

Øystein Grøvlen Senior Principal Software Engineer MySQL Optimizer Team, Oracle February 25, 2015





# Program Agenda

- Introduction to MySQL cost-based optimizer
- Selecting data access method
- Join optimizer
- 4 Sorting
- Tools for monitoring, analyzing, and tuning queries
- Influencing the Optimizer





# Program Agenda

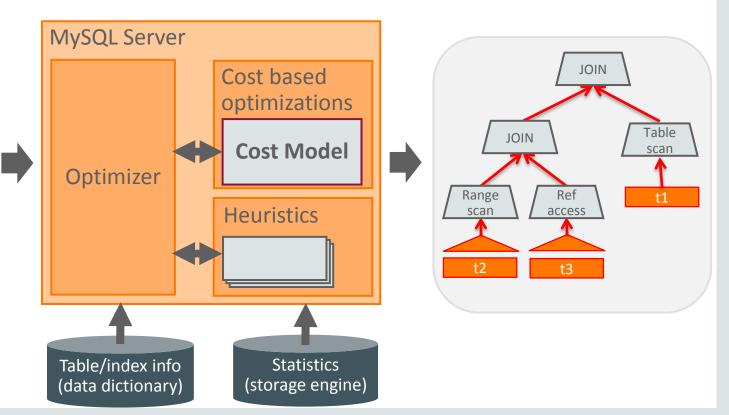
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# MySQL Optimizer



SELECT a, b
FROM t1, t2, t3
WHERE t1.a = t2.b
AND t2.b = t3.c
AND t2.d > 20
AND t2.d < 30;





# Cost-based Query Optimization General idea

- Assign cost to operations
- Computes cost of partial or alternative plans
- Search for plan with lowest cost

Cost-based optimizations:

JOIN Table JOIN scan Range Ref scan access

Access method

Join order

Subquery strategy



# Input to Cost Model

#### IO-cost:

- Estimates from storage engine based on number of pages to read
- Both index and data pages

#### Schema:

- Length of records and keys
- Uniqueness for indexes
- Nullability

#### • Statistics:

- Number of rows in table
- Key distribution/Cardinality:
  - Average number of records per key value
  - Only for indexed columns
  - Maintained by storage engine
- Number of records in an index range



# Cost Model Example

SELECT SUM(o\_totalprice) FROM orders
WHERE o\_orderdate BETWEEN '1994-01-01' AND '1994-12-31';

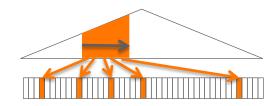
#### Table scan:

- IO-cost: #pages in table
- CPU cost: #rows \* ROW\_EVALUATE\_COST

#### Range scan (on secondary index):

- IO-cost: #pages to read from index + #rows\_in\_range
- CPU cost: #rows\_in\_range \* ROW\_EVALUATE\_COST







#### Cost Model

#### **Example**

EXPLAIN SELECT SUM(o\_totalprice) FROM orders
WHERE o\_orderdate BETWEEN '1994-01-01' AND '1994-12-31';

id	select type	table	type	possible keys	key	key len	ref	rows	extra
1	SIMPLE	orders	ALL	i_o_orderdate	NULL	NULL	NULL	15000000	Using where

# EXPLAIN SELECT SUM(o\_totalprice) FROM orders WHERE o\_orderdate BETWEEN '1994-01-01' AND '1994-06-30';

Id	select type	table	type	possible keys	key	key len	ref	rows	extra
1	SIMPLE	orders	range	i_o_orderdate	i_o_orderdate	4	NULL	2235118	Using index condition





