Swarup Chandra, Vishal Karande, Zhiqiang Lin, Latifur Khan, Murat Kantarcioglu and Bhavani Thuraisingham

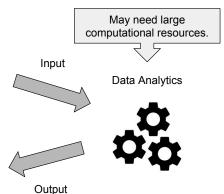
The University of Texas at Dallas

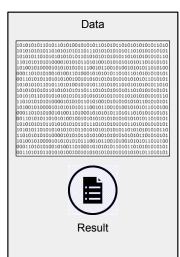
ESORICS 2017

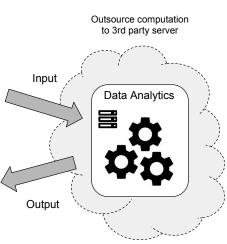
Data

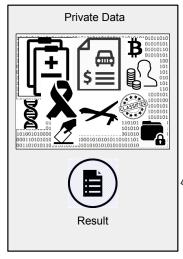


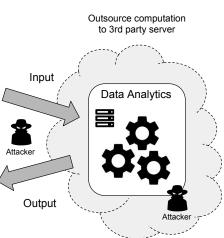
Result

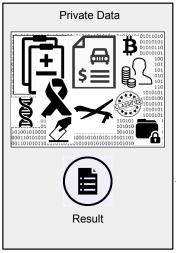


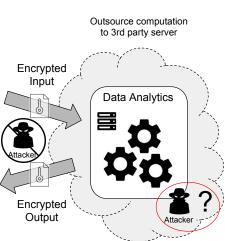












Evaluation

Existing Efforts

Privacy Preserving Analytics

Modifications to data before sharing it to the third-party server [AP08].

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Homomorphic Encryption

Using fully homomorphic encryption scheme (not practical) [LWN+15].

Privacy Preserving Analytics

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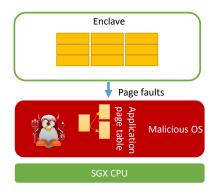
Homomorphic Encryption

Using fully **homomorphic encryption** scheme (not practical) [LWN⁺15].

Leveraging Trusted Execution Environment (TEE)

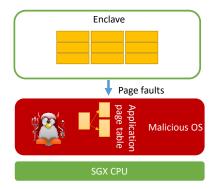
- Recent TEEs: Intel SGX [MAB+13], ARM TrustZone.
- Perform analytics on private data within secure environment, isolated from adversary [SCF⁺15].

Issues w/ TEE Approaches
Information leakage via (controlled) side-channels harms data privacy



Issues w/ TEE Approaches

Information leakage via (controlled) side-channels harms data privacy



Attacker can use page-fault attack to observe execution flow and guess parameter values [XCP15].



State-of-the-Art

- **Balanced Execution** [SCNS16]: Execute each conditional branch.
- Data oblivious execution [OSF+16]: Manipulate execution sequence to be fully data independent.

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Execution overhead impractical for large analytics.

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State-of-the-Art

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Evaluation

 Data oblivious execution [OSF+16]: Manipulate execution sequence to be fully data independent.

Execution overhead impractical for large analytics.

 When executing every path for each input, this creates bottleneck due to unnecessary path execution especially when large parameters are involved.

What are large parameters?

- Decision Tree: Large number of nodes
- Naive Bayes: Large domain size
- K-Means: Large domain size

Limtations in State-of-the-Art **Example Decision Tree**

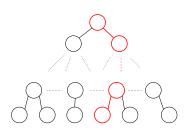
 Given a trained decision tree.

Limitations in State-of-the-Art Example Decision Tree

- Given a trained decision tree.
- **Input**: Data instance *X* for evaluation.

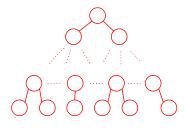
Limitations in State-of-the-Art Example Decision Tree

- Given a trained decision tree.
- Input: Data instance X for evaluation.
- Defense: Hide structure using dummy nodes.



Limitations in State-of-the-Art Example Decision Tree

- Given a trained decision tree.
- Input: Data instance X for evaluation.
- Defense: Hide structure using dummy nodes.
- But, evaluation of X will explore all branches causing computational bottleneck.



