

# ECE 220: **Computer Systems and Programming**

## Lecture 12: Recursion, recursive sorting and Recursion with Backtracking

Instructor: Bruce X.B. Yu

Time: Monday, 9:00-10:20 AM (EE), 10:30-11:50 AM (CompE)

Friday, 13:00-14:20 PM (EE), 14:30-15:50 PM (CompE)

Location: LT Building North A 418/420



A **recursive function** is one that solves its task by **calling itself** on smaller pieces of data.

- Similar to recurrence function in mathematics.
- Like iteration -- can be used interchangeably; sometimes recursion results in a simpler solution.
- Must have at least 1 **base case** (terminal case) that ends the recursive process.

Example:  $n!$

# Factorial:



$$n! = n \cdot (n-1) \cdot (n-2) \cdot \dots \cdot 3 \cdot 2 \cdot 1$$

$$n! = \begin{cases} n \cdot (n-1)! & , n > 0 \\ 1 & , n = 0 \end{cases}$$

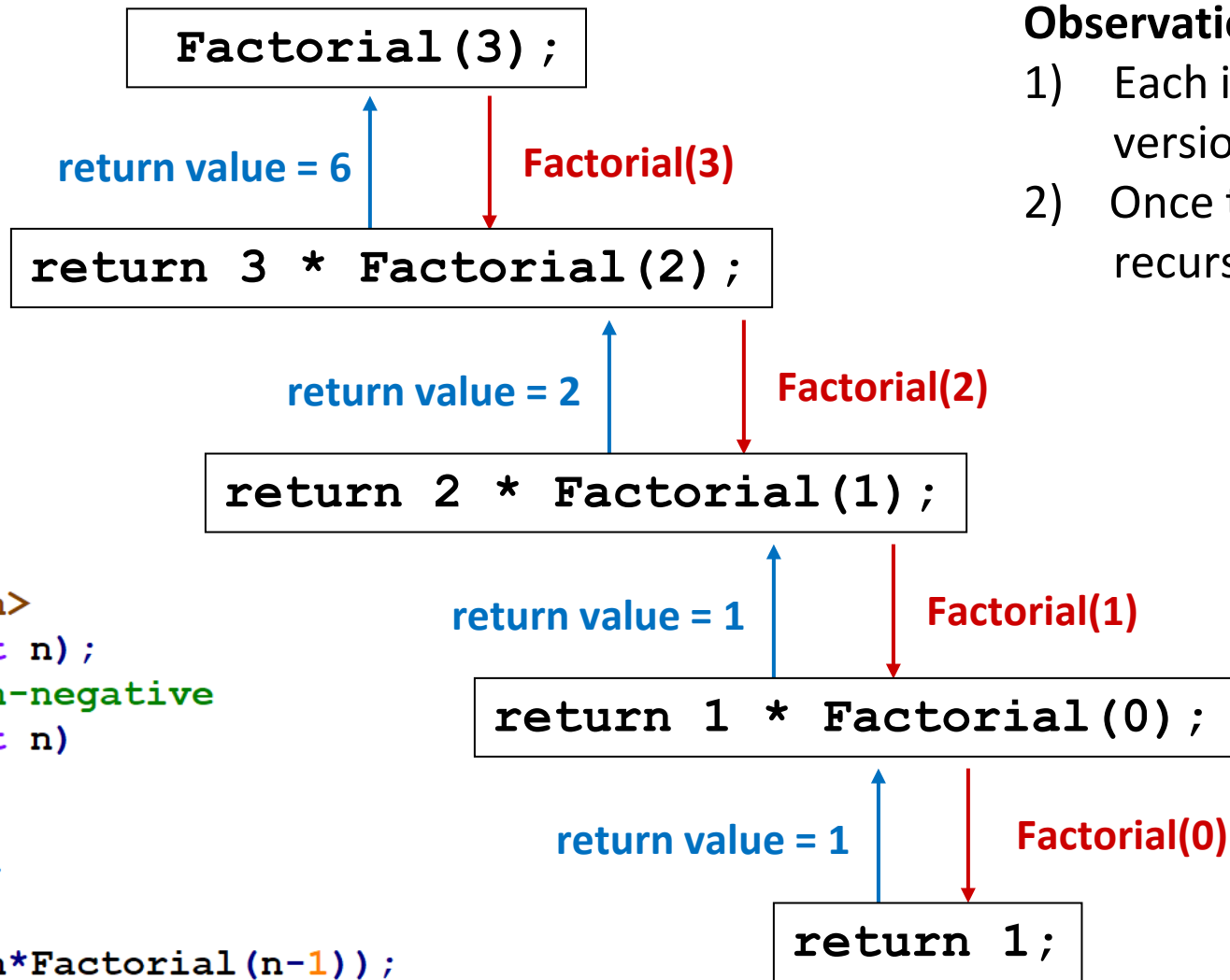
```
int Factorial(int n)
{
    if
        Return ....

    else

    return

}
```

```
1  #include <stdio.h>
2  int Factorial(int n);
3  //assume n is non-negative
4  int Factorial(int n)
5  {
6      if(n == 0)
7          return 1;
8      else
9          return (n*Factorial(n-1));
10 }
11
12 int main()
13 {
14     int n=3;
15     int result = Factorial(n);
16     printf("Factorial(%d)=%d \n",n,result);
17
18     return 0;
19 }
```



## Observation:

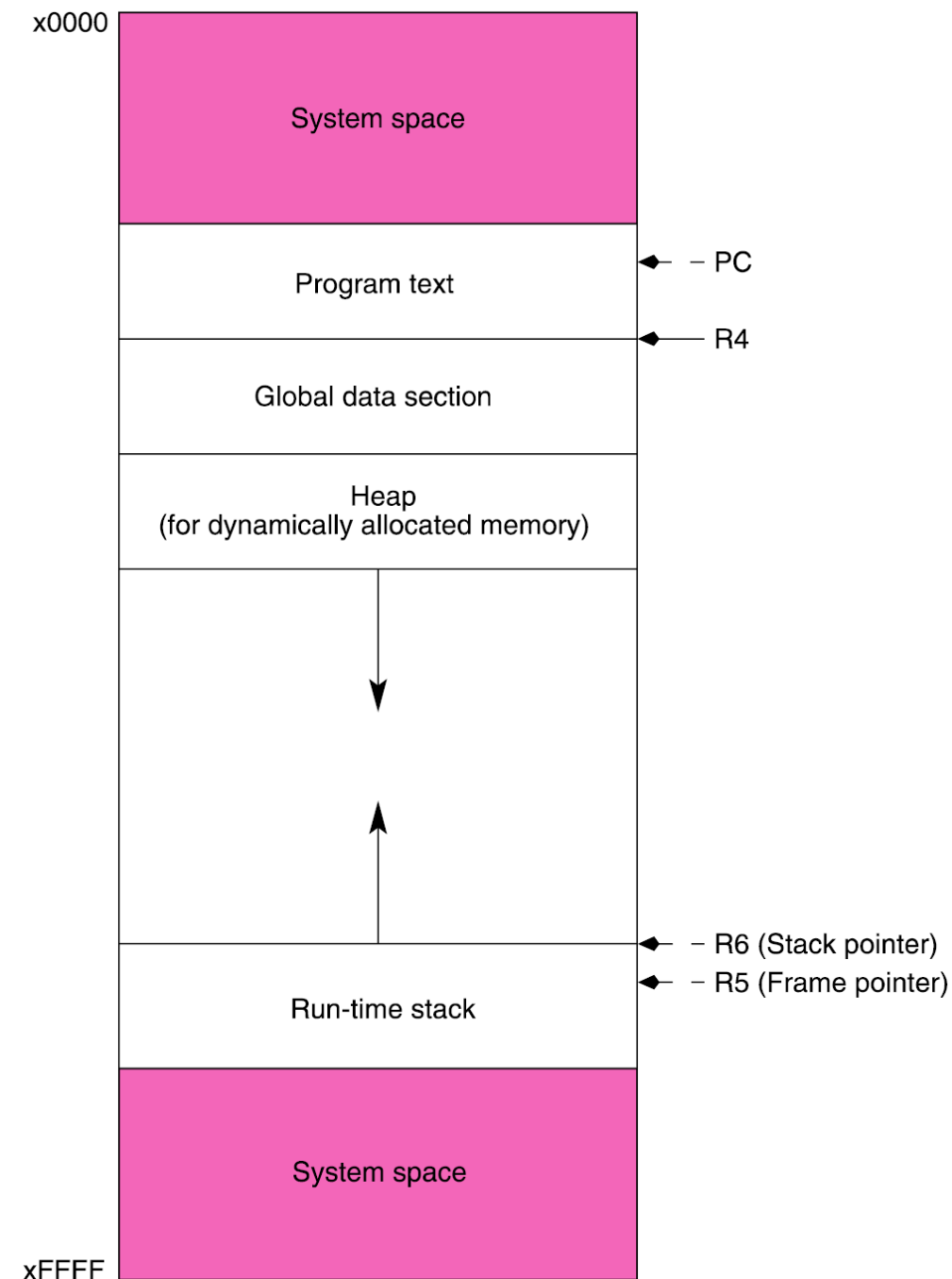
- 1) Each invocation solves a smaller version of the problem;
- 2) Once the base case is reached, recursive process stops.

```
1  #include <stdio.h>
2  int Factorial(int n);
3  //assume n is non-negative
4  int Factorial(int n)
5  {
6      if(n == 0)
7          return 1;
8      else
9          return (n*Factorial(n-1));
10 }
```

