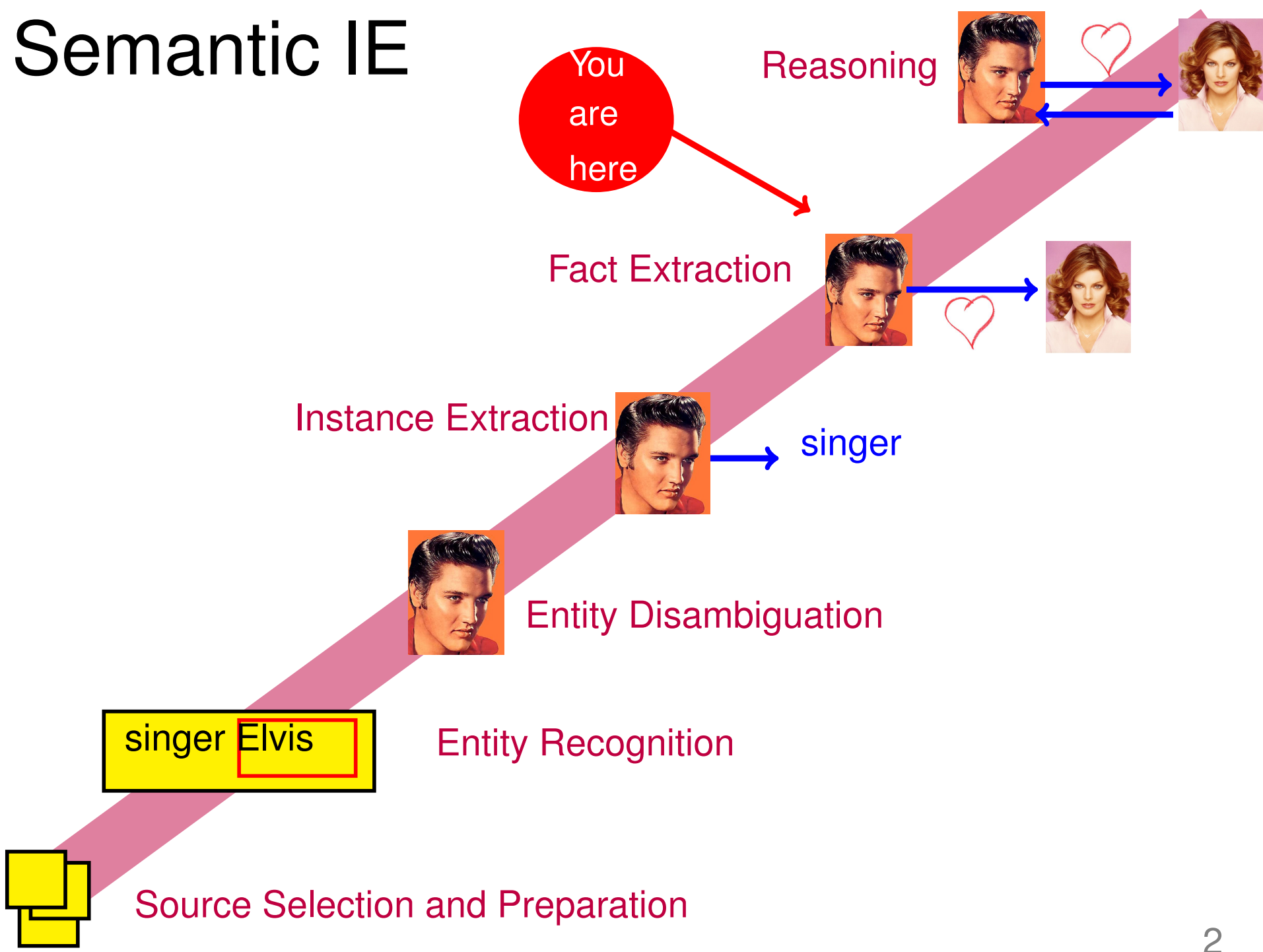


Information Extraction by Reasoning

Fabian M. Suchanek

Semantic IE

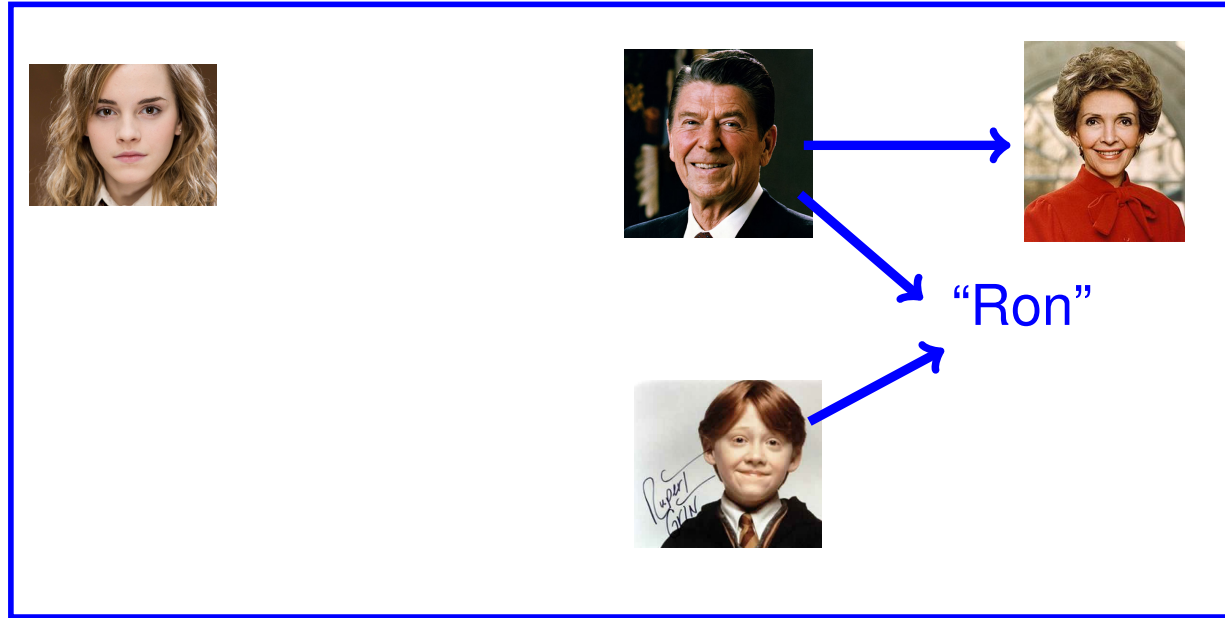
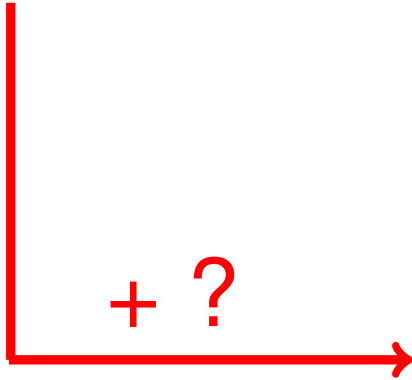