### Cost Model Example: Optimizer Trace

#### join\_optimization / row\_estimation / table : orders / range\_analysis

```
"table_scan": {
 "rows": 15000000,
 "cost": 3.12e6
}/* table scan */,
"potential range indices": [
  "index": "PRIMARY",
  "usable": false,
  "cause": "not applicable"
  "index": "i o orderdate",
  "usable": true.
  "key parts": ["o orderDATE", "o orderkey"]
]/* potential range indices */,
```

```
"analyzing range alternatives": {
 "range scan alternatives": [
   "index": "i o orderdate",
   "ranges": [ "1994-01-01 <= o orderDATE <= 1994-12-31"
   "index dives for eq ranges": true,
   "rowid ordered": false,
   "using mrr": false,
   "index only": false,
   "rows": 4489990,
   "cost": 5.39e6.
   "chosen": false.
   "cause": "cost"
 /* range scan alternatives */,
}/* analyzing range alternatives */
```



#### Cost Model vs Real World

#### **Measured Execution Times**

	Data in Memory	Data on Disk	Data on SSD
Table scan	6.8 seconds	36 seconds	15 seconds
Index scan	5.2 seconds	2.5 hours	30 minutes

#### Force Index Scan:

**SELECT SUM(o\_totalprice)** 

FROM orders FORCE INDEX (i\_o\_orderdate)

WHERE o\_orderdate BETWEEN '1994-01-01' AND '1994-12-31';





# Performance Schema Disk I/O

SELECT event\_name, count\_read, avg\_timer\_read/1000000000.0 "Avg Read Time (ms)", sum\_number\_of\_bytes\_read "Bytes Read"
FROM performance\_schema.file\_summary\_by\_event\_name
WHERE event\_name='wait/io/file/innodb/innodb\_data\_file';

#### Table Scan

event_name	count_read	Avg Read Time (ms)	Bytes Read
wait/io/file/innodb/innodb_data_file	115769	0.0342	1896759296

#### **Index Scan**

event_name	count_read	Avg Read Time (ms)	Bytes Read
wait/io/file/innodb/innodb_data_file	2188853	4.2094	35862167552



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### Selecting Access Method

#### Finding the optimal method to read data from storage engine

- For each table, find the best access method:
  - Check if the access method is useful
  - Estimate cost of using access method
  - Select the cheapest to be used
- Choice of access method is cost based

#### Main access methods:

- Table scan
- Index scan
- Ref access
- Range scan
- Index merge
- Loose index scan





#### **Ref Access**

#### **Single Table Queries**

#### **EXPLAIN SELECT \* FROM customer WHERE c\_custkey = 570887;**

id	select type	table	type	possible keys	key	key len	ref	rows	extra
1	SIMPLE	customer	const	PRIMARY	PRIMARY	4	const	1	NULL

#### **EXPLAIN SELECT \* FROM orders WHERE o\_orderdate = '1992-09-12';**

id	select type	table	type	possible keys	key	key len	ref	rows	extra
1	SIMPLE	orders	ref	i_o_orderdate	i_o_orderdate	4	const	6271	NULL





# Ref Access

**Join Queries** 

EXPLAIN SELECT \*
FROM orders JOIN customer ON c\_custkey = o\_custkey
WHERE o\_orderdate = '1992-09-12';

ld	select type	table	type	possible keys	key	key len	ref	rows	extra
1	SIMPLE	orders	ref	i_o_orderdate, i_o_custkey	i_o_orderdate	4	const	6271	Using where
1	SIMPLE	customer	eq_ref	PRIMARY	PRIMARY	4	dbt3.orders. o_custkey	1	NULL



### **Ref Access**

Join Queries, continued

EXPLAIN SELECT \*
FROM orders JOIN customer ON c\_custkey = o\_custkey
WHERE c\_acctbal < -1000;

Id	select type	table	type	possible keys	key	key len	ref	rows	extra
1	SIMPLE	customer	ALL	PRIMARY	NULL	NULL	NULL	1500000	Using where
1	SIMPLE	orders	ref	i_o_custkey	i_o_custkey	5	dbt3.customer. c_custkey	7	NULL