Intel SGX

Intel Software Guard Extension (SGX)

Capabilities

- Offering secure enclave
- Confidentiality, Integrity



Evaluation

Intel SGX

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- New extensions to Intel x86/64 architecture
- Applications keep secret data and code inside enclave
- Minimum attack surface

0000000

Intel Software Guard Extension (SGX)

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Assumptions

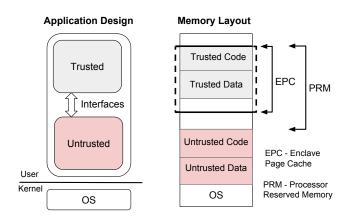
All privileged software (kernel, hypervisor) is malicious.

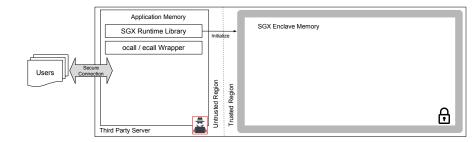


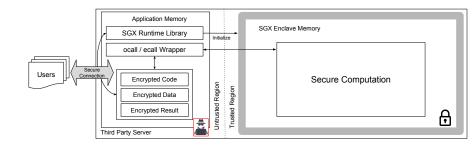
Intel SGX

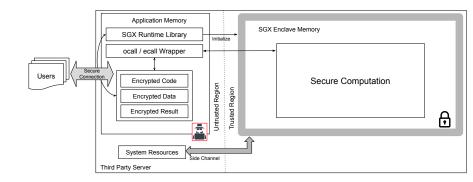
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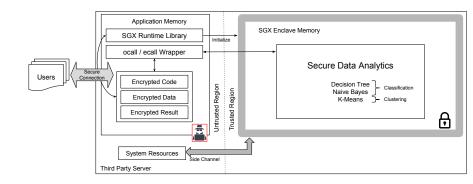
Intel Software Guard Extension (SGX)











Objective: reduce leakage of private information during execution of analytic algorithms

What should we secure?

Model and data specific to each analytics

- Trained parameters such as posterior probability in Naive Bayes or internal node values in decision tree.
- Learned structure such as decision tree.

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- Learned structure such as decision tree.

What need not be secure?

Public information in Analytics.

- Number of data instances.
- Total number of class labels.
- User-defined parameters such as number of clusters in K-Means clustering.

How to reduce computational cost?

- Create **dummy data instances**, equivalent to input data.
- Randomly mix dummy data instances with input test instances to form a contaminated test dataset.
- Evaluate with the newly generated dataset
- Obliviously ignore results from dummy data instances.

SGX-Rand Key Idea

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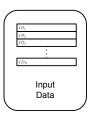
Attacker's perspective

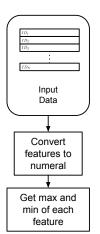
- Obtain traces from side-channel during evaluation.
- Cannot distinguish traces corresponding to real vs. dummy data instances.

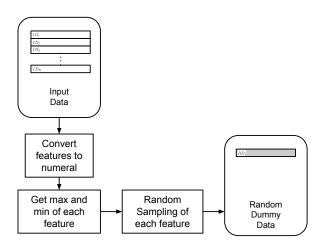


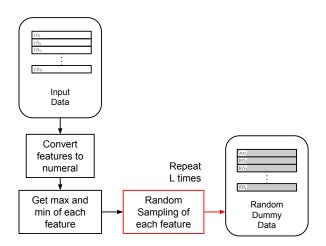
Background

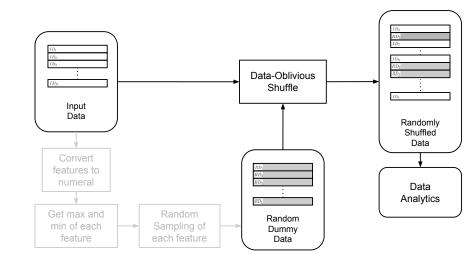
- Dummy data generation
- Data oblivious methods
 - shuffling dummy data with input data to create contaminated dataset.
 - ignoring results of dummy data.











SGX-Rand: **Data Oblivious Comparison** Primitives

- Example function foo.
- Call function bar on different variables depending on condition.

```
int foo(int x, int y)
{
    ...
    if (x > y) {
         bar(x);
    } else {
        bar(y);
    }
    ...
}
```

SGX-Rand: Data Oblivious Comparison Primitives

- Example function foo
- Call function bar on different variables depending on condition
- Instead, we define a temporary variable (t) to capture condition
- Use it to pass appropriate value to bar, making the function data oblivious

```
int foo(int x, int y)
  int t:
  if (x > y) {
        t=1:
  } else {
        t = 0:
  bar(t*x + (1-t)*y);
```

SGX-Rand: Secure Data Analytics

How to Use SGX-Rand

- Train model using confidential training data.
- Provide input to third-party server
 - Trained model
 - Unknown data set for evaluation.
- Obtain output from third-party server.
- Discard dummy output.

