

Computing Optimal Bayesian Decisions for Rank Aggregation via MCMC Sampling

David Hughes, Kevin Hwang and Lirong Xia

Presented by Zhibing Zhao

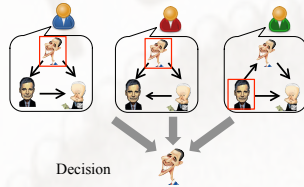


Rensselaer

Introduction

Social Choice (Rank Aggregation)

- Presidential elections
- Alternatives: candidates
- Agents: voters
- Profile: preferences of agents
- Mechanism: plurality rule
- Decision: president



Democracy



Classical social choice

The axiomatic approach [Arrow-50]

Truth



Statistical approaches

MLE; MAP

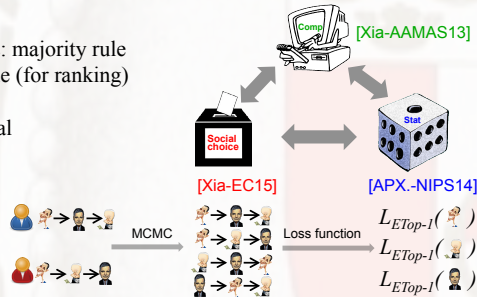
Our Goal and Approach

- Mechanisms
- 2 alternatives: majority rule
- Kemeny's rule (for ranking)

Goal

- Computational
- Axiomatic
- Statistical

Our approach



Statistical decision-theoretic framework for social choice

Inputs

- statistical model: $\Theta, S, \Pr_\theta(s)$
- decision space: D
- loss function: $L(\theta, d) \in \mathbb{R}$

unknown ground truth decision to make

Complexity

Theorem: if BayesianLoss for Mallows' model w.r.t. the exact Top-1 loss function and uniform prior has a polynomial-time approximation algorithm with constant approximation ratio, then $P=NP$

The mechanism

f : Profiles $\rightarrow D$ with minimum Bayesian expected loss:
 $f(R) \in \argmin_d E_{\theta \sim R} L(\theta, d)$

Mallows' Model [Mallows-1957]

- Fix the dispersion $\phi < 1$
- Parameter space
 - all full rankings over alternatives
- Sample space
 - i.i.d. generated full rankings
- Probabilities: given a ground truth ranking W , generate a ranking V
 $\Pr_\mu(V) \propto \phi^{\text{Kendall}(V, W)}$

Decision Space Examples

- All alternatives (resolute voting)
- Power set of alternatives excluding the empty set (irresolute voting)
- All full rankings over alternatives

Loss Function Examples

- Top- k
 - $L_{E\text{Top-}k}(W, d) = 0$, when d is ranked within top k
 - $L_{E\text{Top-}k}(W, d) = 1$, otherwise
- Other smooth functions

References

[Mallows 1957] Colin L. Mallows. Non-null ranking model. *Biometrika*, 44(1/2): 114–130, 1957.
[Sinclair 1995] Alistair Sinclair. Improved Bounds for Mixing Rates of Markov Chains and Multicommodity Flow. *Combinatorics, Probability and Computing*, 14(1):351–370, 1995.
[Sinclair 1999] Alistair Sinclair and Mark Jerrum. Approximate counting, uniform generation and rapidly mixing Markov chains. *Information and Computation*, 82(1):53–133, 1999.
[Jerrum 1998] Mark Jerrum and Alistair Sinclair. The Markov chain Monte Carlo method: an approach to approximate counting and integration. In Dori S. Hochbaum, editor, *Approximation algorithms for NP-hard problems*, pages 462–519. PWS Publishing Company, 1998.
[Liu 1996] Jun S. Liu. Metropolisized independent sampling with comparisons to rejection sampling and importance sampling. *Statistics and Computing*, 6(2):113–119, 1996.

Our Approach

Metropolis-Hastings MCMC Algorithms

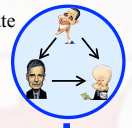
MCMC Algorithms

- Given profile R
- Generate N samples Q from a Markov chain with stationary distribution $\Pr(|R)$
- Compute $\argmin_d \sum_{\theta \in Q} L(\theta, d)$

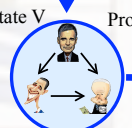
Metropolis-Hastings Sampling

- Initial state – a ranking
- Given proposal distributions $\{p_v(\cdot) : V \in \Theta\}$
 - $p_v(W) = p_w(V)$
- In each step
 - Current state V
 - Proposal next state W
 - The next state is W with probability $\frac{\Pr(W|R)p_w(V)}{\Pr(V|R)p_v(W)}$
 - Otherwise the next state is V

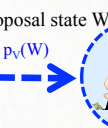
Initial state



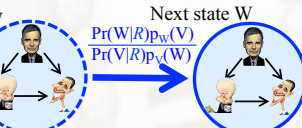
Current state V



Proposal state W



Next state W



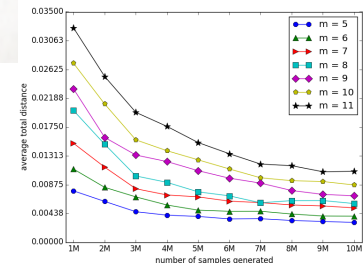
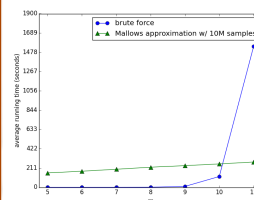
Theoretical Guarantee

- The performance of the algorithm closely depends on the mixing time of the MC
- Theorem.** The mixing time of the M-H algorithm for Mallows' model for any starting state is $\phi^{-k_{\max}} \text{Poly}(\text{other inputs})$
- Profile: $\{ \begin{matrix} \text{Candidate 1} > \text{Candidate 2} > \text{Candidate 3} \\ \text{Candidate 2} > \text{Candidate 1} > \text{Candidate 3} \\ \text{Candidate 3} > \text{Candidate 1} > \text{Candidate 2} \end{matrix} \times 2, \begin{matrix} \text{Candidate 1} > \text{Candidate 2} > \text{Candidate 3} \\ \text{Candidate 2} > \text{Candidate 3} > \text{Candidate 1} \\ \text{Candidate 3} > \text{Candidate 2} > \text{Candidate 1} \end{matrix} \times 3, \begin{matrix} \text{Candidate 1} > \text{Candidate 2} > \text{Candidate 3} \\ \text{Candidate 2} > \text{Candidate 1} > \text{Candidate 3} \\ \text{Candidate 3} > \text{Candidate 1} > \text{Candidate 2} \end{matrix} \times 4 \}$
- WMG: $k_{\max} = 14$

- k_{\max} : maximum cut in the weighted majority graph
- ϕ : dispersion parameter of the model
- Proof uses techniques from
 - novel analysis of a set of canonical paths [Sinclair, 1992; Jerrum and Sinclair 1989, 1996; Liu 1996]

Experiments

- Mallows' model
- D : single winner
- L : Top-1 loss function
- Data: from Preflib.org



Summary & Future Research

Summary

- Mallows' model and Condorcet's model
- Bayesian loss is NP-hard to approximate
- New MCMC algorithms
- Mallows' model: sample complexity/running time only exponential in k_{\max}
- Condorcet's model: polynomial sample complexity

Future Research

- Performance improvement of Markov chain sampler in practice
- Design and analysis of other Markov chain samplers
- Evaluation of the Markov chain approach