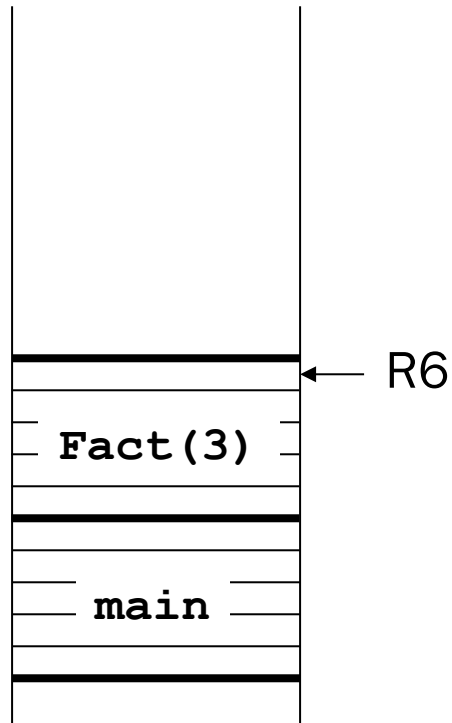
# Space for Variables

1. Global data section  
(global variables)
2. Run-time stack  
(local variables)

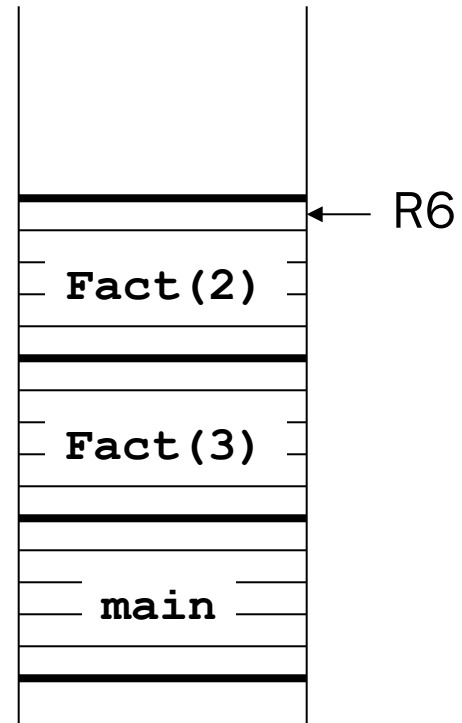
- **R4** (global pointer) points the first global variable
- **R5** (frame pointer) points the first local variable
- **R6** (stack pointer) points the top of run-time stack



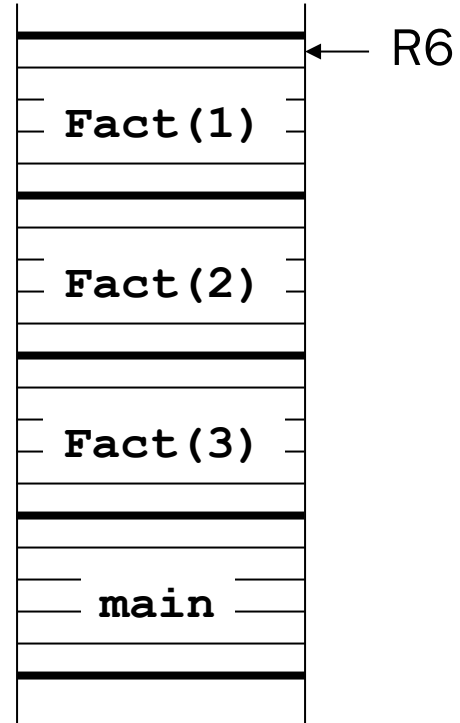
**main calls  
Factorial(3)**



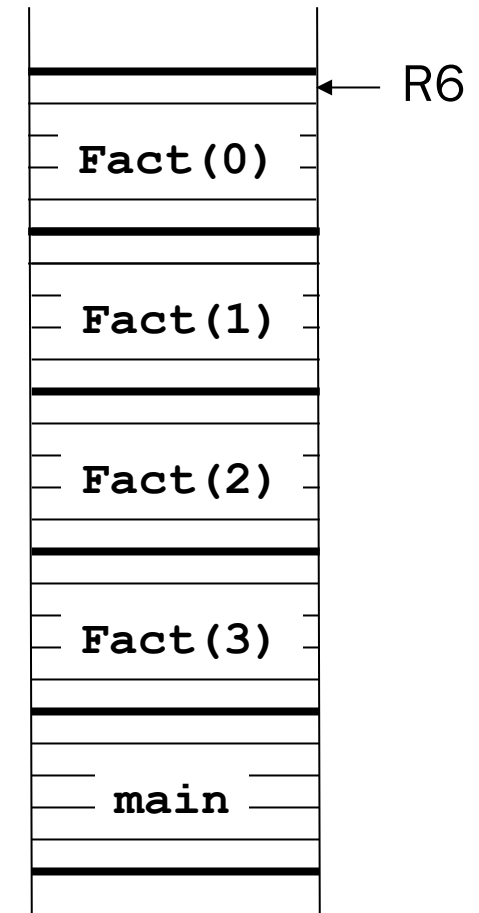
**Factorial(3) calls  
Factorial(2)**



**Factorial(2) calls  
Factorial(1)**



**Factorial(1) calls  
Factorial(0)**



```
1 .ORIG x3000
2 ; push argument
3     LD R6, STACK_TOP
4     AND R0,R0,#0
5     ADD R0,R0,#3
6     ADD R6,R6,#-1 ;R6 <- R6-1;
7     STR R0,R6,#0 ;push argument n
8 ; call subroutine
9     JSR FACTORIAL
10 ; pop return value from run-time stack (to R0)
11     LDR R0,R6,#0
12     ADD R6, R6, #2
13 ;Store the result
14     STR R0,R6,#0 ;dump the result at STACK_TOP
15     HALT
16
```

```

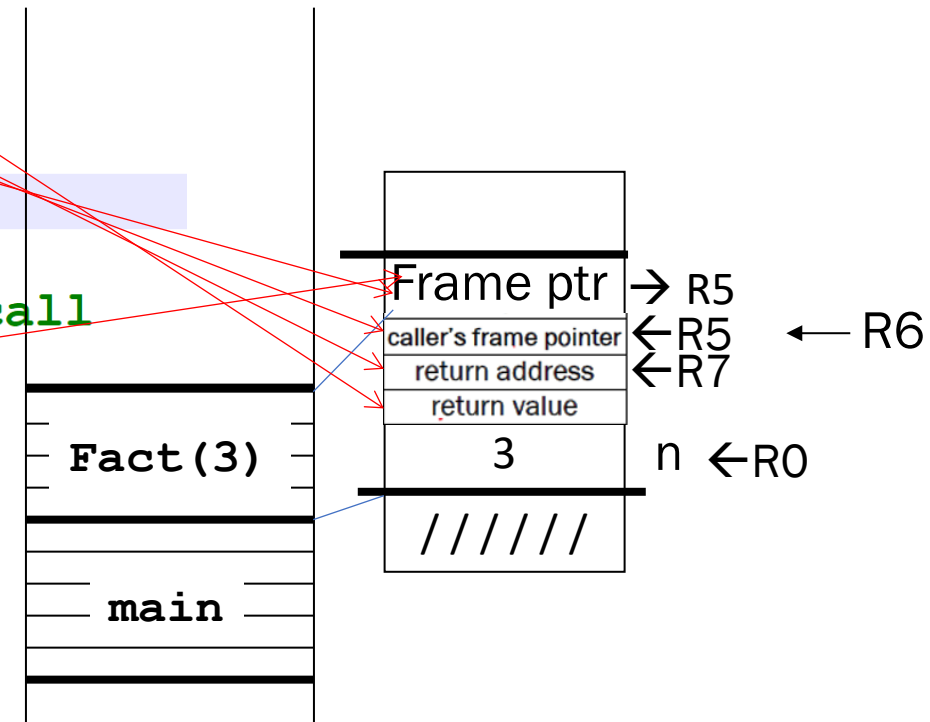
18 FACTORIAL:
19 ; push callee's bookkeeping info onto the run-time stack
20 ; allocate space in the run-time stack for return value
21     ADD R6, R6, #-1
22 ; store caller's return address and frame pointer
23     ADD R6, R6, #-1
24     STR R7, R6, #0
25     ADD R6, R6, #-1
26     STR R5, R6, #0
27 ; Update frame pointer for the callee
28     ADD R5, R6, #-1
29
30 ; if (n>0)
31     LDR R1, R5, #4
32     ADD R2, R1, #-1    R2 ← 2n-1
33     BRn ELSE
34 ; compute fn = n * factorial(n-1)
35 ; caller-built stack for factorial(n-1) function call
36 ; push n-1 onto run-time stack
37     ADD R6, R6, #-1
38     STR R2, R6, #0
39 ; call factorial subroutine
40     JSR FACTORIAL
41 ; pop return value from run-time stack (to R0)
42     LDR R0, R6, #0
43     ADD R6, R6, #1

```

```

1 #include <stdio.h>
2 int Factorial(int n);
3 //assume n is non-negative
4 int Factorial(int n)
5 {
6     if(n == 0)
7         return 1;
8     else
9         return (n*Factorial(n-1));
10 }
11
12 int main()
13 {
14     int n=3;
15     int result = Factorial(n);
16     printf("Factorial(%d)=%d \n",n,result);
17
18     return 0;
19 }