Evaluation

SGX-Rand: Secure Data Analytics

How to Use SGX-Rand

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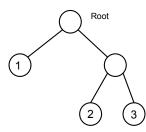
Securing popular analytics during evaluation

- Classification
 - Decision Tree
 - Naive Bayes
- Clustering
 - K-Means

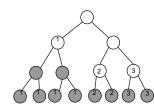


Private Information	Public Information		
Tree Structure	Number of input data instances		
Node values	Number of class labels.		
Data distribution	Number of class labels.		

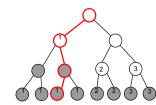
• Train a decision tree offline.



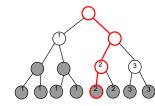
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- Load it into SGX enclave.



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- Within enclave, tree evaluated on contaminated dataset.
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- Train a decision tree offline.
- Balance the tree with dummy nodes such that leaf has same label as real-tree parent leaf. (extra step to secure tree structure).
- Load it into SGX enclave.
- Within enclave, tree evaluated on contaminated dataset.
- Obliviously ignore results from dummy data instances.
- Attacker cannot differentiate between 2 traces.



Empirical Evaluation

Objectives

- At minimum privacy, how much reduction in execution overhead can randomization achieve?
- How much is the trade-off between computation time and privacy-guarantee?

Empirical Evaluation

Objectives

- At minimum privacy, how much reduction in execution overhead can randomization achieve?
- How much is the trade-off between computation time and privacy-guarantee?

Dataset

Dataset	Size	Feature	Classes
Arrhythmia (A)	452	280	13
Defaulter (D)	30,000	24	2
ForestCover (F)	50,000	55	7
Synthetic (S)	50,000	71	7

Experiment Setup

Baseline

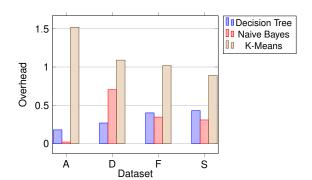
- **SGX**: Analytics within SGX enclave, *without* any defense.
- SGX+Obliv: Analytics within SGX enclave, with fully data-oblivious defense.
- SGX+Rand: Analytics within SGX enclave, with our randomization defense.

Measurement

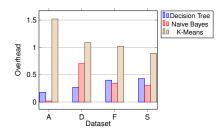
- Execution time $e = \frac{time(SGX + x)}{time(SGX)}$, where $x \in \{Obliv, Rand\}$
- Classification accuracy

Results

- With minimum privacy, L = N.
- Overhead $o = \frac{e(SGX + Rand)}{e(SGX + Obliv)}$
- Lower o is better, with o < 1 is gain in computation time.



Results



Good

Decision Tree and **Naive Bayes** have significantly large gain in execution time.

Not Good

K-Means for SGX+Rand has larger overhead than SGX+Obliv since every cluster is accessed even for dummy data instances.

Objective

- Is there noise in access traces?
- Are traces of input and dummy data instances indistinguishable?

Security Evaluation

Objective

- Is there noise in access traces?
- Are traces of input and dummy data instances indistinguishable?

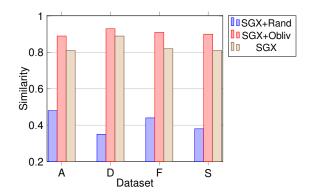
Methodology

- Using Intel PIN to obtain traces during execution.
- Measure Levenshtein distance between memory access sequence.
 - Greater distance is dissimilar traces, i.e., less similarity.
 - More similarity is similar traces.



Security Evaluation

Measuring similarity of execution traces obtained from two independent datasets



Smaller value indicates greater randomization.



Conclusion

- Randomization is effective in reducing execution time overhead in data analytics compared to fully data oblivious solution.
- We quantitatively measure privacy and provide a choice in trade-off between privacy and efficiency.
- Our evaluation demonstrates both advantages and disadvantages in employing randomization.

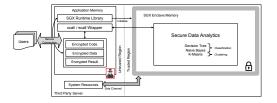
Open-Source

Project code available at https:

//github.com/utds3lab/secure-analytics-sqx



Thank You





Thanks also to the AFOSR, NSA, and NSF for their sponsorship of this research.

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Background

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