
Database Systems

Lecture #2 E/R Model

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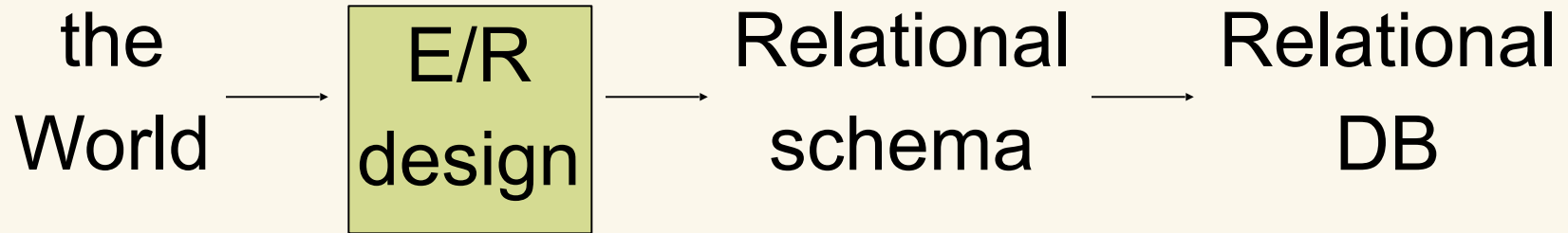


Agenda

- Last time: intro
- This time: E/R model
 1. Identify entity sets, relations and attributes
 2. One-one, one-many, many-many relations
 3. Simple ER diagrams to model a situation
 4. 3-way relationships, multiple roles, subclasses
- Design issues
 1. Simplicity
 2. Redundancy
 3. Replacing a relationships with entity sets

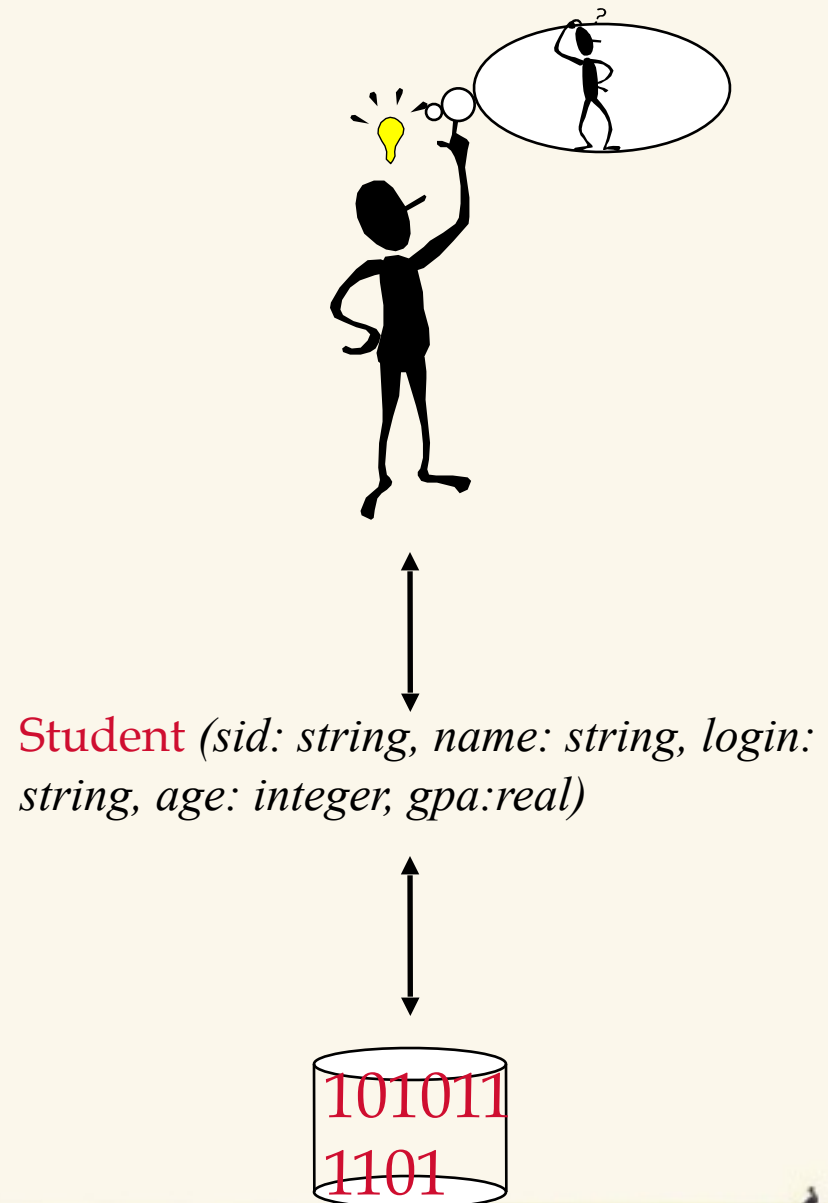


DB development path



Data Models

- DBMS models real world
- *Data Model* is link between user's view of the world and bits stored in computer
- Many models exist
- We will ground ourselves in the Relational Model
 - clean and common
- But use the Entity-Relationship model as a middle ground for design



Entity Relationship (E/R) Model

- A popular data model – useful to database designers
 - Graphical representation of miniworld
 - E/R design translated to a relational design
 - then implemented in an RDBMS
 - Elements of model
 - Entities
 - Entity Sets
 - Attributes
 - Relationships (*!= relations!*)
- /* 注意：联系 (*!=关系*) */



E/R Model: Entity Sets

- Entity: like an object
 - Particular instance of a concept
- *Entity set*: set of one sort of entities or a concept
 - All with same attributes
- Represented by a rectangle:

World Leader

- A “good” entity set
 - Common properties
 - Correspond to class of phys. or bus. objects
 - E.g., Employees, products, accounts, grades, campaigns, etc.

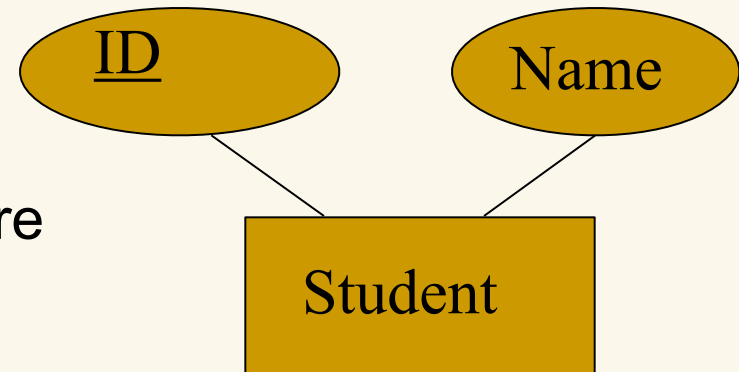


E/R Model: Attributes

- Properties of entities in entity set
 - ❑ Like fields in a struct
 - ❑ Like columns in a table/spreadsheet
 - ❑ Like data members in an object
- Values in some domain (e.g., ints, strings)

```
struct student
{
    int id;
    char* name;
};
```

- Represented by ovals:
- Assumed atomic
 - ❑ But could have limited structure
 - ❑ ints, strings, etc.



- Each entity set has a *key* (*underlined attribute*).

