

CSE215

Foundations of Computer Science

State University of New York, Korea

Some slides taken from Prof. Pramod Ganapathi (Stony Brook). Thanks!

Plan

- Review exercises on Mathematical Induction
- Recursive sequences

Use mathematical induction to prove the following identities.

(a) [5 points] For all natural numbers n ,

$$1^2 \times 2 + 2^2 \times 3 + 3^2 \times 4 + \cdots + n^2 \times (n + 1) = \frac{n(n + 1)(n + 2)(3n + 1)}{12}$$

Use mathematical induction to prove the following identities.

(a) [5 points] For all natural numbers n ,

$$1^2 \times 2 + 2^2 \times 3 + 3^2 \times 4 + \cdots + n^2 \times (n + 1) = \frac{n(n + 1)(n + 2)(3n + 1)}{12}$$

- Proof.
 - We will prove this proposition with mathematical induction.
 - Let $P(n)$ be the predicate $1^2 \times 2 + \dots + n^2 \times (n+1) = n(n+1)(n+2)(3n+1)/12$
 - Base step: We first prove $P(1)$ holds.
 - $LHS = 1^2 \times 2 = 2$. $RHS = 1 \times 2 \times 3 \times 4 / 12 = 2$.
 - Inductive step: Then, we prove $P(k) \rightarrow P(k+1)$ for all $k \geq 1$
 - Let k be an arbitrary integer and $k \geq 1$.
 - Assume $P(k)$ holds. That is, $1^2 \times 2 + \dots + k^2 \times (k+1) = k(k+1)(k+2)(3k+1)/12$
 - We want to prove $P(k+1)$, that is, $1^2 \times 2 + \dots + k^2 \times (k+1) + (k+1)^2 \times (k+2) = (k+1)(k+2)(k+3)(3k+4)/12$
 - $LHS = k(k+1)(k+2)(3k+1)/12 + (k+1)^2 \times (k+2)$ following assumption $P(k)$
 - The latter can be further reduced to $(k+1)(k+2)/12 \times (3k^2 + k + 12k + 12) = RHS$
- QED.

Use mathematical induction to prove the following identities.

(b) [5 points] For all natural numbers n ,

$$\frac{1}{1 \times 2} + \frac{1}{2 \times 3} + \frac{1}{3 \times 4} + \cdots + \frac{1}{n \times (n+1)} = \frac{n}{n+1}$$

Use mathematical induction to prove the following identities.

