

第七章 死锁

1. Consider the traffic deadlock depicted in Figure 1.

- Show that the four necessary conditions for deadlock indeed hold in this example.
- State a simple rule for avoiding deadlocks in this system.

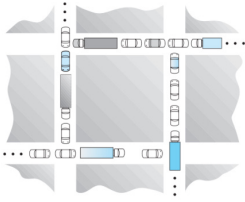


Figure 1. Traffic deadlock

a.

- ① 互斥：路口的每个区域只能被一辆车占用。
- ② 持有并等待：每个车都已经“持有”了路口的一块区域，同时“等待”下一个区域。
- ③ 不可抢占：没有外力可以把车从占据的路口区域“挪出”，只能自己离开。
- ④ 循环等待：车辆形成了循环的等待链，A车等待B车的区域，B等C，C等D，D等A，互相等待对方的资源。

b.

打破持有并等待

车辆在进入交叉路口，必须能一次性通过整个路口。

即执行前申请所有资源

2. Consider the following snapshot of a system:

	<u>Allocation</u>	<u>Max</u>	<u>Available</u>
	A B C D	A B C D	A B C D
P ₀	0012	0012	1520
P ₁	1000	1750	
P ₂	1354	2356	
P ₃	0632	0652	
P ₄	0014	0656	

Answer the following questions using the banker's algorithm:

- What is the content of the matrix Need?
- Is the system in a safe state?
- If a request from process P₁ arrives for (0,4,2,0), can the request be granted immediately?

a. $need_0 = [0, 0, 0, 0]$

$$need_1 = [0, 7, 5, 0]$$

$$need_2 = [1, 0, 0, 2]$$

$$need_3 = [0, 0, 2, 0]$$

$$need_4 = [0, 6, 4, 2]$$

b. ① $need_0 \leq available$

$$P_0 \text{ 释放, } available = [1, 5, 2, 0] + [0, 0, 1, 2] = [1, 5, 3, 2]$$

② $need_2 \leq available$

$$P_2 \text{ 释放, } available = [1, 5, 3, 2] + [1, 3, 5, 4] = [2, 8, 8, 6]$$

③ $need_3 \leq available$

$$P_3 \text{ 释放, } available = [2, 8, 8, 6] + [0, 6, 3, 2] = [2, 14, 11, 8]$$

④ $need_4 \leq available$

$$P_4 \text{ 释放, } available = [2, 14, 11, 8] + [0, 0, 1, 4] = [2, 14, 12, 12]$$

⑤ $need_1 \leq available$

$$P_1 \text{ 释放, } available = [2, 14, 12, 12] + [1, 0, 0, 0] = [3, 14, 12, 12]$$

系统处于安全状态 序列为 $P_0 \rightarrow P_2 \rightarrow P_3 \rightarrow P_4 \rightarrow P_1$

c.

$$R_1 \leq need_1 \quad [0, 4, 2, 0] < [0, 7, 5, 0]$$

$$R_1 \leq Available \quad [0, 4, 2, 0] < [1, 5, 2, 0]$$

$$Allocation_1 = [1, 0, 0, 0] + [0, 4, 2, 0] = [1, 4, 2, 0]$$

$$need_1 = [0, 7, 5, 0] - [0, 4, 2, 0] = [0, 3, 3, 0]$$

$$Available = [1, 5, 2, 0] - [0, 4, 2, 0] = [1, 1, 0, 0]$$

$$\text{need}_0 < \text{Available} \Rightarrow [0, 0, 0, 6] < [1, 1, 0, 0]$$

$$\text{Available} = [1, 1, 0, 0] + [0, 0, 1, 2] = [1, 1, 1, 2]$$

$$\text{need}_2 < \text{Available} \Rightarrow [1, 0, 0, 2] < [1, 1, 1, 2]$$

$$\text{Available} = [1, 1, 1, 2] + [1, 3, 5, 4] = [2, 4, 6, 6]$$

$$\text{need}_1 < \text{Available} \Rightarrow [0, 3, 3, 0] < [2, 4, 6, 6]$$

$$\text{Available} = [2, 4, 6, 6] + [1, 4, 2, 0] = [3, 8, 8, 6]$$

$$\text{need}_3 < \text{Available} \quad [0, 0, 2, 0] < [3, 8, 8, 6]$$

$$\text{Available} = [3, 8, 8, 6] + [0, 6, 3, 2] = [3, 14, 11, 8]$$

$$\text{need}_4 < \text{Available} \quad [0, 6, 4, 2] < [3, 14, 11, 8]$$

$$\text{Available} = [3, 14, 11, 8] + [0, 0, 1, 4] = [3, 14, 12, 12]$$

是安全的序列为 $P_0 \rightarrow P_2 \rightarrow P_1 \rightarrow P_3 \rightarrow P_4$

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