

Krylov method

$$Ax = b. \quad A \in \mathbb{R}^{n \times n}. \quad \text{sym. positive definite.}$$

equivalent to minimization problem.

$$\varphi(x) = \frac{1}{2} x^T A x - b^T x$$

$$\nabla \varphi(x) = Ax - b = 0. \quad \text{Hessian is } A \Rightarrow \text{minimizer.}$$

residual $r = b - AX = -\nabla \varphi(x)$ negative gradient.

$$\langle u, v \rangle_A = u^T A v. \quad A\text{-inner product.}$$

$$\langle u, u \rangle_A = u^T A u \geq 0. \quad \langle u, u \rangle_A = 0 \Rightarrow u = 0.$$

$$\|u\|_A^2 = \langle u, u \rangle_A.$$

steepest descent

$$x_{k+1} = x_k + \alpha_k p_k, \quad p_k = r_k = b - A x_k, \quad \alpha_k \in \mathbb{R}.$$

$$\alpha_k = \arg \min_{\alpha} \varphi(x_k + \alpha_k P_k)$$

$$\varphi(x_k + \alpha_k P_k) = \frac{1}{2} (x_k + \alpha_k P_k)^T A (x_k + \alpha_k P_k) - b^T (x_k + \alpha_k P_k)$$

$$\frac{\partial \varphi(x_k + \alpha_k P_k)}{\partial \alpha_k} = (P_k^T A P_k) \alpha_k + P_k^T \underbrace{(A x_k - b)}_{-r_k} = 0$$

$$\alpha_k = \frac{P_k^T r_k}{P_k^T A P_k} .$$

Alg. for $K=1$, maxit

$$r_k = b - Ax_k$$

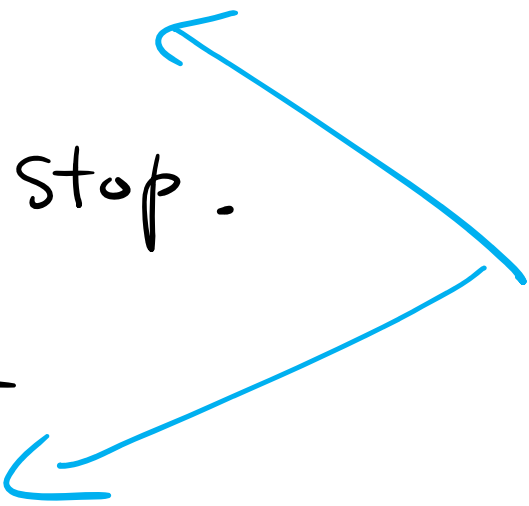
if $\|r_k\| < \tau$. Stop .

$$\alpha_k = \frac{r_k^T r_k}{r_k^T A r_k}$$

$$x_{k+1} = x_k + \alpha_k r_k$$

end .

A is applied
twice .



Reformulate

$$r_{k+1} = b - A x_{k+1}$$

$$= b - A (x_k + \alpha_k r_k) = r_k - \alpha_k \underline{\underline{A r_k}}.$$

store this.

Alg. for $k = 1, \text{maxit}$

$w = A r_k$. if $\|r_k\| < \bar{\epsilon}$ stop

$$\alpha_k = \frac{r_k^T r_k}{r_k^T w}$$

$$x_{k+1} = x_k + \alpha_k r_k, \quad r_{k+1} = r_k - \alpha_k w$$

Convergence of SD.

Thm. A sym. pos. def. eigenvalues $0 < \lambda_1 \leq \lambda_2 \leq \dots \leq \lambda_n$

Cond. number $K(A) = \frac{\lambda_n}{\lambda_1}$

$$\|x_k - x_*\|_A \leq \left(\frac{K(A) - 1}{K(A) + 1} \right)^k \|x_0 - x_*\|_A.$$

1. SD converges. linear rate.

$$2. \quad \gamma = \frac{1 - \frac{1}{K(A)}}{1 + \frac{1}{K(A)}} \approx 1 - \frac{2}{K(A)}$$

$$Pf: \quad \varphi(x) + \frac{1}{2} x_{\star}^T A x_{\star} = \frac{1}{2} (x - x_{\star})^T A (x - x_{\star})$$

$$SD: \quad \varphi(x_k) \leq \varphi(x_{k-1} + \alpha r_{k-1}) \quad \forall \alpha.$$

opt. step
length 

$$\Rightarrow \cancel{\frac{1}{2}} (x_k - x_{\star})^T A (x_k - x_{\star}) \leq \cancel{\frac{1}{2}} (x_{k-1} + \alpha r_{k-1})^T A (x_{k-1} + \alpha r_{k-1})$$