(b) [5 points] For all natural numbers n ,

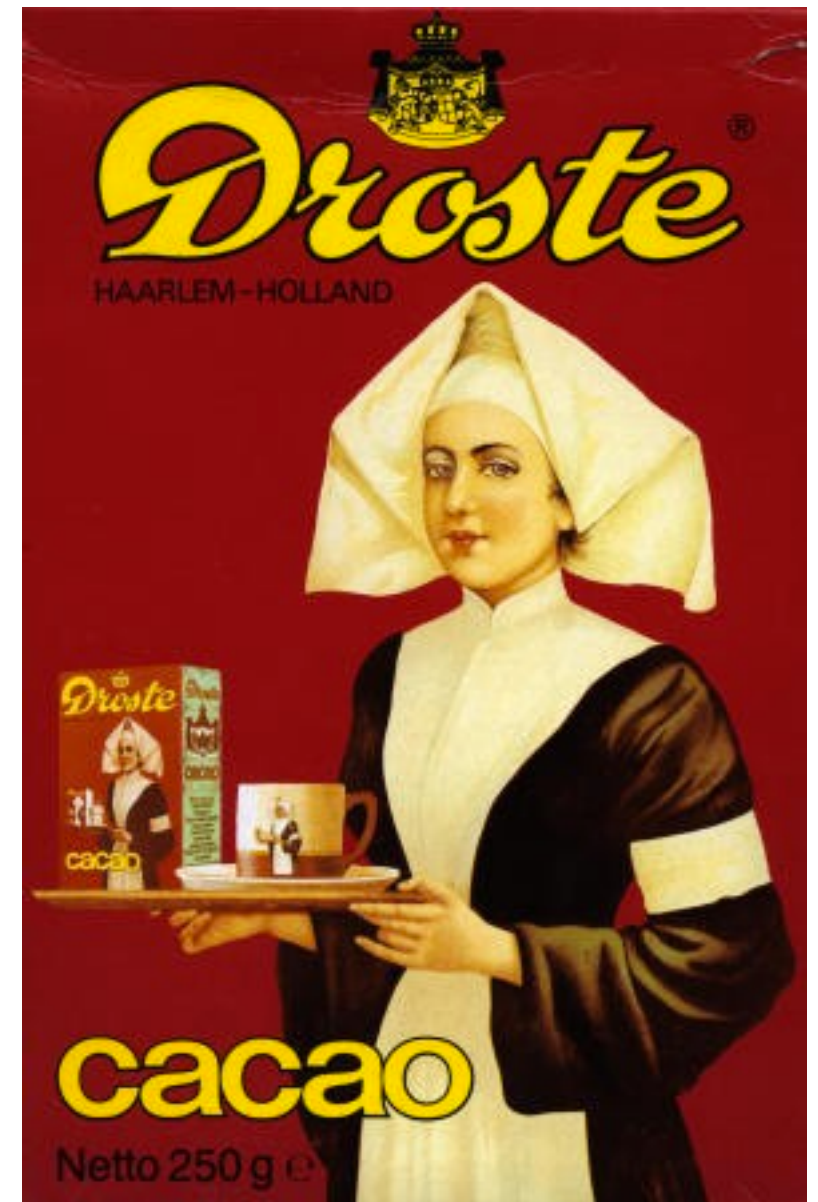
$$\frac{1}{1 \times 2} + \frac{1}{2 \times 3} + \frac{1}{3 \times 4} + \cdots + \frac{1}{n \times (n+1)} = \frac{n}{n+1}$$

- Proof.
 - We will prove this proposition with mathematical induction.
 - Let $P(n)$ be the predicate $1/(1 \cdot 2) + 1/(2 \cdot 3) + \dots + 1/(n \cdot (n+1)) = n/(n+1)$
 - Base step: We first prove $P(1)$ holds.
 - LHS is $1/2$, and RHS is $1/(1+1) = 1/2$
 - Inductive step: Then, we prove $P(k) \rightarrow P(k+1)$ for all integer $k \geq 1$
 - Let k be an arbitrary integer and $k \geq 1$.
 - Assume $P(k)$ holds. That is $1/(1 \cdot 2) + 1/(2 \cdot 3) + \dots + 1/(k \cdot (k+1)) = k/(k+1)$
 - We want to prove $P(k+1)$, namely, $1/(1 \cdot 2) + 1/(2 \cdot 3) + \dots + 1/(k \cdot (k+1)) + 1/((k+1) \cdot (k+2)) = (k+1)/(k+2)$
 - LHS can be reduced to $k/(k+1) + 1/((k+1) \cdot (k+2))$ following the assumption $P(k)$
 - The latter can be further reduced to $(k^2 + 2k + 1)/((k+2)(k+1))$, which equals to RHS.
- QED.

Recursive sequences

Recursion = Repeating itself

- recursive sequences
- recursive functions
- recursive data structures



Example

Examples

- Suppose $f(n) = n!$, where $n \in \mathbb{W}$. Then,

$$f(n) = \begin{cases} 1 & \text{if } n = 0, \\ n \cdot f(n-1) & \text{if } n \geq 1. \end{cases}$$

Closed-form formula: $f(n) = n \cdot (n-1) \cdot \dots \cdot 1$

Examples

Examples

- Suppose $f(n) = n!$, where $n \in \mathbb{W}$. Then,

$$f(n) = \begin{cases} 1 & \text{if } n = 0, \\ n \cdot f(n-1) & \text{if } n \geq 1. \end{cases}$$

