

# **CSE216**

# **Programming Abstraction**

**State University of New York, Korea**

# Why Harvard, MIT, Stanford, Cambridge all teach functional programming



<https://youtu.be/6APBx0WsgeQ>

# Another good explanation (to watch at home)

## Functional languages predict the future

- Garbage collection  
*Java [1995], LISP [1958]*
- Generics  
*Java 5 [2004], ML [1990]*
- Higher-order functions  
*C#3.0 [2007], Java 8 [2014], LISP [1958]*
- Type inference  
*C++11 [2011], Java 7 [2011]*
- What's next?





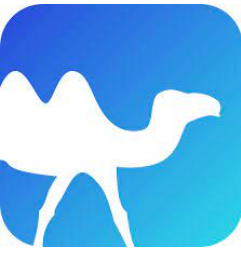
# Install Ocaml

- Official guide: <https://ocaml.org/docs/up-and-running>
- SBU Guide: <https://sites.google.com/cs.stonybrook.edu/cse216/lectures?authuser=0> (may need SBU credentials)
- Once installed, get into the toplevel by running “ocaml”. In the toplevel, run: *print\_endline “hello”;;*

```
OCaml version 4.14.0  
Enter #help;; for help.
```

```
# print_endline "hello";;  
hello  
- : unit = ()  
_
```

- If nothing works, use TryOcaml for now: <https://try.ocamlpro.com/>.



# Toplevel Demo

```
# 42;;
```

```
- : int = 42
```

```
# let f x y = x + y;;
```

```
val f : int -> int -> int = <fun>
```

```
# f 3 ;;
```

```
- : int -> int = <fun>
```

```
# f 3 4 ;;
```

```
- : int = 7
```

```
# #use "hello.ml";;
```

```
hello world!
```

```
- : unit = ()
```

# Ocaml Basics

# Ocaml program = Definitions + Expressions

```
let x = 1 + 2;;
```

**Definition**

```
print_int x;;
```

**Expression**

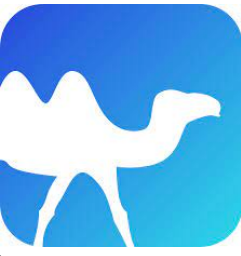
```
let y = x * x;;
```

**Definition**

```
print_int y;;
```

**Expression**





# Let-binding is a definition

```
# let x = 5 ;;  
val x : int = 5
```

```
# let f x y = x + y ;;  
val f : int -> int -> int = <fun>
```

```
# let f = (fun x y -> x+y) ;;  
val f : int -> int -> int = <fun>
```

```
# let a = f 2 3 ;;  
val a : int = 5
```

```
# let g = f 2 ;;  
val g : int -> int = <fun>
```



# A note on “let x = e”

Java

```
final int x = 42;
```

OCaml

```
let x = 42
```

- **let x = e** defines a global variable
- Differences with regard to Java variables:
  - 1 necessarily initialized
  - 2 type not declared but inferred
  - 3 cannot be modified in-place

# A note on functions

```
# let f = (fun x y -> x+y);;
```

```
val f : int -> int -> int = <fun>
```

```
# let f x y = x + y;;
```

```
val f : int -> int -> int = <fun>
```

```
# f 2 3;;
```

```
- : int = 5
```

```
# f 2 ;;
```

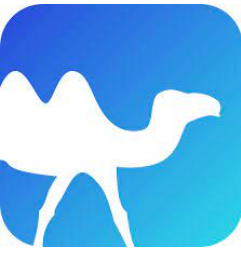
```
- : int -> int = <fun>
```

- Similar to  $\lambda x. \lambda y. x + y$

**Everything other than the  
definition is an expression**

# Constants and their operations are expressions

```
# 42;;  
- : int = 42  
# 42.8;;  
- : float = 42.8  
# "hello";;  
- : string = "hello"  
# true;;  
- : bool = true  
# 'a';;  
- : char = 'a'  
# 1+1;;  
- : int = 2  
# 1+1=2;;  
- : bool = true
```