```





```

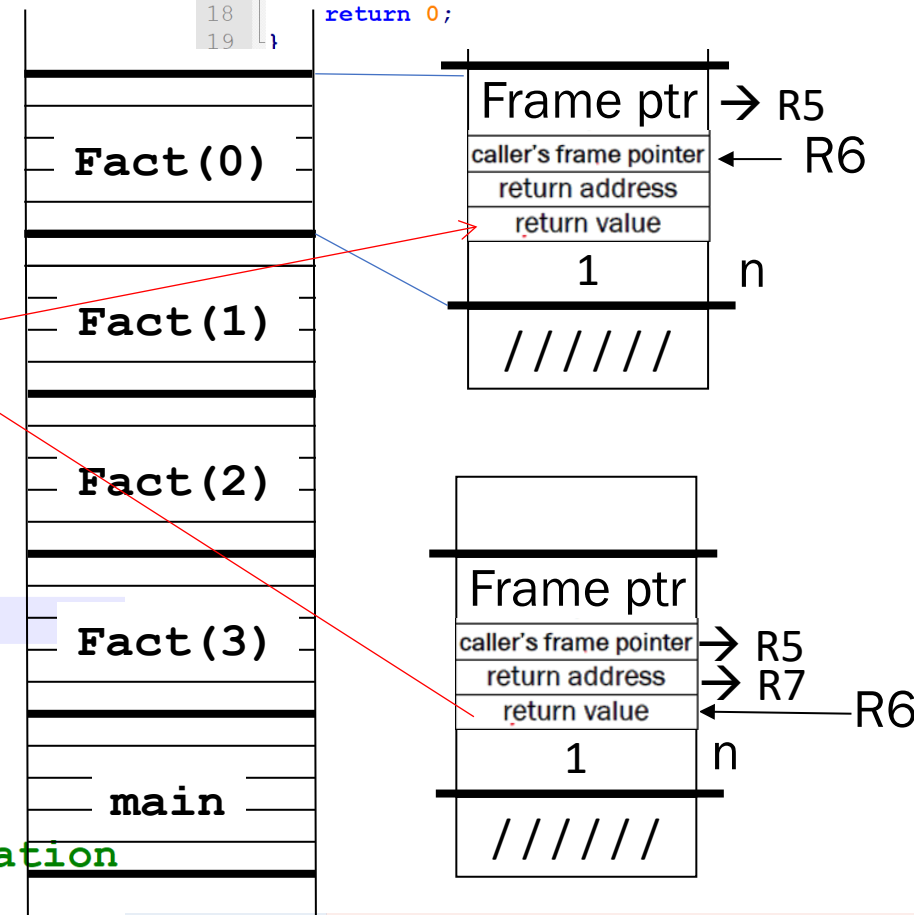
39 ; call factorial subroutine
40 JSR FACTORIAL
41 ; pop return value from run-time stack (to R0)
42 LDR R0, R6, #0
43 ADD R6, R6, #1
44 ; pop function argument from the run-time stack
45 ADD R6, R6, #1
46 ; multiply n by the return value (already in R0)
47 LDR R1, R5, #4
48 ; MUL R2, R0, R1 ; R2 <- n * factorial(n-1)
49 ST R7, SAVE_R7
50 JSR MULT
51 LD R7, SAVE_R7
52 ADD R0, R2, #0
53 BRnzp RETURN
54 ELSE:
55 ; store value of 1 in to the memory of return value
56 AND R0, R0, #0
57 ADD R0, R0, #1
58 ; tear down the run-time stack and return
59 RETURN:
60 ; write return value to the return entry
61 STR R0, R5, #3
62 ; pop local variable(s) from the run-time stack
63 ;no local variable for this implementation
64 ; restore caller's frame pointer and return address
65 LDR R5, R6, #0
66 ADD R6, R6, #1
67 LDR R7, R6, #0
68 ADD R6, R6, #1 ;stack pointer is at the return value location
69 ; return control to the caller function
70 RET

```

```

1 #include <stdio.h>
2 int Factorial(int n);
3 //assume n is non-negative
4 int Factorial(int n)
5 {
6     if(n == 0)
7         return 1;
8     else
9         return (n*Factorial(n-1));
10 }
11
12 int main()
13 {
14     int n=3;
15     int result = Factorial(n);
16     printf("Factorial(%d)=%d \n",n,result);
17
18     return 0;
19 }

```





```
71 ; multiply subroutine
72 ; input should be in R0 and R1
73 ; output should be in R2
74 MULT
75     ; save R3
76     ST R3, SAVE_R3
77     ; reset R2 and initialize R3
78     AND R2, R2, #0
79     ADD R3, R0, #0
80     ; perform multiplication
81     MULT_LOOP
82     ADD R3, R3, #-1
83     BRn MULT_DONE
84     ADD R2, R2, R1
85     BRnzp MULT_LOOP
86     MULT_DONE
87     ; restore R0
88     LD R3, SAVE_R3
89     RET
90
91 SAVE_R3                .BLKW #1
92 SAVE_R7                .BLKW #1
93 STACK_TOP              .FILL x4000
94 .END
```

# Recursive Binary Search



```
1 // C program to implement binary search using recursion
2 #include <stdio.h>
3
4 // A recursive binary search function. It returns location
5 // of x in given array arr[l..r] if present, otherwise -1
6 int binarySearch(int arr[], int l, int r, int x)
7 {
8     // checking if there are elements in the subarray
9     if (r >= l) {
10
11         // calculating mid point
12         int mid = (l + r) / 2;
13
14         // If the element is present at the middle itself
15         if (arr[mid] == x)
16             return mid;
17
18         // If element is smaller than mid, then it can only
19         // be present in left subarray
20         if (arr[mid] > x) {
21             return binarySearch(arr, l, mid - 1, x);
22         }
23
24         // Else the element can only be present in right
25         // subarray
26         return binarySearch(arr, mid + 1, r, x);
27     }
28
29     // We reach here when element is not present in array
30     return -1;
31 }
```

Ref: <https://www.geeksforgeeks.org/binary-search/#>

# Fibonacci Series

$$f(n) = f(n - 1) + f(n - 2)$$

$$f(1) = 1$$

$$f(0) = 1$$

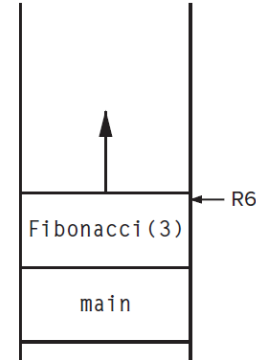
```
1  #include <stdio.h>
2
3  int Fibonacci(int n);
4
5  int main(void)
6  {
7      int in;
8      int number;
9
10     printf("Which Fibonacci number? ");
11     scanf("%d", &in);
12
13     number = Fibonacci(in);
14     printf("That Fibonacci number is %d\n", number);
15 }
16
17 int Fibonacci(int n)
18 {
19     int sum;
20
21     if (n == 0 || n == 1)
22         return 1;
23     else {
24         sum = (Fibonacci(n-1) + Fibonacci(n-2));
25         return sum;
26     }
```

```

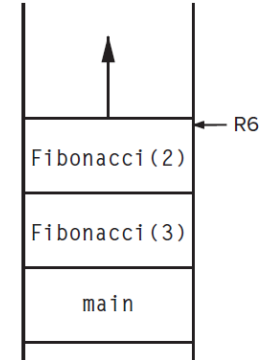
1  #include <stdio.h>
2
3  int Fibonacci(int n);
4
5  int main(void)
6  {
7      int in;
8      int number;
9
10     printf("Which Fibonacci number? ");
11     scanf("%d", &in);
12
13     number = Fibonacci(in);
14     printf("That Fibonacci number is %d\n", number);
15 }
16
17 int Fibonacci(int n)
18 {
19     int sum;
20
21     if (n == 0 || n == 1)
22         return 1;
23     else {
24         sum = (Fibonacci(n-1) + Fibonacci(n-2));
25         return sum;
26     }

```

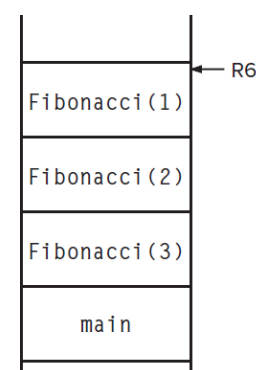
Consider, n=3



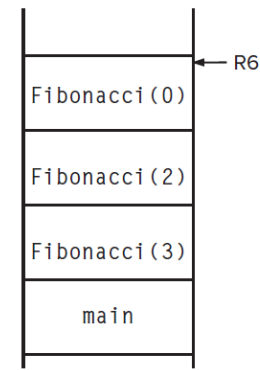
Step 1: Initial call



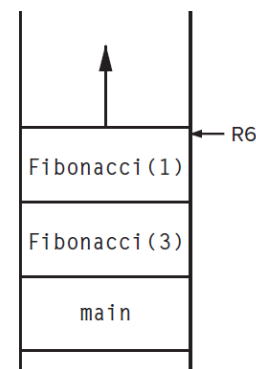
Step 2: Fibonacci(3) calls Fibonacci(2)



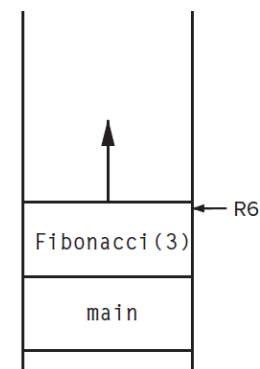
Step 3: Fibonacci(2) calls Fibonacci(1)