$\|x_k - x_{\star}\|_A^2$

$$\hookrightarrow \leq (x_{k-1} - x_{\star})^T (I - \alpha A) A (I - \alpha A) (x_{k-1} - x_{\star})$$

$$\|P(A)(x_k - x_*)\|_A^2 \leq \max_{1 \leq i \leq n} |P(\lambda_i)|^2 \|x_{k-1} - x_*\|_A^2 \quad \left| \begin{array}{l} P(A) = I - \alpha A \end{array} \right.$$

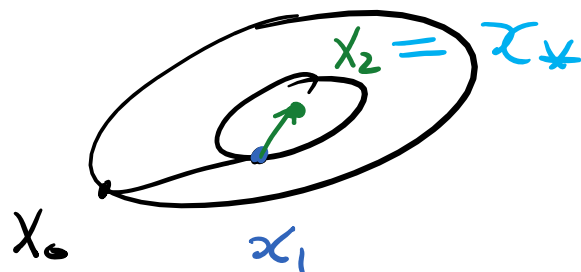
$$\leq \min_{\alpha} \max_{\lambda_1 \leq t \leq \lambda_n} |1 - \alpha t|^2 \|x_{k-1} - x_*\|_A^2$$

sol. $\alpha_* = \frac{2}{\lambda_1 + \lambda_n}$

val: $|1 - \alpha_* t| = \frac{\lambda_n - \lambda_1}{\lambda_n + \lambda_1} \quad \cdot \square$

CG.

2D.



First step. SD.

$$\alpha_0 = \frac{r_0^T P_0}{P_0^T A P_0}$$

$$x_1 = x_0 + \alpha_0 P_0, \quad P_0 = r_0, \quad r_0 = b - Ax_0.$$

$$x_2 = x_1 + \alpha_1 P_1$$

$$\alpha_1 = \frac{r_1^T P_1}{P_1^T A P_1} \leftarrow \text{universal. opt.}$$

choose $P_1^T A P_0 = 0$. A -ortho to P_0 .

Compute $r_2 = b - A(x_1 + \alpha_1 p_1)$

$$= b - Ax_1 - \alpha_1 A p_1$$

$$= (I - \alpha_0 A) p_0 - \alpha_1 A p_1$$

$\{p_0, p_1\}$ form a basis in \mathbb{R}^2

$$p_0^T r_2 = (p_0^T p_0 - \underbrace{\alpha_0}_{=0} p_0^T A p_0) - \alpha_1 \underbrace{p_0^T A p_1}_{=0} = 0$$

$$p_1^T r_2 = (p_1^T p_0 - \alpha_1 p_1^T A p_1) - \alpha_0 \underbrace{p_1^T A p_0}_{=0}$$

$$P_1^T P_0 - \alpha_1 P_1^T A P_1 = P_1^T P_0 - P_1^T r_0 = \alpha_0 P_1^T A P_0 = 0.$$

$$\Rightarrow r_2 = 0 \Rightarrow x_2 = x_4.$$

CG for \mathbb{R}^n .

$$P_1 = r_1 + \beta_1 P_0$$

$$P_0^T A P_1 = P_0^T A r_1 + \beta_1 P_0^T A P_0 \Rightarrow \beta_1 = - \frac{P_0^T A r_1}{P_0^T A P_0}$$

$$x_{k+1} = x_k + \alpha_k P_k$$

$$\alpha_k = \frac{r_k^T P_k}{P_k^T A P_k} \text{ minimization}$$

$$r_{k+1} = b - A x_{k+1}$$

$$P_{k+1} = r_{k+1} + \beta_k P_k$$

$$\beta_k = - \frac{P_k^T A r_{k+1}}{P_k^T A P_k} \text{ ortho.}$$

Conceptual implementation of CG.

2 applications of A each iteration.

practical implementation only requires

1.

In general.

Krylov subspace

$$K_m(A, r_0) \equiv K_m = \text{span} \{ r_0, A r_0, \dots, A^{m-1} r_0 \}.$$

① For all m , $p_m \in K_m$.

$$x_{m+1} \in x_0 + K_m$$

② $p_{m+1} \in K_{m+1}$, $p_{m+1} \perp_A K_m$.

$\{P_0, P_1, \dots, P_{m+1}\}$ mutually A -orthogonal.

→ short (3-term) recurrence.

