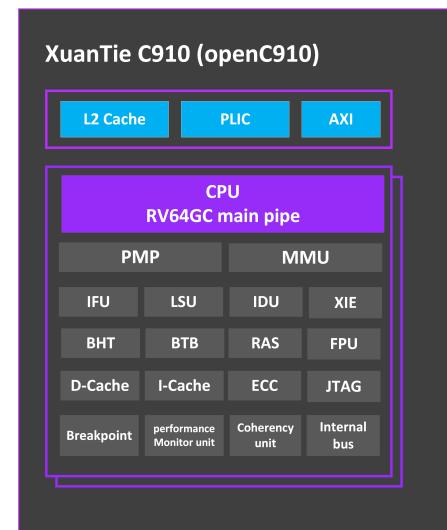




XuanTie C910 Datasheet

Overview

C910 is a RISC-V compatible 64-bit high performance processor developed by XuanTie Semiconductor Co., Ltd. It delivers industry-leading performance in control flow, computing and frequency through architecture and micro-architecture innovations. The C910 processor is based on the RV64GC instruction set and implements the XIE (XuanTie Instruction Extension) technology. C910 adopts a state of the art 12 stages out-of-order multiple issue superscalar pipeline with high frequency, IPC, and power efficiency. C910 supports hardware cache coherency. Each cluster contains 2 cores. The C910 supports the AXI4 bus interface. The C910 uses the Sv39 virtual address system with XMAE (XuanTie Memory Attributes Extension) technology. In addition, C910 includes the standard CLINT and PLIC interrupt controllers and performance monitors. C910 implements the XuanTie debug protocol.