Step 4: Fibonacci(2) calls Fibonacci(0)



Step 5: Fibonacci(3) calls Fibonacci(1)

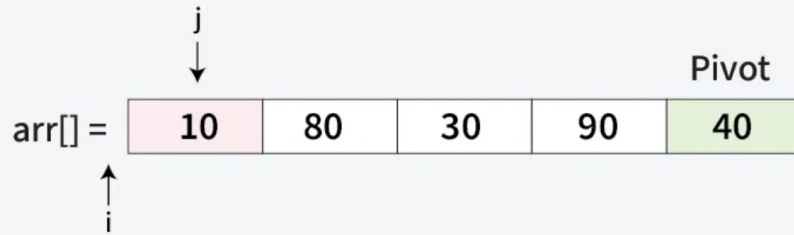


Step 6: Back to the starting point

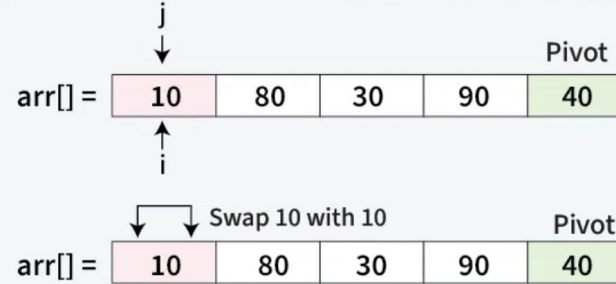
# int partition(int array[], int l, int h)



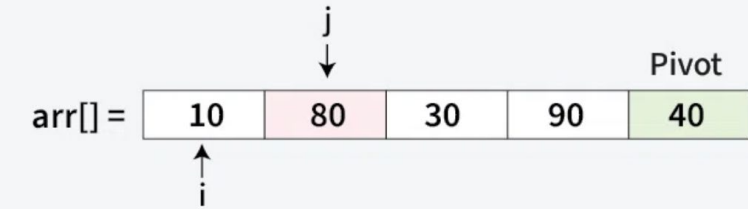
**01** Step | Pivot Selection: The last element  $\text{arr}[4] = 40$  is chosen as the pivot. Initial Pointers:  $i = -1$  and  $j = 0$ .



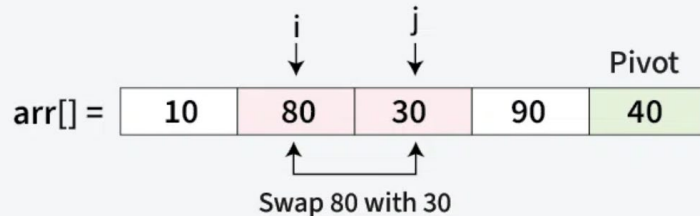
**02** Step | Since,  $\text{arr}[j] < \text{pivot}$  ( $10 < 40$ ) Increment  $i$  to 0 and swap  $\text{arr}[i]$  with  $\text{arr}[j]$ . Increment  $j$  by 1



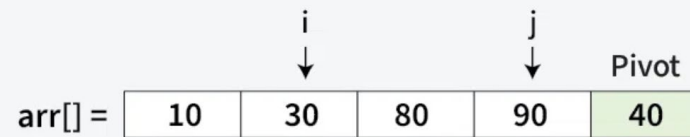
**03** Step | Since,  $\text{arr}[j] > \text{pivot}$  ( $80 > 40$ ) No swap needed. Increment  $j$  by 1



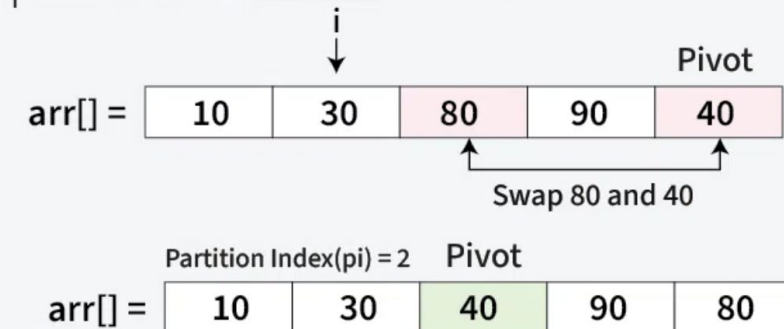
**04** Step | Since,  $\text{arr}[j] < \text{pivot}$  ( $30 < 40$ ) Increment  $i$  by 1 and swap  $\text{arr}[i]$  with  $\text{arr}[j]$ . Increment  $j$  by 1



**05** Step | Since,  $\text{arr}[j] > \text{pivot}$  ( $90 > 40$ ) No swap needed. Increment  $j$  by 1



**06** Step | Since traversal of  $j$  has ended. Now move pivot to its correct position, Swap  $\text{arr}[i + 1] = \text{arr}[2]$  with  $\text{arr}[4] = 40$ .



The logic is simple, we start from the leftmost element and keep track of the index of smaller (or equal) elements as  $i$ . While traversing, if we find a smaller element, we swap the current element with  $\text{arr}[i]$ . Otherwise, we ignore the current element.

```
void swap(int *a, int *b)
```

```
{  
    int tmp;  
    tmp=*a;  
    *a=*b;  
    *b=tmp;  
}
```

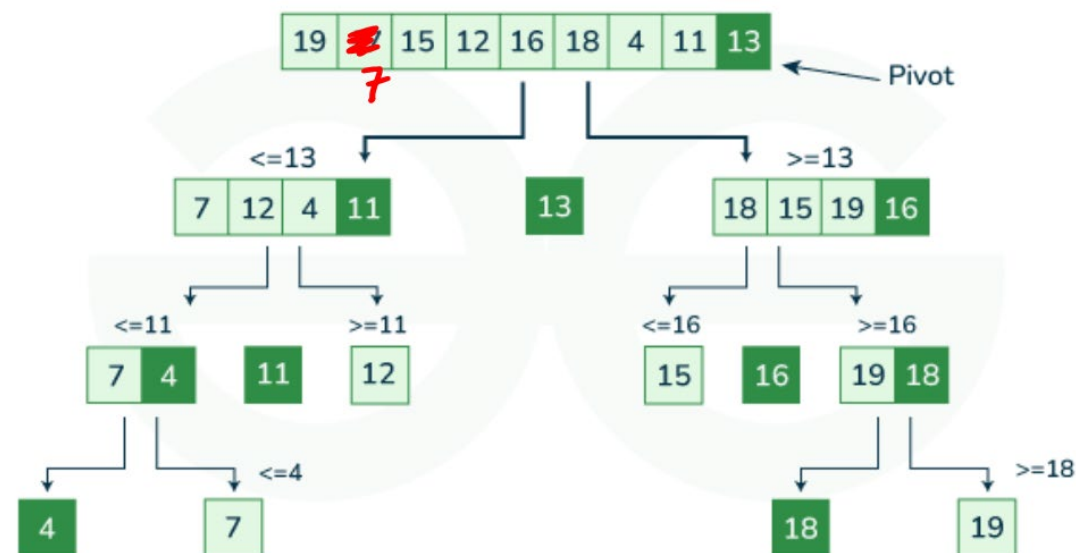
```
void print(int array[], int size)
```

```
{  
    int i;  
    for(i=0;i<size;i++)  
        printf("%d ", array[i]);  
    printf("\n");  
}
```

```
int main()
```

```
{  
    int size=8;  
    int l=0;  
    int h=size-1;  
    int array[8]={10, 20, 80, 30, 100, 90, 15, 40};  
    quickSort(array, l, h);  
    print(array, size);  
    return 0;  
}
```

## Quick Sort Algorithm



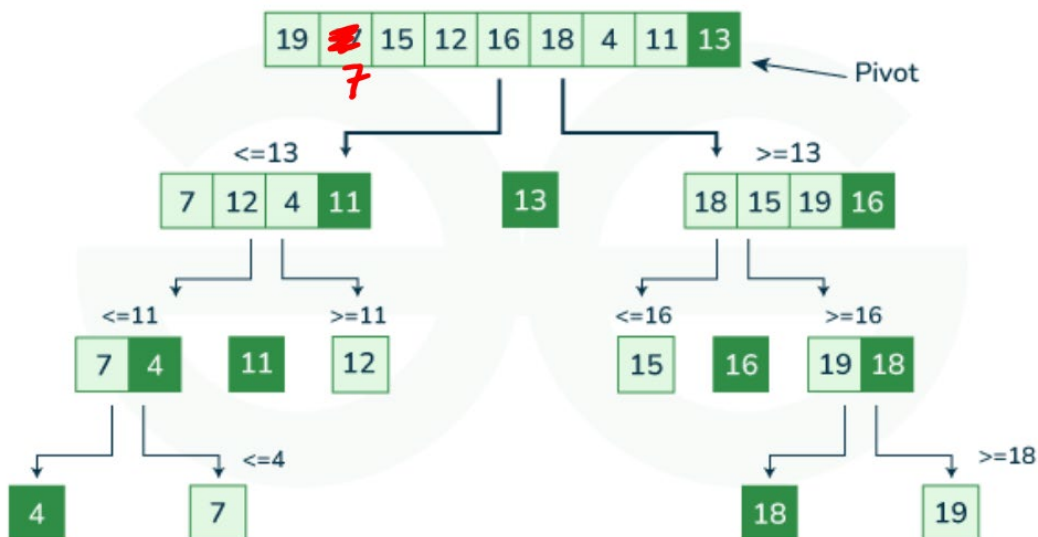
<https://www.geeksforgeeks.org/quick-sort-in-c/>



```
void quickSort(int array[], int l, int h)
{
    if (l < h) {
        // Call partition() to get the partition index
        int p = partition(array, l, h);
        // Call partition() to partition the left side
        quickSort(array, l, p-1);
        // Call partition() to partition the right side
        quickSort(array, p+1, h);
    }
}
```

```
int partition(int array[], int l, int h)
{
    int i = l-1;
    int x = array[h];
    int j;
    for (j = l; j <= h-1; j++) {
        if (array[j] <= x) {
            i++;
            swap(&array[i], &array[j]);
        }
    }
    swap(&array[i+1], &array[h]);
    return (i+1);
}
```