③ x_{m+1} achieves the global minimum of $\varphi(x)$
within all $x_0 + K_m$

Convergence of CG.

$$\forall x \in x_0 + K_m.$$

$$x_* - x = x_* - (x_0 + C_1 r_0 + \dots + C_m A^{m-1} r_0)$$

$$= A^{-1} (r_0 + C_1 A r_0 + \dots + C_m A^m r_0)$$

$$= A^{-1} P_m(A) r_0.$$

$$P_m(z) = 1 + C_1 z + \dots + C_m z^m. \quad P_m(0) = 1$$

$$\|x_* - x_{m+1}\|_A \leq \left(\max_{\lambda_i} |P_m(\lambda_i)| \right) \underbrace{\|A^{-1}r_0\|_A}_{\|x_0 - x_*\|_A}.$$

$$\leq \left(\min_{\substack{P \in \mathcal{P}_m \\ P(0)=1}} \max_{\lambda_1 \leq t \leq \lambda_n} |p(t)| \right) \|x_0 - x_*\|_A.$$

sol. \downarrow Chebyshev poly.

$$\leq 2 \left(\frac{\sqrt{\kappa} - 1}{\sqrt{\kappa} + 1} \right)^{m+1} \|x_0 - x_*\|_A$$

SD : κ . CG : $\sqrt{\kappa}$.

General non-sym matrices.

$$K_m = \text{span}\{r_0, Ar_0, \dots, A^{m-1}r_0\}$$

Ortho normalize \rightarrow Arnoldi process.

$$\beta v_1 = r_0, \quad \beta = \|r_0\|, \quad \|v_1\| = 1.$$

$$A v_1 = h_{21} v_2 + h_{11} v_1$$

\uparrow
normalization.

\uparrow

orthogonal.

$$v_2^T v_1 = 0, \quad v_2^T v_2 = 1$$

$$\Rightarrow h_{11} = v_1^T A v_1$$

$$\Rightarrow h_{21} = \|A v_1 - h_{11} v_1\|$$

$$A v_1 = [v_1 \ v_2] \begin{bmatrix} h_{11} \\ h_{21} \end{bmatrix}$$

$$A [v_1 \ v_2] = [v_1 \ v_2 \ v_3] \begin{bmatrix} h_{11} & h_{12} \\ h_{21} & h_{22} \\ 0 & h_{32} \end{bmatrix}$$

orthogonalization.

normalization

$$A [v_1, \dots, v_m] = [v_1, \dots, v_{m+1}] \begin{bmatrix} h_{11} & \dots & h_{1m} \\ h_{21} & \dots & h_{2m} \\ \vdots & \ddots & \vdots \\ 0 & \dots & h_{mm} \\ \hline 0 & \dots & h_{m+1,m} \end{bmatrix} \rightarrow \begin{bmatrix} H_m \\ \hline 0 \dots 0 \ h_{m+1,m} \end{bmatrix}$$

m.

Arnoldi:

$$A [v_1 \cdots v_m] = [v_1 \cdots v_m] H_m + \underbrace{v_{m+1} e_{m+1}^T}_{\text{"error"}} h_{m+1,m}.$$

↑
upper Hessenberg

Alg . Arnoldi process .

Pick V_1 , $\|V_1\|=1$.

for $k = 1, \dots, m$.

$$w = A V_k$$

for $j = 1, \dots, k$

$$h_{j,k} = V_j^T w$$

$$w = w - h_{j,k} V_j$$

end

$$h_{k+1,k} = \|w\|$$

if $|h_{k+1,k}| < \tau$. stop

$$V_{k+1} = w / h_{k+1,k}$$

end .

Arnoldi: for linear sys.

Start from $r_0 = b - Ax_0$.

Arnoldi: $\bar{V}_m = [v_1, \dots, v_m]$

$$A\bar{V}_m = \bar{V}_m H_m + v_{m+1} e_{m+1}^T h_{m+1,m}.$$

Project to the space \bar{V}_m

$$\bar{V}_m^T A \bar{V}_m = H_m..$$

Find $y_m \in \mathbb{R}^m$

$$x_m = x_0 + V_m y_m.$$

$$r_m \perp K_m. \Rightarrow V_m^* (b - Ax_m) = 0$$

$$\Downarrow$$

$$V_m^* (b - Ax_0 - AV_m y_m) = 0.$$

$\beta e_i = H_m y_m \rightarrow m \times m$ eq. can solve.

Full orthogonalization method (FOM)

Problems of FOM

- ① Error can be very large. in the middle.
- ② Keep entire V_m in memory.

technical. \rightarrow restarting.



Arnoldi for fixed # steps.