Problems in extending a KB

“Hermione is married to Ron”

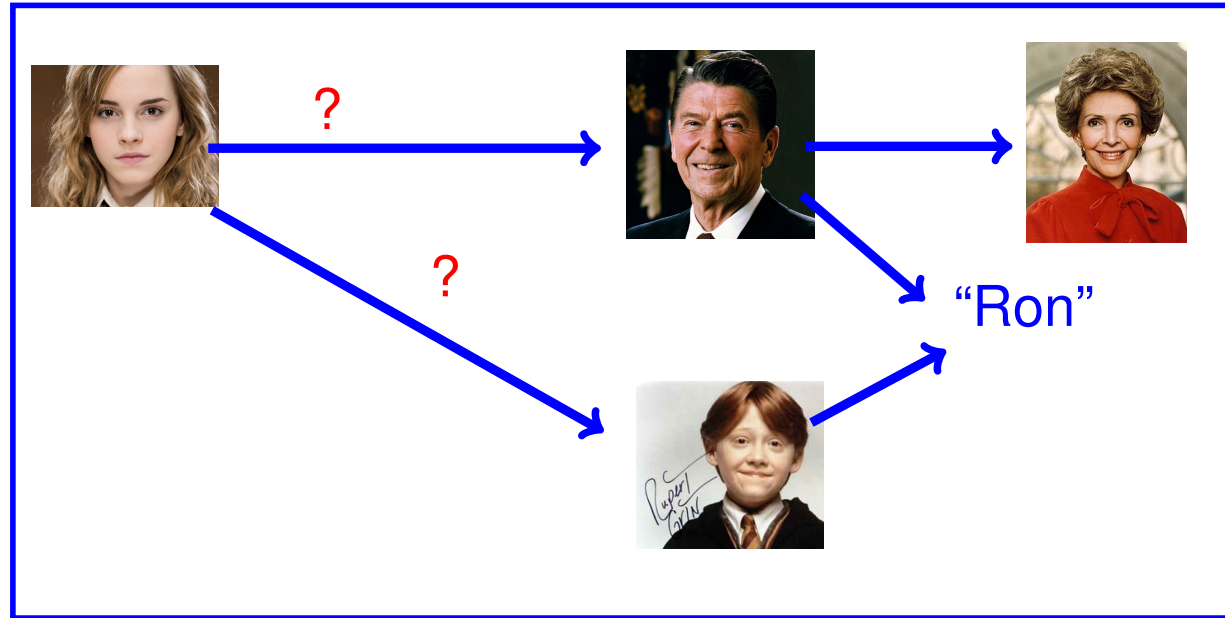
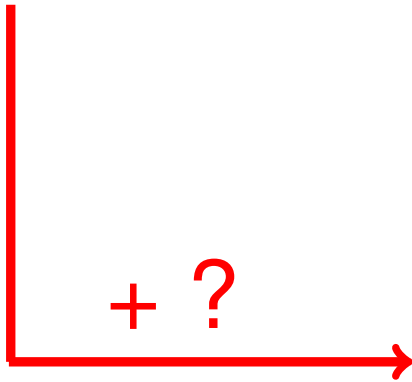
Problems in extending a KB

“Hermione is married to Ron”



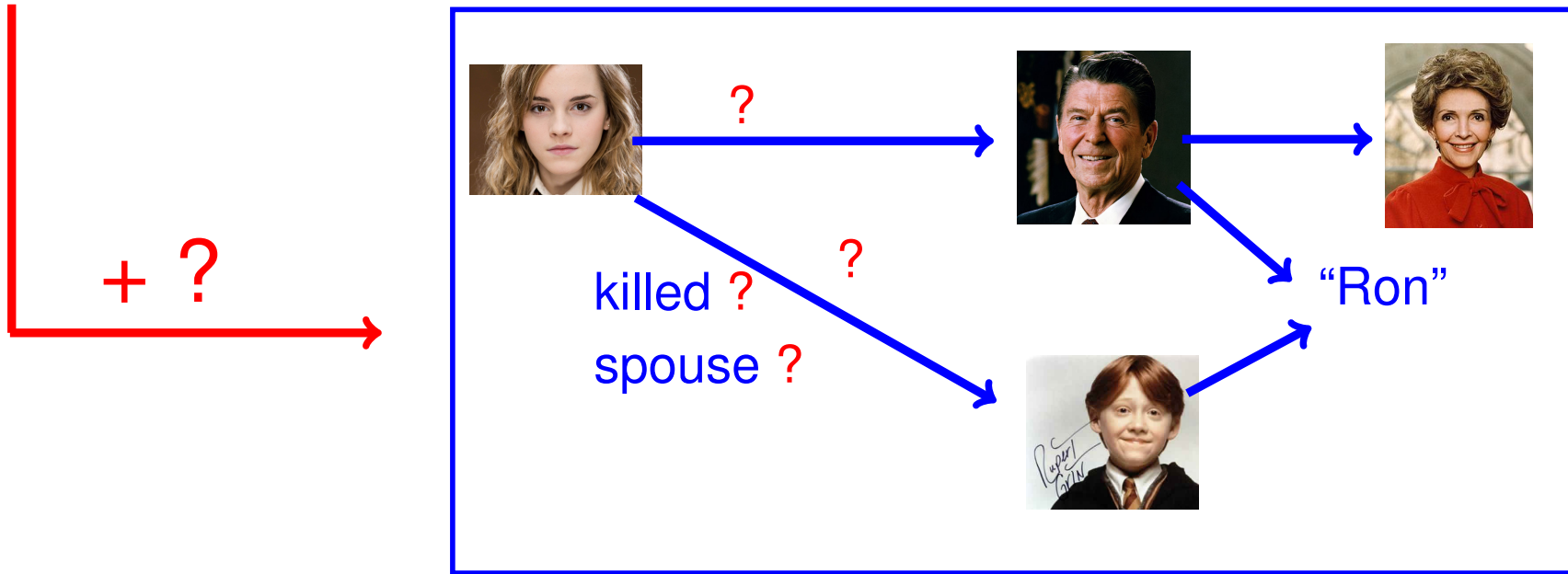
Problems in extending a KB

“Hermione is married to Ron”



