

CSE216

Programming Abstraction

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Clarification

- Attendance checks will take place weekly
- Quizzes, or in-class assignments, will take place weekly, unless we are behind schedule.
- Each quiz taken at week N will be reviewed at week $N+1$.
- Next quiz will be on Wednesday, covering everything we will have learned before the quiz. Things that require hard memorization, if any, will be reminded in the quiz.
- Homework, or after-class assignments, will directly relate to coding, to be done after we have learned more about Ocaml/Java/Python

Today

- Quiz week 02 review
- Context-free grammars, derivation and ambiguity
- Lambda calculus 1

Quiz week 02 review

Quiz Week 02 statistics

Number of submitted grades: 24 / 24



- Clarification: The total score is 100.
- The last four exercises are counted as 5 points each, instead of 10.

Programming paradigms

Consider the following program in Python:

```
numbers = [1, 2, 3, 4, 5, 6]
sum = 0
for number in numbers:
    if number % 3 == 0:
        sum += number
print(sum)
```



1. (points = 10) What will be output if we run this program?

- Answer: **9**

2. (points = 10) What is the major paradigm used in the code? Choose from (a-d) below:

(a) functional (b) object-oriented (c) imperative

- Answer: **c: imperative**

Lambda functions

In Python, lambda functions are defined using the lambda keyword followed by a comma-separated list of arguments (if any), followed by a colon and an expression. Here's an example:

```
add = lambda x, y: x + y
print(add(9, 3))
```

In this example, we define a lambda function `add` that takes two arguments `x` and `y`, and returns their sum. We then call the `add` function with arguments 9 and 3, which returns the sum 12.

3. (points = 10) Define an lambda function in python that takes a two input numbers and returns their distance. You can use the python function `abs` for the absolute value function. For example, `abs(-4.2)` returns 4.2. Write out the lambda expression. It should start with the key word "lambda".

- Answer: **`lambda x, y: abs(x-y)`**

A starter for onject-oriented programming

Consider the following code in Java

```
public class Customer {  
    private String name;  
    private SeniorityLevel level;  
  
    public Customer(String name, SeniorityLevel level) {  
        this.name = name;  
        this.level = level;  
    }  
  
    public String getName() {  
        return name;  
    }  
  
    public SeniorityLevel getLevel() {  
        return level;  
    }  
  
    public void setLevel(SeniorityLevel level) {  
        this.level = level;  
    }  
}  
  
enum SeniorityLevel {  
    NEW, REGULAR, VIP  
}
```



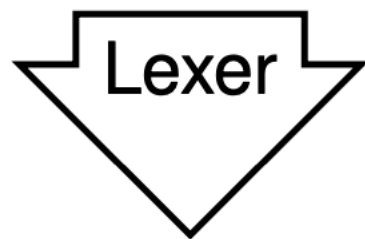
Now, your co-worker Ji-Ho is working on web development to design a frontend for customers to operate on their accounts. Assume Ji-Ho's code has access to Custom objects.

7. (points = 10) Is there any possibility that Ji-Ho's code can change a customer's name? **No**
8. (points = 10) Is there any possibility that Ji-Ho's code can change a customer's seniority level? **Yes**

Context-free grammar

Parsing

"x + y * 10"