Best known fix for ①

Generalized Minimal Residual Method (GMRES)

[Saad. Schultz. 1986].

$$H_m = \begin{pmatrix} h_{11} & \cdots & \cdots \\ h_{21} & \ddots & \vdots \\ 0 & \ddots & h_{m,m-1} & h_{m,m} \end{pmatrix} \quad \tilde{H}_m = \begin{pmatrix} H_m & \\ \cdots & \cdots \\ 0 & \cdots & h_{m+1,m} \end{pmatrix}$$

Find $x_m = x_0 + V_m y_m$ s.t.

$$\min_y \|b - Ax_m\|_2$$

$$\begin{aligned}
\|b - Ax_m\| &= \|b - Ax_0 - AV_m y_m\| \\
&= \|\beta v_1 - V_{m+1} \tilde{H}_m y_m\| \xrightarrow[\text{exact}]{\text{Project to } V_{m+1}} \\
&= \|\beta \tilde{e}_1 - \tilde{H}_m y_m\|
\end{aligned}$$

$${}^{(m+1)} \begin{bmatrix} \tilde{H}_m^m \\ \tilde{H}_m \end{bmatrix} y_m = \rho \begin{bmatrix} 1 \\ 0 \\ \vdots \\ 0 \end{bmatrix}^{(m+1)} \quad \begin{array}{l} \text{least squares.} \\ \text{solved via QR decomp.} \end{array}$$

$$A = QR$$

$$Q^T Q = I_n.$$

$${}^m_m \begin{bmatrix} \\ \\ \end{bmatrix} = {}^m_n \begin{bmatrix} \\ \\ \end{bmatrix} {}^n_n \begin{bmatrix} \times & \times \\ & \times \\ & & \times \end{bmatrix}$$

↑
ortho

↑
upper triangular.

$$\min \| \beta \tilde{e}_1 - \tilde{H}_m y_m \|. \quad \tilde{H}_m = QR.$$

$$\Downarrow$$

$$\min \| \beta Q^T \tilde{e}_1 - R y_m \|.$$

If \tilde{H}_m has full col rank.

$$R y_m = \beta Q^T \tilde{e}_1 \Rightarrow y_m = R^{-1} (\beta Q^T \tilde{e}_1)$$

$$x_{m+1} = x_0 + \tilde{V}_m y_m.$$

Doing GMRCS more efficiently.

Remove redundant operations in QR decomposition.

Given's rotation.

$$\min \| Q_m^T \cdots Q_2^T Q_1^T \tilde{H}_m y_m - Q_m^T \cdots Q_1^T \beta e \| \text{eq. size } (m+1) \times (m+1)$$

$$\min \left\| \begin{bmatrix} R \\ 0 \cdots 0 \end{bmatrix} y_m - \begin{bmatrix} r_1 \\ s \end{bmatrix} \right\| = |s|. \text{ when } Ry_m = r_1$$

Efficient because

① More efficient QR. Reuse previous inf

② Automatically keep track of error

w.o. forming x_m .

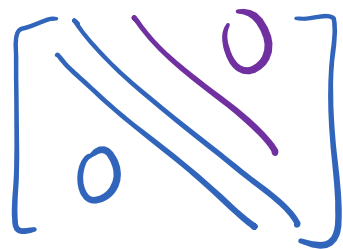
Go back to Pos. Def. matrices $A \in \mathbb{R}^{N \times N}$.

Assume do Arnoldi.

$$A V = V H. \rightarrow \text{upper Hess.} \quad V^* V = I.$$

$$V^* A V = H$$

sym



Tridiagonal. 3-term
recurrence.

Efficient Arnoldi + Sym \rightarrow Lanczos.

Alg. (Lanczos. *conceptual. no reorthogonalization*)

Pick v_1 , $\|v_1\|=1$, $\beta_1=0$, $v_0=0$.

for $k=1, \dots, m$.

$$\tilde{v}_{k+1} = A v_k - \beta_k v_{k-1} \leftarrow \text{ortho. w. } v_{k-1}$$

$$\alpha_k = (v_k, \tilde{v}_{k+1})$$

$$\tilde{v}_{k+1} = \tilde{v}_{k+1} - \alpha_k v_k \leftarrow \text{orth. w. } v_k$$

$$\beta_{k+1} = \|\tilde{v}_{k+1}\|, \quad v_{k+1} = \tilde{v}_{k+1} / \beta_{k+1} \leftarrow \text{normal.}$$

end

$$A V = V T \quad , \quad T = \begin{pmatrix} \alpha_1 & \beta_1 & & 0 \\ & \ddots & \ddots & \\ 0 & & \beta_{N-1} & \alpha_N \end{pmatrix}$$

Lanczos + FOM \rightarrow CG.