

TREE OF ORIENTATIONS ON A NESTED COLLECTION OF SETS

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Let $H \subseteq 2^X$ be a sub-pocset for some fixed set X (so that, in particular, H is closed under complements). With the definitions in Section 1, we prove the following

Theorem A (Propositions 2.5, 2.6, 2.7). *If H is nested, then the graph \mathcal{T}_H , whose:*

- *Vertices are finitely-based orientations $p \subseteq H$;*
- *Edges are pairs $\{p, q\}$ such that $q = p \triangle \{h, \neg h\}$ for some \subseteq -minimal $h \in H$;*

is acyclic. Furthermore, if H is finitely-separating (or more generally, chain-vanishing), then \mathcal{T}_H is a tree.

In particular, this applies to when (X, G) is a graph and H is a nested collection of cuts on X . No further assumptions on X (like local-finiteness) is needed.

1. PRELIMINARIES

Let $H \subseteq 2^X$ be a sub-pocset for some fixed set X , whose elements $h \in H$ we call *half-spaces*.

Definition 1.1. Two elements $h, k \in H$ are *nested* if $h^i \cap k^j = \emptyset$ for some $i, j \in \{1, -1\}$, where $h^i := h$ for $i = 1$ and $h^i := h^c$ otherwise. We say that H is *nested* if every pair $h, k \in H$ are nested.

1.1. Orientations. We give the standard definition of orientations on H and characterize them as ‘consistent assignments of half-spaces to hyperplanes’.

Definition 1.2. An *orientation* on H is a subset $U \subseteq H$ such that

1. (*Upward-closure*). If $h \in U$ and $k \in H$ contains h , then $k \in U$.
2. (*Ultra*). For each $h \in H$, exactly one of h, h^c is contained in U .

Consider the equivalence relation \sim on H generated by $h \sim h^c$ for all $h \in H$, whose classes are called *hyperplanes* $\delta h := \{h, h^c\}$ where $\delta : H \rightarrow H/\sim$ is the projection. We show that an orientation on H is just a choice $\varphi : H/\sim \rightarrow H$ of a half-space for each hyperplane, that is consistent in the sense below.

Proposition 1.3. *An orientation $p \subseteq H$ is exactly the data of a function $\varphi : H/\sim \rightarrow H$ such that $\varphi(\delta h) \in \delta h$ and $\varphi(\delta h) \not\subseteq \varphi(\delta k)^c$ for every $h, k \in H$.*

Proof. Given an orientation $p \subseteq H$, let $\varphi_p(\delta h) := h^i \in U$ for the unique $i \in \{1, -1\}$. That $\varphi_p(\delta h) \in \delta h$ is clear, and if $\varphi_p(\delta h) \subseteq \varphi_p(\delta k)^c$, then U contains both $\varphi_p(\delta k)$ and $\varphi_p(\delta k)^c$ by upward-closure, a contradiction.

Conversely, given such a function $\varphi : H/\sim \rightarrow H$, let $p_\varphi := \text{im } \varphi \subseteq H$. This is ultra since if $h \in H$ and $h^c \notin p_\varphi$, then $\varphi(\delta h) \in \delta h = \{h, h^c\}$ implies $h \in p_\varphi$. Furthermore, if $p_\varphi \ni h \subseteq k$, then $k^c \in p_\varphi$ implies $\varphi(\delta h) = h \subseteq k = \varphi(\delta k)^c$, a contradiction, so $k \in p_\varphi$ by the above.

Finally, given an orientation $p \subseteq H$, we have $h \in p$ iff $\varphi_p(\delta h) = h$, which occurs iff $h \in \text{im } \varphi_p = p_{\varphi_p}$. Thus $p_{\varphi_p} = p$. Conversely, given $\varphi : H/\sim \rightarrow H$, and $h \in H$, we have $\varphi_{p_\varphi}(\delta h) = h^i$ iff $h^i \in p_\varphi = \text{im } \varphi$, which occurs iff $\varphi(\delta h) = h^i$. Thus $\varphi_{p_\varphi} = \varphi$ too, as desired. ■

Definition 1.4. A *base* for an orientation $p \subseteq H$ is a \subseteq -minimal subset $p_0 \subseteq p$ such that $p = \uparrow p_0$, where

$$\uparrow p_0 := \bigcup_{h \in p_0} \uparrow h := \bigcup_{h \in p_0} \{k \in H : k \supseteq h\}.$$

We say that p is *finitely-based* if it admits a finite basis.