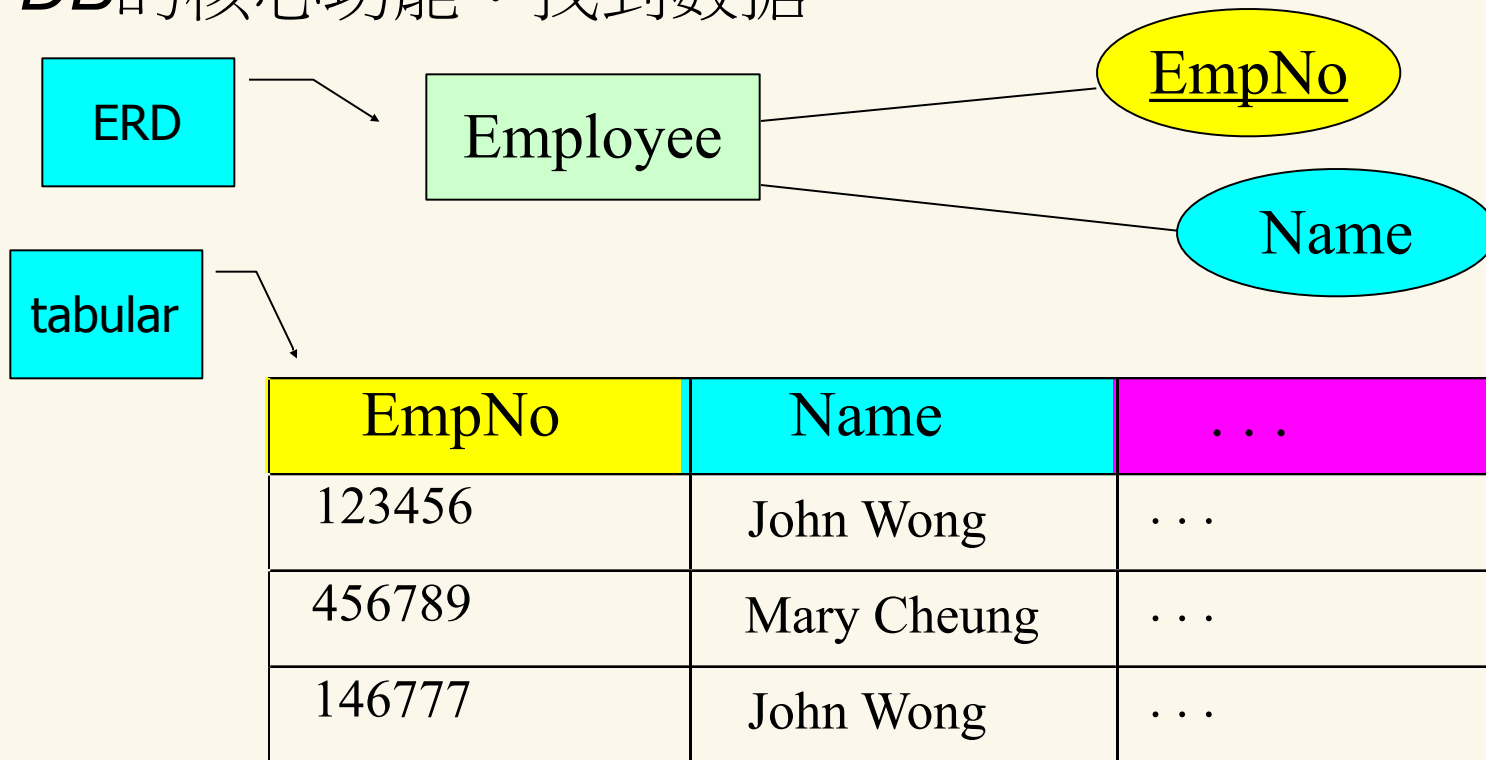


Key Attributes

Super Key

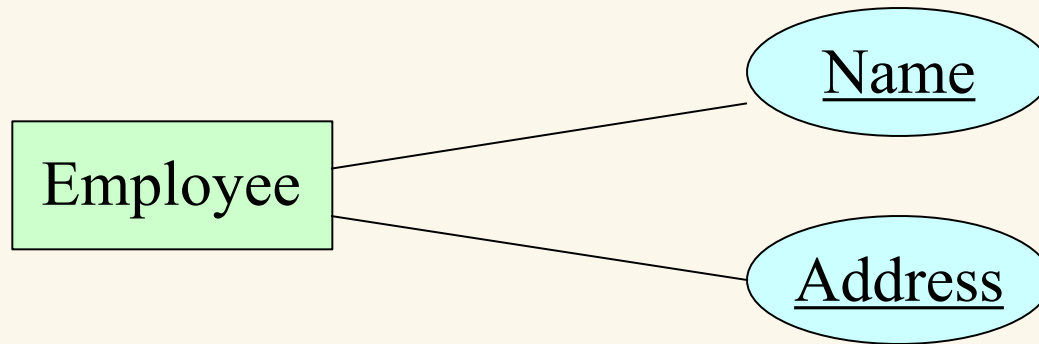
- A set of **attributes** that can uniquely identify an entity (唯一标识一个实体)

DB的核心功能：找到数据



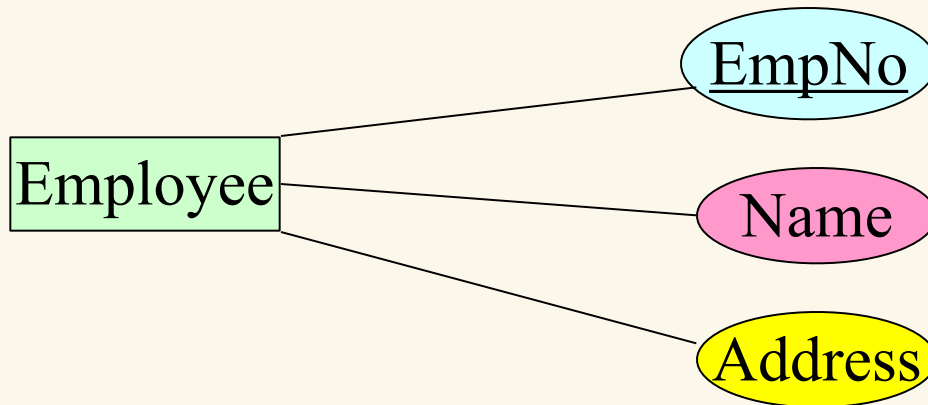
Key Attributes

- **Composite key:** *Name* or *Address* alone cannot uniquely identify an employee, but together they can!



Key Attributes

- An entity may have more than **one** key
 - ❑ e.g., EmpNo, (Name, Address)
 - ❑ only one is selected as the key. (sometimes called the **Primary key**)

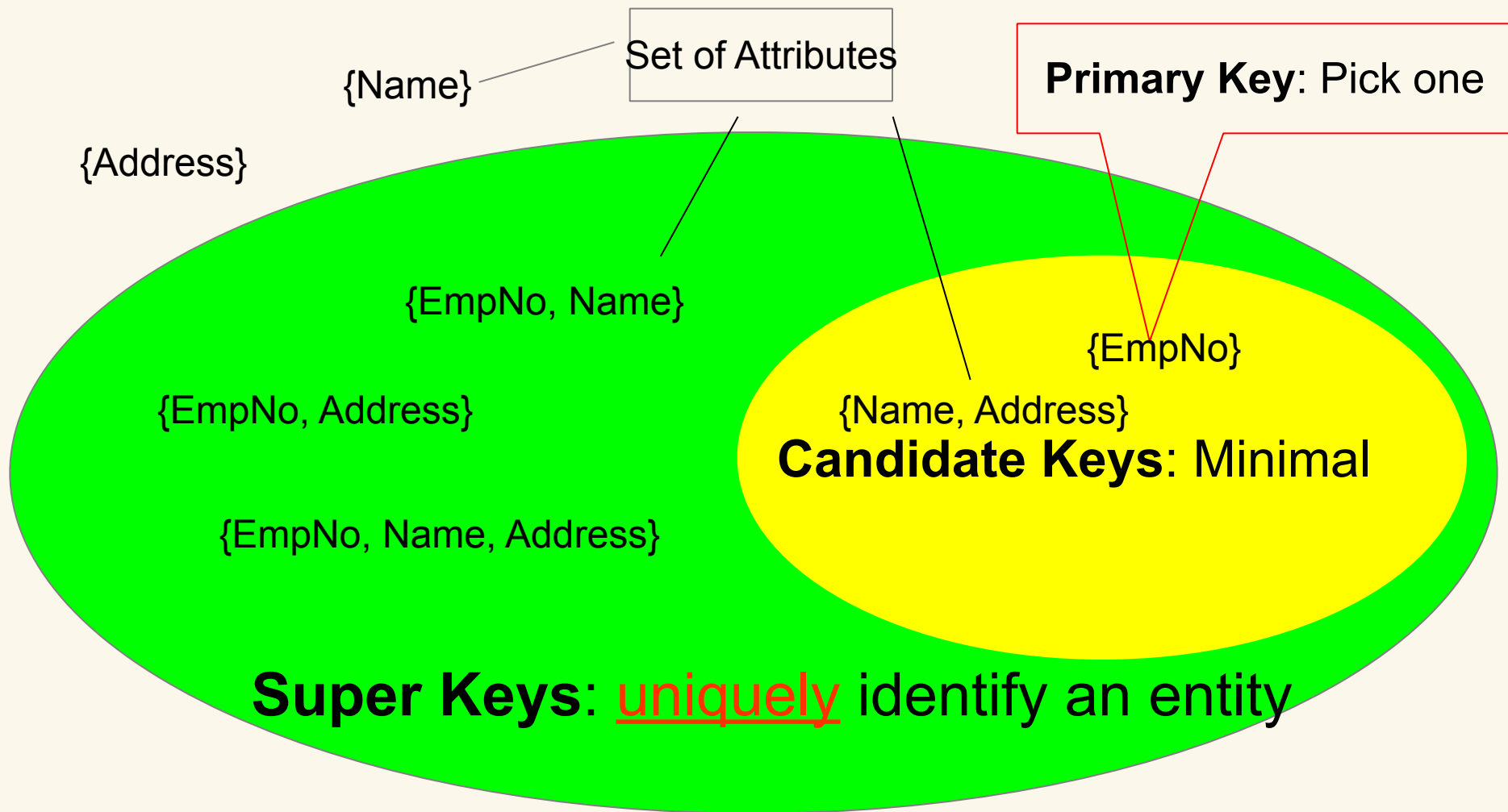


In many cases, a key is artificially introduced (e.g., EmpNo) to make applications more efficient.

