CSC3150 Assignment 4 Report

Environment

I use the centos OS on the computer of TC301

How to compile my program

- 1. "make" to makefiles
- 2. "./fs" to execute
- 3. "makeclean" to clean

How did I Design My Program

My FCB structure:

```
/*
    * my FCB structure
    * |0|1|2|3|4|5|6|7|8|9|10|11|12|13|14|15|16|17|18|19| 20|21 | 22|23 | 24 | 25|26 | 27 |28 |29 | 30 | 31 |
    * | file name | location | size | create_t|modify_t|
    */
struct My_FCB
{
    char file_name[20];
    u32 location;
    u32 size;
    int create_time;
    int modified_time;
}Current_FCB;
```

function: fs_open

in my fs_open function:

- 1. search for the location of file name s if it exists, if it does not exist in the FCB we return -1 to denote.
- 2. in the read mode if we find the match of file s we simply record the location of its block in FCB by variable file_exist and find its start location by looking at its 20th to 23rd bits, in the end we return its pointer start_location in the superblock. recall that in my program the info of start_location is stored from 20-13 in each FCB.
- 3. in the write mode if no such file name can be found, we will just create a new zero byte file. If at the end of the day we find the match we clear the old content in the storage as will as modified the superblock to 0. return the start_location pointer in the FCB

function: fs_read:

in my fs_read function:

- 1. we want to check if fp is a valid. if no, return an error.
- 2. so for the starting area to the size we put the content of fs->volume[fp * 32 + i + fs->FILE_BASE_ADDRESS] to the output buffer

```
__device__ void fs_read(FileSystem *fs, uchar *output, u32 size, u32 fp)
{
    if(fp == -1){
        printf("%d", size);
        printf("error\n");
    }

    for (int i = 0; i < size; i++) {
        output[i] = fs->volume[fp * 32 + i + fs->FILE_BASE_ADDRESS];
    }
}
```

function: fs write:

in my write function we can divide this problem into two cases:

1.the original file space is larger than the size of input written file:

we clear the file contents and directly input the things in input buffer to the storage, interate from 0 to size, then update the start location in the super block.

2. the space in the storage is not enough: allocate a new area for the input file, load the data to the new place as well as update the FCB and superblock

in the end if externel segmentation occured during this process, remove them as will as update the FCB and Superblock

function: fs_gsys LS_D LS_S:

first we search where the file system is going to end

then according to the option we choose we do bubble sort, then we display the file name by the cretirea of the sort

function: fs_gsys rm:

- 1. we find wether the file is exist in the file system. if yes we are trying to find its location in the file system.
- 2. find where its file start in the storage and its size
- 3. clear it from the storage as well as FCB and Superblock
- 4. Then we manage the externel segementation

```
| Implement rm operation here */
```

Sample Output:

Input:

```
cuhksz@TC-301-34:~/Desktop/zhouliang
File Edit View Search Terminal Help
)000040: 6363 6363 6363 6363 6363 aa13 68df cccccccccc.h.
)000050: 1309 5756 d319 c270 5b39 ab09 lac2 6fad ..WV...p[9....o.
)000060: 6852 8014 bab6 12b9 f2fc da9d 10c2 fc23 hR.....#
)000070: ccd4 f820 6d3b 10c8 743b d18f fdc0 bce5 ... m;..t;.....
)000080: 133d fa4e f30d 0866 89e2 8399 2581 3c71 .=.N...f....%.<q
)000090: 5635 91c3 ef21 0c65 5c5d 745a 1f31 4132 V5...!.e\]tZ.1A2
)0000a0: 6fbb 80e2 4808 4ad1 6acd ea90 ce27 0225 o...H.J.j....'.%
)0000b0: db14 68cb 3574 b092 5225 ec71 57ad 2346 ..h.5t..R%.qW.#F
)0000c0: 69a4 2931 acf2 0318 41ed 2810 942a b570 i.)1....A.(..*.p
)0000d0: 3ele bcf3 136d 8665 93f2 d66a alf9 b08a >...m.e...j....
)0000e0: 1e59 bc4b 4d3f e28e 2e0b 9ec2 b4d3 b373 .Y.KM?....s
)0000f0: f270 6706 5d6d 6bf0 60c1 da81 3b8b 8c5a .pg.]mk.`...;..Z
)000100: e549 25b2 0808 4136 925f 7947 342d 3aa6 .I%...A6. yG4-:.
)0000110: 9da1 acfa 0f97 6cef d847 7114 53fd edb8 .....l..Gq.S...
)0000120: c613 6bcf 9a2c 852d 8bfe f43f 2c2f e549 ..k..,.-...?,/.I
)0000130: 5112 c4df 29b0 cf02 77c1 96ca 3f04 8386 Q...)...w...?...
)000140: 976e d532 9a5b dfa6 da54 e507 034c d054 .n.2.[...T...L.T
)000150: dd95 3508 c584 893d 4620 0986 a40c 0d3c ..5....=F .....<
)000160: fae2 ed95 bd4d bb98 a122 2025 6ef0 794c .....M..." %n.yL
                                                  1,1
                                                              Top
```

Output: test case1:

```
t.txt
b.txt
===sort by file size===
t.txt 32
b.txt 32
===sort by file size===
t.txt 32
b.txt 32
b.txt 12
===sort by modified time===
b.txt
t.txt
===sort by modified time===
```