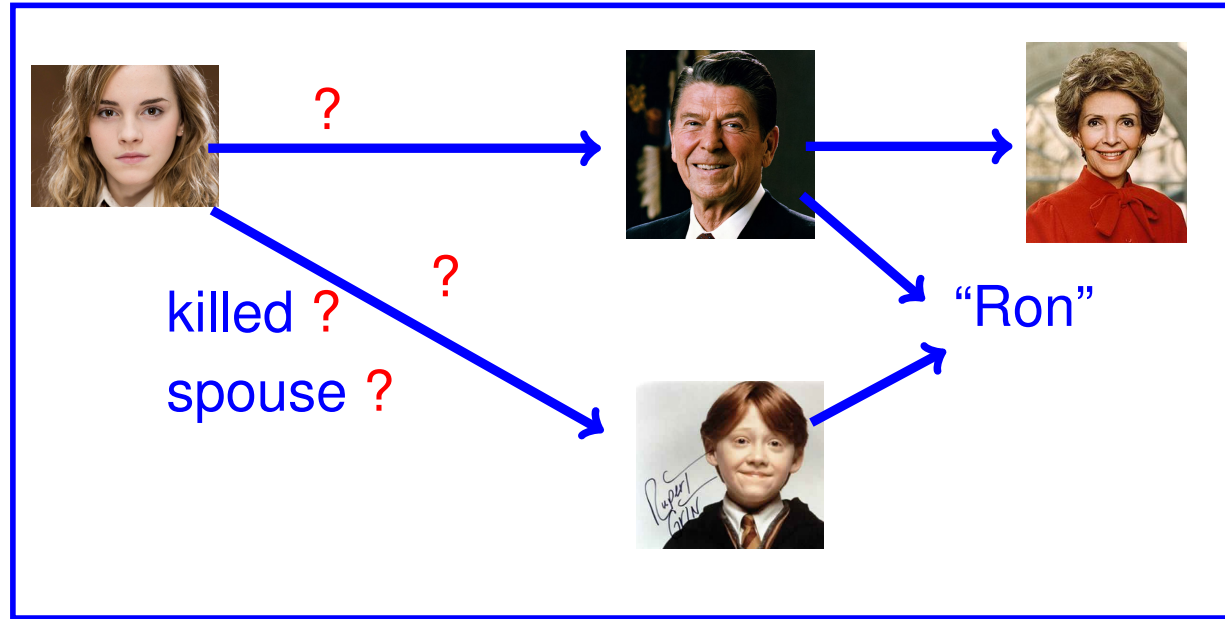
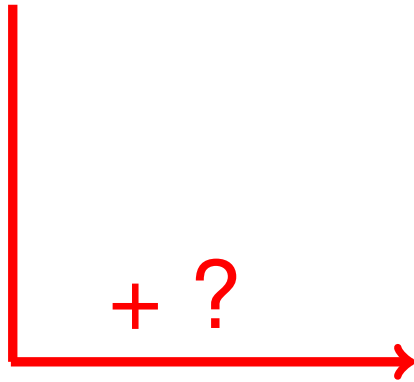
Problems in extending a KB

“Hermione is married to Ron”



Problems in extending a KB

“Hermione is married to Ron”

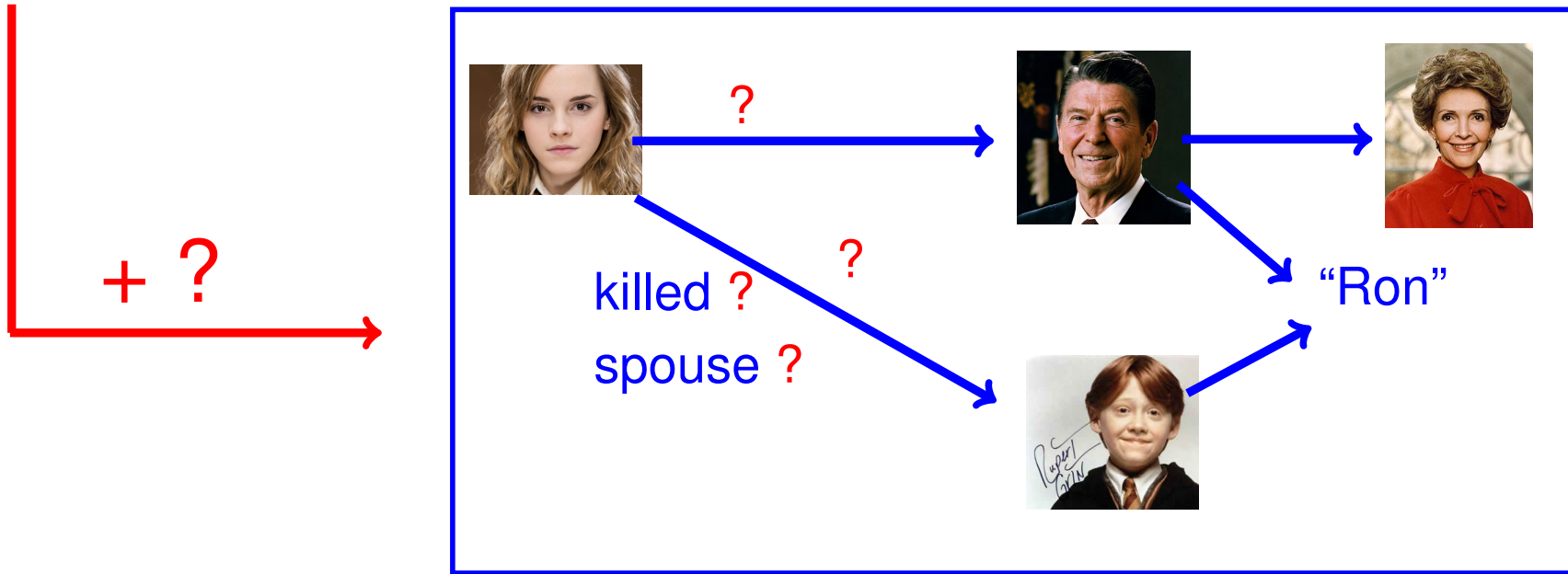


IE faces at least 3 problems:

- Understand patterns (“X is married to Y” = killed(X,Y)?)
- Disambiguate entities (“Ron”= Ronald Reagan?)
- Resolve inconsistencies (Reagan married to 2 women?)

Problems in extending a KB

“Hermione is married to Ron”



- Disambiguation avoids inconsistency
- Pattern helps disambiguation
- Consistency helps finding pattern

=> Solve all 3 problems together!

