

# **CSE216**

# **Programming Abstraction**

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# Today

- Capture avoiding substitution in beta reduction
- Extending lambda calculus core

# Capture avoiding substitutions

- The specific names of bound variables in the lambda calculus are meaningless
- $\lambda x. x$  same as  $\lambda y. y$
- $(\lambda x. (\lambda y. yx))$  is equivalent to  $(\lambda a. (\lambda b. ba))$
- Lambda terms that differ only by bound variable names are called alpha equivalent

# Exercise: alpha equivalence?

- $\lambda x. xy \quad ? \quad \lambda z. zy$
- $\lambda x. xy \quad ? \quad \lambda z. xz$
- $\lambda x. x \lambda y. y \quad ? \quad \lambda z. z \lambda p. p$
- $\lambda x. x \lambda y. y \quad ? \quad \lambda y. y \lambda y. y$
- $\lambda x. x \lambda y. xy \quad ? \quad \lambda y. y \lambda y. yy$

# Exercises: beta reduction

- $(\lambda x. \lambda y. x)y$

# Exercises: beta reduction

- $(\lambda x. x \lambda y. x)y$

# Exercises: beta reduction

- $(\lambda x. \lambda y. x y) y$

# Exercises: beta reduction

- $(\lambda x. \lambda y. xy)(\lambda x. \lambda y. xy)$



# Extending lambda calculus core

# “extending lambda calculus core”?

- The core has only three constructs
- It does not have numbers, conditions, logic, loops
- These things can all be encoded by the core
- We will write  $\| \langle \text{language syntax} \rangle \| = \langle \text{lambda-term} \rangle$  for the encoding
- *"The integers were created by God, everything else is the work of man"*

# TRUE, FALSE, and IF

- IF c e1 e2 returns e1 if c is TRUE, or e2 if c is FALSE
- So we encode TRUE by  $\lambda x. \lambda y. x$
- We encode FALSE by  $\lambda x. \lambda y. y$
- We encode IF by  $\lambda c. \lambda x. \lambda y. c x y$

## ► Examples

- if true then b else c =  $(\lambda x. \lambda y. x) b c \rightarrow (\lambda y. b) c \rightarrow b$
- if false then b else c =  $(\lambda x. \lambda y. y) b c \rightarrow (\lambda y. y) c \rightarrow c$

# Logic AND

- $\| \text{AND } x \ y \| = \| \text{IF } x \ y \ \text{FALSE} \| = x \ y \ \text{FALSE} = x \ y \ x$
- Thus  $\| \text{AND} \| = \lambda x. \lambda y. x \ y \ x$
- Exercise:  $\| \text{OR} \| = ?$
- Exercise:  $\| \text{NOT} \| = ?$

# Logic OR

- Exercise:  $\|OR\|=?$

# Logic NOT

- Exercise:  $\| \text{NOT} \| = ?$

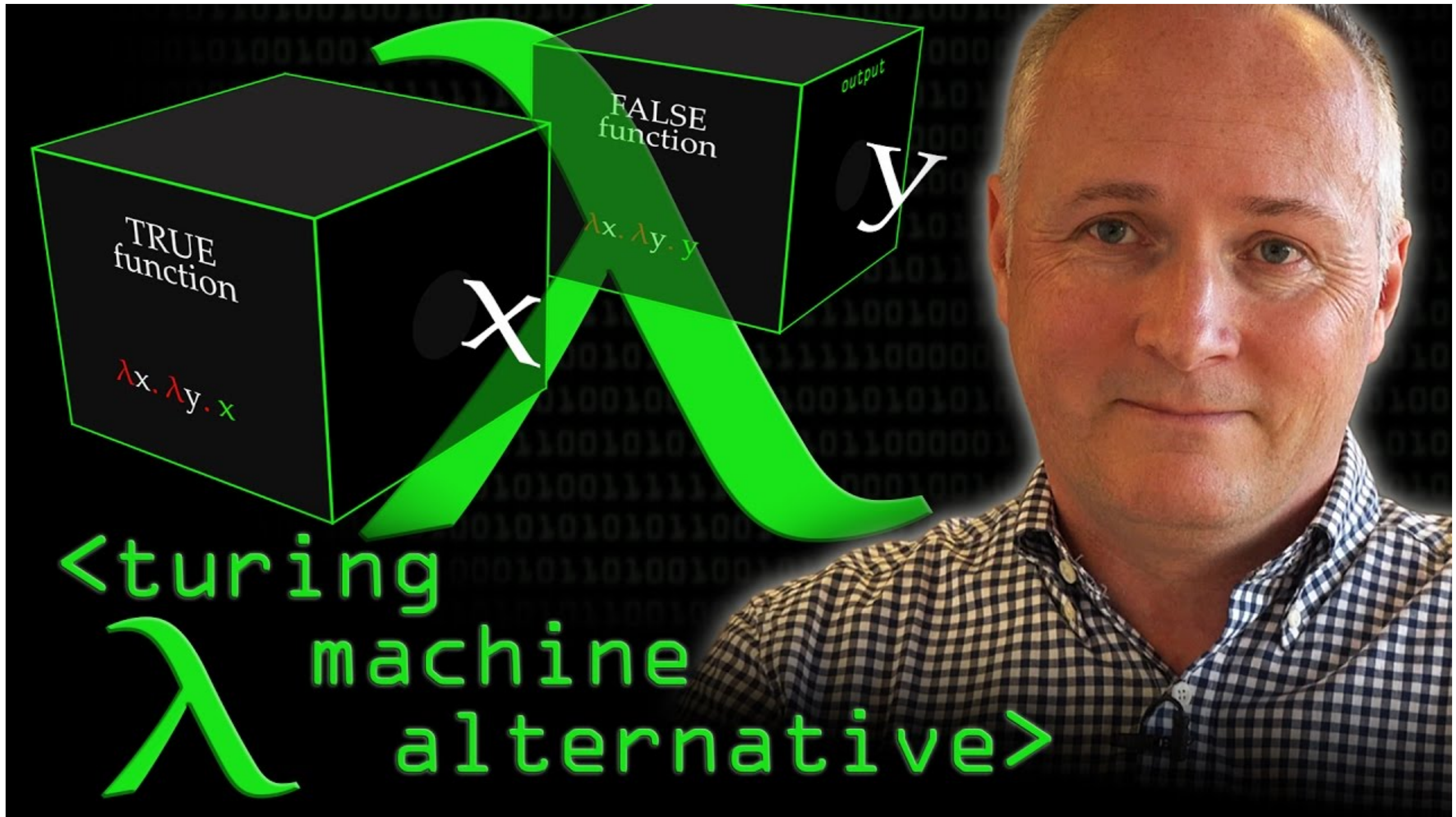
# Numbers

- Any counting system that makes sense would work
- We want  $n \ f \ x = f( f( f( f( \dots f(x) ) ) ) ) )$
- Since  $0 \ f \ x = x$ , we have  $\|0\| = \lambda f. \lambda x. x$
- Since  $1 \ f \ x = f \ x$ , we have  $\|1\| = \lambda f. \lambda x. f \ x$
- ... (called Church numerals)

# Things important in CS but not essential for next parts

- SUM
- PROD
- ....
- Recursion





6'-10'

- [https://youtu.be/eis11j\\_iGMs](https://youtu.be/eis11j_iGMs)