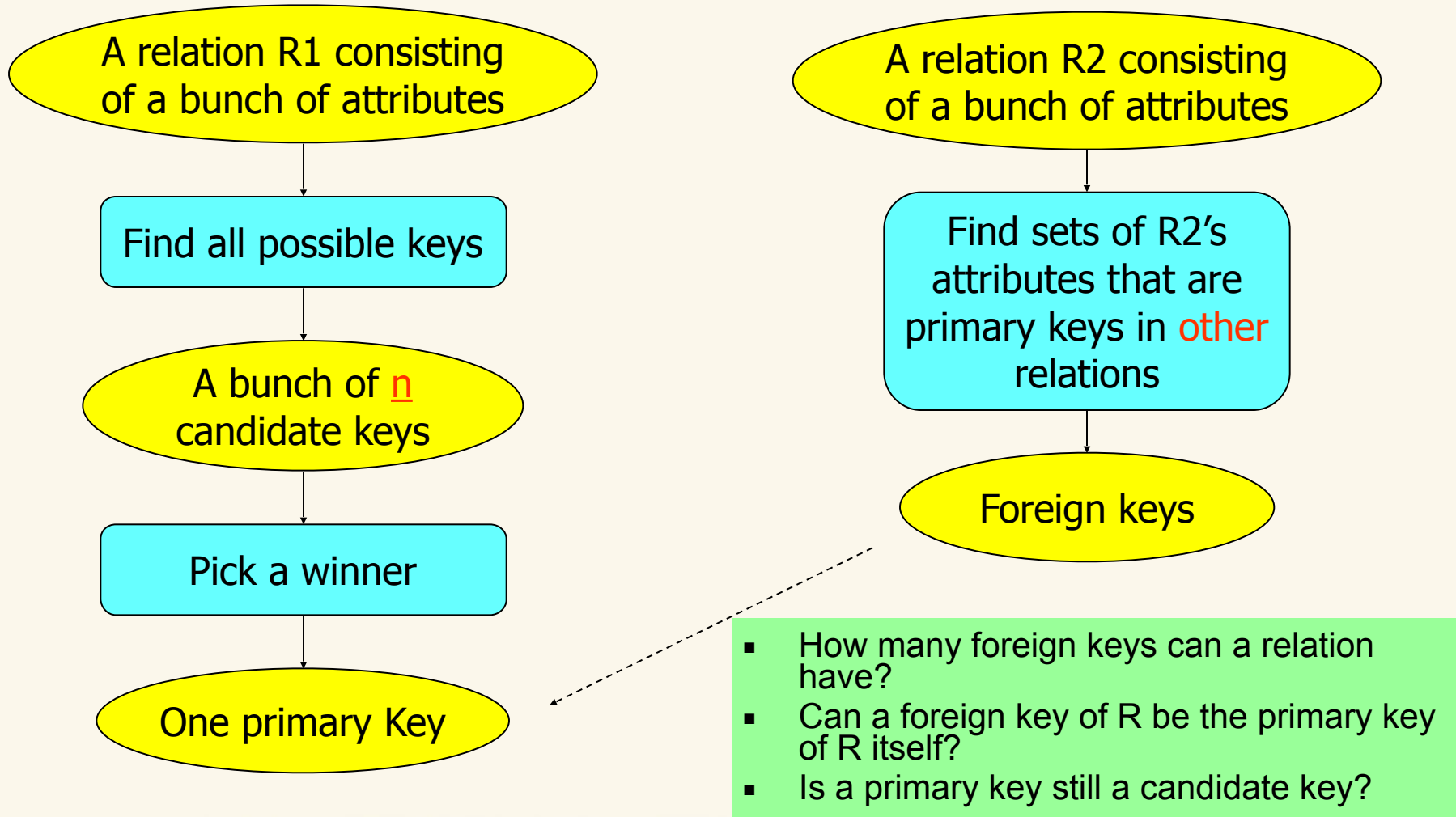
Question: does a desk has a key?



Overview of Various Kinds of Keys



Overview of Various Kinds of Keys



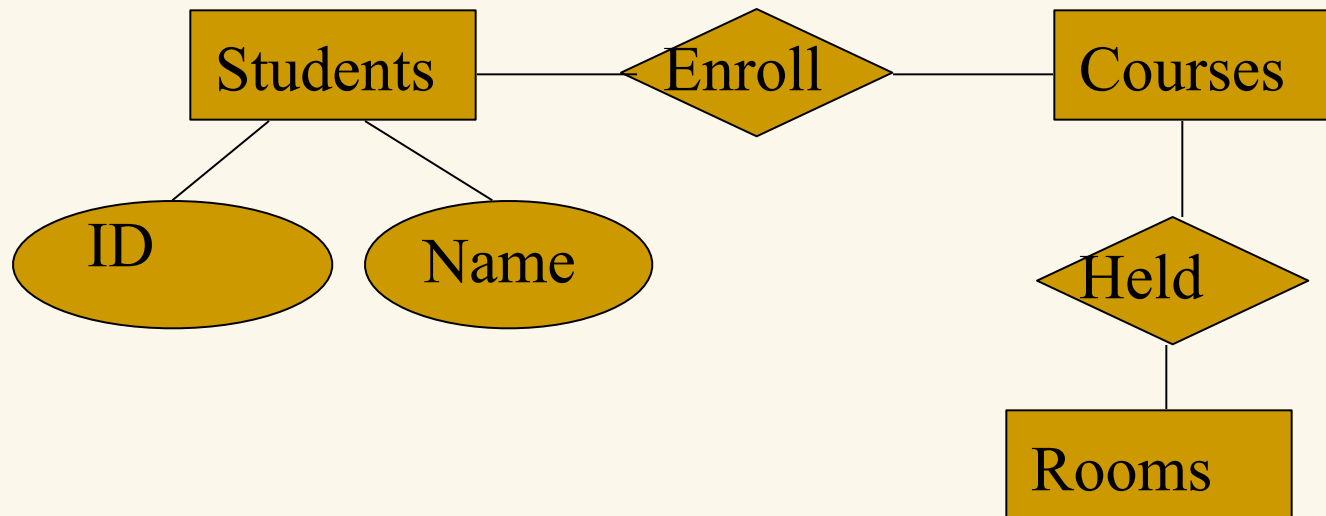
E/R Model: Relationships

- Connect two or more entity sets
 - e.g. students *enroll* in courses
 - Binary relationships: connect two entity sets
 - most common
 - Multiway relationships: connect several ESs
- Represented by diamonds:



E/R Model: Relationships

- Students *Enroll* in courses
- Courses are *Held* in rooms
- The E/R data model:



Set Theory

- Invented by Georg Cantor
 - Great 19th-C German mathematician
- Big set theory results in 1870s-1890s
- Controversial at the time
 - Kronecker: “humbug”
 - First rigorous math of the “actual infinite”
- we’ll mostly deal with *finite* sets



A little set theory

- A mathematical *set* is a collection of *members*
- A set is defined by its members
 - “Are you in or are you out?”
 - No other structure, no order, no duplicates allowed
- Sets specified by listing:
 - $\{1, 2, 3, \dots\} = \mathbf{N}$
 - $\{1, 2, \text{George Bush}\}$ (tho usually homogeneous sets in DBMS...)
- Or by “set-builder” notation:
 - $\{x \text{ in } \mathbf{N} : 2 \text{ divides } x\} = ?$
 - $\{x \text{ in Presidents} \mid \text{reelected}(x)\} = ?$
 - $\{2x : x \text{ in } \mathbf{N}\} = ?$



A little set theory

- One set can be a *subset* of another (which is a *superset* of it)
 - ReelectedPresidents is a subset of Presidents
 - Also, RP is a *proper subset* (真子集) of Pres – some lost reelection
- Given two sets X and Y, the *cross product* or *Cartesian product* is
$$X \times Y = \{(x,y): x \text{ in } X, y \text{ in } Y\}$$

= the set of all *ordered pairs*
- Important: $(x,y) \neq \{x,y\}$
- In an *order pair* or tuple
 - Order matters; duplicates are allowed

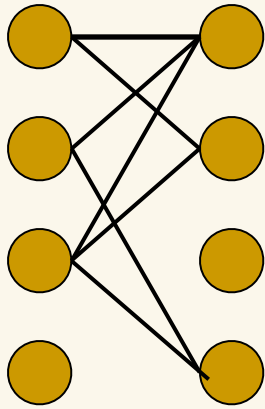


A little set theory

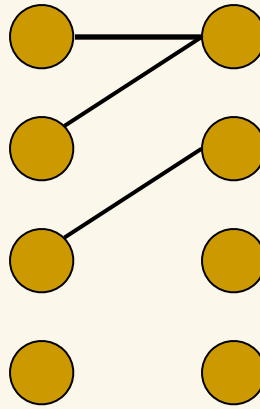
- Mathematically, a *relation* between X and Y is just a *subset of* $X \times Y$ = all those pairs (x,y) s.t. x is related to y
- Example: owner-of O on People, Cats
 - $O(\text{MPJ}, \text{Gödel the Cat})$ holds
- The equals relation E on N, N :
 - $E(3,3)$ holds because $3 = 3$
 - $E(3,4)$ does not hold
 - E is still a set: $E = \{(1,1), (2,2), (3,3), \dots\}$
- Father-of relation F on People, People:
 - $F(\text{GHWB}, \text{GWB})$ holds
 - $F(\text{GWB}, \text{GHWB})$ does not hold
 - a Relations aren't necessarily *symmetric*