Quick Sort Algorithm





# Recursive Bubble Sort



```
21 void bubble_recursion(int *array, int n)
22 {
23     if(n==0)
24         return;
25     int i, temp, swap = 0;
26
27     //sort number in ascending order
28
29     swap = 0;
30     for(i=0;i<n;i++)
31     {
32         //swap the two numbers if order is incorrect
33         if(array[i]>array[i+1])
34         {
35             temp = array[i];
36             array[i] = array[i+1];
37             array[i+1] = temp;
38             //set the swap flag
39             swap = 1;
40         }
41     }
42     n--;
43     if (swap==0)
44         return;
45     bubble_recursion(array, n);
46 }
```

```
1  #include <stdio.h>
2  #define SIZE 8
3  void bubble_recursion(int *array, int n);
4
5  int main()
6  {
7      int n = SIZE-1;
8      int array[] = {6,5,3,1,8,7,2,4};
9      int i;
10     bubble_recursion(array, n);
11
12     printf("sorted array: \n");
13     for(i=0;i<SIZE;i++){
14         printf("%d ", array[i]);
15     }
16     printf("\n");
17
18     return 0;
19 }
```

# Recursive Merge-Insertion Sort

```
19 void merge_recursion(int *array, int i)
20 {
21     if (i==SIZE)
22         return;
23     int j, temp;
24     temp = array[i];
25     for(j=i-1; (j>=0 && (temp < array[j])); j--)
26     {
27         //shift element to the right
28         array[j+1] = array[j];
29     }
30     //insert at the proper location
31     array[j+1] = temp;
32     i++;
33     merge_recursion(array, i);
34 }
```

```
1 #include <stdio.h>
2 #define SIZE 8
3 void merge_recursion(int *array, int i);
4 int main()
5 {
6     int i=1;
7     int array[] = {6,5,3,1,8,7,2,4};
8     merge_recursion(array, i);
9     printf("sorted array: \n");
10    for(i=0;i<SIZE;i++){
11        printf("%d ", array[i]);
12    }
13    printf("\n");
14    return 0;
15 }
```

## Recursive Backtracking:

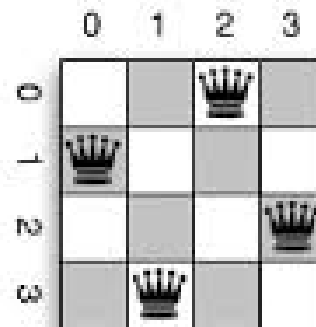
Backtracking is an algorithmic-technique for solving problems recursively by trying to build a solution incrementally, one piece at a time, removing those solutions that fail to satisfy the constraints of the problem statement.

# N queens problem using recursive Backtracking

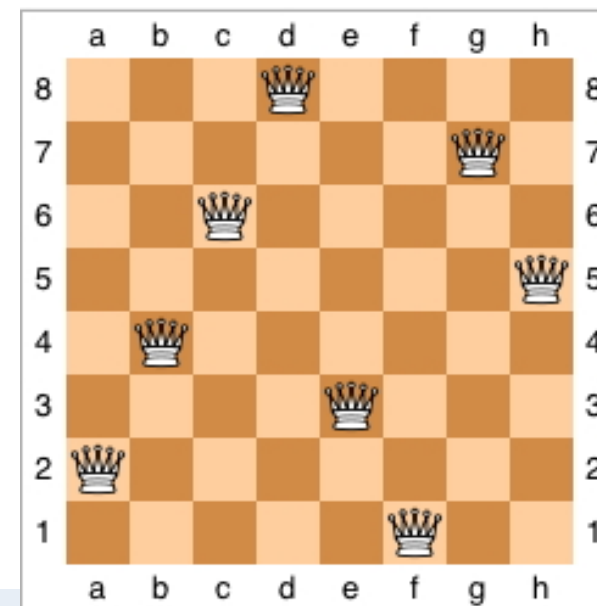


ZJUI

- Place N queens on an NxN chessboard so that none of the queens are **under attack**;
- Brute force: total number of possible placements:
- $\sim N^2$  Choose N  $\sim 4.4 B$  (N=8)



0	0	1	0
1	0	0	0
0	0	0	1
0	1	0	0





**Row 0**

0	1	2	3

	0	1	2	3
0				
1				
2				
3				


QueenPosition[]


	0	1	2	3
Row 0				

	0	1	2	3
0				
1				
2				
3				

QueenPosition[]

Place **0**<sup>th</sup> Queen on the **0**<sup>th</sup> Column of **0**<sup>th</sup> Row

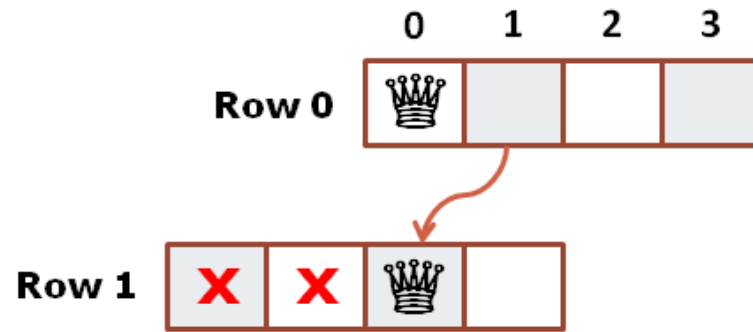
	0	1	2	3
Row 0				

	0	1	2	3
0				
1				
2				
3				

QueenPosition[]

Row 0 (0,0)

Add **0**<sup>th</sup> Queen's position to position array



	0	1	2	3
0	♀			
1	X	X	♀	
2				
3				

QueenPosition[]

Row 0 (0,0)

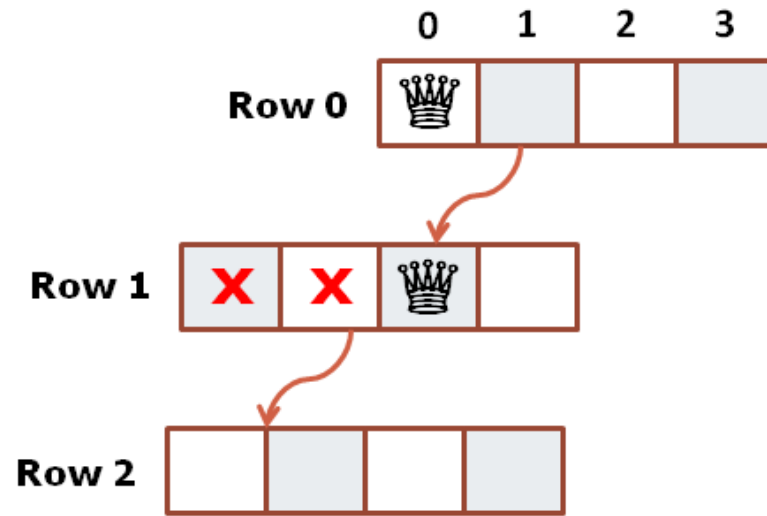
Row 1 (1,2)

Go to the next level of recursion.

Place the 1<sup>st</sup> queen on the 1<sup>st</sup> row such that she does not attack the 0<sup>th</sup> queen and add that to Positions.







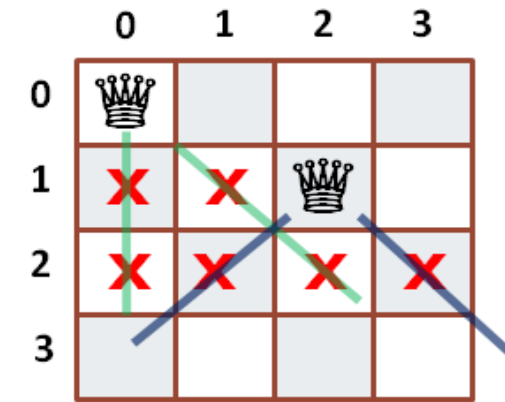
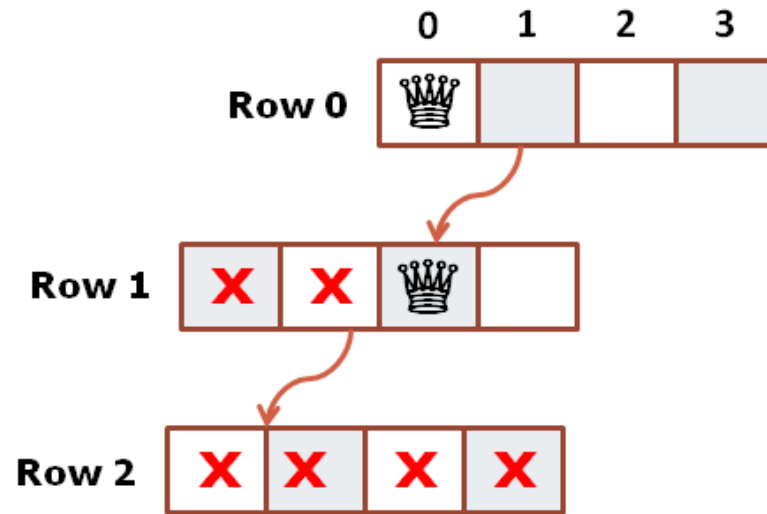
	0	1	2	3
0				
1	X	X		
2				
3				

QueenPosition[]

**Row 0** (0,0)

**Row 1** (1,2)

In the next level of recursion, find the cell on **2<sup>nd</sup>** row such that it is not under attack from any of the available queens.