Idea: Solve all problems together

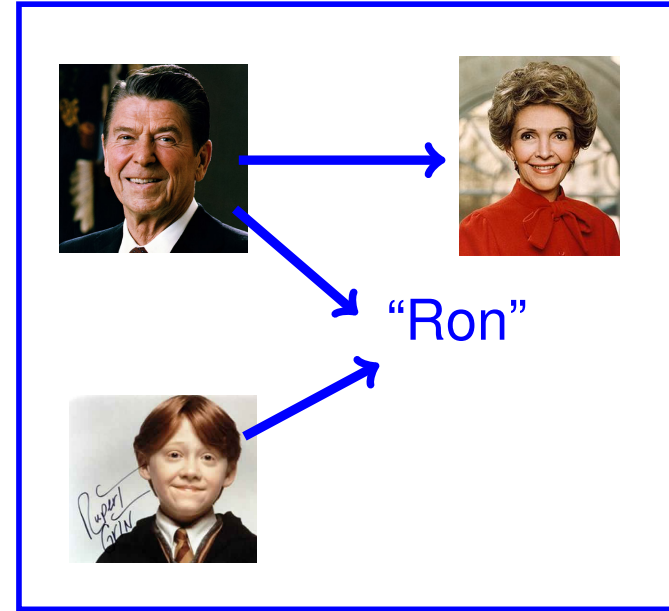
“Hermione is married to Ron”

transform
everything
to logical
formulas

$$A \wedge B \Rightarrow C$$

Find best conclusion

hasSpouse(Hermione, RonWeasley)



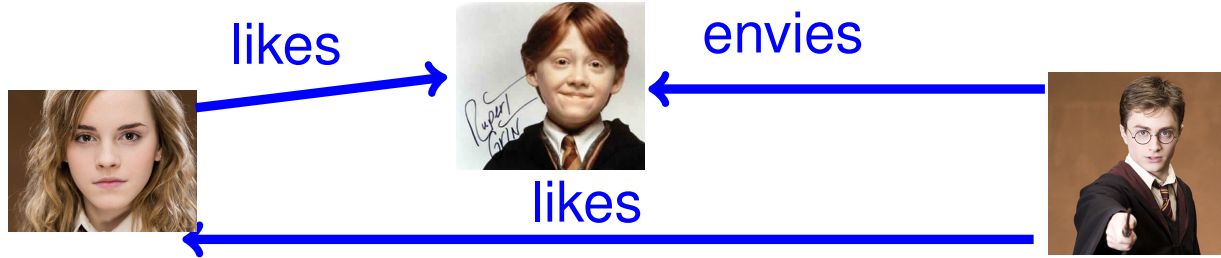
Implications > 13

Refresh: Atoms and KBs

An atom **holds** (“is true”) in a KB, if it appears in the KB.

A negated atom $\neg A$ **holds** in a KB if A does not hold.

A conjunction $A \wedge B \wedge \dots \wedge Z$ **holds** in a KB, if all of its elements hold.



likes(Hermione, Ron) ?

\neg envies(Ron, Harry) ?

\neg envies(Ron, Harry) \wedge likes(Harry, Hermione)

envies(Harry, Ron) \wedge likes(Hermione, Ron) \wedge likes(Harry, Elvis)

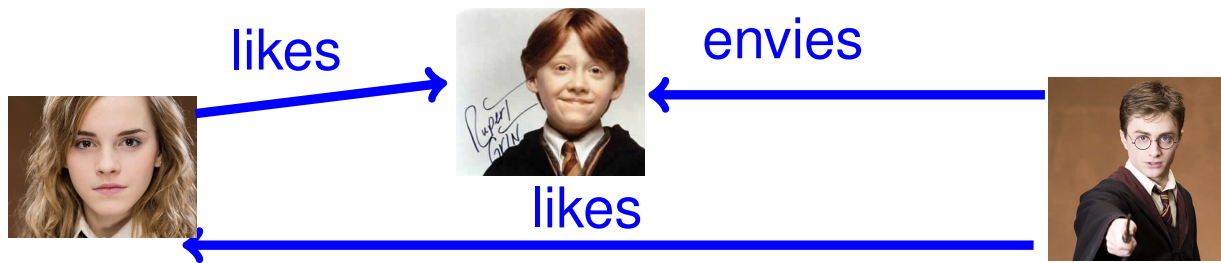
Refresh: Implications

An atom **holds** (“is true”) in a KB, if it appears in the KB.

A negated atom $\neg A$ **holds** in a KB if A does not hold.

A conjunction $A \wedge B \wedge \dots \wedge Z$ **holds** in a KB, if all of its elements hold.

An implication $\vec{B} \Rightarrow H$ **holds** in a KB if \vec{B} does not hold or H holds.



$likes(Hermione, Ron) \Rightarrow likes(Harry, Ron)$

$likes(Harry, Ron) \Rightarrow hasSpouse(Harry, Hermione)$

$\neg likes(Hermione, Harry) \Rightarrow envies(Harry, Ron)$

$\neg likes(Hermione, Harry) \Rightarrow hasSpouse(Harry, Ron)$

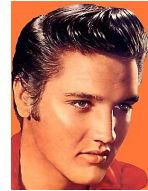
$likes(Herm., Ron) \wedge \neg envies(Harry, Ron) \Rightarrow ffe(Harry, Ron)$

Refresh: Universally quantified formulas

A universally quantified formula **holds** in a KB, if all of its instantiations hold.



hasSpouse



hasSpouse



hasSpouse



$hasSpouse(x, y)$

$hasSpouse(x, y) \Rightarrow hasSpouse(y, x)$

$hasSpouse(Elvis, y) \Rightarrow hasSpouse(y, Elvis)$

$hasSpouse(Ron, y) \Rightarrow hasSpouse(y, Ron)$

Def: Weighted Rule

A **weighted rule** is a rule with an associated real-valued weight.

$$hasSpouse(X, Y) \wedge different(Y, Z) \Rightarrow \neg hasSpouse(X, Z) [3.14]$$

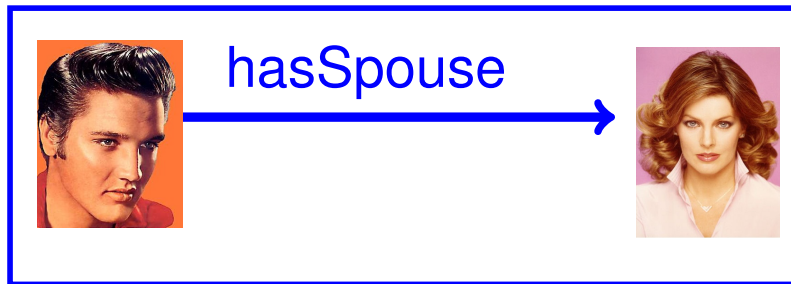
Def: Weight of a KB

Given a KB (a “possible world”) and a set of instantiated rules with weights, the **weight of the KB** is the sum of the weights of all true rules.