### Range Optimizer

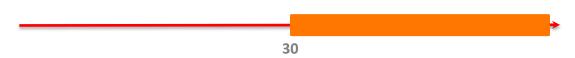
- Goal: find the "minimal" ranges for each index that needs to be read
- Example:

SELECT \* FROM t1 WHERE (key1 > 10 AND key1 < 20) AND key2 > 30

Range scan using INDEX(key1):



Range scan using INDEX(key2):



# Range Optimizer Optimizer Trace show ranges

SELECT a, b FROM t1

WHERE a > 10

AND a < 25

AND a NOT IN (11, 19))

AND (b < 5 OR b > 10);

```
"analyzing range alternatives": {
  "range scan alternatives": [
      "index": "i a",
      "ranges": [
        "10 < a < 11",
        "11 < a < 19",
        "19 < a < 25"
      "index dives for eq ranges": true,
      "rowid ordered": false,
      "using mrr": false,
      "index only": false,
      "rows": 3,
      "cost": 6.61,
      "chosen": true
      "index": "i b",
      "ranges": [
        "NULL < b < 5",
        "10 < b"
      "index dives for eq ranges": true,
      "rowid ordered": false,
```



Why table scan?

SELECT \* FROM orders
WHERE YEAR(o\_orderdate) = 1997 AND MONTH(o\_orderdate) = 5
AND o\_clerk LIKE '%01866';

id	select type	table	type	possible keys	key	key len	ref	rows	extra
1	SIMPLE	orders	ALL	NULL	NULL	NULL	NULL	15000000	Using where

#### Index not considered

mysql> SELECT \* FROM orders WHERE year(o\_orderdate) = 1997 AND MONTH(...
...
15 rows in set (8.91 sec)



Rewrite query to avoid functions on indexed columns

SELECT \* FROM orders
WHERE o\_orderdate BETWEEN '1997-05-01' AND '1997-05-31'
AND o\_clerk LIKE '%01866';

id	select type	table	type	possible keys	key	key len	ref	rows	extra
1	SIMPLE	orders	range	i_o_orderdate	i_o_orderdate	4	NULL	376352	Using index condition; Using where

```
mysql> SELECT * FROM orders WHERE o_orderdate BETWEEN '1997-05-01' AND ...
15 rows in set (0.91 sec)
```



Adding another index

CREATE INDEX i\_o\_clerk ON orders(o\_clerk);

SELECT \* FROM orders
WHERE o\_orderdate BETWEEN '1997-05-01' AND '1997-05-31'
AND o\_clerk LIKE '%01866';

id	select type	table	type	possible keys	key	key len	ref	rows	extra
1	SIMPLE	orders	range	i_o_orderdate	i_o_orderdate	4	NULL	376352	Using index condition; Using where

New index not considered





Rewrite query, again

SELECT \* FROM orders
WHERE o\_orderdate BETWEEN '1997-05-01' AND '1997-05-31'
AND o\_clerk = 'Clerk#000001866';

i	d	select type	table	type	possible keys	key	key len	ref	rows	extra
	1	SIMPLE	orders	range	i_o_orderdate, i_o_clerk	i_o_clerk	16	NULL	1504	Using index condition; Using where

```
mysql> SELECT * FROM orders WHERE o_orderdate BETWEEN '1997-05-01' AND ...
15 rows in set (0.01 sec)
```





# Range Access for Multi-part Index

#### **Example table with multi-part index**

• Table:



INDEX idx(a, b, c);

Logical storage layout of index:





Create multi-column index

**CREATE INDEX i\_o\_clerk\_date ON orders(o\_clerk, o\_orderdate)**;

SELECT \* FROM orders
WHERE o\_orderdate BETWEEN '1997-05-01' AND '1997-05-31'
AND o\_clerk = 'Clerk#000001866';

id	select type	table	type	possible keys	key	key len	ref	row s	extra
1	SIMPLE	orders	range	i_o_orderdate, i_o_clerk, i_o_clerk_date	i_o_clerk_date	20	NULL	14	Using index condition

mysql> SELECT \* FROM orders WHERE o\_orderdate BETWEEN '1997-05-01' AND  $\dots$ 

. . .