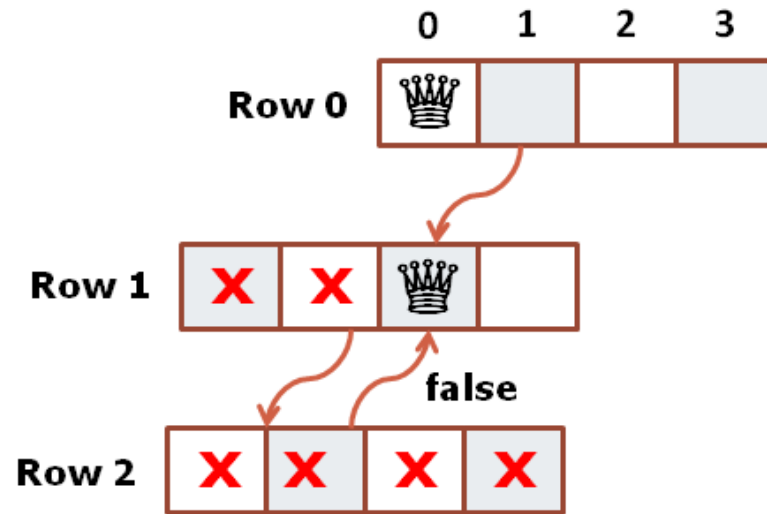




QueenPosition[ ]

**Row 0** (0,0)

**Row 1** (1,2)

But cell **(2,0)** and **(2,2)** are under attack from **0<sup>th</sup>** queen and cell **(2,1)** and **(2,3)** are under attack from **1<sup>st</sup>** queen.



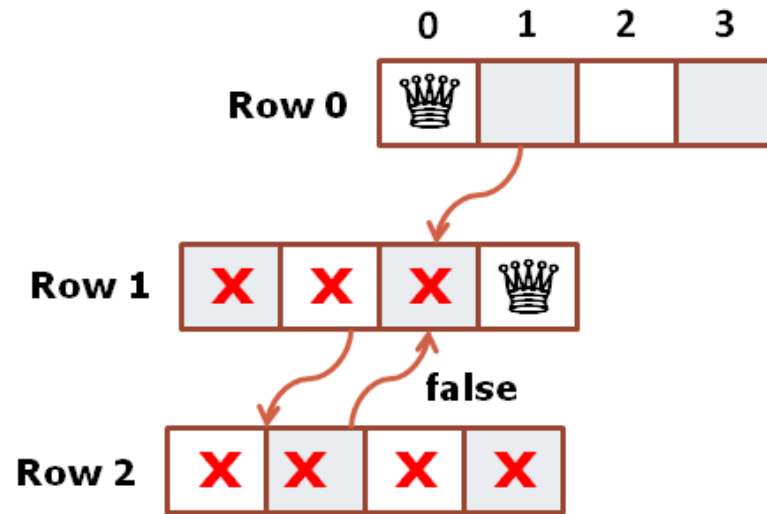
	0	1	2	3
0				
1	X	X		
2	X	X	X	X
3				

QueenPosition[]

**Row 0** (0,0)

**Row 1** (1,2)

So function will return false to the calling function.



	0	1	2	3
0				
1	X	X	X	
2				
3				

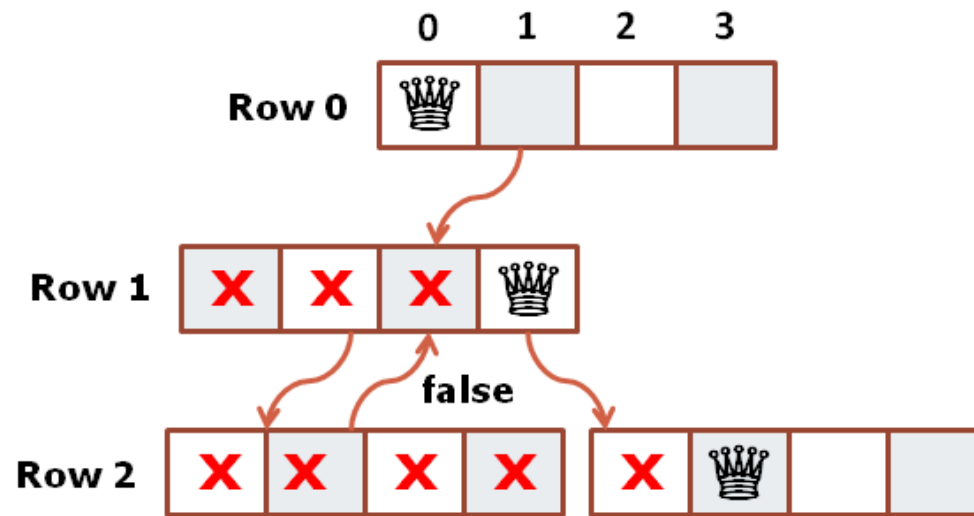
QueenPosition[]

Row 0 (0,0)

~~Row 1~~ (1,2)

Row 1 (1,3)

Calling function will try to find next possible place for the 1<sup>st</sup> queen on 1<sup>st</sup> row and update the queen position in position array.



	0	1	2	3
0	Queen			
1	X	X	X	Queen
2	X	Queen		
3				

QueenPosition[ ]

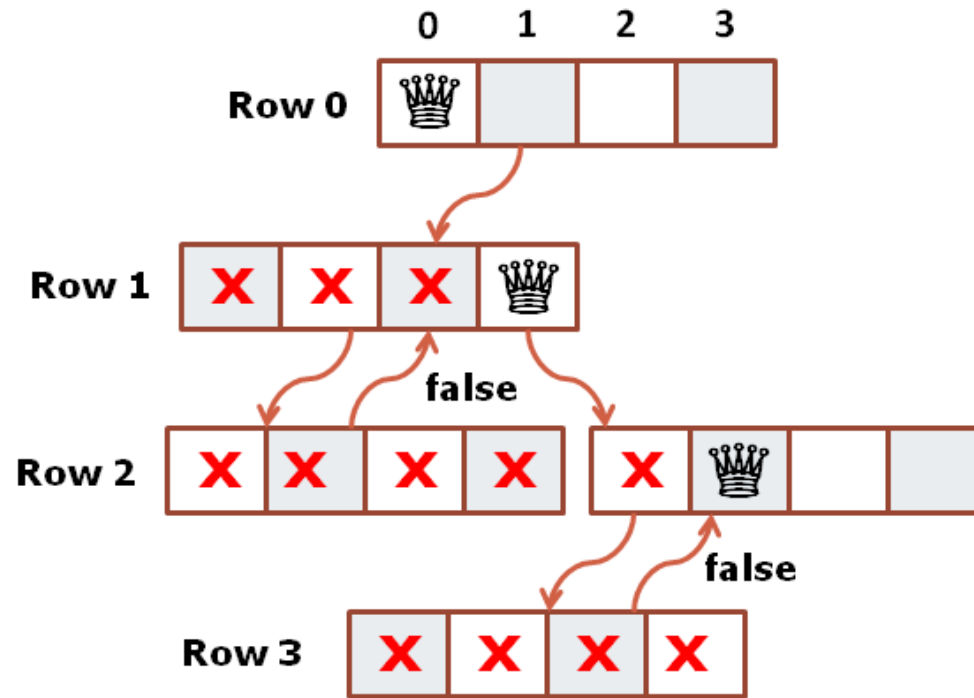
Row 0 (0,0)

Row 1 (1,3)

Row 1 (2,1)

Again find the cell on **2<sup>nd</sup>** row such that it is not under attack from any of the available queens.

Placing the queen in cell **(2,1)** as it is not under attack from any of the queen.