Closed-form formula: $f(n) = n \cdot (n-1) \cdot \dots \cdot 1$

- Suppose $F(n) = n$ th Fibonacci number. Then,

$$F(n) = \begin{cases} 1 & \text{if } n = 0 \text{ or } 1, \\ F(n-1) + F(n-2) & \text{if } n \geq 2. \end{cases}$$

Closed-form formula: $F(n) = ?$

Example: Arithmetic sequence

Let a_0, a_1, a_2, \dots be the sequence defined recursively as follows: For all integers $k \geq 1$,

$$(1) \quad a_k = a_{k-1} + 2 \quad \text{recurrence relation}$$

$$(2) \quad a_0 = 1 \quad \text{initial condition.}$$

- Write out a_1, a_2, a_3, a_4 , and a_5
- Derive an explicit formula of the sequence
- Confirm the explicit formula satisfies the recursive definition

Example: Geometric sequence

You deposit \$100,000 in a bank account for a 3% interest compounded annually. How much will you get after 21 years?

Solution

- Suppose $A_k = \text{Amount in your account after } k \text{ years}$. Then,
$$A_k = \begin{cases} 100,000 & \text{if } k = 0, \\ (1 + 3\%) \times A_{k-1} & \text{if } k \geq 1. \end{cases}$$
- Solving the recurrence by the method of iteration, we get

$$A_k = ((1.03)^k \cdot 100,000) \text{ dollars}$$

- Homework: **Confirm the explicit formula**
- **satisfies the recursive definition**

$$A_{21} = ((1.03)^{21} \cdot 100,000) \approx 186,029.46 \text{ dollars}$$

Exercise 1

- Consider the recursive sequence below. Find its explicit form and prove your answer, namely Confirm the explicit formula satisfies the recursive definition

$$a_k = ka_{k-1}, \text{ for all integers } k \geq 1$$

$$a_0 = 1$$

Solution

We have: $a_k = k \cdot a_{k-1} = k \cdot (k-1) \cdot a_{k-2} = k(k-1)(k-2) \cdots 1 \cdot a_0 = k!$

- To prove it indeed satisfies the recursive definition, we plug it to the recursive definition:
 - The explicit form clearly satisfies $a_0=1$
 - The explicit form also satisfies $a_k = k a_{k-1}$, since the left is $k!$, the right is $k \cdot (k-1)!$ which is also $k!$

Exercise 2

- Consider the recursively defined sequence below. Find its explicit form and prove your answer.

$$b_k = \frac{b_{k-1}}{1 + b_{k-1}}, \text{ for all integers } k \geq 1$$
$$b_0 = 1$$

Solution

- We have $b_0=1$, $b_1=1/2$, $b_2=1/3$, $b_3=1/4$... So a possible explicit form is $b_n=1/(n+1)$
- To prove it indeed satisfies the recursive definition, we plug it to the recursive definition:
 - The explicit form clearly satisfies $b_0=1$
 - The explicit form also satisfies $b_k = b_{k-1} / (1+b_{k-1})$, since LHS= $1/(k+1)$,
RHS = $1/(k+1) / (1 + 1/(k+1)) = 1/(k+1)$

Exercise 3

- Consider the recursively defined sequence below. Find its explicit form and prove your answer.

$$c_k = 3c_{k-1} + 1, \text{ for all integers } k \geq 2$$

$$c_1 = 1$$

Solution

- We have $c_1=1$, $c_2=4$, $c_3=13$, $c_4=40$... So a possible explicit form is $c_n=(3^n-1)/2$
- To prove it indeed satisfies the recursive definition, we plug it to the recursive definition:
 - The explicit form clearly satisfies $c_1=1$
 - The explicit form also satisfies_____, since $LHS= \underline{\hspace{1cm}}$ $RHS = \underline{\hspace{1cm}} = \underline{\hspace{1cm}}$