Regular
expression
Specification

x

+

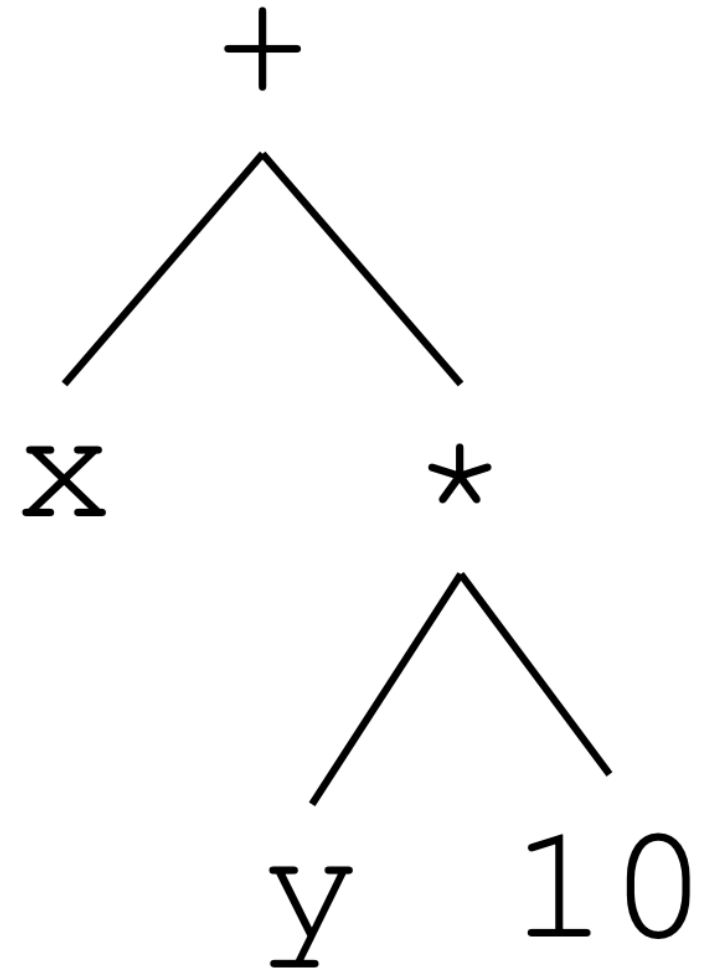
y

*

10

Parser

Context-
free grammar
specification



The need for a grammar

"Afternoon good, I'd room a like."



Mr. Men and Little Miss Series

Formal grammar

- Formal grammars are typically expressed using **a set of symbols and production rules** that define how those symbols can be combined to form valid sentences or expressions.
- There are **different types of formal grammars, including context-free grammars, regular grammars, and context-sensitive grammars.**
- Formal grammar is an essential tool for **analyzing and understanding the structure of languages.** It has applications in fields such as computational linguistics, natural language processing, and artificial intelligence.

Context-free grammar: Example 1

$$1. S \rightarrow aSb$$

$$2. S \rightarrow ba$$

- We start with S , and can choose a rule to apply to it.
- If we choose rule 1, we obtain the string aSb . If we then choose rule 1 again, we replace S with aSb and obtain the string $aaSbb$. If we now choose rule 2, we replace S with ba and obtain the string $aababb$, and are done.
- We can write this series of choices more briefly, using the **derivation**: $S \Rightarrow aSb \Rightarrow aaSbb \Rightarrow aababb$
- **The language of the grammar** is the infinite set $\{a^n ba b^n \mid n > 0\}$ where n is repeated times.
- This grammar is **context-free**, with only single nonterminals appear as left-hand sides)

Example 2: arithmetic expressions (ambiguous)

$\langle \text{expr} \rangle ::= \langle \text{expr} \rangle + \langle \text{expr} \rangle$

$| \langle \text{expr} \rangle - \langle \text{expr} \rangle$

$| \langle \text{expr} \rangle * \langle \text{expr} \rangle$

$| \langle \text{expr} \rangle / \langle \text{expr} \rangle$

$| (\langle \text{expr} \rangle)$

$| \langle \text{number} \rangle$

$\langle \text{number} \rangle ::= 0 \mid 1 \mid 2 \mid \dots \mid 9$

This grammar is **ambiguous** because it allows for multiple parse trees to represent the same expression. For example, the expression $2 + 3 * 4$ can be parsed as $(2 + 3) * 4$ or $2 + (3 * 4)$

Ambiguous grammar

- An ambiguous grammar is a formal grammar that can produce **multiple parse trees** or interpretations for the same input sentence or sequence of symbols.
- This can be problematic in various contexts because it can **make it difficult to determine the correct meaning** or parse tree of a sentence or sequence of symbols.
- To avoid ambiguity, it is often necessary to **use unambiguous grammars** or to add rules or constraints to the ambiguous grammar to disambiguate the interpretations.