	0	1	2	3
0	♔			
1	✗	✗	✗	♔
2	✗	♔		
3				

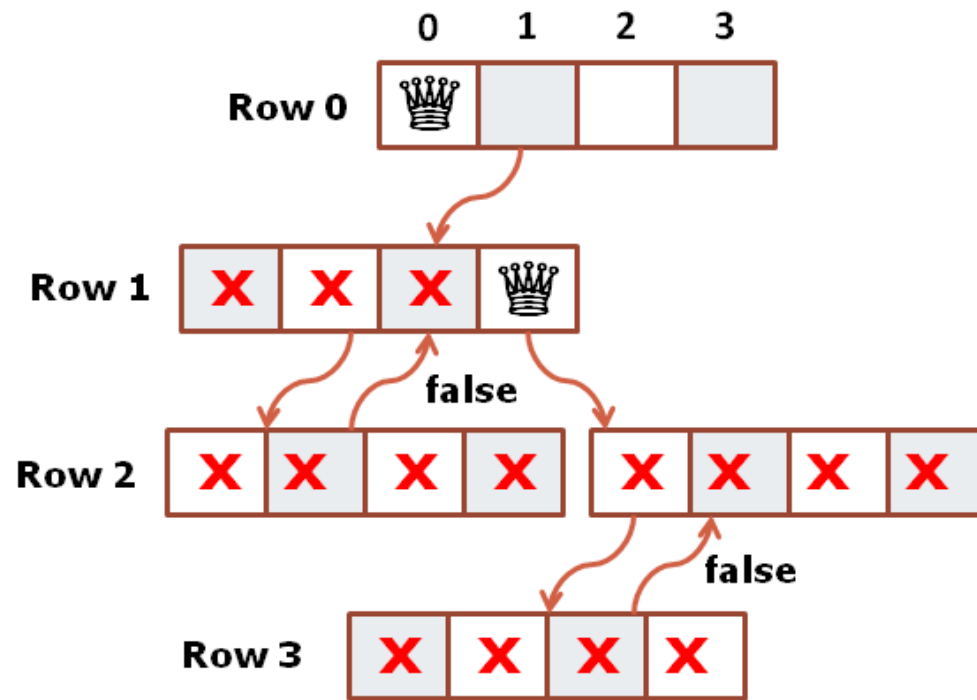
QueenPosition[ ]

Row 0 (0,0)

Row 1 (1,3)

Row 1 (2,1)

For **3<sup>rd</sup>** queen, no safe cell is available on **3<sup>rd</sup>** row.  
So function will return false to calling function.



	0	1	2	3
0				
1	X	X	X	
2	X	X	X	X
3				

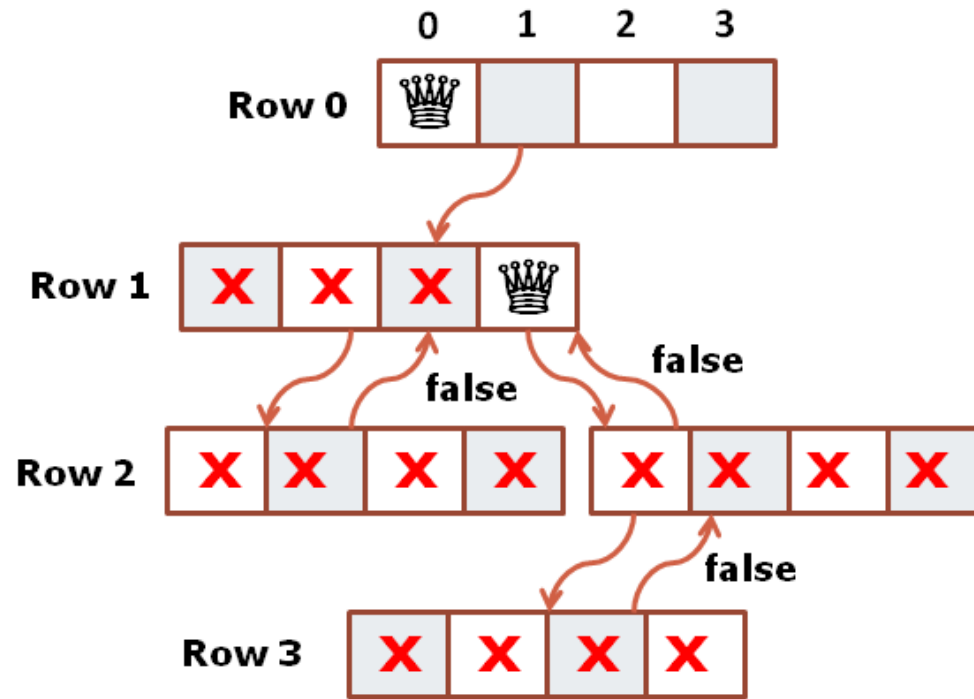
QueenPosition[]

Row 0 (0,0)

Row 1 (1,3)

~~Row 2 (2,1)~~

Queen at the 2<sup>nd</sup> row tries to find next safe cell.



	0	1	2	3
0				
1	X	X	X	
2				
3				

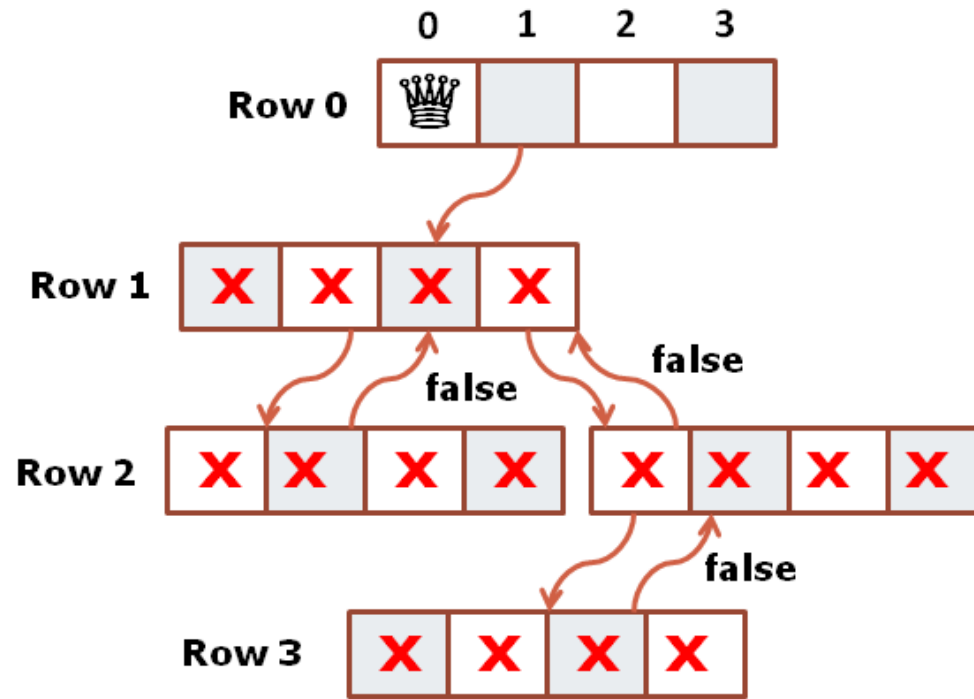
QueenPosition[ ]

Row 0 (0,0)

Row 1 (1,3)

But as both remaining cells are under attack from other queens, this function also returns false to its calling function.





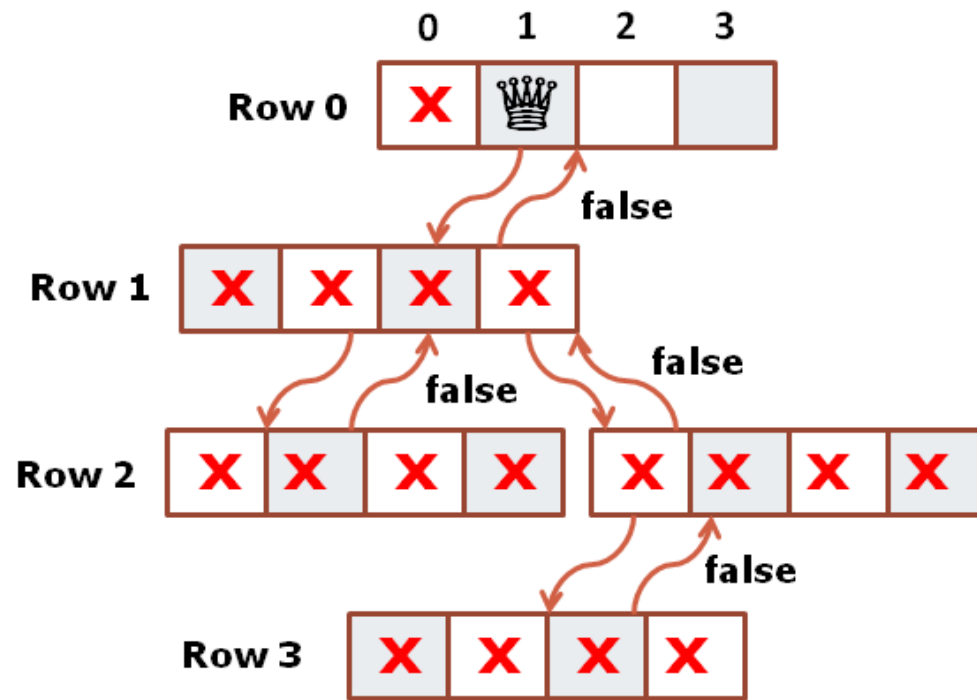
	0	1	2	3
0				
1	X	X	X	X
2				
3				

QueenPosition[ ]

Row 0 (0,0)

~~Row 1~~ (1,3)

Queen at the **1<sup>st</sup>** row tries to find next safe cell.  
 But as queen is in the last cell, it will return false to  
 its calling function.



	0	1	2	3
0	<b>X</b>			
1				
2				
3				


QueenPosition[ ]

~~Row 0~~ (0,0)

Row 0 (0,1)

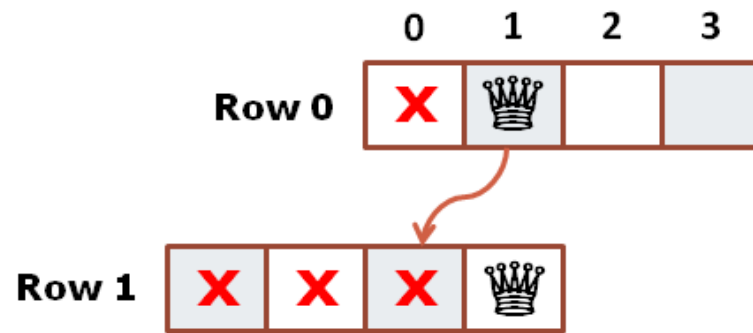
Queen at the **1<sup>st</sup>** row tries to find next safe cell.  
 Let us remove these failed recursion calls from the screen.

