

“数据库技术” 考试答题纸

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一. Database Design

1.1. ODL

```
interface Country (key name)
{
    attribute string name;
    attribute int population;
    attribute string size;
    attribute string socialSystemType;
    relationship Set<Province> have
        inverse Province:: belongTo
}

interface Province (key name)
{
    attribute string name;
    attribute int population;
    attribute string size;
    attribute Struct Range
        { string start, string end; latitudeRange;
    relationship Country belongTo
        inverse Country:: have;
    relationship Set<Deceased> govern
        inverse Deceased:: comeFrom;
}
```

interface Deceased (key ID)

```
{  
    attribute string ID;  
    attribute string name;  
    attribute string birthdate;  
    attribute string dateOfDeath;  
    relationship Province ComeFrom  
        inverse Province :: govern  
}
```

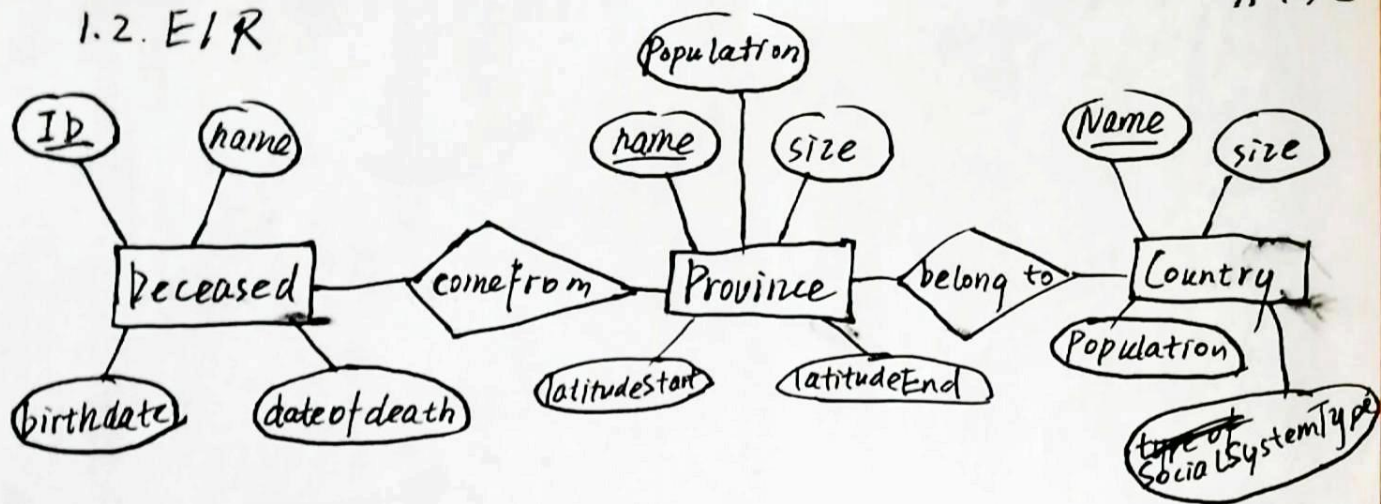
- Relational Schemas:

Country ( name, population, size, socialSystemType)

Province ( name, population, size, latitudeStart, latitudeEnd,  
countryName )

Deceased ( ID, name, birthdate, dateOfDeath, ProvinceName )

## 1.2. EIR



- Relational Schemas:

Country ( Name, population, size, socialSystemType )

Province ( Name, population, size, latitudeStart, latitudeEnd )

Deceased ( ID, ~~name~~, birthdate, date of Death )

ComeFrom ( ID, ProvinceName )

BelongTo ( ProvinceName, CountryName )

## 1.3

E<sub>1s</sub> ( A, B, C )

E<sub>2s</sub> ( D, E, F )

E<sub>3s</sub> ( G, H, I, A, B )

E<sub>4s</sub> ( D, J, K )

R<sub>1</sub> ( A, B, D )



## 二. Normal Form of Relational Database

2.1. Key of  $R$ :  $\{A, C\}$

4NF violations:  $A \twoheadrightarrow C$ ,  $AB \twoheadrightarrow D$ ,  $B \twoheadrightarrow D$ ,

• Decomposition Process:

1. use  $A \twoheadrightarrow C$  to decompose  $R$  into:

$R_1(A, C)$  and  $R_2(A, B, D)$

$R_1$  is in 4NF

Key of  ~~$R_2$~~   $R_2$ :  $\{A, B\}$

4NF violations:  $B \twoheadrightarrow D$

2. use  $B \twoheadrightarrow D$  to decompose  $R_2$  into:

$R_3(B, D)$  and  $R_4(A, B)$

$R_3, R_4$  are in 4NF.