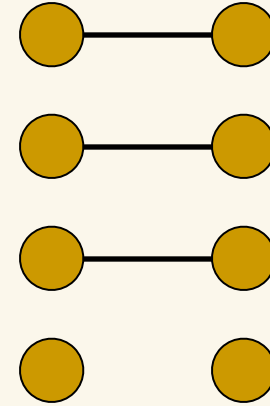
Multiplicity of Relationships



Many-many



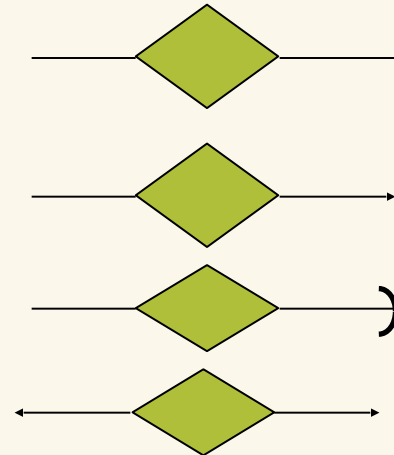
Many-one



One-one

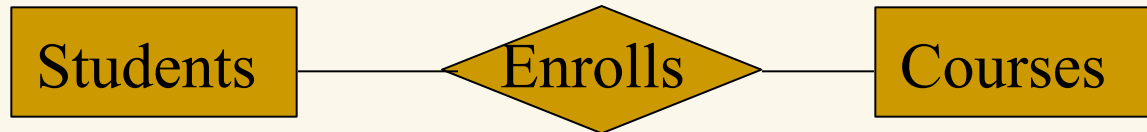
Representation of relationships

- No arrow: many-to-many
- Sharp arrow: many-to-one
- Rounded arrow: “exactly one”
 - “key constraint”
- One-one:

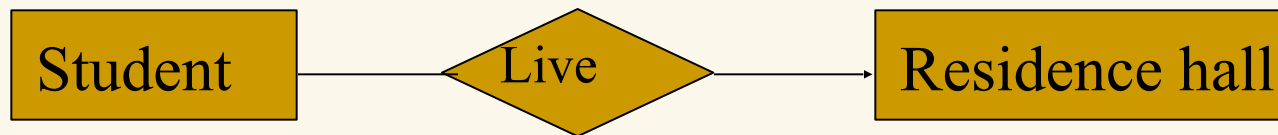


Multiplicity of Relationships

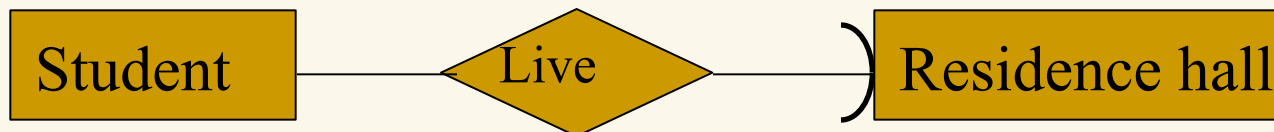
Many-to-many:



Many-to-one: a student living in a residence hall



Many-to-exactly-one: a student must live in a residence hall

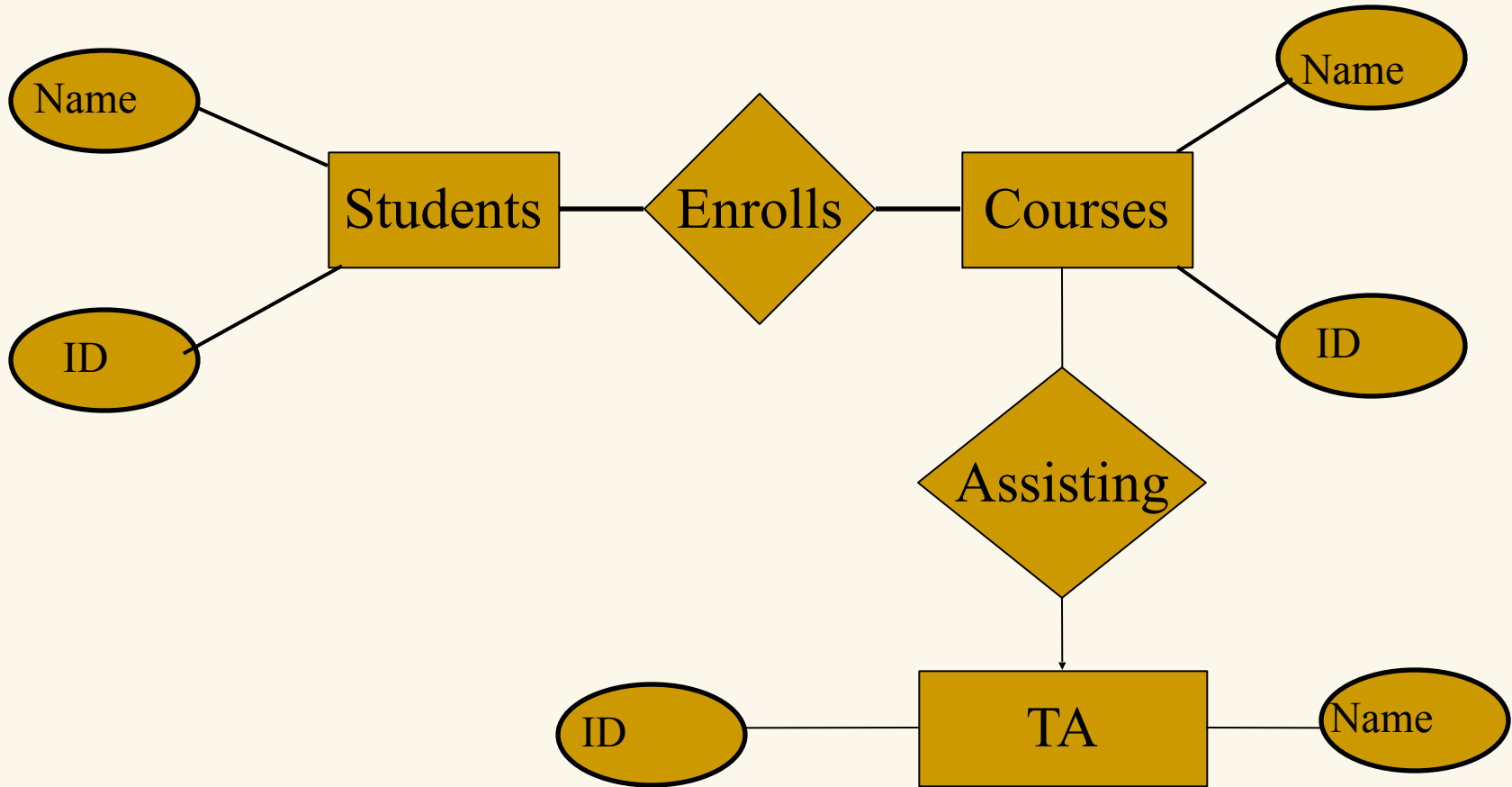


Multiplicity, set-theoretically

- Assume *no vars below are equal*
- Many-one means:
 - if (x_1, y_1) in R then (x_1, y_2) cannot be in R
- One-many means:
 - (Y, X) is many-one
- One-one means:
 - if (x_1, y_1) in R , then *neither* (x_2, y_1) nor (x_1, y_2) can be in R
- Notice: one-one is stronger than many-one
- One-one *implies* both many-one and one-many



E/R Diagram

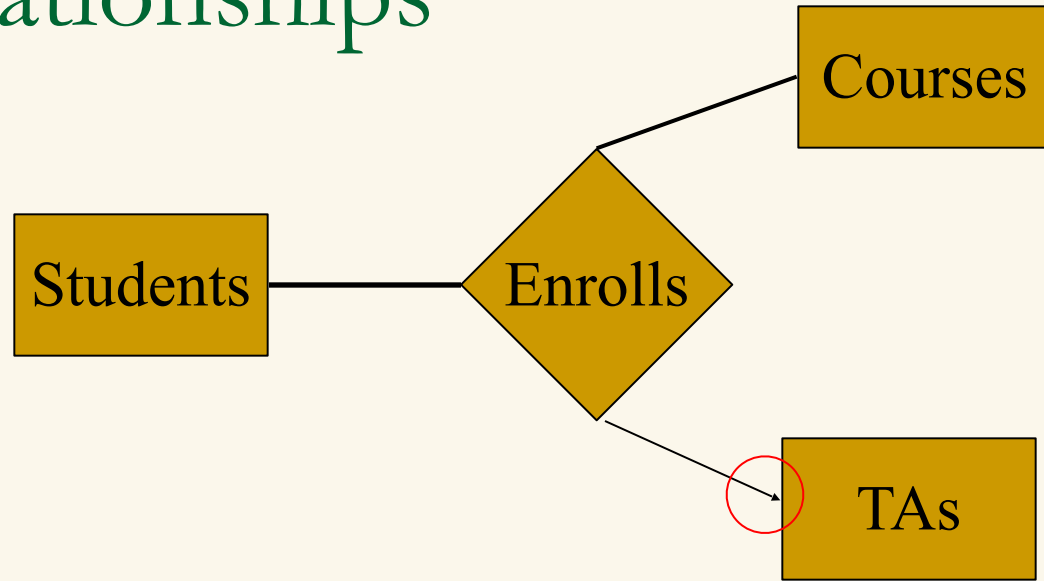


E/R Diagrams

- OK if each TA is a TA of all students
 - Student and TA connected only through Course
- But what if students were divided among multiple TAs?
 - Then a student in SE-304 would be related to only one of the TA's for SE-304—which one?
 - Schema doesn't store enough info
- 3-way relationship is helpful here



