

# **CSE216**

# **Programming Abstraction**

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# Today

- Capture avoiding substitution in beta reduction
- Extending lambda calculus core

# Capture avoiding substitutions

- The specific names of bound variables in the lambda calculus are meaningless
- $\lambda x. x$  same as  $\lambda y. y$
- $(\lambda x. (\lambda y. yx))$  is equivalent to  $(\lambda a. (\lambda b. ba))$
- Lambda terms that differ only by bound variable names are called alpha equivalent

# Exercise: alpha equivalence?

- $\lambda x. xy \quad ? \quad \lambda z. zy$
- $\lambda x. xy \quad ? \quad \lambda z. xz$
- $\lambda x. x \lambda y. y \quad ? \quad \lambda z. z \lambda p. p$
- $\lambda x. x \lambda y. y \quad ? \quad \lambda y. y \lambda y. y$
- $\lambda x. x \lambda y. xy \quad ? \quad \lambda y. y \lambda y. yy$

# Exercises: beta reduction

- $(\lambda x. \lambda y. x)y$

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- $(\lambda x. x \lambda y. x)y$

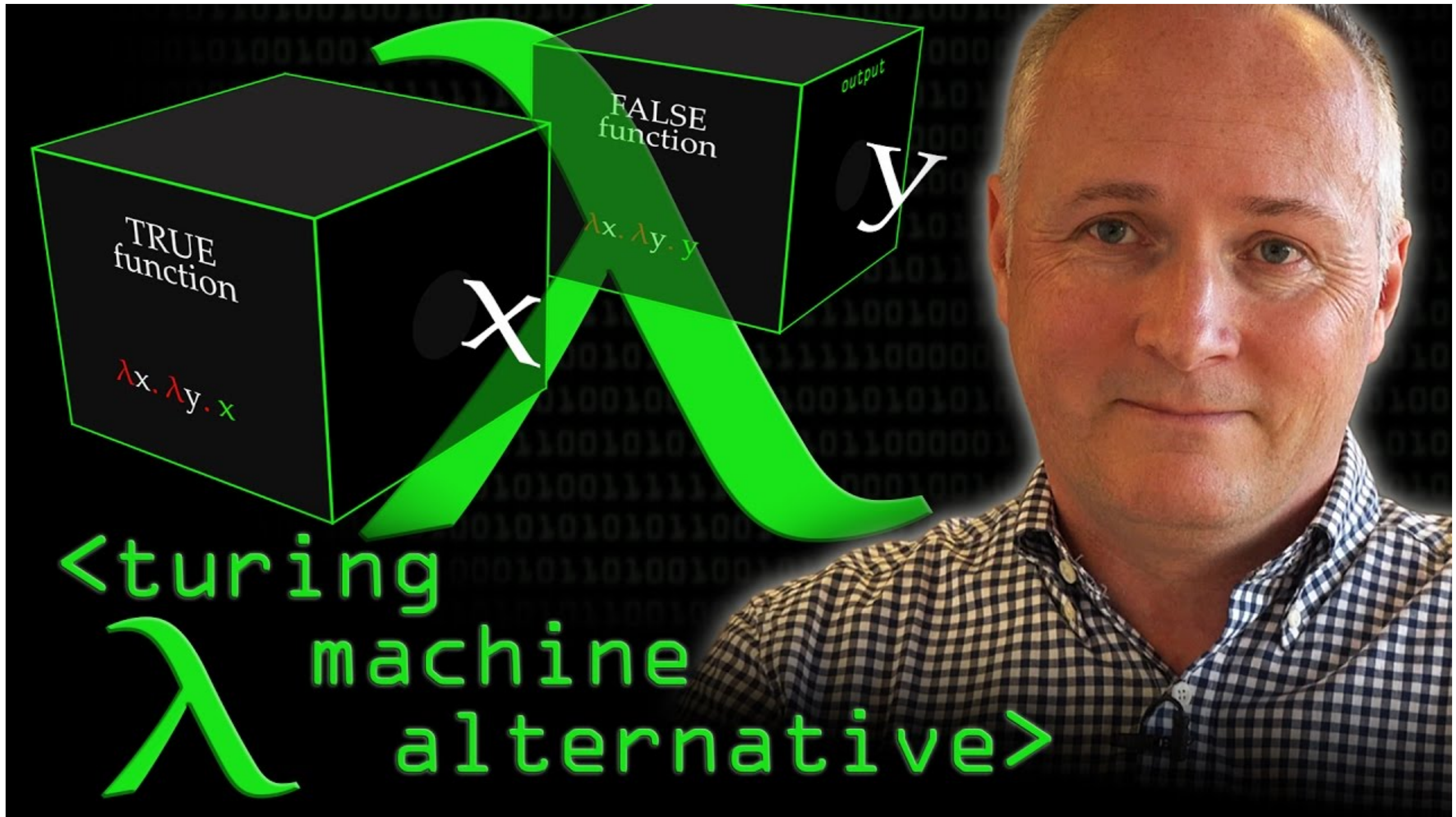
# Exercises: beta reduction

- $(\lambda x. \lambda y. x y) y$

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- $(\lambda x. \lambda y. xy)(\lambda x. \lambda y. xy)$





6'-10'

- [https://youtu.be/eis11j\\_iGMs](https://youtu.be/eis11j_iGMs)