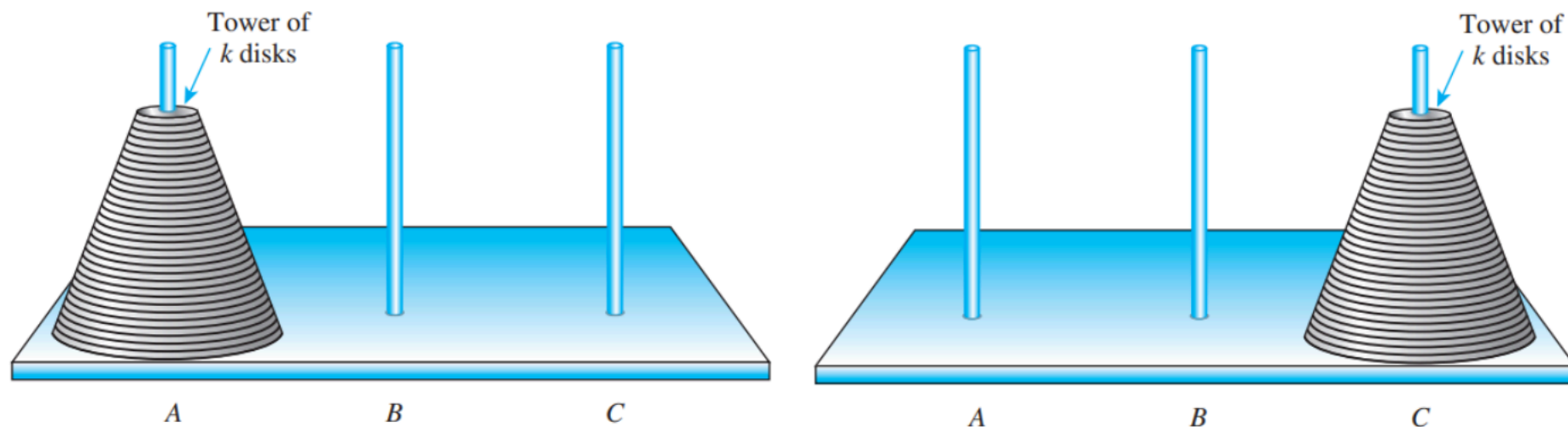
Application

Example: Towers of Hanoi

Problem

- There are k disks on peg 1. Your aim is to move all k disks from peg 1 to peg 3 with the minimum number of moves. You can use peg 2 as an auxiliary peg. The constraint of the puzzle is that at any time, you cannot place a larger disk on a smaller disk.

What is the minimum number of moves required to transfer all k disks from peg 1 to peg 3?



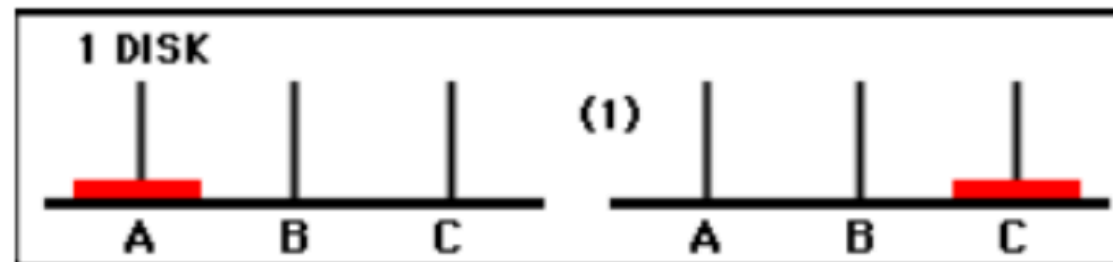
Demo: <https://www.mathsisfun.com/games/towerofhanoi.html>

Example: Towers of Hanoi

Solution

Suppose $k = 1$. Then, the 1-step solution is:

1. Move disk 1 from peg A to peg C .

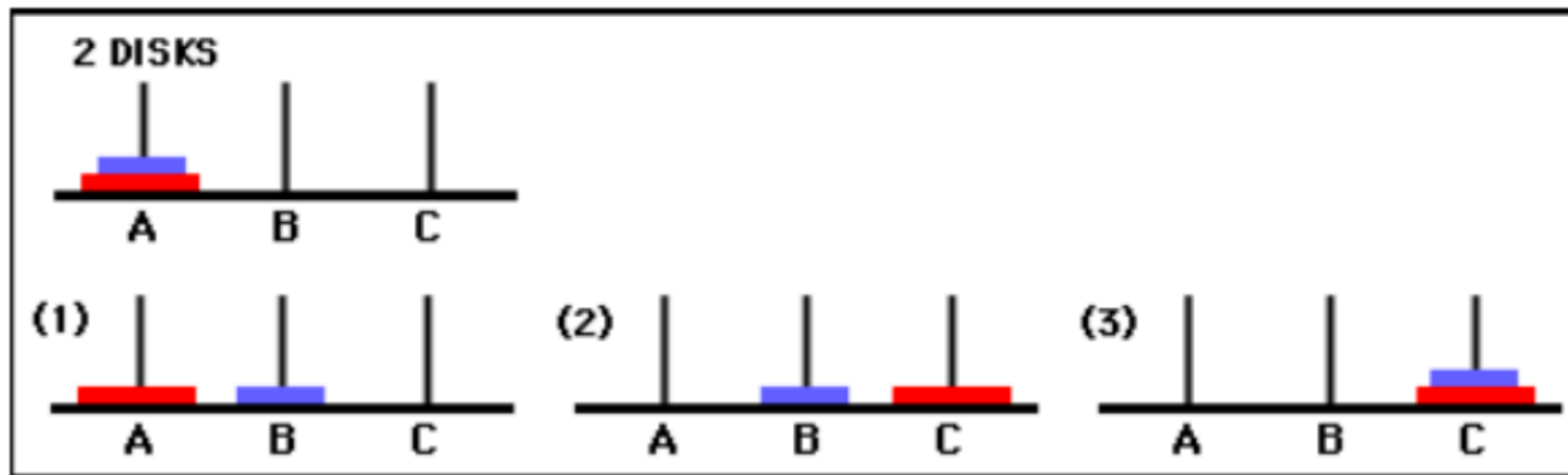


Example: Towers of Hanoi

Solution

Suppose $k = 2$. Then, the 3-step solution is:

1. Move disk 1 from peg A to peg B .
2. Move disk 2 from peg A to peg C .
3. Move disk 1 from peg B to peg C .

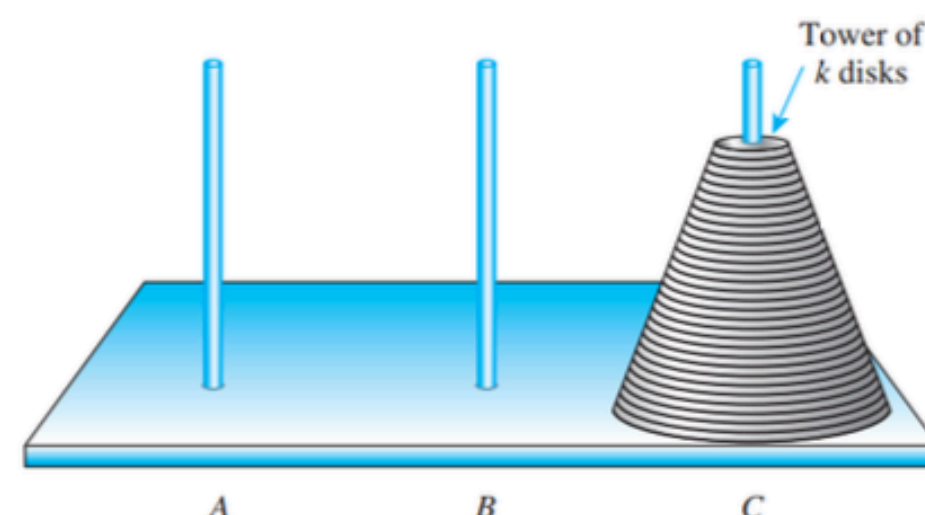
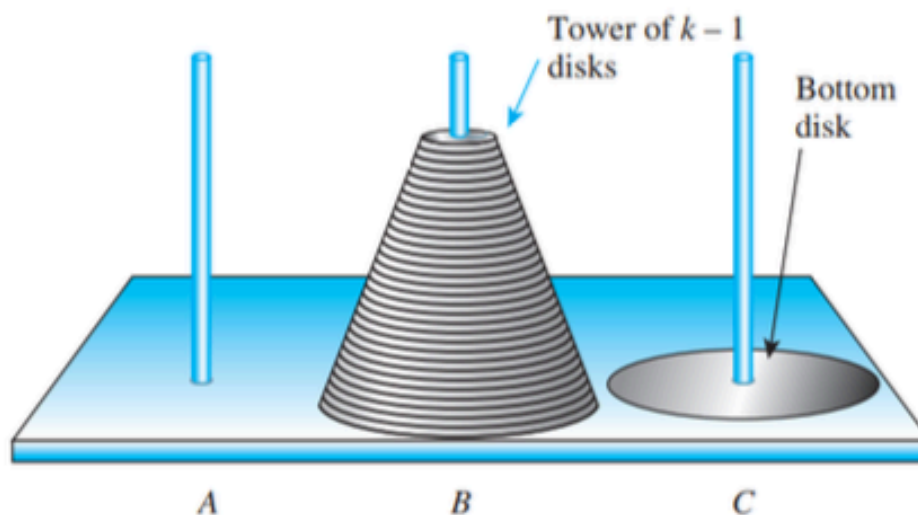
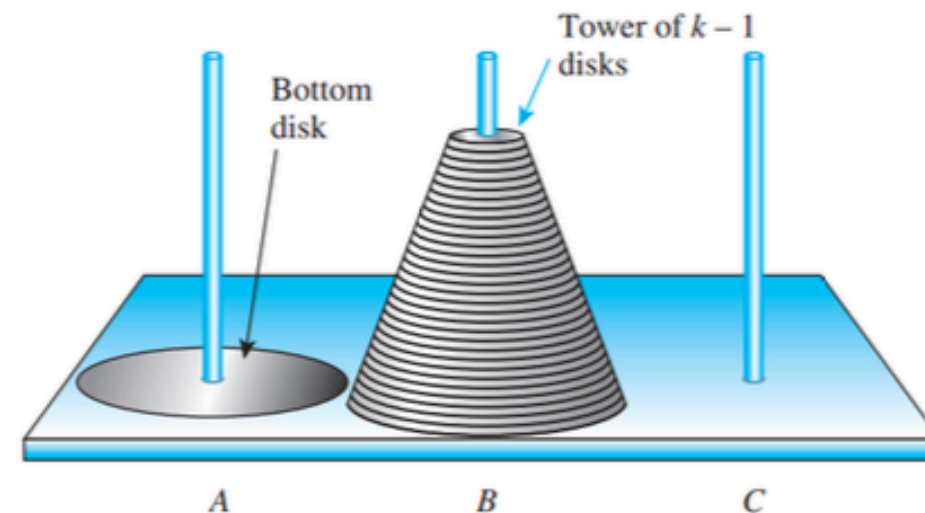
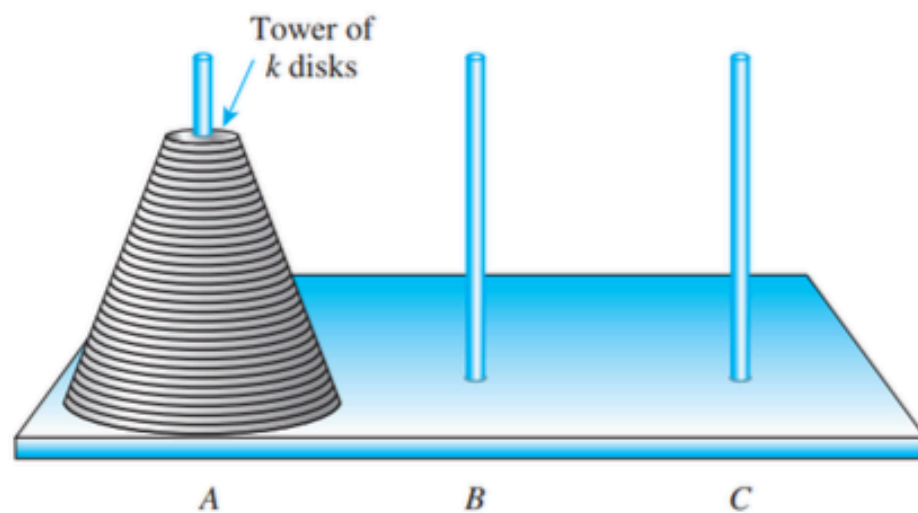