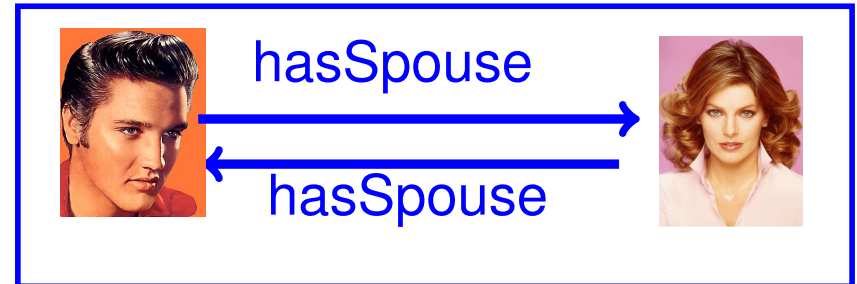
$hasSpouse(Elvis, Priscilla) \Rightarrow hasSpouse(Priscilla, Elvis)[3]$

$hasSpouse(cat, dog) \Rightarrow hasSpouse(dog, cat)[2]$

KB1



KB2



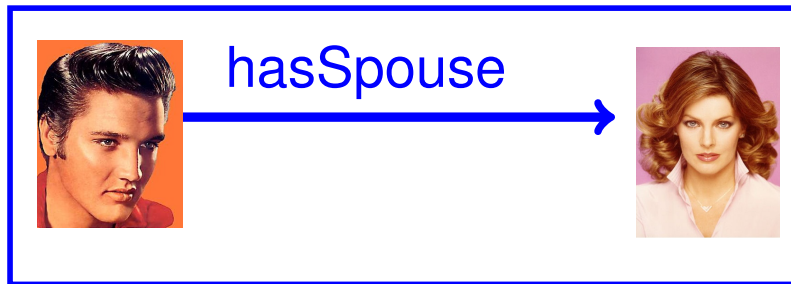
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$hasSpouse(Elvis, Priscilla) \Rightarrow hasSpouse(Priscilla, Elvis)[3]$

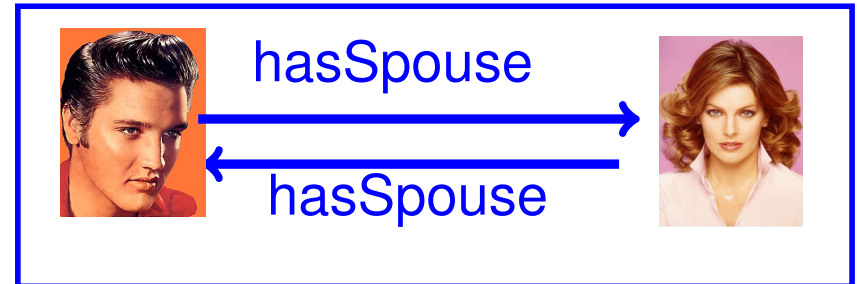
$hasSpouse(cat, dog) \Rightarrow hasSpouse(dog, cat)[2]$

KB1



Weight: 2

KB2



Weight: 5

Def: Weighted MAX SAT

Given a set of instantiated rules with weights, **weighted MAX SAT** is the problem of finding the KB with the highest weight.

(Since SAT is NP-complete, so is MAX SAT and Weighted MAX SAT. There may be multiple such worlds. We are interested in minimal worlds.)

is(Ron, immature)[10]

is(Ron, immature) \wedge type(H., sorceress) \Rightarrow likes(H., Ron)[3]

type(Hermione, sorceress)[4]

Def: Weighted MAX SAT

Given a set of instantiated rules with weights, **weighted MAX SAT** is the problem of finding the KB with the highest weight.

(Since SAT is NP-complete, so is MAX SAT and Weighted MAX SAT. There may be multiple such worlds. We are interested in minimal worlds.)

is(Ron, immature)[10]

is(Ron, immature) \wedge type(H., sorceress) \Rightarrow likes(H., Ron)[3]

type(Hermione, sorceress)[4]

is(Ron, immature)

Best world:

type(Hermione, sorceress)

weight: 17

likes(Hermione, Ron)

Task: Weighted MAX SAT

Find the KB with the highest weight:

is(Hermione, smart)[1]

is(Herm., smart) \wedge is(Harry, smart) \Rightarrow likes(Herm., Harry)[3]

likes(Hermione, Ron) \Rightarrow \neg likes(Hermione, Harry)[100]

is(Harry, smart)[10]

likes(Hermione, Ron)[20]



Hint: Start by satisfying disjunctions with one literal and high weight.

Back to our problem

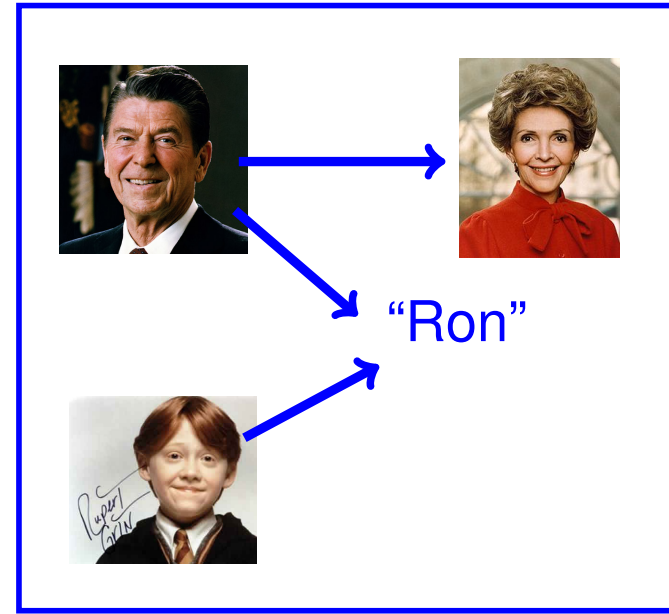
“Hermione is married to Ron”

transform
everything
to logical
formulas

$$A \wedge B \Rightarrow C$$

Find best conclusion

hasSpouse(Hermione, RonWeasley)



Consistency

Consistency constraints can be expressed by rules:

hasSpouse(*X*, *Y*) \wedge *different*(*Y*, *Z*) \Rightarrow \neg *hasSpouse*(*X*, *Z*)[10]

hasSpouse(*X*, *Y*) \Rightarrow *type*(*X*, *person*)[20]

loves(*X*, *Y*) \wedge \neg *hasSpouse*(*X*, *Y*) \Rightarrow *type*(*X*, *stupidPerson*)[3]

...

Rules and weights can be designed manually. Such rules will guide our information extraction process.

