

# Active GNN

---

- S2: An efficient graph based active learning algorithm with application to nonparametric classification
  - Problem setup
    - active learning for binary label prediction on a graph
    - nonparametric active learning, S2 sequentially select vertices to be labeled
      - *cut-set*:  $C = \{\{x, y\} \in E : f(x) \neq f(y)\}$
      - *boundary*:  $\partial C = \{x \in V : \exists e \in C \text{ with } x \in e\}$
      - goal is to identify  $\partial C$
    - algorithm assume a noiseless oracle that return label of a multiset of vertices, noisy oracle version algorithms can be transferred from noiseless
    - can be extended to multi-class
  - Datasets
    - Digits:
      - Cedar Buffalo binary digits database
      - construct symmetrized 10-nearest-neighbor graph
    - Congressional Voting Records (CVR):
      - 380 vertices, boundary size of 234
    - Grid:
      - synthetic example of a 15x15, positive core in the center
  - Methods
    - S2: Shortest Shortest Path

---

**Algorithm 1**  $S^2$ : Shortest Shortest Path

---

**Input** Graph  $G = (V, E)$ , BUDGET  $\leq n$ 

```
1:  $L \leftarrow \emptyset$ 
2: while 1 do
3:    $x \leftarrow$  Randomly chosen unlabeled vertex
4:   do
5:     Add  $(x, f(x))$  to  $L$ 
6:     Remove from  $G$  all edges whose two ends have different labels.
7:     if  $|L| = \text{BUDGET}$  then
8:       Return LABELCOMPLETION( $G, L$ )
9:     end if
10:    while  $x \leftarrow \text{MSSP}(G, L)$  exists
11:  end while
```

---

- LABELCOMPLETION: Any off-the-shelf graph prediction algorithms
- MSSP: return midpoint on the shortest among all the shortest-paths that connect oppositely labeled vertices in  $L$

---

**Sub-routine 2** MSSP

---

**Input** Graph  $G = (V, E)$ ,  $L \subseteq V$ 

```
1: for each  $v_i, v_j \in L$  such that  $f(v_i) \neq f(v_j)$  do
2:    $P_{ij} \leftarrow$  shortest path between  $v_i$  and  $v_j$  in  $G$ 
3:    $\ell_{ij} \leftarrow$  length of  $P_{ij}$  ( $\infty$  if no path exists)
4: end for
5:  $(i^*, j^*) \leftarrow \arg \min_{v_i, v_j \in L: f(v_i) \neq f(v_j)} \ell_{ij}$ 
6: if  $(i^*, j^*)$  exists then
7:   Return mid-point of  $P_{i^*j^*}$  (break ties arbitrarily).
8: else
9:   Return  $\emptyset$ 
10: end if
```

---

- Can be seen as: random sampling + aggressive search
  - aggressive search: like binary search to find the cut-edge, then unzip the cut-edge

- Baselines

- measure query complexity
- AFS On the complexity of finding an unknown cut via vertex queries
- ZLG Combining active learning and semi- supervised learning using Gaussian fields and harmonic functions
- BND Towards active learning on graphs: An error bound minimization approach

---

- Active Learning for Networked Data

- Problem setup

- rr