Output: test case2:

```
(base) [cuhksz@TC-301-10 zhouliang]$ ./fs
---sort by time---
b.txt
t.txt
---sort by file size---
t.txt 32
b.txt 32
---sort by file size---
t.txt 32
b.txt 12
---sort by time---
b.txt
t.txt
---sort by file size---
b.txt 12
---sort by file size---
*ABCDEFGHIJKLMNOPQR 33
)ABCDEFGHIJKLMNOPQR 32
(ABCDEFGHIJKLMNOPQR 31
'ABCDEFGHIJKLMNOPQR 30
&ABCDEFGHIJKLMNOPQR 29
%ABCDEFGHIJKLMNOPQR 28
$ABCDEFGHIJKLMNOPQR 27
#ABCDEFGHIJKLMNOPQR 26
"ABCDEFGHIJKLMNOPQR 25
!ABCDEFGHIJKLMNOPQR 24
b.txt 12
---sort by time---
*ABCDEFGHIJKLMNOPQR
)ABCDEFGHIJKLMNOPQR
(ABCDEFGHIJKLMNOPQR
'ABCDEFGHIJKLMNOPQR
&ABCDEFGHIJKLMNOPQR
```

Output: test case3:

```
ABCDEFGHIJKLMNOPQR
)ABCDEFGHIJKLMNOPOR
(ABCDEFGHIJKLMNOPQR
'ABCDEFGHIJKLMNOPQR
&ABCDEFGHIJKLMNOPQR
b.txt
===sort by file size===
~ABCDEFGHIJKLM 1024
ABCDEFGHIJKLM 1023
ABCDEFGHIJKLM 1022
{ABCDEFGHIJKLM 1021
zABCDEFGHIJKLM 1020
VABCDEFGHIJKLM 1019
xABCDEFGHIJKLM 1018
wABCDEFGHIJKLM 1017
VABCDEFGHIJKLM 1016
uABCDEFGHIJKLM 1015
tABCDEFGHIJKLM 1014
sABCDEFGHIJKLM 1013
rABCDEFGHIJKLM 1012
qABCDEFGHIJKLM 1011
pABCDEFGHIJKLM 1010
oABCDEFGHIJKLM 1009
nABCDEFGHIJKLM 1008
mABCDEFGHIJKLM 1007
1ABCDEFGHIJKLM 1006
kabcdefghijklm 1005
jABCDEFGHIJKLM 1004
iABCDEFGHIJKLM 1003
hABCDEFGHIJKLM 1002
gABCDEFGHIJKLM 1001
fabcdefghijklm 1000
eABCDEFGHIJKLM 999
dabcdefghijklm 998
cABCDEFGHIJKLM 997
babcdefghijklm 996
aABCDEFGHIJKLM 995
ABCDEFGHIJKLM 994
ABCDEFGHIJKLM 993
^ABCDEFGHIJKLM 992
ABCDEFGHIJKLM 991
```

The Problem I met

- 1. How should I do free space management we can do remove those externel segmentation while they occured, in my program:
 - 1. I remove a certain portion of file storage by setting corresponding volume[i] to be 0, and move the predessessor file upward.
 - 2. we set the corresponding area in superblock to be 0

- 3. update the file_start_location and get its quotient and reminder.
- 4. update the superblock and change the contents in FCB

Free Space Management

- File system maintains free-space list to track available blocks/clusters
 - (Using term "block" for simplicity)
- Bit vector or bit map (n blocks)



For example, consider a disk where blocks 2, 3, 4, 5, 8, 9, 10, 11, 13, 17, 18 are free and the rese blocks are allocated.

The free-space bit map would be 0011110011111100011...

- 2. I stuck in the problem when the size are the same, sort its modified time for those files with same sizes do the sort algorithm again for these files, but keep other files static
- 3. How to construct a tree-structure directory in bonus task:

In my program I denote the file is a directory or file in the 28th and 29th bit to denote who is the directory/file's parent. this makes the tree structure feasible the structure of tree directory is:

```
struct directory{
   char name[20];
   u32 sibling;
   u32 parent;
   u32 child;
   u32 size;
   u32 modified date;
}directory[4];

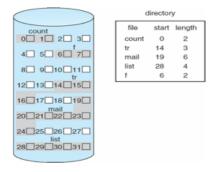
/*
   * my FCB structure of tree directory
   * |0|1|2|3|4|5|6|7|8|9|10|11|12|13|14|15|16|17|18|19| 20|21 | 22|23 | 24 | 25|26 | 27 |28 | 29 | 30| 31 |
   * file name | location | size | index | parent|modify_t |
   */
```

What I learnt from the task

- 1. The relationship between Superblock and contents of file:
 each bit in the volume[i] in superblock correpsonding to a unit block in the contents of files.
 in my work in the superblock we use 1 to denote the file exist, 0 to denote it doesn't exist.
- 2. How to do contiguous allocation

Contiguous Allocation

- An allocation method refers to how disk blocks are allocated for files:
- Contiguous allocation each file occupies set of contiguous blocks



basicly we should know where the file start, and allocate the corresponding slots according to its size. when we are going to remove the file we put the predecessor file upwards have the same start_block with the file to be deleted.