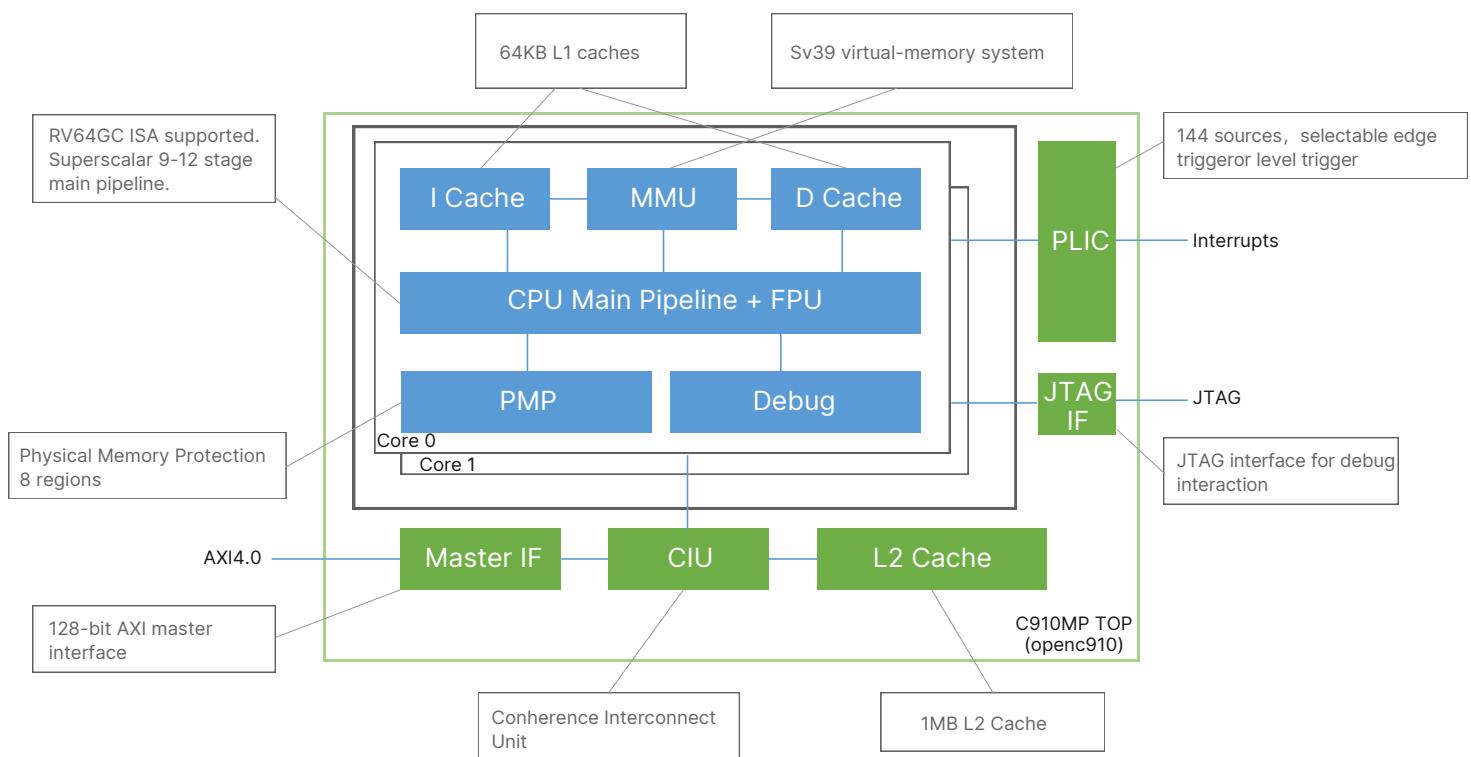
Features

Feature	Description
Architecture	RV64GC
SMP	2 cores in each cluster
Pipeline	12 stages (Integer)
Floating-point Unit	Support RISC-V F, D instruction extension Support IEEE 754-2008 standard
Bus interface	AXI4-128 master
Instruction Cache	64KB
Data Cache	64KB

Feature	Description
L2 Cache	1MB Supporting parallel access with multi-bank
XuanTie extensions	XuanTie Instruction Extension (XIE) XuanTie Memory Attributes Extension (XMAE)
Memory Management Unit (MMU)	Sv39 virtual memory translation
PMP	8 regions
Interrupt Controller	Platform-Level Interrupt Controller (PLIC) for supporting wide range of system event scenarios

XuanTie C910 Components

- Processor Overview



- Multi-Core

- ◇ Support dual core homogeneous multi-core system.
- ◇ MOESI coherency protocol.
- ◇ 2-way centralized snoop buffer.
- ◇ Exclusive memory access instructions.
- ◇ Integrates multi-core interrupt controllers, timers, and debuggers.

• Memory sub-system

The C910 has 64K instruction and data caches with cache coherency support. Hardware cache coherency ensures the consistency of all caches efficiently. The shared L2 Cache is 1MB. Software and hardware collaborative optimization of data consistency between TLB, I-Cache and D-Cache.

◇ The L1 instruction memory system has the following key features:

- VIPT, two-way set-associative instruction cache.
- Fixed cache line length of 64 bytes.
- 128-bit read interface from the L2 memory system.

◇ The L1 data memory system has the following features:

- PIPT, two-way set associative L1 data cache.
- Fixed cache line length of 64 bytes.
- 128-bit read interface from the L2 memory system.
- Up to 128-bit read data paths from the data L1 memory system to the data path.
- Up to 128-bit write data path from the data path to the L1 memory system.

◇ The L2 Cache has the following features:

- Size: 1MB.
- PIPT, 16-way set-associative structure.
- Fixed line length of 64 bytes.
- Support data prefetch.

• Memory Management Unit (MMU)

- ◇ Sv39 virtual memory systems supported.
- ◇ 32/17-entry fully associative I-uTLB/D-uTLB.
- ◇ 1024-entry 4-way set-associative shared TLB.
- ◇ Hardware page table walker.
- ◇ Virtual memory support for full address space and easy hardware for fast address translation.
- ◇ Code/data sharing.
- ◇ Support for full-featured OS such as Linux.
- ◇ XMAE (XuanTie Memory Attributes Extension) technology extends page table entries for additional attributes.

• Physical Memory Protection (PMP)

8 regions basic read/write/execute memory protection with low cost.

• Performance Monitor Unit (PMU)

Program code performance tuning.

• Platform-Level Interrupt Controller (PLIC)

- ◊ Support multi-core interrupt control.
- ◊ 144 PLIC interrupt sources.
- ◊ Selectable edge trigger or level trigger.

