15-213 Recitation 11Processes, Signals, Tshlab

4 November 2019

Outline

- Logistics
- Process Lifecycle
- Error Handling
- Signal Handling

Logistics

- Malloc Final due tomorrow (11/5)
 - Can use up to 2 late days!
 - Style grading mm.c (not checkheap)
- Midterm regrades released
 - Review exam in Professor OH

Shell Lab

- **Due date:** next Thursday (November 14th)
- Simulate a Linux-like shell with I/O redirection
- Review the writeup carefully.
 - Review once before starting, and again when halfway through
 - This will save you a lot of style points and a lot of grief!
- Read Chapter 8 in the textbook:
 - Process lifecycle and signal handling
 - How race conditions occur, and how to avoid them
 - Be careful not to use code from the textbook without understanding it first.

Process "Lifecycle"

- fork()
 Create a duplicate, a "child", of the process
- execve()
 Replace the running program
- ... [Complete Work]
- exit()
 End the running program
- waitpid()Wait for a child process to terminate

Notes on Examples

- Full source code of all programs is available
 - TAs may demo specific programs
- In the following examples, exit() is called
 - We do this to be explicit about the program's behavior
 - Exit should generally be reserved for terminating on error
- Unless otherwise noted, assume all syscalls succeed
 - Error checking code is omitted.
 - Be careful to check errors when writing your own shell!

Processes are separate

- How many lines are printed?
- If pid is at address 0x7fff2bcc264c, what is printed?

```
int main(void) {
    pid_t pid;
    pid = fork();
    printf("%p - %d\n", &pid, pid);
    exit(0);
}
```

Processes are separate

- How many lines are printed?
- If pid is at address 0x7fff2bcc264c, what is printed?

```
int main(void) {
   pid_t pid;
   pid = fork();
   printf("%p - %d\n", &pid, pid);
   exit(0);
```

0x7fff2bcc264c - 24750
0x7fff2bcc264c - 0

The order and the child's PID (printed by the parent) may vary, but the address will be the same in the parent and child.

What does this program print?

```
int main(void) {
    char *args[3] = {
        "/bin/echo", "Hi 18213!", NULL
    };
    execv(args[0], args);
    printf("Hi 15213!\n");
    exit(0);
}
```

What does this program print?

```
int main(void) {
    char *args[3] = {
        "/bin/echo", "Hi 18213!", NULL
    };
    execv(args[0], args);
    printf("Hi 15213!\n");
    exit(0);
}
```

Hi 18213!

What about this program? What does it print?

```
int main(void) {
    char *args[3] = {
        "/bin/blahblah", "Hi 15513!", NULL
    };
    execv(args[0], args);
    printf("Hi 14513!\n");
    exit(0);
}
```

What about this program? What does it print?

```
int main(void) {
    char *args[3] = {
        "/bin/blahblah", "Hi 15513!", NULL
    };
    execv(args[0], args);
    printf("Hi 14513!\n");
    exit(0);
}
```

Hi 14513!

On Error

What should we do if malloc fails?

```
const size_t HUGE = 1 * 1024 * 1024 * 1024;
int main(void) {
   char *buf = malloc(HUGE * HUGE);
```

```
printf("Buf at %p\n", buf);
free(buf);
exit(0);
```

On Error

What should we do if malloc fails?

```
const size_t HUGE = 1 * 1024 * 1024 * 1024;
int main(void) {
    char *buf = malloc(HUGE * HUGE);
    if (buf == NULL) {
        fprintf(stderr, "Failure at %u\n", LINE );
       exit(1);
    printf("Buf at %p\n", buf);
    free(buf);
   exit(0);
```

Exit values can convey information

Two values are printed. Are they related?

Exit values can convey information

Two values are printed. Are they related?

```
int main(void) {
    pid t pid = fork();
    if (pid == 0) { exit(getpid()); }
    else {
        int status = 0;
        waitpid(pid, &status, 0);
        printf("0x%x exited with 0x%x\n", pid,
               WEXITSTATUS(status));
    exit(0);
```

0x7b54 exited with 0x54

They're the same!... almost. Exit codes are only one byte in size.