Multiway Relationships



Enrolls entries:

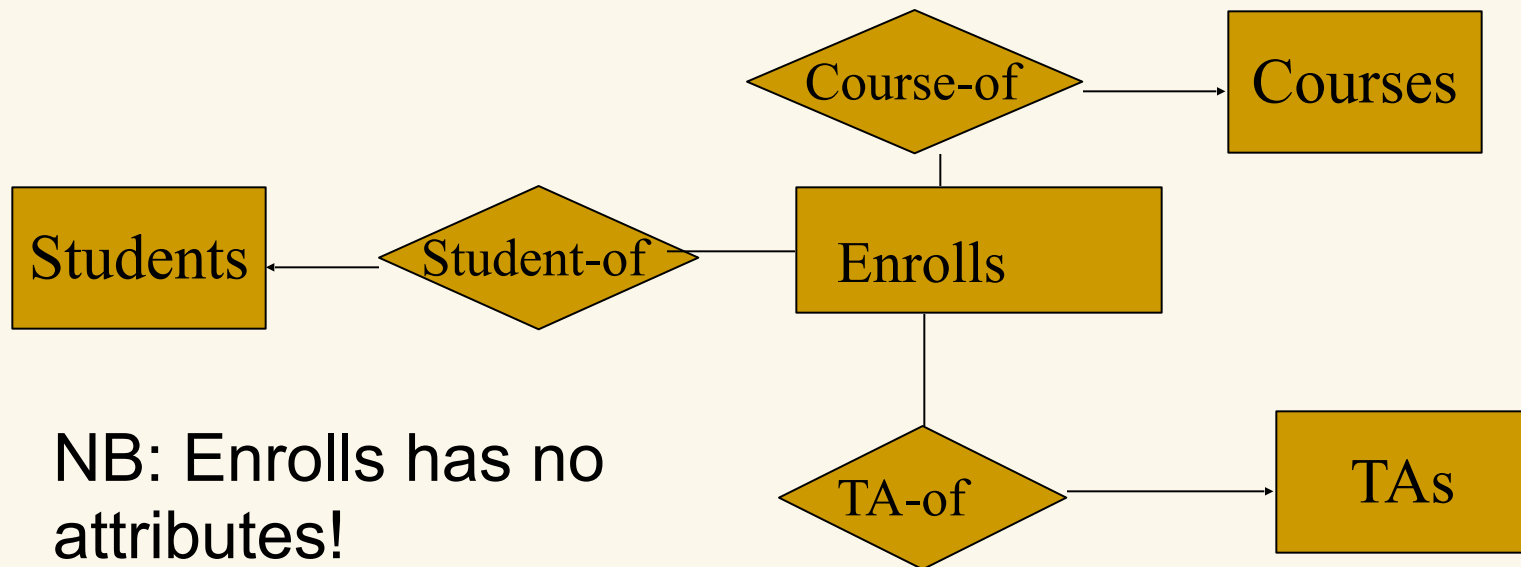
Student	Course	TA
John	SE-304	Chen
Mary	SE-304	Li
Alice	SE-304	Zhang
Mary	SE-304	Wang
...

NB: *Enrolls* determines *TA*:
(student, course) \longrightarrow at most one TA



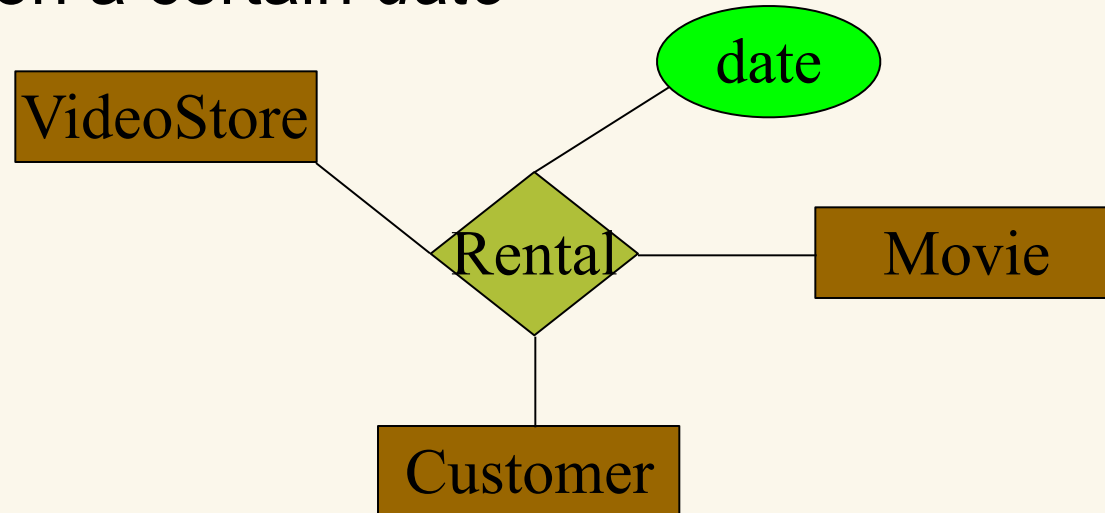
Converting multiway relships to binary

- Binary relationships are as strong as multiway
- Replace relationship with *connecting entity set* and multiple binary relationships



Second multiway e.g.: renting movies

- Scenario: a *Customer Rents a Movie* from a *VideoStore* on a certain *date*

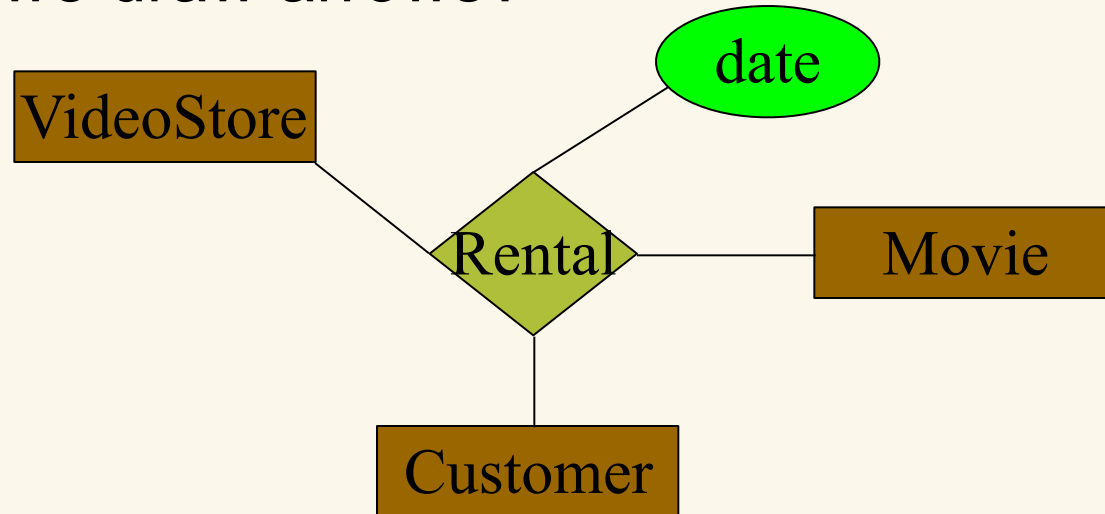


- date should belong to the *fact* of the renting
 - Relationship attribute



Second multiway e.g.: renting movies

- Where can we draw arrows?

