Algorithms and Data Structures Part 4

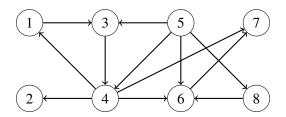
Lecture 7: Depth-First Search

George Mertzios

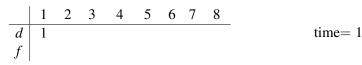
george.mertzios@durham.ac.uk

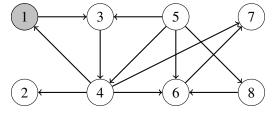
Depth-first search

- Like BFS, Depth-first search explores the graph (but does not find distances to the source).
- In contrast to BFS, when a vertex is discovered it is immediately explored.
- Two timestamps are recorded for each vertex, *d* and *f*; the discovery and finish times. We can also record predecessors again.
- Again colours are used:
 - white for undiscovered,
 - grey for discovered but not finished, and
 - black for finished.

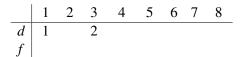


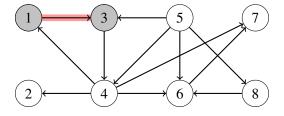
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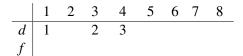


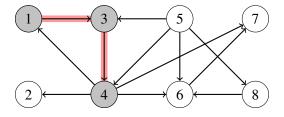
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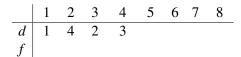


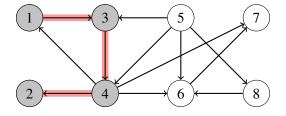
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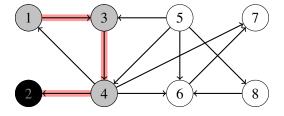
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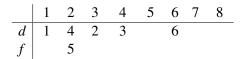


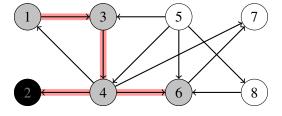
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	1	2	3	4	5	6	7	8
\overline{d}	1	4	2	3				
f		5						



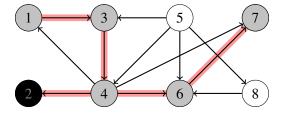
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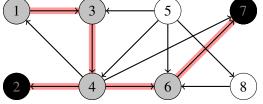
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	_	_	_	4	_	6	7	8
\overline{d}	1	4	2	3		6	7	
f		5						



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	1	2	3	4	5	6	7	8			
				3		6	7		-	ti	me = 8
f		5					8				



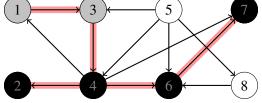
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	1	2	3	4	5	6	7	8	
\overline{d}	1	4	2	3		6	7		time=
f		5				9	8		
			1			3) —		5 7

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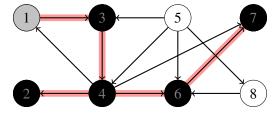
9

	1	2	3	4	5	6	7	8		
\overline{d}	1	4	2	3		6	7		-	time = 10
f		5		10		9	8			
			(1			$\sqrt{2}$	\ .		(5)	



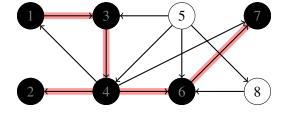
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	1	2	3	4	5	6	7	8
\overline{d}	1	4	2	3		6	7	
f		5	11	10		9	8	



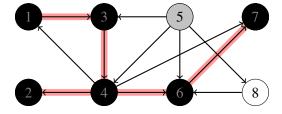
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				4	6	7	8	
\overline{d}	1	4	2	3	6	7		
f	12	5	11	3 10	9	8		



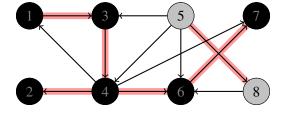
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				4				8
\overline{d}	1	4	2	3	13	6	7	
f	12	5	11	3 10		9	8	



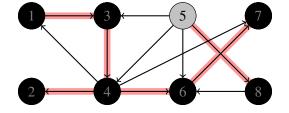
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				4				
\overline{d}	1	4	2	3	13	6	7	14
f	12	5	11	3 10		9	8	



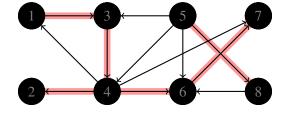
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				4				
\overline{d}	1	4	2	3 10	13	6	7	14
f	12	5	11	10		9	8	15



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\overline{d}	1	4	2	3	13	6	7	14
f	12	5	11	10	16	9	8	15



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Depth-first search

```
DFS (G)
1 for each vertex u \in V[G]
       do colour[u] \leftarrow WHITE
           \pi[u] \leftarrow \text{NIL}
4 time \leftarrow 0
5 for each vertex u \in V[G]
       do if colour[u] = WHITE
6
           then DFS-VISIT(u)
DFS-VISIT(u)
1 \operatorname{colour}[u] \leftarrow \operatorname{GREY}
                                              [vertex u has just been discovered]
2. time \leftarrow time + 1
3 d[u] \leftarrow time
4 for each vertex v \in Adj[u]
                                                                [explore edge (u, v)]
       do if colour[v] = WHITE
           then \pi[v] \leftarrow u
6
                DFS-VISIT(v)
8 \operatorname{colour}[u] \leftarrow \operatorname{BLACK}
                                                             [u has been processed]
9f[u] \leftarrow time \leftarrow time + 1
```

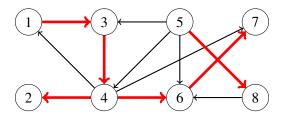
Analysis

- Initialization takes time O(V).
- Time O(V) is spent on incrementing time, colouring vertices and updating d and f.
- Each vertex in each adjacency list is considered at most once. This takes time O(E).
- Total time is O(V + E).

Analysis

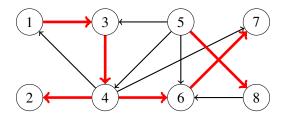
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- Total time is O(V + E).

The edges used for discovering new vertices form the depth-first tree (or forest). Again, we can find this with a predecessor array.



Once we have run DFS on a graph we can construct the predecessor subgraph. This has the same vertex set as the graph, and for each vertex v there is an edge from the predecessor of v to v.

The predecessor subgraph is a depth-first forest.



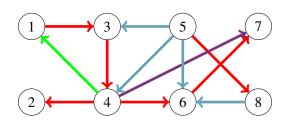
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How many trees do we have in the above forest?

Once we have obtained a DFS-forest for a graph G, we can classify the edges of G.

- Tree edges are those edges in the DFS-forest.
- Back edges are edges that join a vertex to an ancestor.
- Forward edges are edges not in the tree that join a vertex to its descendant.
- Cross edges: all other edges.



Tree edges -----

Forward edges -----

Back edges ———

Cross edges -----

The classification is ambiguous for undirected graphs (back edges and forward edges are the same thing),

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Let us refine the definition: suppose that e is an edge that joins a vertex u to its descendant v.

- \blacksquare e is a forward edge if DFS first considers e from u.
- \blacksquare *e* is a back edge if DFS first considers *e* from *v*.

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Theorem

In an undirected graph, every edge is a tree edge or a back edge.

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Theorem

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Questions: Which types of edges can be found in a directed graph? What if BFS is used?