

## Question 1

1)

No

Because:  $D^+ = \{D\}$

Hence  $D \rightarrow A$  not  $\in F^+$

2)

$X = \{A, B, C, D, I, J\}$  if the left-hand side of  $F$  is a super key.

$A$  can be removed because  $BCDIJ^+ = ABCDEGHIJ = R$

$B$  cannot be removed because  $CDIJ^+ = ACDEGHIJ \neq R$

$C$  can be removed because  $BDIJ^+ = ABCDEGHIJ = R$

$D$  can be removed because  $BIJ^+ = ABCDEGHIJ = R$

$I$  cannot be removed because  $BJ^+ = BDGHJ \neq R$

$J$  cannot be removed because  $BI^+ = BI \neq R$

Because  $C \rightarrow IJ$

Hence  $BC^+ = ABCDEGHIJ = R$

.....

$AB/BC/BDI/BIJ$  are the candidate keys for  $R$

3)

64

ABC

ABD

ABI

ABG

ABIJ

4)

$F_m = \{AB \rightarrow C, C \rightarrow I, C \rightarrow J, J \rightarrow D, J \rightarrow H, J \rightarrow G, DI \rightarrow A, DI \rightarrow E\}$

5)

2NF.

Because non-prime attribute  $E, G, H$  is fully functionally determined by prime attribute.

6)

Not dependency-preserving

Because  $R1 = \{ACE\}$ ,  $R2 = \{BCDE\}$ ,  $R3 = \{DJGHI\}$ , their FD set  $F1, F2, F3$ .  $(F1 \cup F2 \cup F3)^+ \text{ not } = F^+$

7)

Not lossless-join

decomposition	A	B	C	D	E	G	H	I	J
$R_1 = \{ACE\}$	a		a		a				
$R_2 = \{BCDE\}$		a	a	a	a				
$R_3 = \{DJGHI\}$				a		a	a	a	a

decomposition	A	B	C	D	E	G	H	I	J
$R_1 = \{ACE\}$	a		a	b	a	a	a	b	b
$R_2 = \{BCDE\}$	a	a	a	a	a	a	a	b	b
$R_3 = \{DJGHI\}$	b			a	b	a	a	a	a

8)

Step 1:

$F_m = \{AB \rightarrow C, C \rightarrow I, C \rightarrow J, J \rightarrow D, J \rightarrow H, J \rightarrow G, DI \rightarrow A, DI \rightarrow E\}$

Step 2:

because  $AB/BC/BDI/BIJ$  is candidate key, hence  $AB \rightarrow C$  can be subset  $\{ABC\}$

then because the left-hand of  $J \rightarrow D, J \rightarrow H, J \rightarrow G, DI \rightarrow A, DI \rightarrow E, C \rightarrow I, C \rightarrow J$  contains no  $A/B/C$ ,

hence they can be subset  $\{JD\}, \{JH\}, \{JG\}, \{DIA\}, \{DIE\}$

So the BCNF relations is  $\{ABC\}, \{JD\}, \{JH\}, \{JG\}, \{DIA\}, \{DIE\}, \{CI\}, \{CJ\}$