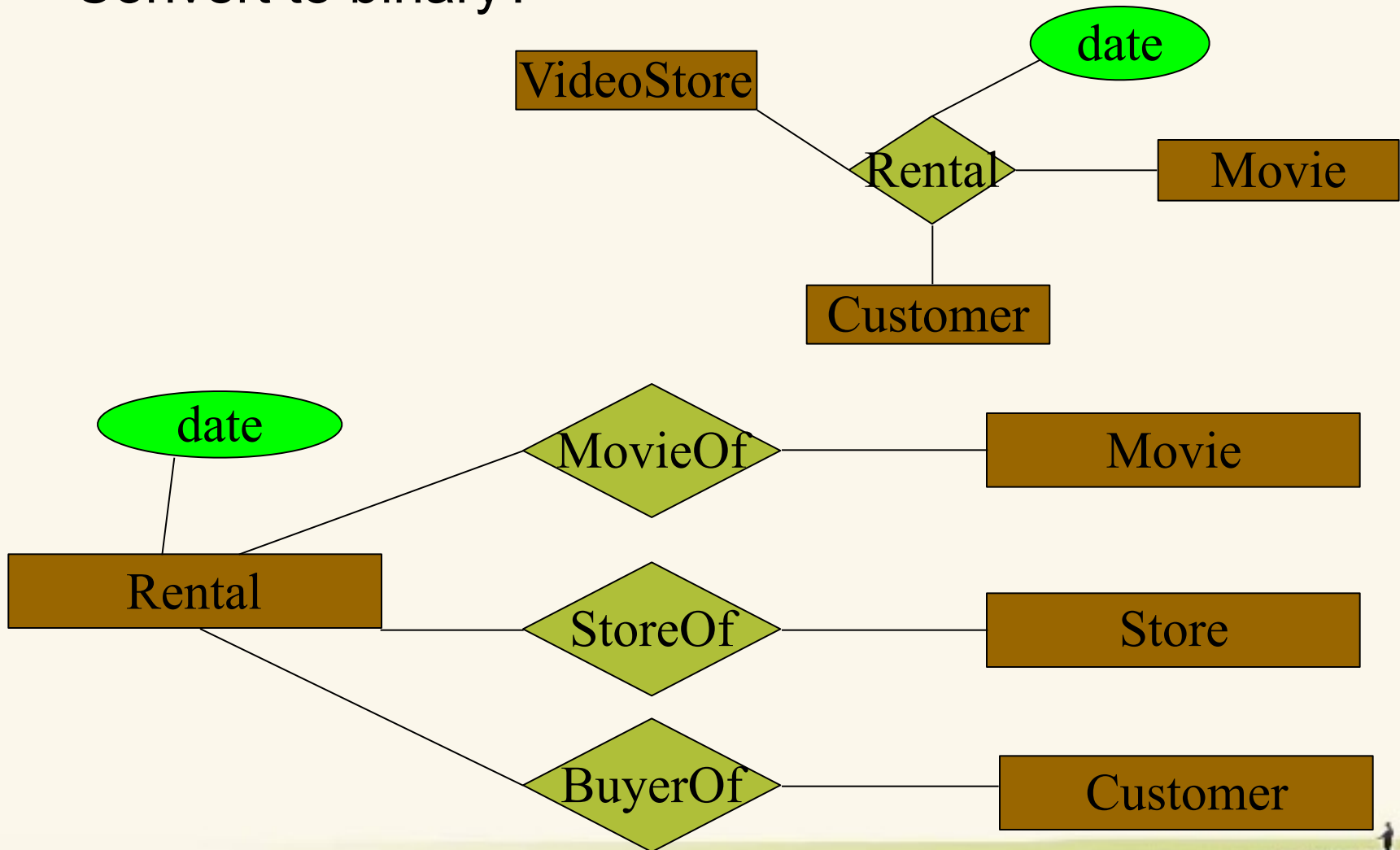


- (store, video, customer) → date ?
- (store, video, date) → customer ?
- (store, date, customer) → video ?
- (video, date, customer) → store ?



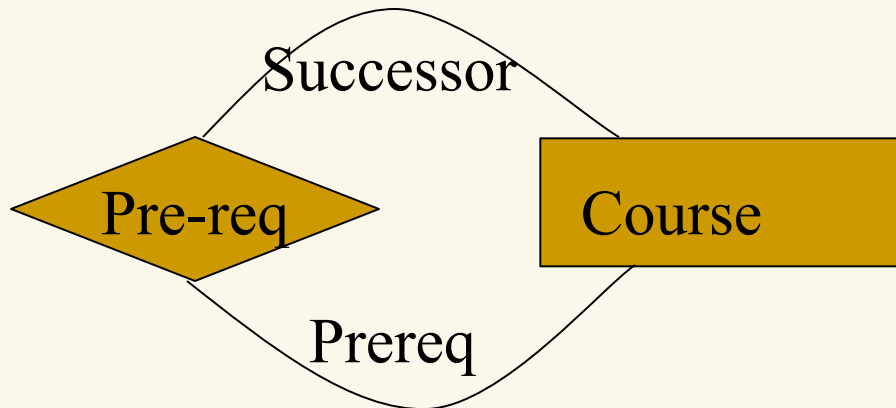
Second multiway e.g.: renting movies

- Convert to binary?



Roles in relationships

- Entity set appears more than once in a relship
 - Generally distinct *entities*
- Each appearance is in a different *role*
- Edges labeled by roles



Course (Pre-req)	Course (Successor)
Accounting	Finance-I
Finance-I	Derivatives
Finance-I	Finance-II
Calculus	Derivatives

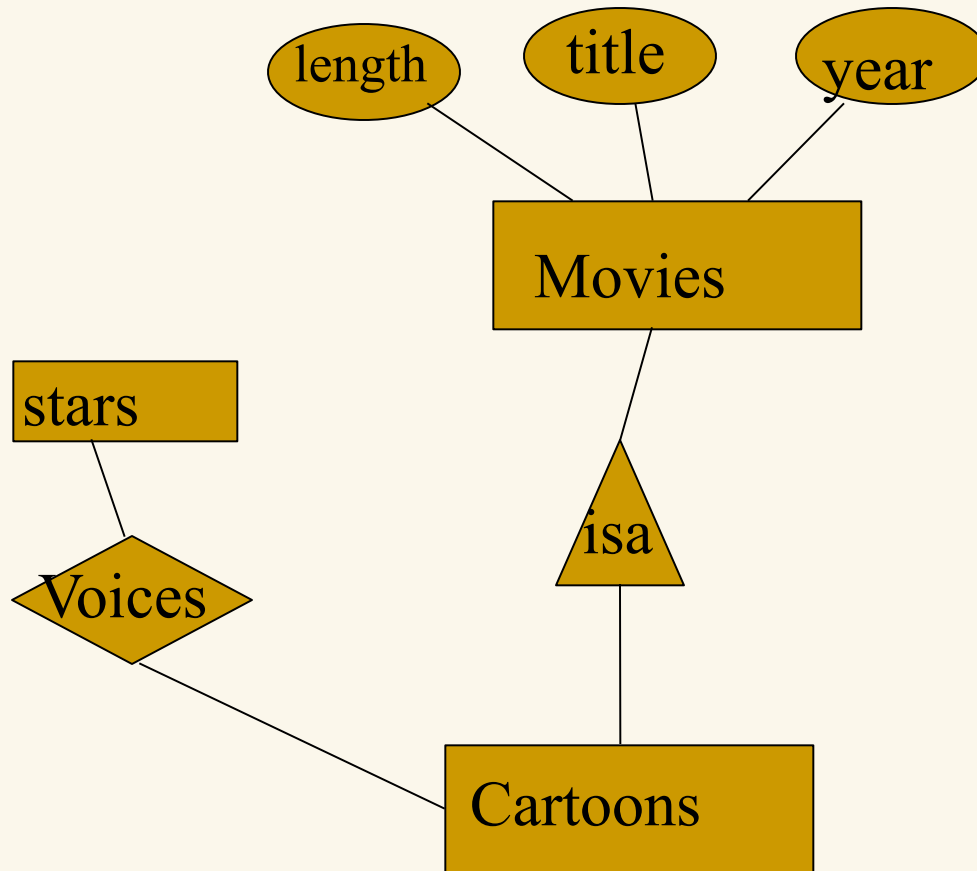


Subclasses in the E/R model

- Some ESs are special cases of others
- Conversely: some are generalizations
 - Mammals, humans, students, grad students
 - NB: These aren't *members* but subclasses
- Subclass A is a B
 - Represented by a triangle
 - Root is more general



Subclasses



New topic: Design Issues

- Faithfulness (如实、正确)
- Avoiding redundancy (避免冗余)
- Simplicity (简单性)
- Choice of relationships
- Picking elements



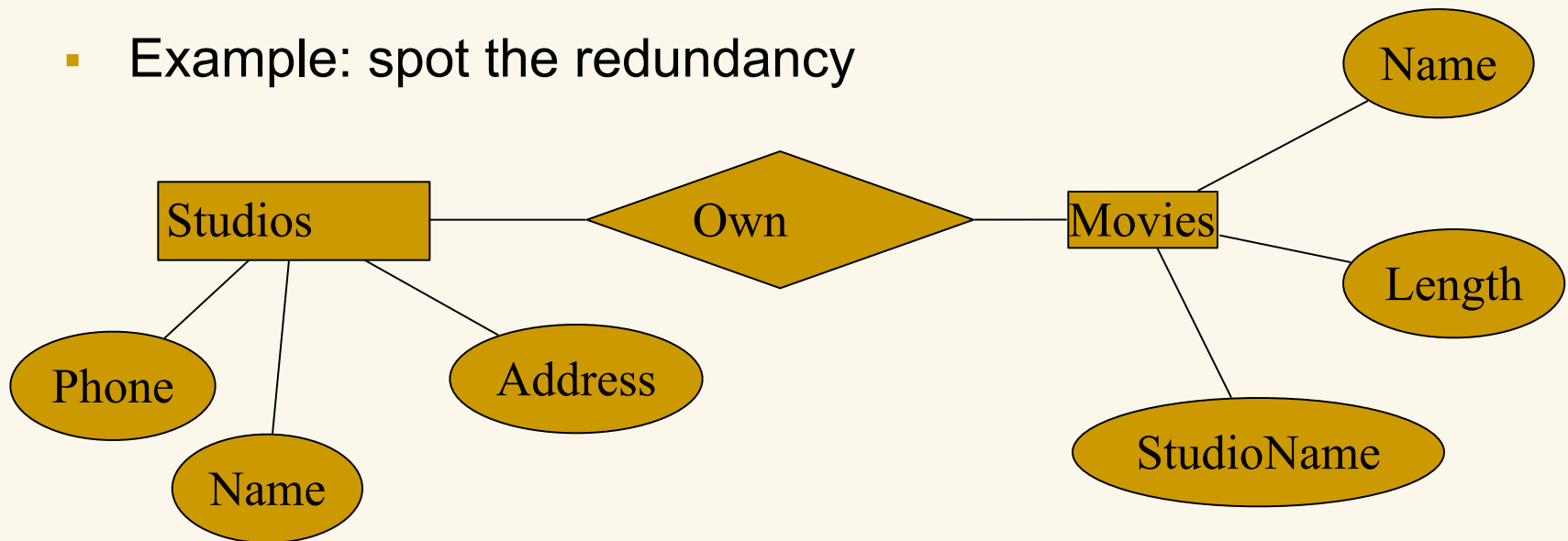
Faithfulness

- Is the relationship many-many or many-one?
- Are the attributes appropriate?
- Are the relationships applicable to the entities?
- Examples:
 - ❑ Courses & instructors
 - maybe many-one, maybe many-many
 - ❑ Bosses & subordinates
 - maybe one-many, maybe many-many



