# CSE215 Foundations of Computer Science

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#### Plan

- Today:
  - Homework 04
  - Review of lambda calculus
  - Mini-mock
- Midterm 1 on Thursday 04/03 12h30-13h50 pm at B103 (this room)
- Covering everything until last lecture. Yes, Ocaml is covered in midterm 1
- Any amount of physical notes are allowed.
- E-devices are not allowed.

#### Homework 04

#### Review on lambda calculus

- (λ x.x) a
- (λ x.y) a
- (λ x.xy) a
- (λ x. yx) a
- (λ x. xx) a
- (λ x. yy) a

- (λ x.x) a b
- (λ x.y) a b
- (λ x.xy) a b
- (λ x. yx) a b
- (λ x. xx) a b
- (λ x. yy) a b

- (λ x.x) λ a. b
- (λ x.y) λ a. b
- (λ x.xy) λ a. b
- (λ x. yx) λ a. b
- (λ x. xx) λ a. b
- (λ x. yy) λ a. b

- (λ x.x) x
- (λ x.y) x
- (λ x.xy) x
- (λ x. yx) x
- (λ x. xx) x
- (λ x. yy) x

- (λ x.x) x y
- (λ x.y) x y
- (λ x.xy) x y
- (λ x. yx) x y
- (λ x. xx) x y
- (λ x. yy) x y

(λz.z) (λz.z z) (λz.z q)

(λs.λq.s q q) (λa.a) b

(λs.λq.s q q) (λq.q) q

((λs.s s) (λq.q)) (λq.q)

(λ x.λy.x) x y

(λ x. λ y. λ z. y (w y x)) λ s. λ z. z

#### Mini-Mock Midterm1

- Regular expressions
- Context-free grammar
- lambda calculus
- Ocaml

### Regular expression

In our class, we have studied the core regular expressions and some abbreviations based on those core regular expressions.

r	Meaning	Language $\mathcal{L}(r)$
a	Character a	{"a"}
arepsilon	Empty string	{""}
$r_1 r_2$	$r_1$ followed by $r_2$	$\{s_1s_2 \mid s_1 \in \mathcal{L}(r_1), s_2 \in \mathcal{L}(r_2)\}$
r*	Zero or more r	$\{s_1 \ldots s_n \mid s_i \in \mathcal{L}(r), n \geq 0\}$
$r_1   r_2$	Either $r_1$ or $r_2$	$\mathcal{L}(r_1)\cup\mathcal{L}(r_2)$

Abbrev.	Meaning	Expansion
[aeiuo]	Set	a e i o u
[0-9]	Range	0 1  8 9
[0-9a-z]	Ranges	0 1  8 9 a b  y z
<i>r</i> ?	Zero or one r	r arepsilon
$r^+$	One or more r	rr*

Give a regular expression over {a, b} that has aab as a substring

## Regular expression

(1) Write a regular expression pattern to match valid music notes according to the criteria below: A music note is represented by a capital letter A to G (inclusive) followed by an optional symbol: sharp (#), flat (b), or natural (n).

```
Example valid inputs: C, D#, Fb, Gn

Example invalid inputs: H, C##, Fm, C#b
```

Note. The sharp symbol ("#") is not a special character in regular expressions. So you do *not* need to escape it with a backslash.

### Context-free grammar

(2)\*\* What is the language generated by the following grammar? Select one answer from the four choices.

```
S -> aSbb | ε
```

- A. The set of all strings that start with 'a' and end with two 'b'.
- B. The set of all strings that contain twice as many 'b's as 'a's.
- C. The set of all strings that contain an odd number of 'a's followed by an even number of 'b's.
- D. The set of all strings that contain n 'a's followed by m 'b's, where m = 2n > = 0

#### lambda calculus

- reduce the lambda term to the normal form. No need to write the reduction in details.
  - $(\lambda x.(x y))(\lambda z.z)$

#### lambda calculus

- reduce the lambda term to the normal form. No need to write the reduction in details.
  - $(\lambda x.(\lambda y.x y)x)(\lambda z.y)$

#### Ocaml

- Give the type of the following OCaml expressions:
  - print\_string
  - Let f = fun x y -> x+y in f 3
  - let  $f x = x^*$ . 3.14 in f 5.6