A KB can be expressed as rules

Every fact from the KB can be expressed as a weighted rule:

type(Hermione, Person)[100]

High weight


This corresponds to the rule

\Rightarrow *type(Hermione, Person)[100]*

Assuming completeness

If we assume the KB to be perfect on a relation, we can create negative rules:

$$\neg type(Hermione, X)[100]$$

$$\forall X : type(Hermione, X) \notin KB$$

i.e.

$$\neg type(Hermione, chicken)[100]$$

$$\neg type(Hermione, cat)[100]$$

$$\neg type(Hermione, mouse)[100]$$

...

Expressing the corpus as rules

D42:

Hermione married Ron.



occurs(Hermione@D42, "married", Ron@D42)

Expressing the corpus as rules

D42:

Hermione married Ron.



occurs(Hermione@D42, “married”, Ron@D42)



Does not talk about Ronald Reagan or Ron Weasley, but about the word “Ron” in document D42.

Ron@D42 is a “word in context” (wic).

Expressing the corpus as rules

D42:

Hermione married Ron.



occurs(Hermione@D42, “married”, Ron@D42)

The word “Ron” in document D42 can mean different entities (from KB):

means(Ron@D42, *RonaldReagan*)

means(RonD42, *RonWeasley*)

But only one entity in practice:

means(*X*, *Y*) \wedge *different*(*Y*, *Z*) \Rightarrow \neg *means*(*X*, *Z*)

Weights for corpus rules

occurs(Hermione@D42, “married”, Ron@D42)[3]

means(Ron@D42, *RonaldReagan*)


means(Ron@D42, *RonWeasley*)

$\text{means}(X, Y) \wedge \text{different}(Y, Z) \Rightarrow \neg \text{means}(X, Z)$

Weights for corpus rules

occurrences

occurs(Hermione@D42, “married”, Ron@D42)[3]



means(Ron@D42, *RonaldReagan*)

means(Ron@D42, *RonWeasley*)

$\text{means}(X, Y) \wedge \text{different}(Y, Z) \Rightarrow \neg \text{means}(X, Z)$

Weights for corpus rules

occurs(Hermione@D42, “married”, Ron@D42)[3]

From disambiguation by context/prior

means(Ron@D42, *RonaldReagan*)[5]

means(Ron@D42, *RonWeasley*)[7]

$\text{means}(X, Y) \wedge \text{different}(Y, Z) \Rightarrow \neg \text{means}(X, Z)$

Weights for corpus rules

occurs(Hermione@D42, “married”, Ron@D42)[3]

means(Ron@D42, *RonaldReagan*)[5]

means(Ron@D42, *RonWeasley*)[7]

“Hard” rule with very high weight

means(*X*, *Y*) \wedge *different*(*Y*, *Z*) \Rightarrow \neg *means*(*X*, *Z*) [100]



Weights for corpus rules

occurs(Hermione@D42, “married”, Ron@D42)[3]

means(Ron@D42, *RonaldReagan*)

means(Ron@D42, *RonWeasley*)

$\text{means}(X, Y) \wedge \text{different}(Y, Z) \Rightarrow \neg \text{means}(X, Z)$ [100]

Let us ignore the weights for a moment.

Deducing patterns

KB



+

Reagan married Davis.

=

“X married Y”
is pattern for
 $hasSpouse(X,Y)$

Deducing patterns

KB



hasSpouse(Reagan, Davis)

+

Reagan married Davis.

=

“X married Y”
is pattern for
hasSpouse(X,Y)

Deducing patterns

KB



+

Reagan married Davis.

=

“X married Y”
is pattern for
 $hasSpouse(X, Y)$

$hasSpouse(Reagan, Davis)$

$occurs(R@1, \text{“married”}, D@1)$

$means(R@1, Reagan)$

$means(D@1, Davis)$

Deducing patterns

KB



+

Reagan married Davis.

=

“X married Y”
is pattern for
 $hasSpouse(X, Y)$

$hasSpouse(Reagan, Davis)$

$occurs(R@1, \text{“married”}, D@1)$

$means(R@1, Reagan)$

$means(D@1, Davis)$

$occurs(X, P, Y)$

$\wedge means(X, X')$

$\wedge means(Y, Y')$

$\wedge R(X', Y')$

$\Rightarrow isPatternFor(P, R)$

Deducing patterns

KB



+

Reagan married Davis.

=

“X married Y”
is pattern for
 $hasSpouse(X, Y)$

$hasSpouse(Reagan, Davis)$

$occurs(R@1, \text{“married”}, D@1)$

$means(R@1, Reagan)$

$means(D@1, Davis)$

$occurs(X, P, Y)$

$\wedge means(X, X')$

$\wedge means(Y, Y')$

$\wedge R(X', Y')$

$\Rightarrow isPatternFor(P, R)$

$isPatternFor(\text{“married”}, hasSpouse)$

Applying patterns

“X married Y”

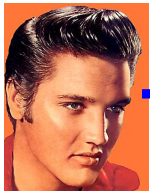
is pattern for

hasSpouse(X,Y)

+

Elvis married Priscilla.

=



hasSpouse



Applying patterns

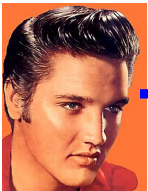
“X married Y”
is pattern for
hasSpouse(X,Y)

isPatternFor(“married”,*hasSpouse*)

+

Elvis married Priscilla.

=



hasSpouse



Applying patterns

“X married Y”

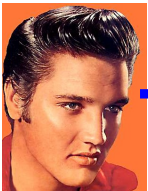
is pattern for

hasSpouse(X,Y)

+

Elvis married Priscilla.

=



hasSpouse



isPatternFor(“married”,*hasSpouse*)

occurs(E@1, “married”, P@1)

means(E@1, *Elvis*)

means(P@1, *Priscilla*)

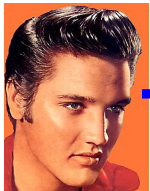
Applying patterns

“X married Y”
is pattern for
hasSpouse(X,Y)

+

Elvis married Priscilla.

=



hasSpouse



isPatternFor(“married”,*hasSpouse*)

occurs(E@1, “married”, P@1)

means(E@1, *Elvis*)

means(P@1, *Priscilla*)

occurs(X, P, Y)

\wedge *means*(X, X')

\wedge *means*(Y, Y')

\wedge *isPatternFor*(P, R)

$\Rightarrow R(X', Y')$

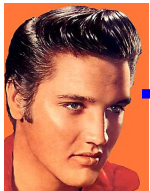
Applying patterns

“X married Y”
is pattern for
hasSpouse(X,Y)

+

Elvis married Priscilla.