In the above correspondence, a base for $\varphi : H/\sim \rightarrow H$ is a function $\varphi_0 \subseteq \varphi$ where $\text{dom } \varphi_0$ is a \subseteq -minimal subset of hyperplanes such that φ_0 extends uniquely to φ . Thus, the finitely-based orientations are the ones determined by a choice of half-spaces from finitely-many hyperplanes.

1.2. Flipping basis elements. We now investigate the behaviour of orientation when its choice on a single half-space is modified. Although the proofs work without the characterization in Proposition 1.3, it makes orientations a lot more intuitive to me, and so will be phrased this way.

Lemma 1.5. *Let $p \subseteq H$ be an orientation and fix $h \in p$. The following are equivalent.*

- (1) $q := p \triangle \{h, h^c\} = p \cup \{h^c\} \setminus \{h\}$ is an orientation on H .
- (2) h is \subseteq -minimal in p .
- (3) h is in every basis $p_0 \subseteq p$.

If any of the above holds and furthermore p is finitely-based, then so is q .

Proof. First, note that $\varphi_q(\partial h) = \varphi_p(h)^c = h^c$ and $\varphi_q = \varphi_p$ away from ∂h .

- (1 \Leftrightarrow 2). If q is an orientation and $p \ni k \subset h$, then we have a contradiction since

$$\varphi_q(\partial k) = \varphi_p(\partial k) = k \subseteq h = \varphi_q(\partial h)^c.$$

Conversely, let $h \in p$ be \subseteq -minimal and suppose that q is not an orientation. Since only $\varphi_q(\partial h) = h^c$ differs from φ_p , this can only occur if $\varphi_q(\partial h) \subseteq \varphi_q(\partial k)^c$ for some $k \in H$. But then

$$h^c = \varphi_q(\partial h) \subseteq \varphi_q(\partial k)^c = k^i \notin p$$

for some unique $i \in \{1, -1\}$, so that $p \ni k^{-i} \subseteq h$ and contradicts that $h \in p$ is \subseteq -minimal.

- (2 \Leftrightarrow 3). We prove the contrapositives. If there is a basis $p_0 \subseteq p$ with $h \notin p_0$, then $p = \uparrow p_0$ gives us some $h_0 \in p_0 \subseteq p$ with $h \supset h_0$. Conversely, if $h \supset k \in p$ but $p_0 \subseteq p$ is a basis containing h , then $k \supseteq k_0$ for some $k_0 \in p_0$, and thus $p = \uparrow(p_0 \setminus \{h\})$, contradicting the minimality of p_0 .

Finally, suppose that p is finitely-based witnessed by some finite $B \subseteq H/\sim$ and $\varphi_0 : B \rightarrow H$. Then $\partial h \in B$ and $\varphi_0(\partial h) = h$, for otherwise φ_0 extends to both φ_p and φ_q . Letting $\varphi'_0 : B \rightarrow H$ agree with φ_0 on $B \setminus \{\partial h\}$ but setting $\varphi'_0(\partial h) := h^c$ then gives a basis for φ_q . **More justification.** ■

2. THE GRAPH \mathcal{T}_P

Let $H \subseteq 2^X$ be a nested sub-pocset as before. Using Lemma 1.5, we construct a graph \mathcal{T}_H whose:

- *Vertices* are finitely-based orientations $p \subseteq H$;
- *Edges* are pairs $\{p, q\}$ such that $q = p \triangle \{h, h^c\}$ for some \subseteq -minimal $h \in H$.

The goal of this section is to establish Theorem A, stating that \mathcal{T}_P is acyclic (Proposition 2.5), and furthermore, that \mathcal{T}_P is a tree when P is finitely-separating (Proposition 2.6), or more generally, when P is chain-vanishing (Proposition 2.7).