Processes have ancestry

What's wrong with this code? (assume that fork succeeds)

```
int main(void) {
    int status = 0, ret = 0;
    pid t pid = fork();
    if (pid == 0) {
        pid = fork();
        exit(getpid());
    }
    ret = waitpid(-1, &status, 0);
    printf("Process %d exited with %d\n", ret, status);
    ret = waitpid(-1, &status, 0);
    printf("Process %d exited with %d\n", ret, status);
    exit(0);
```

Processes have ancestry

What's wrong with this code? (assume that fork succeeds)

```
int main(void) {
    int status = 0, ret = 0;
                                   waitpid will reap only
    pid t pid = fork();
                                   children, not grandchildren,
   if (pid == 0) {
                                   so the second waitpid call
        pid = fork();
        exit(getpid());
                                  will return an error.
    }
    ret = waitpid(-1, &status, 0);
    printf("Process %d exited with %d\n", ret, status);
    ret = waitpid(-1, &status, 0);
    printf("Process %d exited with %d\n", ret, status);
   exit(0);
```

How many different sequences can be printed?

```
int main(void) {
    int status;
    if (fork() == 0) {
        pid_t pid = fork();
        printf("Child: %d\n", getpid());
        if (pid == 0) {
            exit(0);
        // Continues execution...
    pid_t pid = wait(&status);
    printf("Parent: %d\n", pid);
    exit(0);
```

How many different sequences can be printed?

```
int main(void) {
    int status;
    if (fork() == 0) {
        pid t pid = fork();
        printf("Child: %d\n", getpid());
        if (pid == 0) {
            exit(0);
        // Continues execution...
    pid t pid = wait(&status);
    printf("Parent: %d\n", pid);
    exit(0);
                                                              print
                   fork
                                                       exit
                                         exit
                                 print
```

How many different lines are printed?

```
int main(void) {
    char *tgt = "child";
    pid_t pid = fork();
    if (pid == 0) {
        pid = getppid(); // Get parent pid
            tgt = "parent";
    }
    kill(pid, SIGKILL);
    printf("Sent SIGKILL to %s:%d\n", tgt, pid);
    exit(0);
}
```

How many different lines are printed?

```
int main(void) {
    char *tgt = "child";
    pid_t pid = fork();
    if (pid == 0) {
        pid = getppid(); // Get parent pid
            tgt = "parent";
    }
    kill(pid, SIGKILL);
    printf("Sent SIGKILL to %s:%d\n", tgt, pid);
    exit(0);
}
```

Anywhere from 0-2 lines. The parent and child try to terminate each other.

Signals and Handling

- Signals can happen at any time
 - Control when through blocking signals
- Signals also communicate that events have occurred
 - What event(s) correspond to each signal?
- Write separate routines for receiving (i.e., signals)

Counting with signals

Will this code terminate?

```
volatile int counter = 0;
void handler(int sig) { counter++; }
int main(void) {
    signal(SIGCHLD, handler);
    for (int i = 0; i < 10; i++) {
        if (fork() == 0) { exit(0); }
    while (counter < 10) {</pre>
        mine bitcoin();
    return 0;
```

Counting with signals

Will this code terminate?

```
volatile int counter = 0;
void handler(int sig) { counter++; }
int main(void) {
                                              (Don't use signal, use
    signal(SIGCHLD, handler);
                                              Signal or sigaction
    for (int i = 0; i < 10; i++) {
                                              instead!)
        if (fork() == 0) { exit(0); }
    while (counter < 10) {</pre>
        mine bitcoin();
    return 0;
                                            It might not, since
                  (Don't busy-wait, use
```

sigsuspend instead!)

signals can coalesce.

Proper signal handling

How can we fix the previous code?

- Remember that signals will be coalesced, so the number of times a signal handler has executed is **not** necessarily the same as number of times a signal was sent.
- We need some other way to count the number of children.

Proper signal handling

How can we fix the previous code?