15 rows in set (0.00 sec)





# Performance Schema: Query History

# UPDATE performance\_schema.setup\_consumers SET enabled='YES' WHERE name = 'events\_statements\_history';





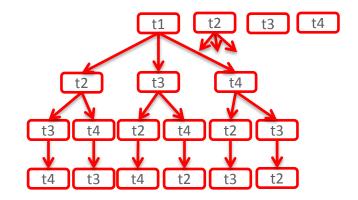
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# Join Optimizer

#### "Greedy search strategy"

- Goal: Given a JOIN of N tables, find the best JOIN ordering
- Strategy:
  - Start with all 1-table plans
  - Expand each plan with remaining tables
    - Depth-first
  - If "cost of partial plan" > "cost of best plan":
    - "prune" plan
  - Heuristic pruning:
    - Prune less promising partial plans
    - May in rare cases miss most optimal plan (turn off with set optimizer\_prune\_level = 0)



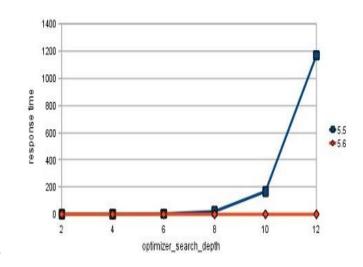


# Complexity and Cost of Join Optimizer

Join of N tables: N! possible plans to evaluate

Heuristics to reduce the number of plans to evaluate:

- Use optimizer\_search\_depth to limit the number of tables to consider
- Pre-sort tables on size and key dependency order (Improved in MySQL 5.6)
- When adding the next table to a partial plan, add all tables that it has an equality reference to (New in MySQL 5.6)





# Join Optimizer: Case study

**DBT-3 Query 8: National Market Share Query** 

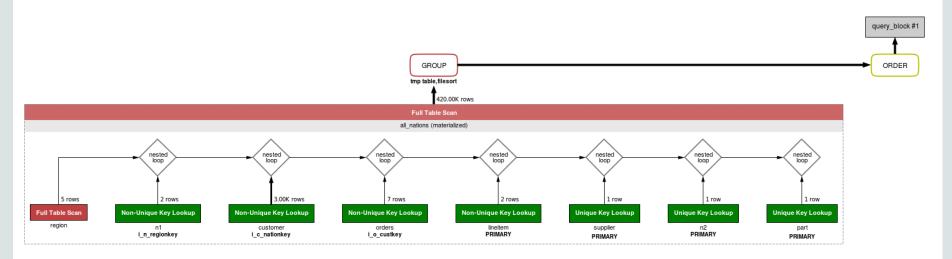
```
SELECT o year, SUM(CASE WHEN nation = 'FRANCE' THEN volume ELSE 0 END) / SUM(volume) AS
       mkt share
FROM (
       SELECT EXTRACT(YEAR FROM o_orderdate) AS o_year,
I_extendedprice * (1 - I_discount) AS volume, n2.n_name AS nation
       FROM part
         JOIN lineitem ON p_partkey = l_partkey
JOIN supplier ON s_suppkey = l_suppkey
JOIN orders ON l_orderkey = o_orderkey
         JOIN customer ON o_custkey = c_custkey
         JOIN nation n1 ON c_nationkey = n1.n_nationkey
       JOIN region ON n1.n_regionkey = r_regionkey
JOIN nation n2 ON s_nationkey = n2.n_nationkey
WHERE r_name = 'EUROPE' AND o_orderdate BETWEEN '1995-01-01' AND '1996-12-31'
            AND p type = 'PROMO BRUSHED STEEL'
 ) AS all_nations GROUP BY o_year ORDER BY o_year;
```



# Join Optimizer: Case Study

**MySQL Workbench: Visual EXPLAIN** 

Execution time: 3 min. 28 sec.