• JTAG Debug

- ◊ Support XuanTie multi-core debug.
- ◊ JTAG debug interface support several triggers.
- ◊ Support software breakpoints.
- ◊ Check and modify CPU register resource.
- ◊ Single step or multi step flexibly supported.
- ◊ High speed program download through JTAG.

• Float Point Unit (FPU)

- ◊ RISC-V F and D extensions.
- ◊ Support half/single/double precision.
- ◊ Fully IEEE-754 compliant.
- ◊ Does not generate floating-point exceptions.
- ◊ User configurable rounding modes.

• Branch Predictor

- ◊ Branch Target Buffer (BTB) and Branch History Table (BHT) to speed up control codes.
- ◊ Return Address Stack (RAS) to speed up procedure returns.
- ◊ Loop buffer to speed up short loops.

• RV Compatibility with Custom Extensions

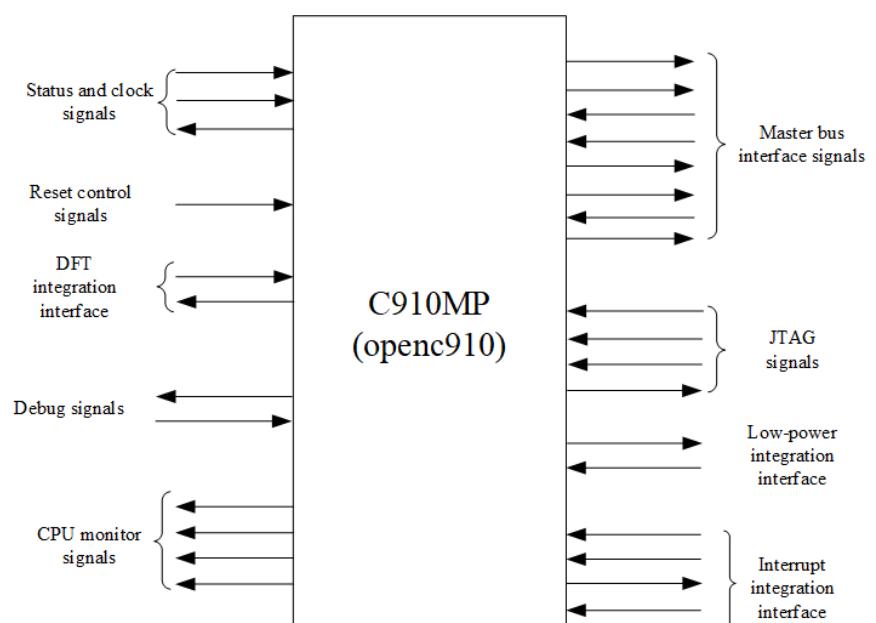
The C910 is fully compatible with the RV64GC instruction set and supports the standard M/S/U privilege program model. The C910 includes a standard 8 region PMP and Sv39 MMU, which is fully compatible with RISC-V Linux. The C910 includes standard CLINT and PLIC interrupt controllers and RV compatible PMU.

• RV Compatibility

Component	RV version
ISA	RV64GC
Privilege	1.10
MMU	Sv39
Interrupt controller	CLINT/PLIC

• Interfaces

- ◊ Master AXI (M-AXI)
- ◊ Debug (JTAG)
- ◊ Interrupts
- ◊ Low power control



Configurations

Config	Options
Core Number	2
L1 D-Cache Size	64K
L1 I-Cache Size	64K
L2-Cache Size	1M

Software Ecosystems

- ◇ Compiler, assembler, linker, debugger and binary tools are contributed to GNU and supported officially.
- ◇ Linux kernel is contributed to Linux foundation and supported officially.
- ◇ QEMU is contributed and supported officially.
- ◇ Chip Integrated Development Environment (CDS), compatible with Eclipse development.
- ◇ Graphical profiling and statistical analysis tools base trace data (simulator) .
- ◇ ICE, CK-Link Pro, high speed (1.1Mbytes/s) JTAG debug hardware.
- ◇ Multi-core JTAG online debug.

Linux System

C910 supports the official RISC-V Linux and its software applications ecosystem. e.g.: GNU toolchain, Fedora, Debian, buildroot and thousands of open source software based on Linux. XuanTie will continue to make contributions to RISC-V architecture porting in Linux ecosystem.