

CS32 Discussion  
Section 1B  
Week 5

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# Recursion

- Function-writing technique where the function refers to itself.
- Recall the following function:

```
int factorial(int n)
{
    if (n <= 1)
        return 1;

    return n * factorial(n - 1);
}
```

- Let us talk about how to come up with such a function.

# Decomposition of the problem

- You're all used to the following technique.

```
int factorial(int n)
{
    int temp = 1;
    for (int i = 1; i <= n; i++)
        temp *= i;
    return temp;
}
```

- $n! = 1 * 2 * 3 * \dots * (n-1) * n$

# Decomposition of the problem

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- $n! = 1 * 2 * 3 * \dots * (n-1) * n$   
     $= \text{factorial}(n-1)!$

# Decomposition of the problem

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- $n! = 1 * 2 * 3 * \dots * (n-1) * n$
- $n! = \text{factorial}(n-1) * n$

# Decomposition of the problem

```
int factorial(int n)
{
    int temp = factorial(n - 1) * n;

    return temp;
}
```

- $n! = 1 * 2 * 3 * \dots * (n-1) * n$
- $n! = \text{factorial}(n-1) * n$

# Power of Belief

- **BELIEVE factorial(n - 1) will do the right thing.**

```
int factorial(int n)
{
    int temp = factorial(n - 1) * n;

    return temp;
}
```

- factorial(n) will believe that factorial(n-1) will return the right value.
- factorial(n-1) will believe that factorial(n-2) will return the right value.
- ...
- factorial(2) will believe that factorial(1) will return the right value.

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- ...
- factorial(2) will believe that factorial(1) will return the right value.
- **AND MAKE factorial(1) return the right value!**



# Base Case

- **BELIEVE factorial(n - 1) will do the right thing.**

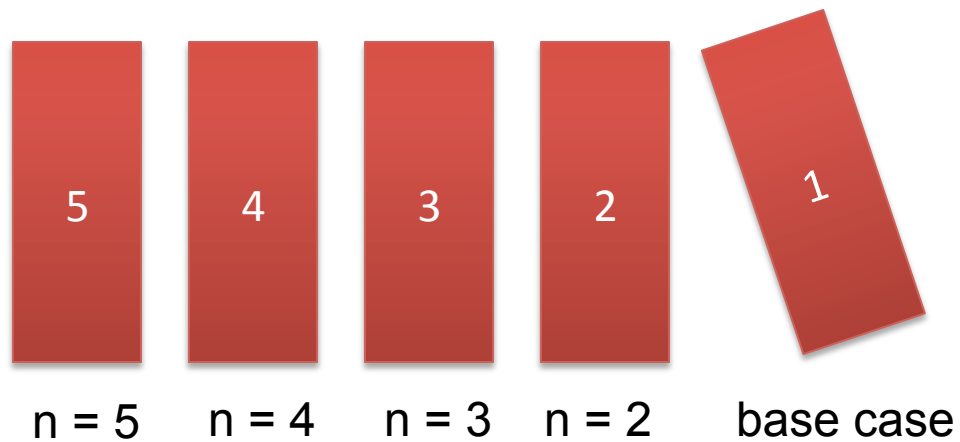
```
int factorial(int n)
{
    if (n <= 1)
        return 1;
    int temp = factorial(n - 1) * n;
    return temp;
}
```



- factorial(n) will believe that factorial(n-1) will return the right value.
- factorial(n-1) will believe that factorial(n-2) will return the right value.
- ...
- factorial(2) will believe that factorial(1) will return the right value.
- **AND MAKE factorial(1) return the right value!**

# Base Case

```
int factorial(int n)
{
    if (n <= 1)
        return 1;
    int temp = factorial(n - 1) * n;
    return temp;
}
```



# Pattern

- How to Write a Recursive Function for Dummies

1. Find the base case(s).

- What are the trivial cases? e.g. empty string, empty array, etc.
- When should the recursion stop?

2. Decompose the problem.

- Example: Tail recursion
  - Take the first (or last) of the  $n$  items of information.
  - Make a recursive call to the rest of  $n-1$  items, believing the recursive call will give you the correct result.
  - Given this result and the information you have on the first (or last) item, conclude about the  $n$  items.

3. Just solve this subproblem!

# Pattern

- Write your function this way:  
    `recursive_function(set of data)`
  1. Take care of all base cases
  2. `x = current data item`
  3. `result = recursive_function(set of data - {x})`
  4. combine `x` and `result`
- Lastly, look back and make sure every call to this function hits some base case eventually.
- There are variations:
  - You may need to make multiple recursive calls.
  - You may make a different recursive call based on `x`.

# Practice I: Average

```
double average(const double arr[], int n)
{
    // assume n > 0

}
```

# Practice 1: Average

- What is the base case?
- What is the relationship between  $(n)$ -step and  $(n-1)$ -step?  
That is, how do I get the average of  $n$  items, knowing the average of  $n-1$  of them?