# Functions and their operations are expressions

```
# fun x y -> x +. y;;
```

```
- : float -> float -> float = <fun>
```

```
# (fun x y -> x +. y) 2.0 3.7;;
```

```
- : float = 5.7
```

```
# (fun x y -> x +. y) 2.0;;
```

```
- : float -> float = <fun>
```

Exercise: What happens if we do `(fun x y -> x +. y) 2 3.7;;`

# Expression `let ... in ...`

`let x = e1 in e2` is an expression

its type and value are those of `e2`,  
in an environment where `x` has the type and value of `e1`

example

```
let x = 1 in (let y = 2 in x + y) * (let z = 3 in x * z)
```

Semantics of “`let...in`” is substitution in lambda calculus

**Exercise:** What is the value and type of above?

# A note on the let ... in expression

in C or Java, the scope of a local variable extends to the **bloc**:

```
{  
  int x = 1;  
  ...  
}
```

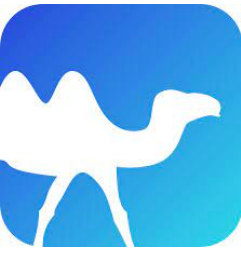
in OCaml, a local variable is introduced with **let in**:

```
let x = 10 in x * x
```

as for a global variable:

- necessarily initialized
- type inferred
- immutable
- but **scope limited to the expression following in**





# Expression

## if c then e1 else e2

```
# let f x y = x + y;;
```

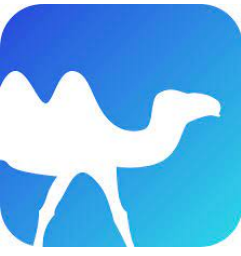
```
val f : int -> int -> int = <fun>
```

```
# if f 4 5 == 9 then "hello" else "world";;
```

```
- : string = "hello"
```

e1 and e2 must be of the same type

- Evaluate the boolean condition c.
- If c is true, evaluate the expression e1 and return its value as the result of the if expression.
- If c is false, evaluate the expression e2 and return its value as the result of the if expression.



# Unit type

Expressions with no meaningful values have type **unit**.

Example: `print_string "okay"`

The type has a single value, written `()`

it is the type given to the **else** branch when it is omitted

```
# let x = 5;;
val x : int = 5
# if x > 0 then print_string "okay";;
okay- : unit = ()
# if x > 0 then 100;
Error: This expression has type int but an expression was expected of type
      unit
      because it is in the result of a conditional with no else branch
.. ■
```

# Expression e1;e2

- The first expressions e1 needs to be of type unit
- The evaluation result of the whole is that of e2

```
# print_string "hello"; 3 ;;
```

```
hello
```

```
- : int = 3
```

```
# 4①; 3 ;;
```

```
- : int = 3
```

```
① Warning : this expression should have type unit.
```

# Exercises

- Evaluate  $\text{let } x = 3 \text{ in let } y = x + 2 \text{ in } y * y$
- Evaluate  $\text{let } x = 3 \text{ in } (\text{let } y = x + 1 \text{ in } y) * (\text{let } z = x * x \text{ in } z)$

# Exercises

- For each of these expressions, what is its value and type. Write the value of a function by “<fun>”:
  - `let x = 3 in x * x`
  - `print_string "hello"`
  - `print_string`
  - `let f x y = x + y in f 5 6`
  - `let x = 5 in (if x > 0 then "pos" else "neg")`
  - `let f x y = x + y in f 5`

# Summary so far

- Compiling/Interpreting an Ocaml program
- Ocaml program = Definitions + Expressions
- `let x = 3` is a definition
- Everything else is an expression, thus has a value and type
- `let x = 3 in x+5` is an expression of value 8, type `int`
- `if x then 3 else 5` is also an expression, of type `int`