Avoiding redundancy

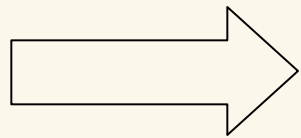
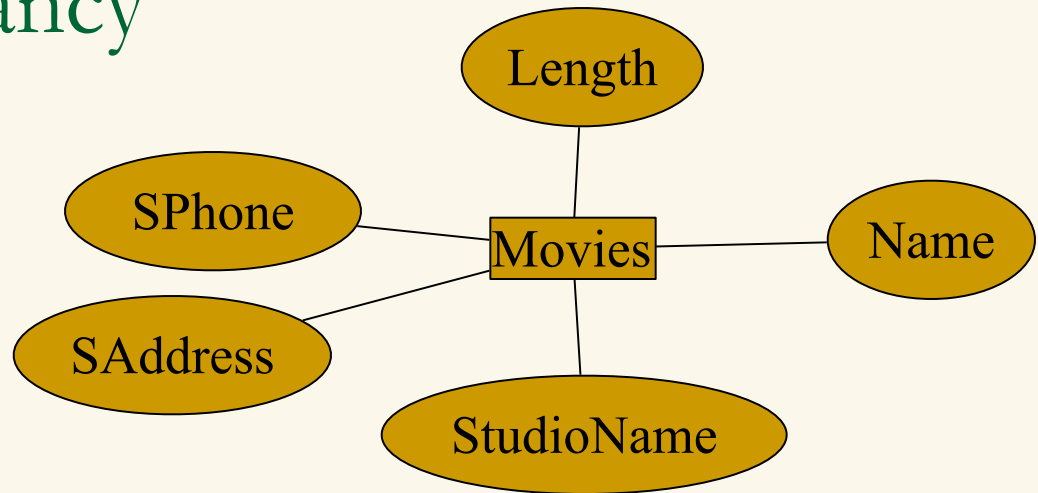
- Say everything once and only once
 - ❑ Minimize database storage requirements
 - ❑ More important: prevent possible update errors
 - One danger: modifying data one place but not the other
- Example: spot the redundancy



Redundancy: Movies “knows” the studio two ways



Spot more redundancy



Name	Length	Studio	SAddress	SPhone
Pulp Fiction	...	Miramax	NYC	212-...
Sylvia	...	Miramax	NYC	212-...
Jay & Sil. Bob	...	Miramax	NYC	212-...
...				

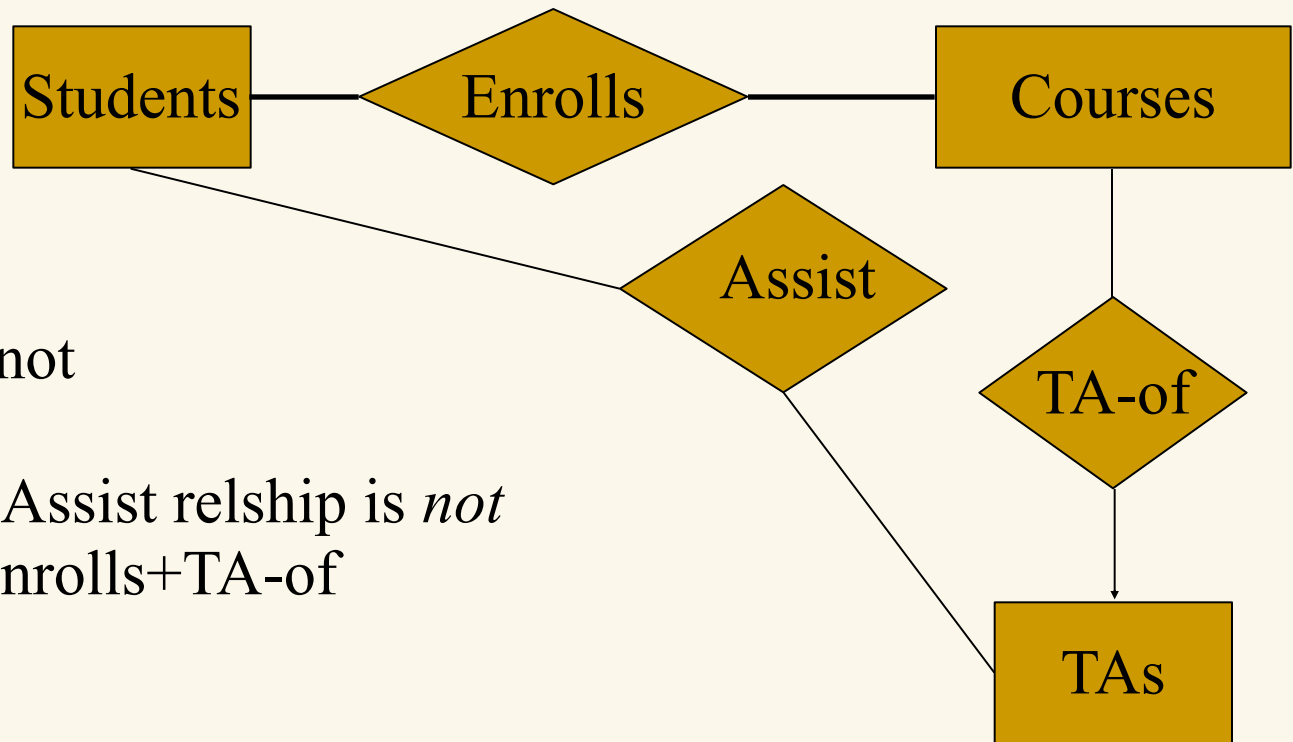
Different redundancy: studio info listed for every movie!



Don't add relships that are implied

Suppose each course again has ≤ 1 TA

Q: Is this good design?



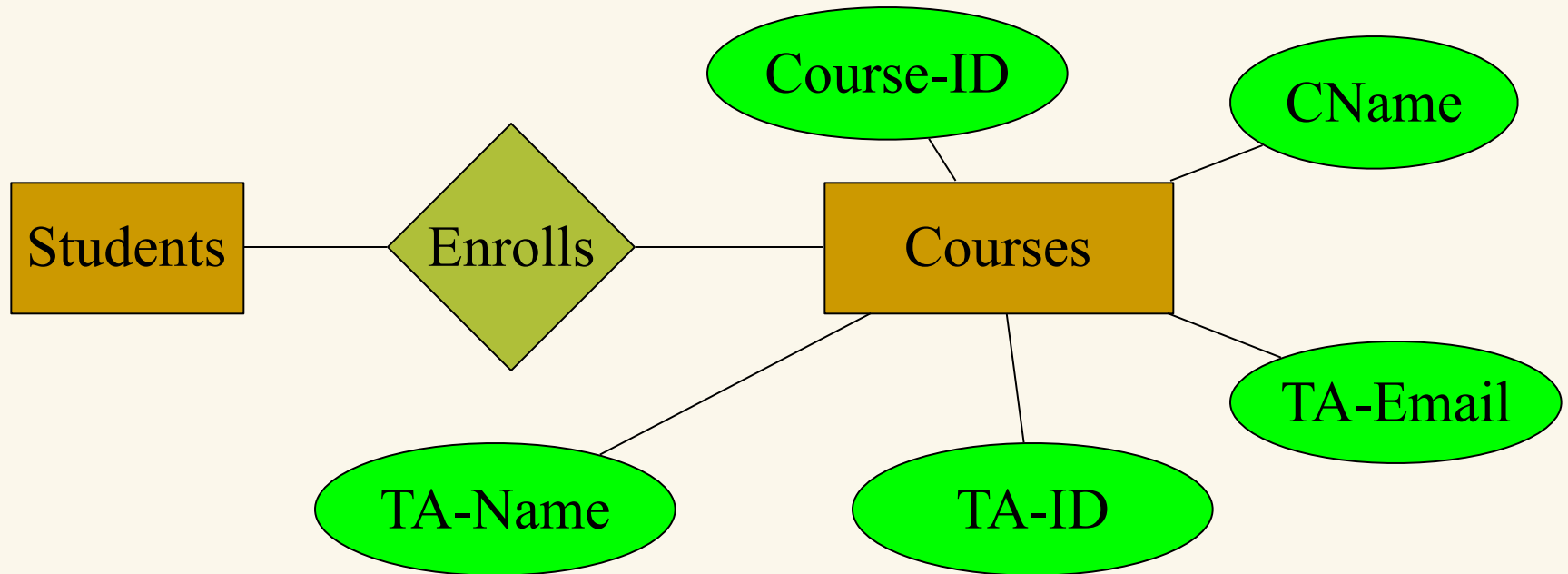
A: probably not

...unless the Assist relship is *not* implied by Enrolls+TA-of



Still more redundancy

Q: What's wrong with this design?