# Practice 1: Average

- What is the base case?
- $n = 1$ , where the average is just the value of the only item.
- What is the relationship between (n)-step and (n-1)-step?  
That is, how do I get the average of n items, knowing the average of n-1 of them?
- $$\begin{aligned}\text{average}(\text{arr}, n) &= \text{total}(\text{all } n \text{ items}) / n \\ &= \{\text{total}(\text{first } n-1 \text{ items}) + n\text{-th item}\} / n \\ &= \frac{\text{average}(\text{arr}, n-1) * (n-1) + n\text{-th item}}{n}\end{aligned}$$

# Practice I:Average

```
double average(const double arr[], int n)
{
    if (n == 1)
        return arr[0];

    double prevAvg = average(arr, n - 1);
    return ((n - 1) * prevAvg + arr[n - 1]) / n;
}
```



# Practice 2: Summing Digits

```
int sumOfDigits(const int n)
{
    // assume n >= 0

}
```

# Practice 2: Summing Digits

- What is the base case?
- What is the relationship between  $(n)$ -step and  $(n-1)$ -step?  
That is, how do I get the sum of  $n$  digits, knowing the sum of digits of  $(n-1)$  digits?

# Practice 2: Summing Digits

- What is the base case?
- $n < 10$  (i.e.  $n$  is a single digit number), where the sum of digits is simply  $n$
- What is the relationship between  $(n)$ -step and  $(n-1)$ -step?  
That is, how do I get the sum of  $n$  digits, knowing the sum of digits of  $(n-1)$  digits?
- Just add the last digit to the sum!

# Practice 2: Summing Digits

```
int sumOfDigits(const int n)
{
    if (n < 10)
        return n;

    return n % 10 + sumOfDigits(n / 10);
}
```

# Practice 3: Deleting characters

```
string deleteChar(const string &s, const char c)
{
```

```
}
```

# Practice 3: Deleting characters

- What is the base case?
- What is the relationship between  $(n)$ -step and  $(n-1)$ -step?

# Practice 3: Deleting characters

- What is the base case?
- $s == ""$ : There is no character to delete! Just return  $""$ .
- What is the relationship between (n)-step and (n-1)-step?
- Suppose the string is of length n, and you make a recursive call on  $s.substr(1)$ . (e.g. If the string is "hello", the recursive call will be made on "ello".)
- What's returned by `deleteChar` must not contain any c. You only need to append  $s[0]$  to it if  $s[0] \neq c$ . If  $s[0] == c$ , just return it.

# Practice 3: Deleting characters

```
string deleteChar(const string &s, const char c)
{
    if (s.empty())
        return s;

    if (s[0] == c)
        return deleteChar(s.substr(1), c);
    else
        return s[0] + deleteChar(s.substr(1), c);
}
```



# Practice 4: Fibonacci numbers

- Fibonacci numbers refer to the sequence of numbers of the following form:

0, 1, 1, 2, 3, 5, 8, 13, 21, 34, 55, ...

- $F(0) = 0$
- $F(1) = 1$
- $F(n) = F(n-1) + F(n-2), n \geq 2$

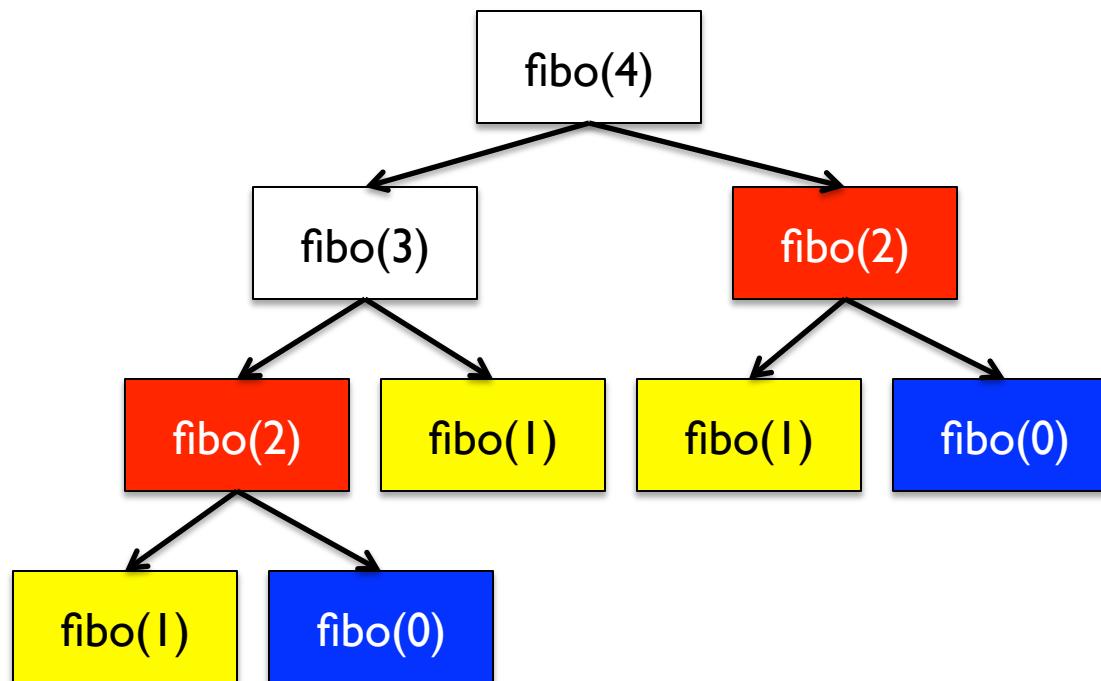
# Practice 4: Fibonacci numbers

```
// A little too obvious now, isn't it?
int fibo(const int n)
{
    if (n == 0)
        return 0;
    if (n == 1)
        return 1;

    return fibo(n - 1) + fibo(n - 2);
}
```