	0	1	2	3
Row 0	<b>X</b>			

	0	1	2	3
0	<b>X</b>			
1				
2				
3				

QueenPosition[]

Row 0 (0,1)

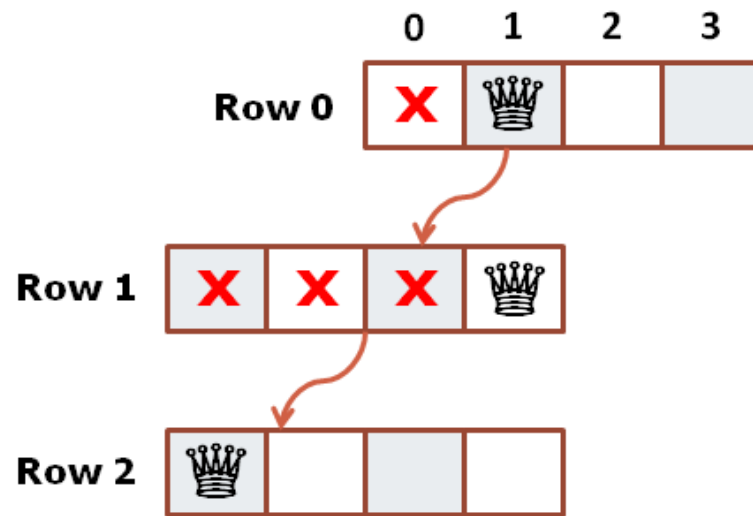


	0	1	2	3
0	X	Queen		
1	X	X	X	Queen
2				
3				

QueenPosition[ ]

**Row 0** (0,1)

**Row 1** (1,3)



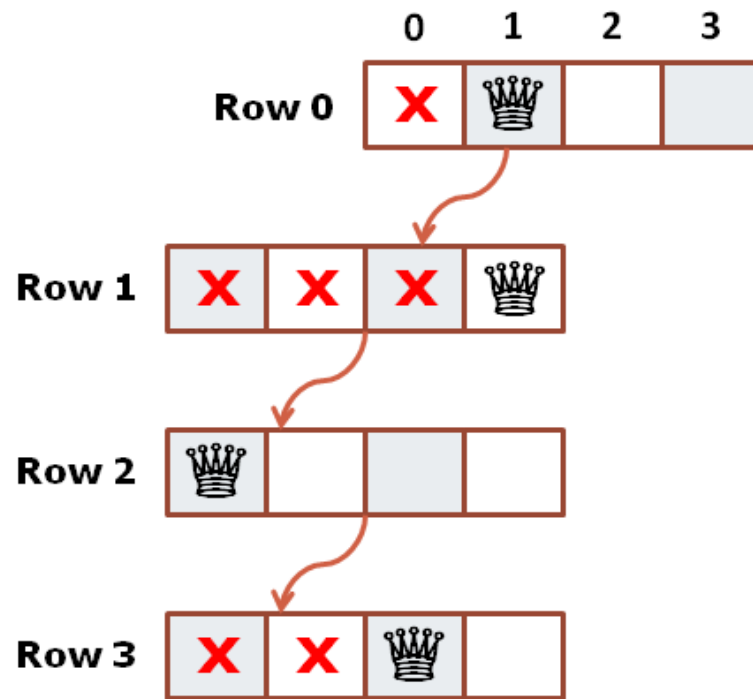
	0	1	2	3
0	X	Queen		
1	X	X	X	Queen
2	Queen			
3				

QueenPosition[]

**Row 0** (0,1)

**Row 1** (1,3)

**Row 2** (2,0)



	0	1	2	3
0	X	Queen		
1	X	X	X	Queen
2	Queen			
3	X	X	Queen	

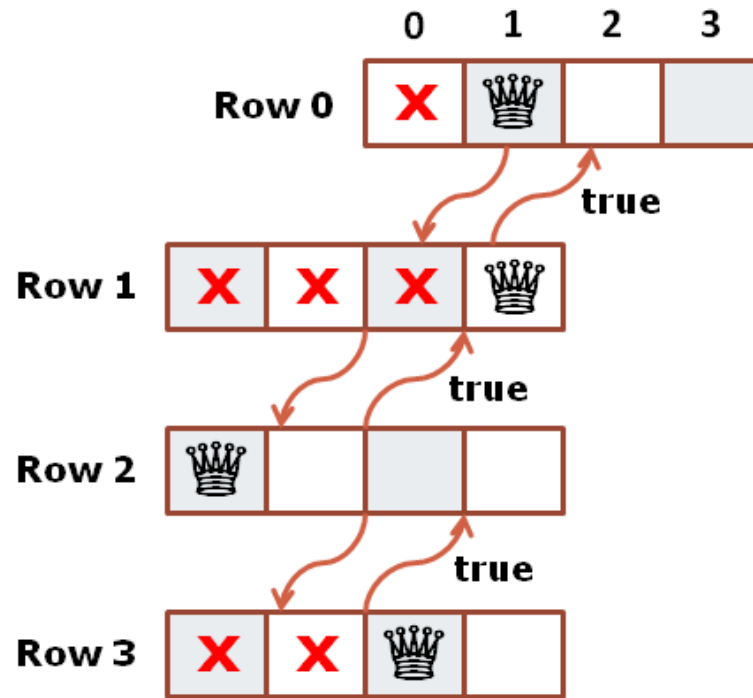
QueenPosition[]

**Row 0** (0,1)

**Row 1** (1,3)

**Row 2** (2,0)

**Row 3** (3,2)



	0	1	2	3
0	X	Queen		
1	X	X	X	Queen
2	Queen			
3	X	X	Queen	

QueenPosition[]

**Row 0** (0,1)

**Row 1** (1,3)

**Row 2** (2,0)

**Row 3** (3,2)

All functions will return true to their calling function.  
It means all queens are placed on the board such that they are not attacking each other.

# N Queens with backtracking



ZJUI

- `int board[N][N]` represents placement of queens
  - `board[i][j] = 0`: no queen at row `i` column `j`
  - `board[i][j] = 1`: queen at row `i` column `j`
- Initialize, for all `i,j` `board[i][j] = 0`
- Functions
  - `PrintBoard(board)`: Prints board on the screen
  - `IsSafe(board, row, col)`: returns 1 if new queen can be placed at (row,col) in board
  - `Solve(board, row)`: recursively attempts to place (N-row) queens; returns 0 if it fails

0	0	0	0
0	0	0	0
0	0	0	0
0	0	0	0

Initial board

`Solve(board,3)` returns 0

1	0	0	0
0	0	0	0
0	1	0	0
0	0	0	0



- N-Queen Problem by **Backtracking**
  1. **Decision**  
Place a queen at a safe place.
  2. **Recursion**  
Explore the solution for the next row.
  3. **Backtrack (Undo)**  
Remove the queen if no solution for the next row.
  4. **Base case**  
Reach the goal.

## N-Queen (4x4) Backtracking – CODE (Main function)

```
1  #include <stdio.h>
2
3  //Solve 4x4 n Queen problem using recursion with backtracking
4
5  #define N 4
6  #define true 1
7  #define false 0
8
9  void printSolution(int board[N][N]);
10 int Solve(int board[N][N], int col);
11 int isSafe(int board[N][N], int row, int col);
12
13 int main()
14 {
15     int board[N][N] = {{0,0,0,0},{0,0,0,0},{0,0,0,0},{0,0,0,0}};
16
17     //game started at row 0
18     if(Solve(board,0) == false)
19     {
20         printf("Solution does not exist.\n");
21         return 1;
22     }
23
24     printf("Solution: \n");
25     printSolution(board);
26     return 0;
27 }
```

## N-Queen (4x4) Backtracking – CODE (Solve function)

```
29 int Solve(int board[N][N], int row)
30 {
31     //base case
32     if(row>=N)
33         return true;
34
35     //find a safe column(j) to place queen
36     int j;
37     for(j=0;j<N;j++)
38     {
39         //column j is safe, place queen here
40         if(isSafe(board, row, j) == true)
41         {
42             board[row][j]=1;
43             printf("Current Play: \n");
44             printSolution(board);
45
46             //increment row to place the next queen
47             if(Solve(board, row+1) == true)
48                 return true;
49             //attempt to place queen at row+1 failed,-
50             //backtrack to row and remove queen
51             board[row][j]=0;
52             printf("Backtrack: \n");
53             printSolution(board);
54         }
55     }
56     return false;
57 }
```

## N-Queen (4x4) Backtracking – CODE (isSafe & PrintSolution functions)

```
59 int isSafe(int board[N][N], int row, int col)
60 {
61     int i, j;
62     for(i=0; i<row; i++)
63     {
64         for(j=0; j<N; j++)
65         {
66             //check whether there's a queen at the same column or the 2 diagonals
67             if(((j==col) || (i-j == row-col) || (i+j == row + col)) && (board[i][j]==1))
68                 return false;
69         }
70     }
71     return true;
72 }
73
74 void printSolution(int board[N][N])
75 {
76     int i, j;
77     for(i=0; i<N; i++)
78     {
79         for(j=0; j<N; j++)
80             printf(" %d ", board[i][j]);
81         printf("\n");
82     }
83 }
84 }
```

1	0	0	0
0	0	0	0
0	1	0	0
0	0	0	0

# Thank You!