Example 3: arithmetic expressions (unambiguous)

$\langle \text{expr} \rangle ::= \langle \text{expr} \rangle + \langle \text{term} \rangle$

$\quad | \langle \text{expr} \rangle - \langle \text{term} \rangle$

$\langle \text{term} \rangle ::= \langle \text{term} \rangle * \langle \text{factor} \rangle$

$\quad | \langle \text{term} \rangle / \langle \text{factor} \rangle$

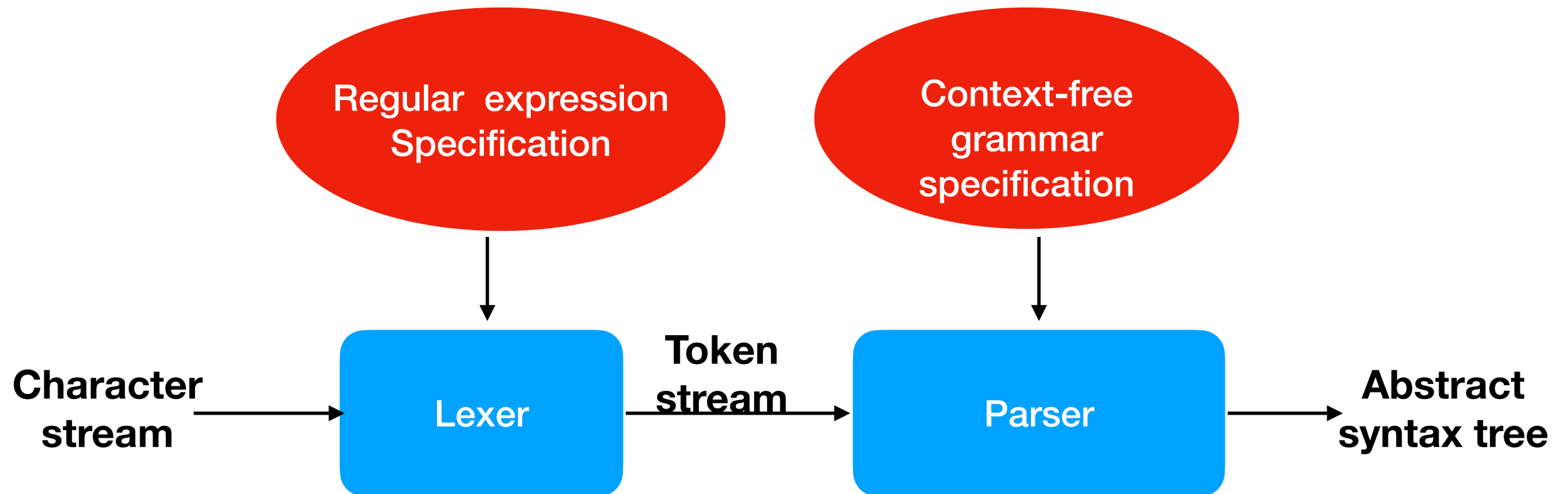
$\quad | \langle \text{factor} \rangle$

$\langle \text{factor} \rangle ::= (\langle \text{expr} \rangle) \mid \langle \text{number} \rangle$

$\langle \text{number} \rangle ::= 0 \mid 1 \mid 2 \mid \dots \mid 9$

In this revised grammar, the $\langle \text{expr} \rangle$ rule now includes a $\langle \text{term} \rangle$ component, which handles multiplication and division. The $\langle \text{term} \rangle$ rule includes a $\langle \text{factor} \rangle$ component, which handles parentheses and numbers. This approach ensures that the order of operations is clear and unambiguous

Summary so far



- A lexer is a DFA transformed from a specification of regular expressions.
- A parser derives a parse tree using a non-ambiguous context-free grammar

Exercise

Give a RE for: $L = \{0^i 1^j \mid i \text{ is even and } j \text{ is odd} \}$

Solution

- $(00)^*1(11)^*$

Exercise

- Write a context-free grammar that generates the language of all strings of a's and b's that of the form $a^n b^n$ where $n > 0$.

Solution

$S \rightarrow 01$

$S \rightarrow 0S1$

Lambda calculus

Many slides adapted from CMU 15-252: More Great Theoretical Ideas in Computer Science
<https://www.cs.cmu.edu/~venkatg/teaching/15252-sp20/notes/lambda-calculus-slides.pdf>

Thanks!

A very nice Intro to lambda calculus