# Practice 4: Fibonacci numbers

- Note that `fibonacci()` makes two recursive calls.



- Look at all this redundancy!!!!

# Trick: Memoization

```
int fibMem[100]; // global
for (int i = 2; i < 100; ++i)
    fibMem[i] = -1;
fibMem[0] = 0;
fibMem[1] = 1;
```

```
int fibo(const int n)
{
    int fib1, fib2;
```

```
    if (fibMem[n-1] != -1)
        fib1 = fibMem[n-1];
    else
        fib1 = fibo(n-1);
```

```
    if (fibMem[n-2] != -1)
        fib2 = fibMem[n-2];
    else
        fib2 = fibo(n-2);

    fibMem[n] = fib1 + fib2;

    return fib1 + fib2;
}
```

Memoization is an optimization technique that helps avoid computing the same value over and over by remembering it.

We like this because memory is cheap, but computing time is not!

# Practice 5: Palindrome

- Examples: “eye”, “racecar”, “deed”

```
bool palindrome(const string &s)
{
```

```
}
```

- Base case?
- General case?

# Practice 5: Palindrome

- Examples: “eye”, “racecar”, “deed”

```
bool palindrome(const string &s)
{
```

```
}
```

- Cases: (1) size = 0, (2) size = 1, (3) the first char differs from the last one, (4) first char = last char

# Practice 5: Palindrome

- Examples: “eye”, “racecar”, “deed”

```
bool palindrome(const string &s)
{
    if (s.size() <= 1)
        return true;
    if (s[0] != s[s.size() - 1])
        return false;
    return palindrome(s.substr(1, s.size() - 2));
}
```

- Cases: (1) size = 0, (2) size = 1, (3) the first char differs from the last one, (4) first char = last char

# More Practice

- Write `reverse()`, which recursively reverses a string and return the reversed version.

```
string reverse(const string &s)
{

}
}
```



# More and More Practice

- Generate mnemonic phone numbers (e.g. 1-800-UCLA-CSD).
- Assume the following function is given:

```
string digitToLetters(char digit)
{
    switch (digit)
    {
        case '0': return "0";
        case '1': return "1";
        case '2': return "ABC";
        case '3': return "DEF";
        ...
        case '9': return "WXYZ";
        default: cout << "ERROR" << endl; abort();
    }
}
```



# More and More Practice

- Given `digitToLetters()`, write a function `mnemonics`, which prints all possible mnemonic numbers given the prefix and digits. (No restrictions here – you can use a loop.)

```
void mnemonics(const string &prefix, const string &digits);
```

```
mnemonics("", "723");
```

PAD  
PAE  
PAF  
PBD  
PBE  
PBF  
...

```
mnemonics("1-800-", "723")
```

1-800-PAD  
1-800-PAE  
1-800-PAF  
1-800-PBD  
1-800-PBE  
1-800-PBF  
...



# Practice helps

- Recursion is somewhat counter-intuitive when confronted for the first time.
- Just do a lot of practice and you will see some patterns.
- Try finding more examples by googling.
- Again, the key to recursion is to “believe”! Do not try to track the call stack down and see what happens until you really have to.

# Advanced Practice

- Given a set of integers, represented by an array `s`, write a function called `exactSum` that checks if the elements of some subset of `s` sums up to exactly `target`.

e.g. If `s[] = {1, 2, 3}`, then `exactSum(s, 3, 6)` returns true.

If `s[] = {-3, 5, 2}`, then `exactSum(s, 3, 2)` returns true.

If `s[] = {-3, 5, 2}`, then `exactSum(s, 3, 8)` returns false.

If `s[] = {}`, then `exactSum(s, 0, 10)` returns false.

If `s[] = {-3, 5, 2, -10}`, then `exactSum(s, 3, 2)` returns true.

```
bool exactSum(int a[], int n, int targetSum) {  
}
```

# Advanced Practice

- Write a `isPrimeNumber` function to test if a number is prime number or not.

How to do it in an un-iterative way? What's the algorithm?

```
Bool isPrimeNumber(int n){  
  
}
```