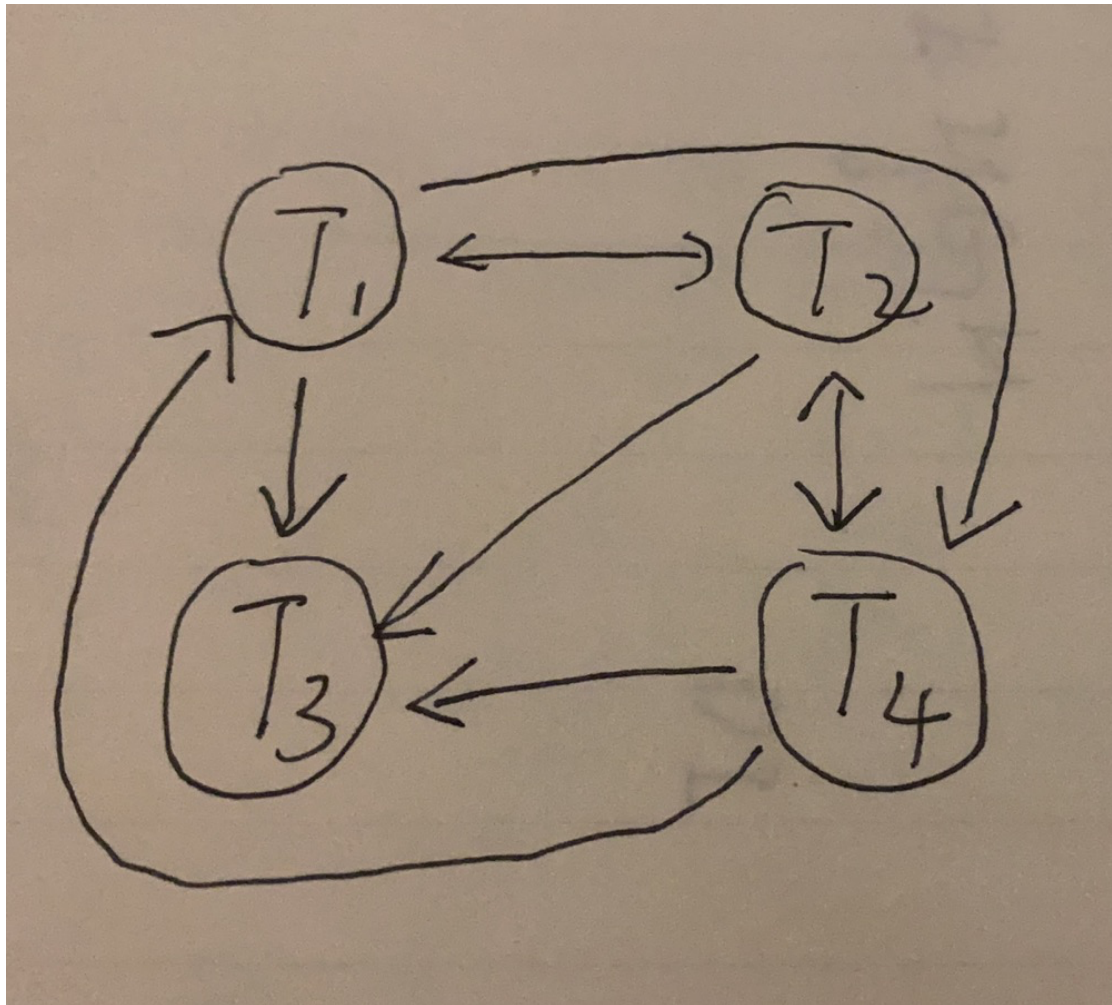
## Question 2

1)

T1 undo T2 redo

T3, T4 do nothing

2)



3)

time	T1	T2	T3	T4
1	R(B)			
2	R(A)			
3	W(B)			
4	W(A)			
5		R(B)		
6		R(A)		
7		W(B)		
8		R(A)		
9		W(A)		
10				R(A)
11				W(A)
12				R(B)
13				W(B)
14		R(B)		
15		W(B)		

4)

There is no way to construct a schedule whose wait-for graph contains cycles.

We have T1 and T2 read and write on B, we have potential to make T2 wait-for T1. We have T1 and T2 read and write on A, we have potential to make T1 wait-for T2. we cannot make both T2 wait-for T1 and T1 wait-for T2.

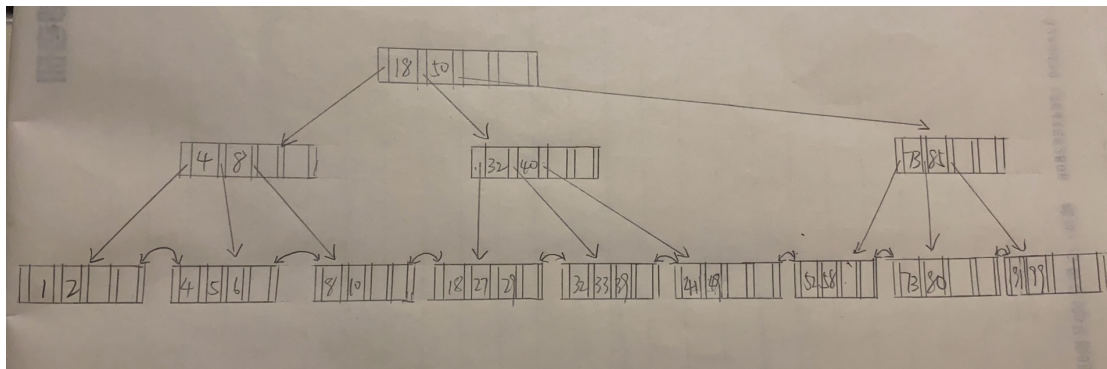
### Question 3

1)

80

Because  $(5 * 5 * 4) - 20 = 80$

2)



3)