=



hasSpouse



isPatternFor(“married”,*hasSpouse*)

occurs(E@1, “married”, P@1)

means(E@1, *Elvis*)

means(P@1, *Priscilla*)

occurs(X, P, Y)

\wedge *means*(X, X')

\wedge *means*(Y, Y')

\wedge *isPatternFor*(P, R)

$\Rightarrow R(X', Y')$

hasSpouse(*Elvis*, *Priscilla*)

task>106

Task: Pattern deduction by rules

Pattern deduction:

$$\begin{aligned} & occurs(X, P, Y) \\ & \wedge means(X, X') \\ & \wedge means(Y, Y') \\ & \wedge R(X', Y') \\ & \Rightarrow isPatternFor(P, R) \end{aligned}$$

Pattern application:

$$\begin{aligned} & occurs(X, P, Y) \\ & \wedge means(X, X') \\ & \wedge means(Y, Y') \\ & \wedge isPatternFor(P, R) \\ & \Rightarrow R(X', Y') \end{aligned}$$

1 : $occurs(P@1, \text{"adores"}, E@1)$
2 : $means(E@1, Elvis)$
3 : $means(P@1, Priscilla)$
4 : $hasSpouse(Priscilla, Elvis)$
5 : $occurs(M@1, \text{"adores"}, E@1)$
6 : $means(M@1, Madonna)$

All rules have weight 1.

Compute facts that will be in the best world.

Life is not easy

- words are ambiguous

“Ron”

- corpora may err

“Madonna is married to Elvis”


- parsing may fail

“Coyote dreams Coyote eats Roadrunner”

- contradictions may occur

Reagan was married twice.

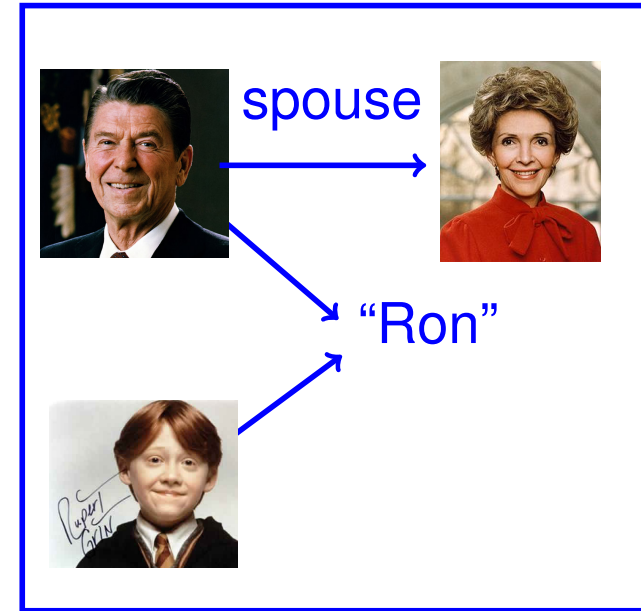
occurs(W@1, “eats”, R@1)



=> we will compute the most plausible world

Finding the most plausible world

“Hermione married Ron”



World1:



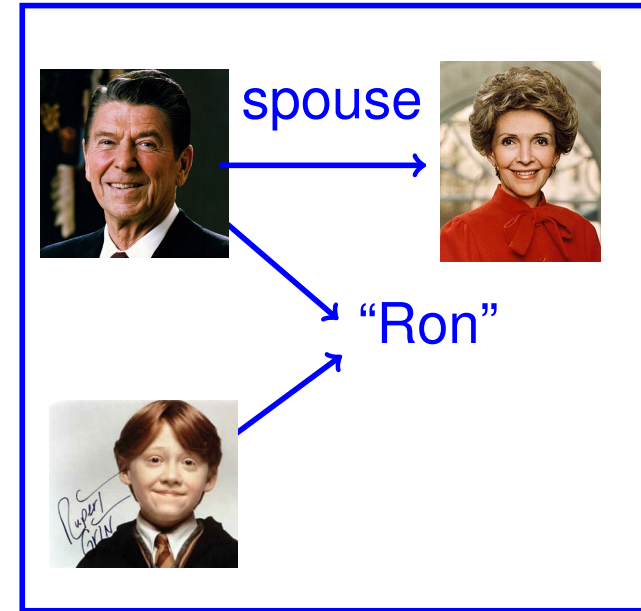
World2:



Finding the most plausible world

“Hermione married Ron”

occurs(H@1, “married”, R@ 1)[1]
isPatternFor(“married”, *spouse*)[1]



World1:



World2:



Finding the most plausible world

“Hermione married Ron”

occurs(H@1, “married”, R@ 1)[1]

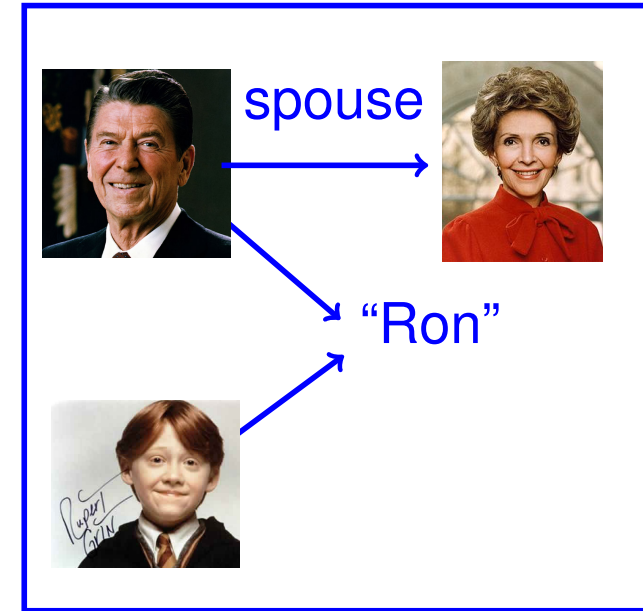
isPatternFor(“married”, *spouse*)[1]

means(H@1, *Hermione*)[5]

means(R@1, *RonWeasley*)[2]

means(R@1, *Reagan*)[3]

means(X, Y) \wedge Y \neq Z $\Rightarrow \neg$ *means*(X, Z)[10]



World1:



World2:



Finding the most plausible world

“Hermione married Ron”

occurs(H@1, “married”, R@1)[1]

isPatternFor(“married”, *spouse*)[1]

means(H@1, *Hermione*)[5]