- Remember that signals will be coalesced, so the number of times a signal handler has executed is **not** necessarily the same as number of times a signal was sent.
- We need some other way to count the number of children.

```
void handler(int sig) {
    pid_t pid;
    while ((pid = waitpid(-1, NULL, WNOHANG)) > 0) {
        counter++;
    }
}
```

(This instruction isn't atomic. Why won't there be a race condition?)

Error in UNIX - return value

- Can System call fail ?
- How to tell the difference?

Error in UNIX - What error?

- Can System call fail ?
- How to tell the difference ?
 - Returned -1
- So, my fantastic system call failed.
- How can I tell what got wrong?

Error handling - What now?

int main() { Can System call fail? How to tell the difference? while (!quit) { Returned -1 So, my fantastic system call failed.int fd = open(userfile,O_RDWR); **if** (fd < 0) { How can I tell what got wrong? printf("Failed\n"); The error is in errno (global) perror("open"); // use errno (only if the syscall fail, what do you get on success ?) exit(-1); What do I do?

Error and signals

- Can System call fail ?
- How to tell the difference ?
 - Returned -1
- So, my fantastic system call failed.
- How can I tell what got wrong?
 - The error is in errno (a global)
 - If success errno may contain anything
- What do I do ?
 - Look at errno and take action
- Hey, here comes a signal...

```
int main() {
 while (!quit) {
    int fd = open(userfile,O RDWR);
    if (fd < 0) {
      if (errno == EACCESS) {
        // tell user he's wrong
        continue;
      } else if(...) {
      } else {
        perror("open"); // use errno
        exit(-1);
```

Error and signals: Recap

- You can't expect people to block signals around all error handling logic
- Hence, your signal handler shouldn't interfere with them
- Solution:
 - Do not make any system call that could set errno
 - Save and restore errno (store at beginning of handler and restore after)

Example Question

What are all possible output values? int main() { int val = 2; printf("%d", 0); fflush(stdout); if (fork() == 0) { val++; printf("%d", val); fflush(stdout); } else { val--; printf("%d", val); fflush(stdout); wait(NULL); val++; printf("%d", val); fflush(stdout); exit(0);

If you get stuck

- Read the writeup!
- Do manual unit testing before runtrace and sdriver!
- Read the writeup!
- Post private questions on Piazza!
- Think carefully about error conditions.
 - Read the man pages for each syscall when in doubt.
 - What errors can each syscall return?
 - How should the errors be handled?

Appendix: Blocking signals

- Surround blocks of code with calls to sigprocmask.
 - Use SIG_BLOCK to block signals at the start.
 - Use SIG_SETMASK to restore the previous signal mask at the end.
- Don't use SIG_UNBLOCK.
 - We don't want to unblock a signal if it was already blocked.
 - This allows us to nest this procedure multiple times.

```
sigset_t mask, prev;
sigemptyset(&mask, SIGINT);
sigaddset(&mask, SIGINT);
sigprocmask(SIG_BLOCK, &mask, &prev);
// ...
sigprocmask(SIG_SETMASK, &prev, NULL);
```

Appendix: Errno

#include <errno.h>

- Global integer variable used to store an error code.
 - Its value is set when a system call fails.
 - Only examine its value when the system call's return code indicates that an error has occurred!
 - Be careful not to call make other system calls before checking the value of errno!
- Lets you know why a system call failed.
 - Use functions like strerror, perror to get error messages.
- Example: assume there is no "foo.txt" in our path

```
int fd = open("foo.txt", O_RDONLY);
if (fd < 0) perror("open");
// open: No such file or directory</pre>
```

Appendix: Writing signal handlers

G1. Call only async-signal-safe functions in your handlers.

- Do not call printf, sprintf, malloc, exit! Doing so can cause deadlocks, since these functions may require global locks.
- We've provided you with sio_printf which you can use instead.

■ G2. Save and restore errno on entry and exit.

- If not, the signal handler can corrupt code that tries to read errno.
- The driver will print a warning if errno is corrupted.

G3. Temporarily block signals to protect shared data.

This will prevent race conditions when writing to shared data.

Avoid the use of global variables in tshlab.

- They are a source of pernicious race conditions!
- You do not need to declare any global variables to complete tshlab.
- Use the functions provided by tsh_helper.