# Join Optimizer: Case study

Force early processing of high selectivity predicates

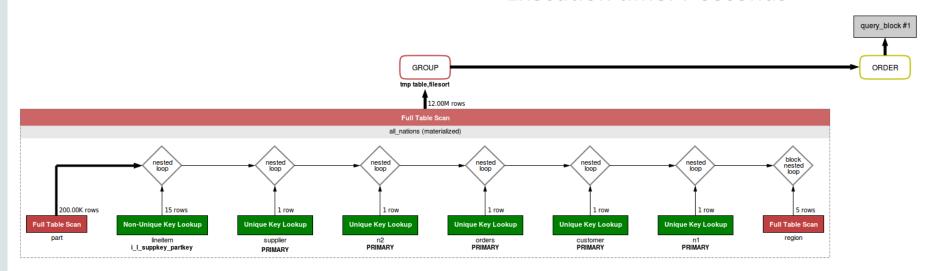
```
SELECT o year, SUM(CASE WHEN nation = 'FRANCE' THEN volume ELSE 0 END) / SUM(volume) AS
       mkt share
                                                                                         part before lineitem
FROM (
       SELECT EXTRACT(YEAR FROM o orderdate) AS o year,
I extendedprice * (1 - I discount) AS volume, n2.n name AS nation
       FROM part \angle
         STRAIGHT_JOIN lineitem ON p_partkey = l_partkey JOIN supplier ON s_suppkey = l_suppkey JOIN orders ON l_orderkey = o_orderkey
         JOIN customer ON o_custkey = c_custkey
         JOIN nation n1 ON c_nationkey = n1.n_nationkey
       JOIN region ON n1.n_regionkey = r_regionkey
JOIN nation n2 ON s_nationkey = n2.n_nationkey
WHERE r_name = 'EUROPE' AND o_orderdate BETWEEN '1995-01-01' AND '1996-12-31'
           AND p_type = 'PROMO BRUSHED STEEL' _
                                                                                                    Highest selectivity
) AS all_nations GROUP BY o_year ORDER BY o_year;
```



# Join Optimizer: Case study

Improved join order

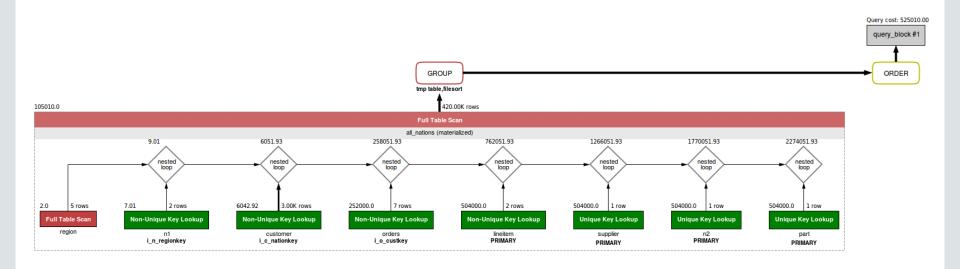
Execution time: 7 seconds







# MySQL 5.7: Cost Information in Structured EXPLAIN







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### **ORDER BY Optimizations**

- General solution; "Filesort":
  - Store query result in temporary table before sorting
  - If data volume is large, may need to sort in several passes with intermediate storage on disk.
- Optimizations:
  - Take advantage of index to generate query result in sorted order
  - For "LIMIT n" queries, maintain priority queue of n top items in memory instead of filesort. (New in MySQL 5.6)



### Filesort

SELECT \* FROM orders ORDER BY o\_totalprice;

id	select type	table	type	possible keys	key	key len	ref	rows	extra
1	SIMPLE	orders	ALL	NULL	NULL	NULL	NULL	15000000	Using filesort