2.1. Acyclicity of \mathcal{T}_P . Essentially, any cycle in \mathcal{T}_P (without backtracking) flips each \subseteq -minimal element

Definition 2.1. Fix $\mathcal{U}_0 \in V(\mathcal{T}_P)$ and $n \in \mathbb{N}$. A sequence $(A_i)_{i < n} \subseteq P$ is said to *induce a path from \mathcal{U}_0* if $(\mathcal{U}_i)_{i < n}$, defined by $\mathcal{U}_i := \mathcal{U}_{i-1} \triangle \{A_{i-1}, A_{i-1}^c\}$ for every $1 \leq i < n$, is a path in \mathcal{T}_P with each $A_i \in \mathcal{U}_i$.

Remark 2.2. Any path in \mathcal{T}_P is induced by its sequence of flipped basis elements.

Lemma 2.3. *Let $n \geq 3$. A path in \mathcal{T}_P from \mathcal{U}_0 induced by $(A_i)_{i < n}$ has no backtracking iff $A_i \neq A_{i-1}^c$ for every $1 \leq i < n$.*

Proof. Take $2 \leq i \leq n$. It suffices to show that $\mathcal{U}_{i-2} = \mathcal{U}_i$ iff $A_{i-1} = A_{i-2}^c$.

(\Rightarrow). We have by definition that $\mathcal{U}_i = \mathcal{U}_{i-2} \cup \{A_{i-1}^c, A_{i-2}^c\} \setminus \{A_{i-1}, A_{i-2}\}$, so since $A_{i-2} \in \mathcal{U}_{i-2} = \mathcal{U}_i$, we have $A_{i-2} = A_{i-1}^c$ as desired.

(\Leftarrow). Again by definition, by noting that the basis-flipping cancels out. ■

Lemma 2.4. *If $(A_i)_{i < n}$ induces a path in \mathcal{T}_P with no backtracking, then $(A_i)_{i < n}$ is strictly increasing.*

Proof. By Lemma 2.3, we have $A_i \neq A_{i-1}^c$ for every $1 \leq i < n$. Thus, since $A_i \in \mathcal{U}_i = \mathcal{U}_{i-1} \cup \{A_{i-1}^c\} \setminus \{A_{i-1}\}$, we see that $A_i \in \mathcal{U}_{i-1}$. Clearly $A_i \neq A_{i-1}$. It suffices to remove the three cases when $A_i \subseteq A_{i-1}$, $A_{i-1} \subseteq A_i^c$, and $A_i^c \subseteq A_{i-1}$, since then nestedness of P gives us $A_{i-1} \subsetneq A_i$, as desired.

- If $A_i \subseteq A_{i-1}$, then $A_{i-1} \in \mathcal{U}_i$, contradicting the definition of \mathcal{U}_i .
- If $A_{i-1} \subseteq A_i^c$, then $A_i^c \in \mathcal{U}_{i-1}$ by upward-closure of \mathcal{U}_{i-1} , a contradiction.
- If $A_i^c \subseteq A_{i-1}$, then $A_{i-1} \in \mathcal{U}_{i+1}$ by upward-closure of $\mathcal{U}_{i+1} \ni A_i^c$. But since $A_{i-1} \neq A_i^c$, we have by definition of \mathcal{U}_{i+1} that $A_{i-1} \in \mathcal{U}_i$, a contradiction. ■

Proposition 2.5. \mathcal{T}_P is acyclic.

Proof. Let $(\mathcal{U}_i)_{i < n}$ be a cycle in \mathcal{T}_P induced by $(A_i)_{i < n}$. Since cycles are non-backtracking, we have $A_0 \subsetneq A_0$ by Lemma 2.4, a contradiction. ■

Proposition 2.6. If P is chain-vanishing, then \mathcal{T}_P is connected (and hence a tree).

Proposition 2.7.

Proposition 2.8. If P is closed under complements, then \mathcal{T}_P is connected (and hence a tree).

Proof. If $\mathcal{U}, \mathcal{U}' \in V(\mathcal{T}_P)$ are two finitely-based orientations on P , then swapping their basis elements one by one as in Remark ? gives us a path between \mathcal{U} and \mathcal{U}' ; the closure of P is needed to ensure that those basis elements induce a path between the vertices, in that their complement lies in P . ■