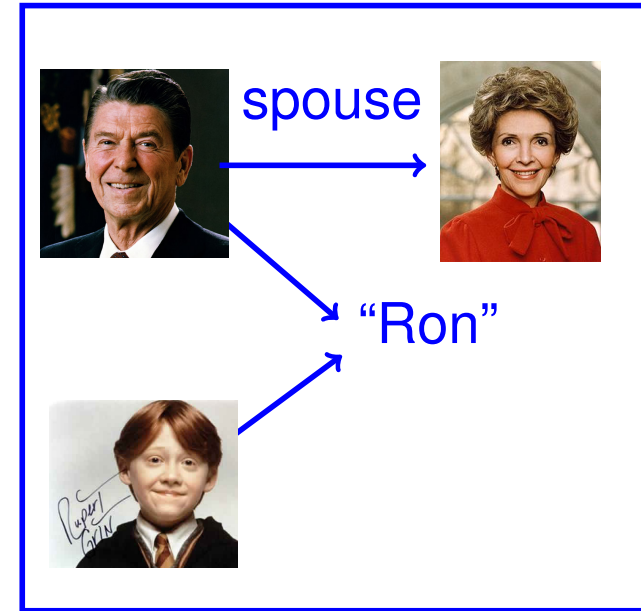
means(R@1, *RonWeasley*)[2]

means(R@1, *Reagan*)[3]

means(X, Y) $\wedge Y \neq Z \Rightarrow \neg means(X, Z)$ [10]

spouse(X, Y) $\wedge Y \neq Z \Rightarrow \neg spouse(Z, X)$ [6]

+ Pattern Application Rule [10]



World1:



World2:



Finding the most plausible world

“Hermione married Ron”

occurs(H@1, “married”, R@1)[1]

isPatternFor(“married”, *spouse*)[1]

means(H@1, *Hermione*)[5]

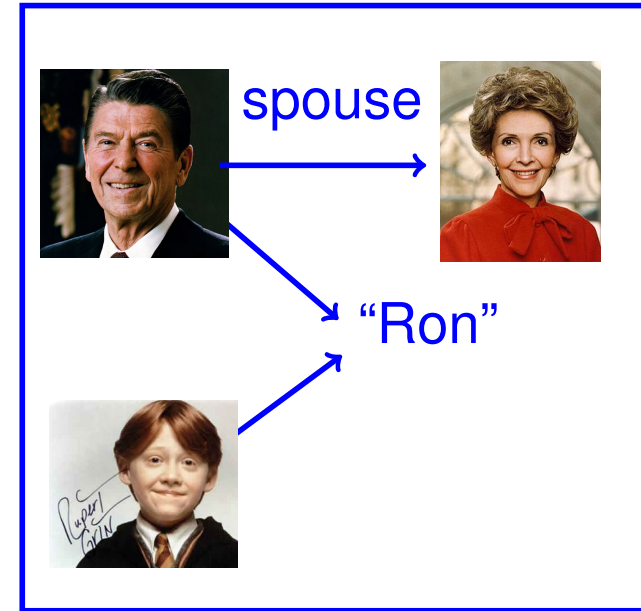
means(R@1, *RonWeasley*)[2]

means(R@1, *Reagan*)[3]

means(*X*, *Y*) \wedge *Y* \neq *Z* \Rightarrow \neg *means*(*X*, *Z*)[10]

spouse(*X*, *Y*) \wedge *Y* \neq *Z* \Rightarrow \neg *spouse*(*Z*, *X*)[6]

+ Pattern Application Rule [10]



World1:



loses 2

wins 3

loses 6

World2:



Finding the most plausible world

“Hermione married Ron”

occurs(H@1, “married”, R@ 1)[1]

isPatternFor(“married”, *spouse*)[1]

means(H@1, *Hermione*)[5]

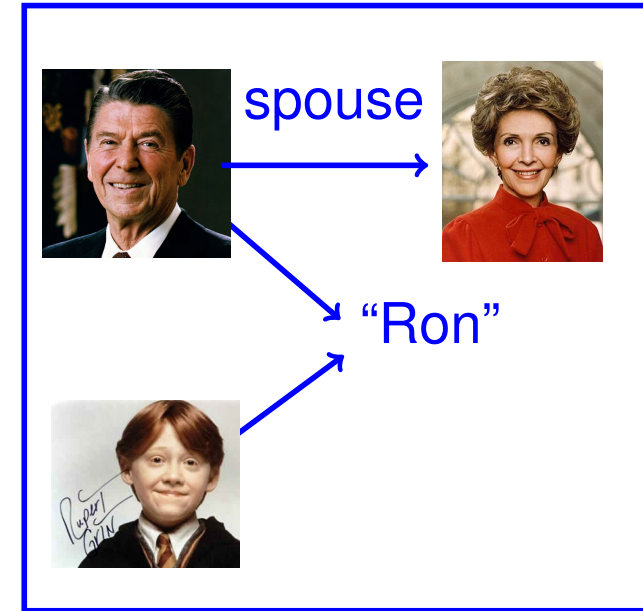
means(R@1, *RonWeasley*)[2]

means(R@1, *Reagan*)[3]

means(X, Y) $\wedge Y \neq Z \Rightarrow \neg means(X, Z)$ [10]

spouse(X, Y) $\wedge Y \neq Z \Rightarrow \neg spouse(Z, X)$ [6]

+ Pattern Application Rule [10]



World1:



loses 2

wins 3

loses 6

World2:



loses 3

wins 6

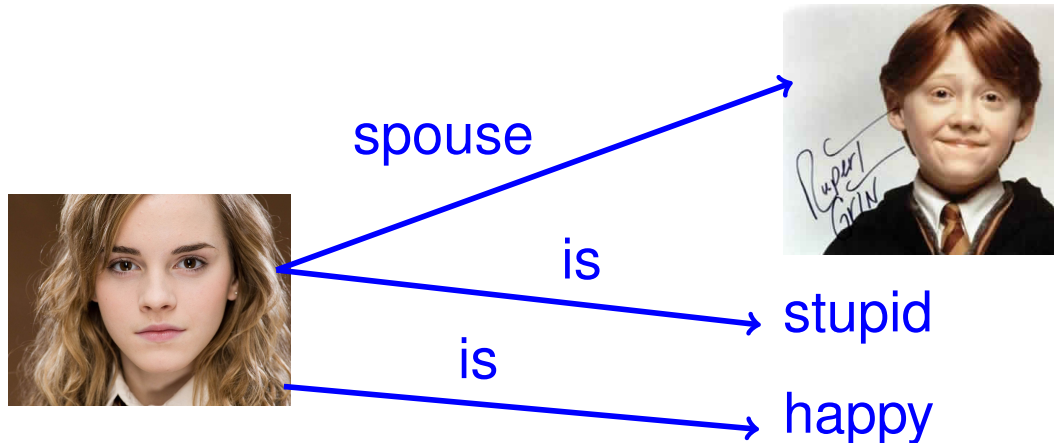
wins 2

Summary: IE by reasoning

IE by reasoning converts

- background knowledge (existing KB)
- logical constraints
- the corpus

into logical rules and aims to find the most plausible possible world given these evidences.



References

SOFIE - a self-organizing framework for IE

->markov-logic

->semantic-web