SELECT c\_name, o\_orderkey, o\_totalprice FROM orders JOIN customer ON c\_custkey = o\_custkey WHERE c\_acctbal < -1000 ORDER BY o\_totalprice;

id	select type	table	type	possible keys	key	key len	ref	rows	extra
1	SIMPLE	customer	ALL	PRIMARY	NULL	NULL	NULL	1500000	Using where; Using temporary; Using filesort
1	SIMPLE	orders	ref	i_o_custkey	i_o_custkey	5	•••	7	NULL



### **Filesort**

#### **Status variables**

Status variables related to sorting:

Number of rows sorted



### Filesort: Case Study

```
mysql> FLUSH STATUS;
Query OK, 0 rows affected (0.00 sec)
mysql> SELECT AVG(o totalprice) FROM (
 SELECT * FROM orders
   ORDER BY totalprice DESC
   LIMIT 100000) td;
 AVG(o_totalprice)
  398185.986158
1 row in set (24.65 sec)
```

Unnecessary large data volume!

```
mysql> SHOW STATUS LIKE 'sort%';
| Variable name
                    | Value
  Sort_merge_passes | 1432
  Sort range
  Sort rows
                  1 100,000
  Sort scan
4 rows in set (0.00 \text{ seq})
```

Many intermediate sorting steps!



# Filesort: Case Study

#### Reduce amount of data to be sorted

```
mysql> FLUSH STATUS;
Query OK, 0 rows affected (0.00 sec)
mysql> SELECT AVG(o totalprice) FROM (
 SELECT o totalprice FROM orders
   ORDER BY o totalprice DESC
   LIMIT 100000) td;
 AVG(o totalprice) |
  398185.986158
1 row in set (8.18 sec)
```



# Filesort: Case Study Increase sort buffer (1 MB)

```
Default is 256 kB
```

```
mysql> set sort buffer size=1024*1024;
mysql> FLUSH STATUS;
Query OK, 0 rows affected (0.00 sec)
mysql> SELECT AVG(o totalprice) FROM
  ( SELECT o totalprice FROM orders
   ORDER BY o totalprice DESC
   LIMIT 100000) td;
  AVG(o_totalprice) |
  398185.986158
1 row in set (7.24 sec)
```



# Filesort: Case Study

Increase sort buffer even more (8 MB)

```
mysql> set
 sort buffer size=8*1024*1024;
mysql> FLUSH STATUS;
Query OK, 0 rows affected (0.00 sec)
mysql> SELECT AVG(o totalprice) FROM
  ( SELECT o totalprice FROM orders
   ORDER BY o totalprice DESC
   LIMIT 100000) td;
 AVG(o totalprice) |
 398185.986158 I
1 row in set (6.30 sec)
```

NB! Bigger sort buffer than needed will give unnecessary overhead



### Use Index to Avoid Sorting

**CREATE INDEX i\_o\_totalprice ON orders(o\_totalprice)**;

SELECT AVG(o\_totalprice) FROM

(SELECT o\_totalprice FROM orders ORDER BY o\_totalprice DESC LIMIT 100000) td;

id	select type	table	Туре	possible keys	key	key len	ref	rows	extra
1	PRIMARY	<derived2></derived2>	ALL	NULL	NULL	NULL	NULL	100000	NULL
2	DERIVED	orders	index	NULL	i_o_totalprice	6	NULL	15000000	Using index

```
mysql> SELECT AVG(o_totalprice) FROM (
    SELECT o_totalprice FROM orders
    ORDER BY o_totalprice DESC LIMIT 100000) td;
...
1 row in set (0.06 sec)
```





