whose:

- *Vertices* of  $\mathcal{T}_{\mathcal{C}}$  are finitely-based orientations on  $\mathcal{C}$ .
- *Neighbors* of  $\mathcal{U} \in V(\mathcal{T}_{\mathcal{C}})$  are  $\mathcal{U} \triangle \{A, A^c\}$  for every minimal  $A \in \mathcal{U}$  with  $A^c \in \mathcal{C}$ .

The goal of this section is to establish Theorem 1.1, stating that  $\mathcal{T}_{\mathcal{C}}$  is acyclic (Proposition 3.5), and furthermore,  $\mathcal{T}_{\mathcal{C}}$  is a tree precisely when  $\mathcal{C}$  is closed under complements (Proposition 3.6).

**3.1. Paths in  $\mathcal{T}_{\mathcal{C}}$ .** To show that  $\mathcal{T}_{\mathcal{C}}$  is acyclic, we characterize backtracking paths in  $\mathcal{T}_{\mathcal{C}}$  as follows.

**Definition 3.1.** Fix  $\mathcal{U}_0 \in V(\mathcal{T}_{\mathcal{C}})$  and  $\alpha \leq \omega$ . A sequence  $(A_n)_{n < \alpha} \subseteq \mathcal{C}$  is said to *represent a path from  $\mathcal{U}_0$*  if  $(\mathcal{U}_n)_{n \leq \alpha}$ , defined by  $\mathcal{U}_n := \mathcal{U}_{n-1} \triangle \{A_{n-1}\}$  for every  $1 \leq n \leq \alpha$ , is a path in  $\mathcal{T}_{\mathcal{C}}$  with each  $A_n \in \mathcal{U}_n$ .

**Remark 3.2.** Any path in  $\mathcal{T}_{\mathcal{C}}$  is represented by its sequence of flipped basis elements.

**Lemma 3.3.** *Let  $\alpha \geq 3$ . A path in  $\mathcal{T}_{\mathcal{C}}$  from  $\mathcal{U}_0$  represented by  $(A_n)_{n \leq \alpha}$  has no backtracking iff  $A_n \neq A_{n-1}^c$  for every  $1 \leq n < \alpha$ .*

*Proof.* Take  $2 \leq n \leq \alpha$ . It suffices to show that  $\mathcal{U}_{n-2} = \mathcal{U}_n$  iff  $A_{n-1} = A_{n-2}^c$ .

( $\Rightarrow$ ). We have by definition that  $\mathcal{U}_n = \mathcal{U}_{n-2} \cup \{A_{n-1}^c, A_{n-2}^c\} \setminus \{A_{n-1}, A_{n-2}\}$ , so since  $A_{n-2} \in \mathcal{U}_{n-2} = \mathcal{U}_n$ , we have  $A_{n-2} = A_{n-1}^c$  as desired.

( $\Leftarrow$ ). Again by definition, by noting that the basis-flipping cancels out.  $\blacksquare$

**Lemma 3.4.** *If  $(A_n)_{n < \alpha}$  represents a path in  $\mathcal{T}_{\mathcal{C}}$  with no backtracking, then  $(A_n)_{n < \alpha}$  is strictly increasing.*

*Proof.* By Lemma 3.3, we have  $A_n \neq A_{n-1}^c$  for every  $1 \leq n < \alpha$ . Thus, since  $A_n \in \mathcal{U}_n = \mathcal{U}_{n-1} \cup \{A_{n-1}^c\} \setminus \{A_{n-1}\}$ , we see that  $A_n \in \mathcal{U}_{n-1}$ . Clearly  $A_n \neq A_{n-1}$ . It suffices to remove the three cases when  $A_n \subseteq A_{n-1}$ ,  $A_{n-1} \subseteq A_n^c$ , and  $A_n^c \subseteq A_{n-1}$ , since then nestedness of  $\mathcal{C}$  gives us  $A_{n-1} \subset A_n$ , as desired.

- If  $A_n \subseteq A_{n-1}$ , then  $A_{n-1} \in \mathcal{U}_n$ , contradicting the definition of  $\mathcal{U}_n$ .
- If  $A_{n-1} \subseteq A_n^c$ , then  $A_n^c \in \mathcal{U}_{n-1}$  by upward-closure of  $\mathcal{U}_{n-1}$ , a contradiction.
- If  $A_n^c \subseteq A_{n-1}$ , then  $A_{n-1} \in \mathcal{U}_{n+1}$  by upward-closure of  $\mathcal{U}_{n+1} \ni A_n^c$ . But since  $A_{n-1} \neq A_n^c$ , we have by definition of  $\mathcal{U}_{n+1}$  that  $A_{n-1} \in \mathcal{U}_n$ , a contradiction.  $\blacksquare$

**Proposition 3.5.**  $\mathcal{T}_{\mathcal{C}}$  is acyclic.

*Proof.* Let  $(\mathcal{U}_i)_{i \leq n}$  be a cycle in  $\mathcal{T}_{\mathcal{C}}$ , say represented by  $(A_i)_{i \leq n}$ . Then  $A_0 = A_n$ , a contradiction since  $(A_i)_{i \leq n}$  is strictly-increasing by Lemma 3.4.  $\blacksquare$

**Proposition 3.6.** *If  $\mathcal{C}$  is closed under complements, then  $\mathcal{T}_{\mathcal{C}}$  is connected (and hence a tree).*

*Proof.* If  $\mathcal{U}, \mathcal{U}' \in V(\mathcal{T}_{\mathcal{C}})$  are two finitely-based orientations on  $\mathcal{C}$ , then swapping their basis elements one by one as in Remark 2.9 gives us a path between  $\mathcal{U}$  and  $\mathcal{U}'$ ; the closure of  $\mathcal{C}$  is needed to ensure that those basis elements represent a path between those vertices.  $\blacksquare$