Example: Towers of Hanoi

Solution

For any $k \geq 2$, the recursive solution is:

1. Transfer the top $k - 1$ disks from peg A to peg B .
2. Move the bottom disk from peg A to peg C .
3. Transfer the top $k - 1$ disks from peg B to peg C .

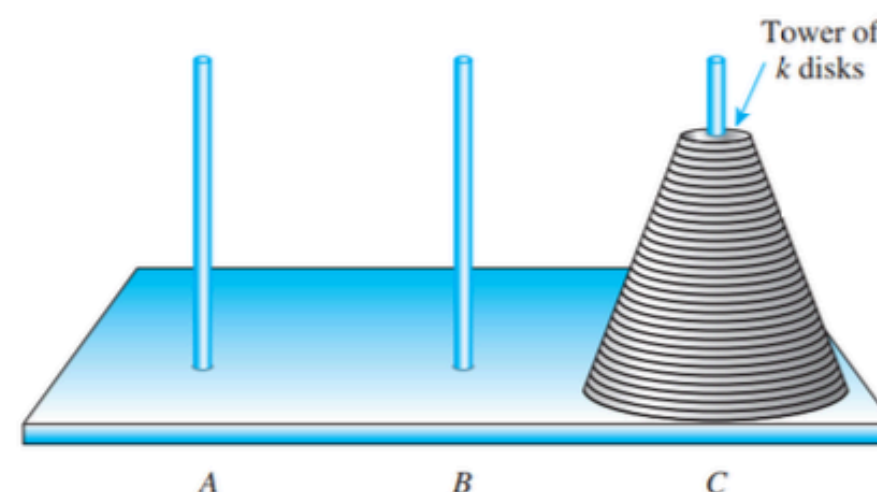
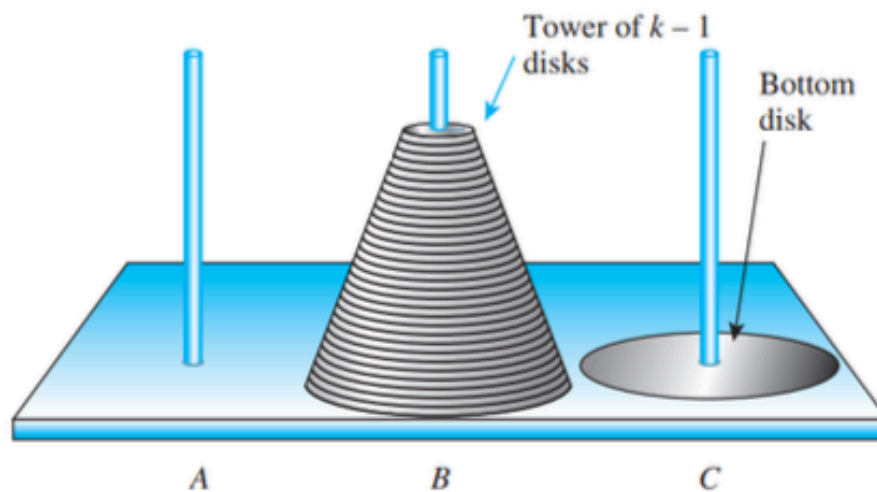
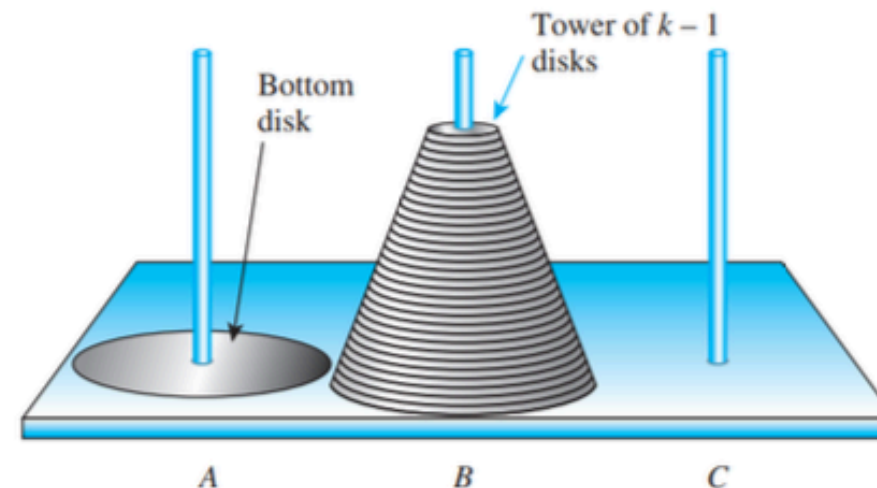
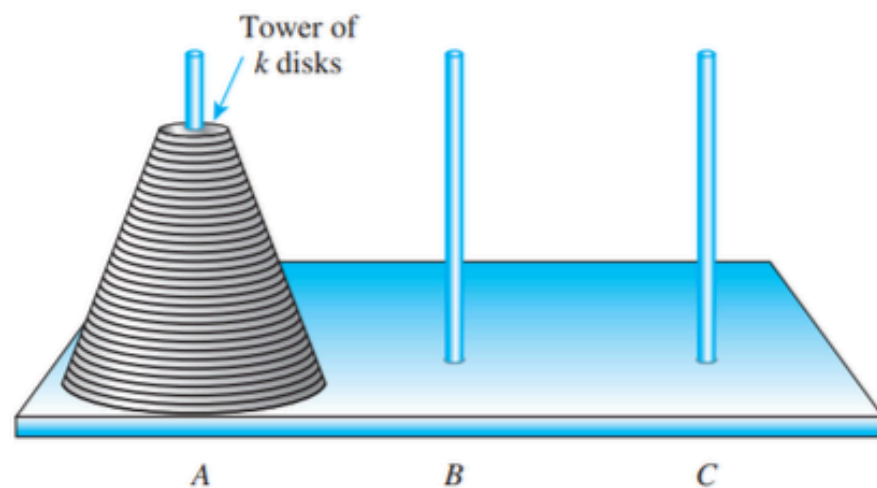


Example: Towers of Hanoi

TOWERS-OF-HANOI(k, A, C, B)

1. if $k = 1$ then
2. Move disk k from A to C .
3. elseif $k \geq 2$ then
4. TOWERS-OF-HANOI($k - 1, A, B, C$)
5. Move disk k from A to C .
6. TOWERS-OF-HANOI($k - 1, B, C, A$)

[Code](#) [Demo](#)



Example: Towers of Hanoi

Solution (continued)

- Let $M(k)$ denote the **minimum number of moves** required to move k disks from one peg to another peg. Then

$$M(k) = \begin{cases} 1 & \text{if } k = 1, \\ 2 \cdot M(k-1) + 1 & \text{if } k \geq 2. \end{cases}$$

- Solving the recurrence by the method of iteration, we get

$$M(k) = 2^k - 1$$

▷ How?

Why minimum?

https://proofwiki.org/wiki/Tower_of_Hanoi