

# **CSE216**

# **Programming Abstraction**

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# Agenda

- Practice on lambda calculus syntax and semantics

# Lambda calculus syntax summary

- $\text{TERM} ::= \text{Var}$  // Variables
- |  $\text{lambda Var. TERM}$  // Definition/Abstraction
- |  $\text{TERM TERM}$  // Application

$\text{Var} ::= x \mid y \mid z \dots$

# Lambda calculus semantics summary

- $(\lambda x.M)N \rightarrow M$  with bound occurrences of  $x$  substituted by  $N$
- Sometimes written  $M [x \rightarrow N]$
- Something to say later, called “capture-avoiding”

# Draw parse trees

- $\lambda x. x \lambda y$

# Draw parse trees

- $\lambda x. \lambda y. x y \lambda z. z y$

# Draw parse trees

- $\lambda x. \lambda y. x y z$

# beta-reduction

- $(\lambda x.xx)(\lambda y.yx)z$



# beta-reduction

- $(\lambda x.xx)(\lambda y.yx)z$

# beta-reduction

- Someone thinks the lambda expression  $((\lambda x. \lambda y. y) y) (\lambda x. x a)$  reduces to  $a$ . Can you spot the issue?