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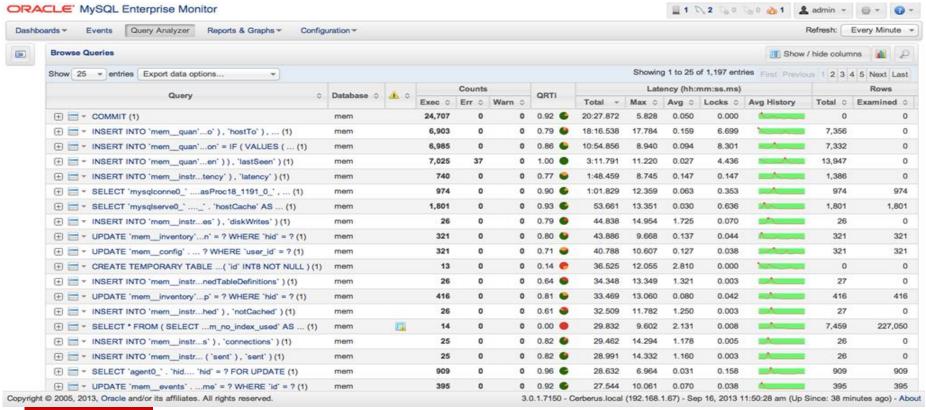


### Useful tools

- MySQL Enterprise Monitor (MEM), Query Analyzer
  - Commercial product
- Performance schema, MySQL sys schema
- EXPLAIN
- EXPLAIN FORMAT=JSON
- Optimizer trace
- Slow log
- Status variables (SHOW STATUS LIKE 'Handler%')

# MySQL Enterprise Monitor, Query Analyzer







### Query Analyzer Query Details







### Performance Schema

#### Some useful tables

- events\_statements\_history\_long
  - Most recent statements executed
- events\_statements\_summary\_by\_digest
  - Summary for similar statements (same statement digest)
- file\_summary\_by\_event\_name
  - Interesting event: wait/io/file/innodb/innodb\_data\_file
- table\_io\_waits\_summary\_by\_table
   table\_io\_waits\_summary\_by\_index\_usage
  - Statistics on storage engine access per table and index





### Performance Schema

#### **Statement digest**

• Normalization of queries to group statements that are similar to be grouped and summarized:

```
SELECT * FROM orders WHERE o_custkey=10 AND o_totalprice>20
SELECT * FROM orders WHERE o_custkey = 20 AND o_totalprice > 100
```

- SELECT \* FROM orders WHERE o\_custkey = ? AND o\_totalprice > ?
- events statements summary by digest

```
DIGEST, DIGEST_TEXT, COUNT_STAR, SUM_TIMER_WAIT, MIN_TIMER_WAIT, AVG_TIMER_WAIT, MAX_TIMER_WAIT, SUM_LOCK_TIME, SUM_ERRORS, SUM_WARNINGS, SUM_ROWS_AFFECTED, SUM_ROWS_SENT, SUM_ROWS_EXAMINED, SUM_CREATED_TMP_DISK_TABLES, SUM_CREATED_TMP_TABLES, SUM_SELECT_FULL_JOIN, SUM_SELECT_FULL_RANGE_JOIN, SUM_SELECT_RANGE, SUM_SELECT_RANGE_CHECK, SUM_SELECT_SCAN, SUM_SORT_MERGE_PASSES, SUM_SORT_RANGE, SUM_SORT_ROWS, SUM_SORT_SCAN, SUM_NO_INDEX_USED, SUM_NO_GOOD_INDEX_USED, FIRST_SEEN, LAST_SEEN
```



### Performance Schema

#### Statement events

• Tables:

```
events_statements_current (Current statement for each thread)
events_statements_history (10 most recent statements per thread)
events_statements_history_long (10000 most recent statements)
```

Columns:

THREAD\_ID, EVENT\_ID, END\_EVENT\_ID, EVENT\_NAME, SOURCE, TIMER\_START, TIMER\_END, TIMER\_WAIT, LOCK\_TIME, SQL\_TEXT, DIGEST, DIGEST\_TEXT, CURRENT\_SCHEMA, OBJECT\_TYPE, OBJECT\_SCHEMA, OBJECT\_NAME, OBJECT\_INSTANCE\_BEGIN, MYSQL\_ERRNO, RETURNED\_SQLSTATE, MESSAGE\_TEXT, ERRORS, WARNINGS, ROWS\_AFFECTED, ROWS\_SENT, ROWS\_EXAMINED, CREATED\_TMP\_DISK\_TABLES, CREATED\_TMP\_TABLES, SELECT\_FULL\_JOIN, SELECT\_FULL\_RANGE\_JOIN, SELECT\_RANGE, SELECT\_RANGE, SELECT\_RANGE\_CHECK, SELECT\_SCAN, SORT\_MERGE\_PASSES, SORT\_RANGE, SORT\_ROWS, SORT\_SCAN, NO\_INDEX\_USED, NO\_GOOD\_INDEX\_USED, NESTING\_EVENT\_ID, NESTING\_EVENT\_TYPE



# MySQL sys Schema / ps\_helper

- Started as a collection of views, procedures and functions, designed to make reading raw Performance Schema data easier
- Implements many common DBA and Developer use cases
- Now bundled within MySQL Workbench
- Available on GitHub
  - https://github.com/MarkLeith/mysql-sys
- Examples of very useful functions:
  - format\_time() , format\_bytes(), format\_statement()



### MySQL sys Schema

#### **Example**

**statement\_analysis:** Lists a normalized statement view with aggregated statistics, mimics the MySQL Enterprise Monitor Query Analysis view, ordered by the total execution time per normalized statement

```
lock_latency: 00:18:29.18

rows_sent: 0

rows_sent_avg: 0

rows_examined: 0

rows_examined_avg: 0

tmp_tables: 0

tmp_disk_tables: 0

rows_sorted: 0

sort_merge_passes: 0

digest: d48316a218e95b1b8b72db5e6b177788!

first_seen: 2014-05-20 10:42:17
```



# Optimizer Trace: Query Plan Debugging

- EXPLAIN shows the selected plan
- TRACE shows WHY the plan was selected:
  - Alternative plans
  - Estimated costs
  - Decisions made
- JSON format





### Optimizer Trace: Example

```
SET optimizer_trace= "enabled=on", end_markers_in_json=on;
SELECT * FROM t1, t2 WHERE f1=1 AND f1=f2 AND f2>0;
SELECT trace INTO DUMPFILE <filename>
    FROM information_schema.optimizer_trace;
SET optimizer_trace="enabled=off";
```

QUERY	SELECT * FROM t1,t2 WHERE f1=1 AND f1=f2 AND f2>0;
TRACE	"steps": [ { "join_preparation": { "select#": 1, } } ]
MISSING_BYTES_BEYOND_MAX_MEM_SIZE	0
INSUFFICIENT_PRIVILEGES	0



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# Influencing the Optimizer

#### When the optimizer does not do what you want

- Add indexes
- Force use of specific indexes:
  - USE INDEX, FORCE INDEX, IGNORE INDEX
- Force specific join order:
  - STRAIGHT JOIN
- Adjust session variables
  - optimizer\_switch flags: set optimizer\_switch="index\_merge=off"
  - Buffer sizes: set sort buffer=8\*1024\*1024;
  - Other variables: set optimizer\_prune\_level = 0;



### More information

- My blog:
  - <a href="http://oysteing.blogspot.com/">http://oysteing.blogspot.com/</a>
- Optimizer team blog:
  - http://mysqloptimizerteam.blogspot.com/
- MySQL Server Team blog
  - http://mysqlserverteam.com/
- MySQL forums:
  - Optimizer & Parser: <a href="http://forums.mysql.com/list.php?115">http://forums.mysql.com/list.php?115</a>
  - Performance: <a href="http://forums.mysql.com/list.php?24">http://forums.mysql.com/list.php?24</a>



#### Safe Harbor Statement

The preceding is intended to outline our general product direction. It is intended for information purposes only, and may not be incorporated into any contract. It is not a commitment to deliver any material, code, or functionality, and should not be relied upon in making purchasing decisions. The development, release, and timing of any features or functionality described for Oracle's products remains at the sole discretion of Oracle.