• Final Result:  $R_1(A, C)$   $R_3(B, D)$   $R_4(A, B)$

2.2. b. compute the closure of each subset of  $\{A, B, C\}$ :

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1.  $\{A\}^+ = \{A\}$      $\{B\}^+ = \{B\}$      $\{C\}^+ = \{C, E, A\}$

• add  $C \rightarrow A$

2.  $\{A, B\}^+ = \{A, B, C, D, E\}$      $\{B, C\}^+ = \{B, C, E, A, D\}$

• add  $AB \rightarrow C$

• add  $BC \rightarrow A$

$\{A, C\}^+ = \{A, C, E\}$

• All the completely nontrivial functional dependencies:

$C \rightarrow A$

$AB \rightarrow C$

$BC \rightarrow A$

• in the form of minimal basis.

$\{C \rightarrow A,$

$AB \rightarrow C\}$

## 三. Relational Algebra

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3.1

$\pi_{Name} (Tourist \bowtie \sigma_{Name='Hubei'} (Province))$

$\pi_{Tourist.Name} (Tourist \bowtie_{Tourist.Pcode=Province.Pcode} \sigma_{Name='Hubei'} (Province))$

3.2.

$\pi_{Name} (Tourist \bowtie \sigma_{Expense > 3000} (TourismRecord))$

## 四. T-SQL.

4.1  $\text{SELECT } \overset{COUNT(*)}{\cancel{SUM(*)}} \text{ AS totalNumber}$   
 $\text{FROM Destination D INNER JOIN Province P}$   
 $\text{ON D.PCode} = \text{P.PCode}$   
 $\text{WHERE P.Name} = 'Hubei'$

4.2  $\text{SELECT } \cancel{SUM(T.ID)} \overset{ArrivalTime}{COUNT(\cancel{T.ID})} \text{ as totalNumber,}$   
 $\text{MAX(Name) as Name,}$   
 $\text{T.ID as ID}$   
 $\text{FROM } \cancel{Tourist} \text{ Tourist T INNER JOIN TourismRecord TR}$   
 $\text{ON T.ID} = \text{TR.ID}$   
 $\text{Group BY T.ID}$   
 $\text{ORDER BY COUNT(ArrivalTime) DESC}$



② -- It show the information of expense 李与怡 第7页

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```
SELECT COUNT(ArrivalTime) as totalNumber,  
       AVG(Expense) as averageExpense,  
       MAX(Name) as Name,  
       T.ID as ID
```

```
FROM Tourist T INNER JOIN TouristRecord TR
```

```
ON T.ID = TR.ID
```

```
GROUP BY T.ID
```

```
ORDER BY COUNT(ArrivalTime) DESC
```

4.3. SELECT ~~MAX(T.ID)~~ as ID,  
 MAX(Name) as Name,  
 AVG(Expense) as averageExpense

```
FROM Tourist T INNER JOIN TouristRecord TR
```

```
ON T.ID = TR.ID
```

```
GROUP BY T.ID
```

```
HAVING AVG(Expense) > (SELECT AVG(Expense)  
                        FROM TouristRecord)
```

## 1. Background

本次大作业正值新冠疫情，积极的队友叶紫璇已率先将自己的初步想法“关于口罩生产”发布于讨论群。我与另外的队友刘又萌认为这个主题切合热点，有实际价值，便一致同意立足于这个主题。

## 2. Requirement Analysis

由于初步想法是一张草图，设计的对象和联系较为冗杂。为了贴合实际生产设计，我认为确立一个立足点较为重要，之后再围绕立足点将重要的信息纳入数据模型。“你们想当卫健委呢？还是国药集团的董事长？”一个戏谑的发问，希望能让我们更快拥有用处的代入感——前者是对医疗资源的宏观调度，后者是需要考虑口罩的实际生产。她们认为“卫健委”比较有前途，我们便在宏观的角度分析口罩资源调配所涉及到的对象和数据：

- ① 口罩 · 型号：疫情期间类别很关键
- 工厂 · 生产能力规模：疫情停工停产/复工时的最大生产力。
- 仓库 · 口罩储量：停工停产状态下可供能力。
- 物流公司 · 运输能力：疫区需求最大，物流业能否支撑。
- 购买单位：口罩的消耗方。 (~~购买单位~~)
- 订单：购买信息。

## 3. Summary

上述的结果实际是在讨论、建模过程中逐渐完善得出的结果。经过大家的修改，基本符合实际需求。

并且在ODL建模时，将工厂、物流公司与购买单位进一步抽象得到基本单位。——包括地址、联系电话等公共属性。



1. 且见我于笔记中写到的-条:

"4.6. 我的姐妹们太强了, 今天题目刚出来, 刘又萌一己之力完成了全部解答, 并写成了word文档, 叹服.  
一个高效率的组织多么重要!"

2. 看来在解题上我是大用武之地了, 于是另辟蹊径.

参考到作业中提到的"代码附上".

何妨具体将题目用 SQL 建立数据库并实现呢?

说干就干, 且说无独有偶, 收到的题目也是60-19大背景下的医疗体系..

需求分析 抄录如下.

1. 导入地区数据

2. 注册医院

• 创建数据库并默认科室

3. 添加医护人员信息.

4. 患者就医 (通过 Procedure 实现)

a. 注册

b. 充值

c. 消费

d. 注销.