

# Network Security Contd

COMP90007

Internet Technologies

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## Public Key Algorithms

- **Fundamentally different** to symmetric key ones
- Diffie & Hellman proposed the new model
  - **Asymmetric key algorithms**
  - **Key used to encrypt and key used to decrypt different**
  - **Not easily derivable from each other**
- Diffie-Hellman key system
  - **Key 1: public key**, usable by anyone **to encrypt** messages to the owner of the key, this key known to all
  - **Key 2: private key**, required **to decrypt** the message and known only by the owner of this key

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## The Process

- C = ciphertext, P = plaintext, E = encryption, D=decryption  
K1, K2 = keys
- $C = E_{K1}(P)$ 
  - Sender knows the public key K1 and the P
- $P = D_{K2}(C)$ 
  - Only receiver knows private K2 which can undo K1's effect
- $D_{K2}(E_{K1}(P)) = P$

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## RSA: An Asymmetric Key Algorithm

- RSA - Rivest, Shamir, Adleman
- Famous and robust algorithm
- Key generation:
  - Choose two large primes, p and q
  - Compute  $n = p \times q$  and  $z = (p - 1) \times (q - 1)$ .
  - Choose d to be relatively prime to z, i.e., no common factors
  - Find e such that
    - $(d \times e) \bmod z = 1$
  - Public key is (e, n), and private key is (d, n)
- Encryption:
  - Cipher =  $\text{Plain}^e \pmod n$
- Decryption:
  - Plain =  $\text{Cipher}^d \pmod n$

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## RSA Security

- RSA's security is **based on the difficulty involved in factoring large numbers in math theory** - approx  $10^{25}$  years to factor a 500 digit number and RSA uses 1024 bits!
- RSA is too slow for encrypting/decrypting large volumes of data, but is widely used for many other things such as **secure key distribution**
- RSA can be used in tandem with symmetric key algorithms

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## RSA Example

- Let  $p=3$ ,  $q=11$ : then  $z$  is  $(3 - 1) \times (11 - 1) = 20$
- What is a potential  $d$ ?
- If  $d = 7$  then  $z$  and  $20$  has no common factors
- What is an  $e$ ?
- If  $e = 3$ , then  $(d \times e)$  is 1 in mod  $z$
- What are the two key tuples then?
- Enc: **3. 33** Dec: **7. 33** (as  $n=3 \times 11=33$  and  $d=7$  and  $e=3$ )

Plaintext (P)		Ciphertext (C)			After decryption	
Symbolic	Numeric	$P^3$	$P^3 \pmod{33}$	$C^7$	$C^7 \pmod{33}$	Symbolic
S	19	6859	28	13492928512	19	S
U	21	9261	21	1801088541	21	U
Z	26	17576	20	1280000000	26	Z
A	01	1	1	1	01	A
N	14	2744	5	78125	14	N
N	14	2744	5	78125	14	N
E	05	125	26	8031810176	05	E

Encryption:  $C = P^3 \pmod{33}$

Decryption:  $P = C^7 \pmod{33}$

S is the 19<sup>th</sup> character in the alphabet...

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## Another Use of Cryptography: Digital Signatures

- Cryptographic approaches can also be used to ensure **authenticity** and allow for **non-repudiation**
- Requirements
  - Receiver can **verify the claimed identity of the sender**
  - **Sender cannot deny she created** contents of the message
  - **Receiver cannot have derived the message themselves**
- Three approaches
  - Using symmetric keys via an intermediary
    - You need a BIG BROTHER to do all the messaging, not good!
  - Using **public keys** as individuals

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## Using Public Keys

- Sender Alice uses **private key on P**
- **Receiver Bob uses her public key to undo and get P**
- RSA can do this as well, as  **$E(D(P)) = P$  in RSA**
- Alice cannot deny signing as she only knows her private key

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## Signatures with Message Digests

- Basic concept of a **message digest is to use a one-way hash function** for an arbitrary length of plaintext, so that it becomes a **"unique" small fixed-length bit string**
- Thus **no need to deal with huge message text and encryption just for authentication** purposes
- A message digest (MD) has four important properties:
  - 1 Given P, it is easy to compute MD(P)
  - 2 Given MD(P) it is effectively impossible to find P
  - 3 Given P, no one can find P' such that MD(P') = MD(P)
  - 4 A change in even a single bit of input produces a very different output

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## Famous Message Digest Algorithms

- MD5
- SHA-1
- Outputs
  - Given "this is a test" (text could have been longer)
  - MD5:  
e19c1283c925b3206685522acfe3e6
  - SHA-1:  
6476df3aac780622368173fe6e768a2edc3932c8

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## Public Key Management

- There is **specific PK infrastructure** to avoid compromising the security of PK's **during the initial distribution process**.
- Certification Authority (CA)
  - A trusted intermediary who uses non-electronic identification to identify users prior to certifying keys and certificates
- X.509
  - An international standard for certificate expression
- PKI (Public Key Infrastructure) is a
  - **Hierarchically structured certificate authorities** allow for the establishment of a chain of trust or certification path
  - *Verisign* was such a company

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## Certificate Issuing

- A Certificate authority (CA) says:

I hereby certify that the public key  
19836A8B03030CF83737E3837837FC3s87092827262643FFA82710382828282A  
belongs to  
Robert John Smith  
12345 University Avenue  
Berkeley, CA 94702  
Birthday: July 4, 1958  
Email: bob@superdupernet.com

SHA-1 hash of the above certificate signed with the CA's private key

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