- A:
- ❑ Repeating TA names & IDs — redundant
 - ❑ TA is not TAing any course now — lose TA's data!
 - ❑ TA should get its own ES

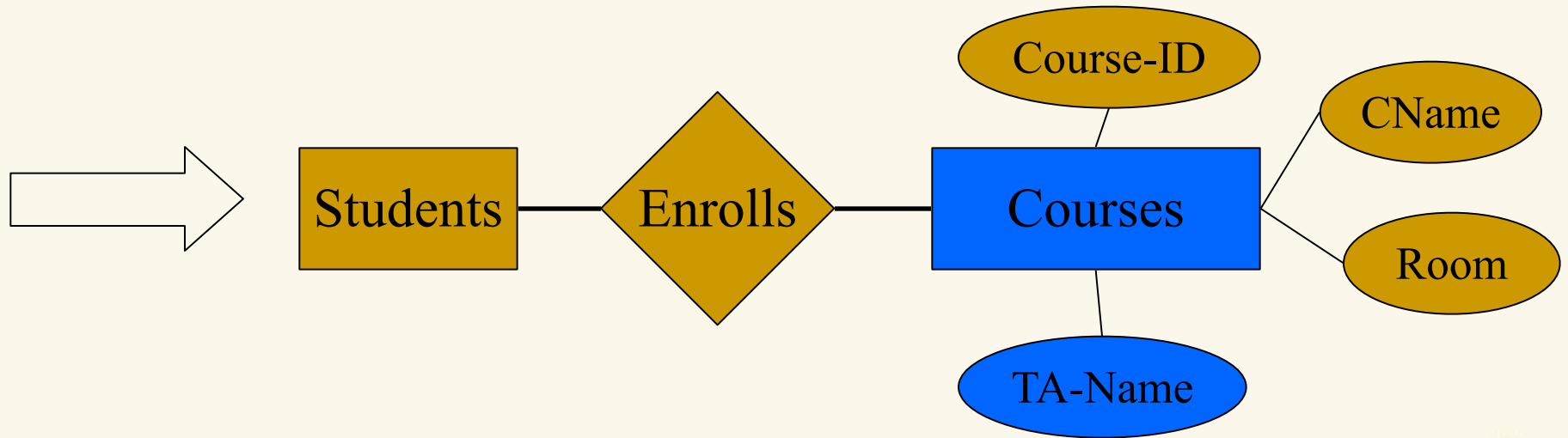
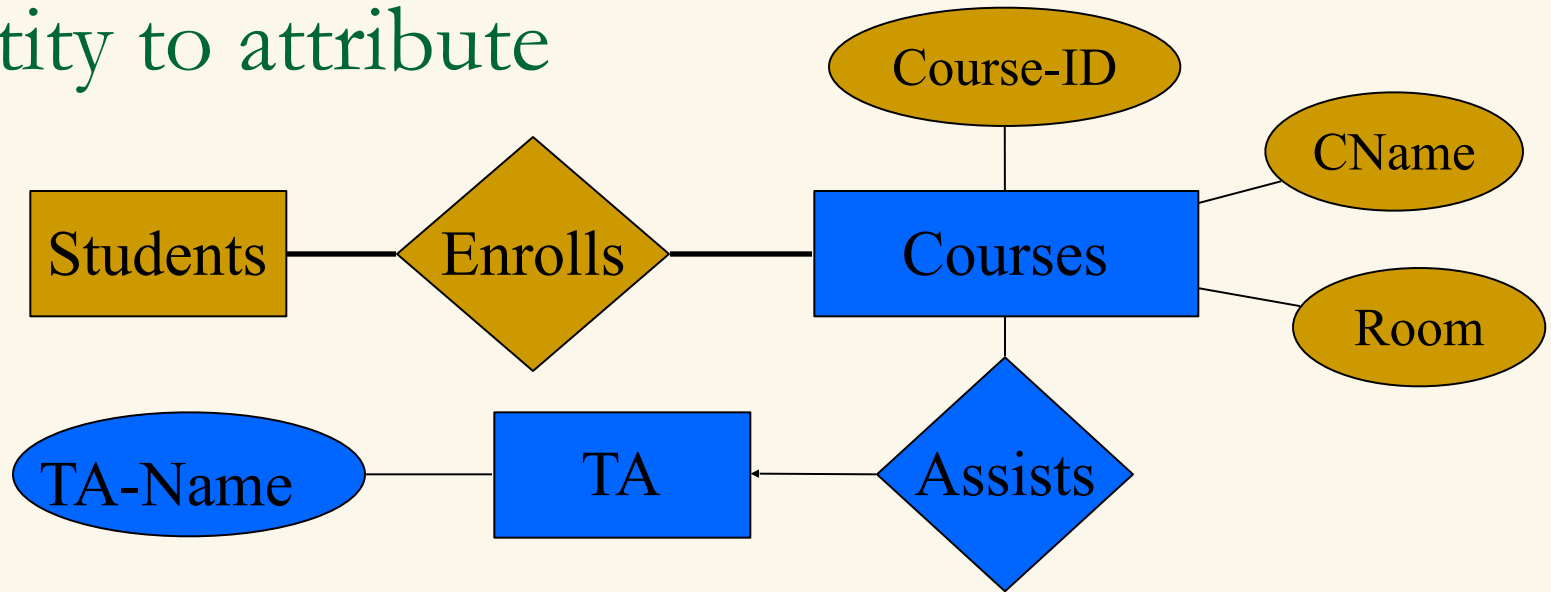


Related issue: entity or attribute?

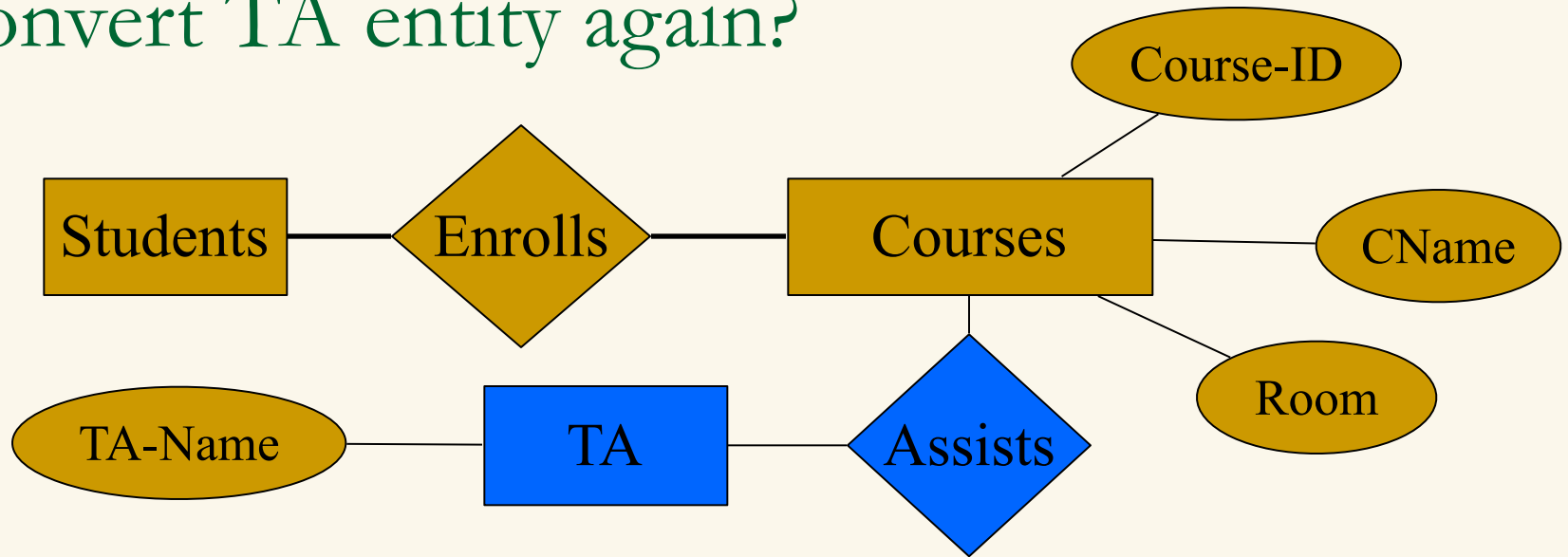
- Some E/Rs improved by removing entities
- Can convert Entity E into attributes of F if
 1. $R:F \rightarrow E$ is many-one (or 1-1)
 2. Attributes for E are *mutually independent*
 - knowing one att val doesn't tell us another att val
- Then
 - remove E
 - add all attributes of E to F



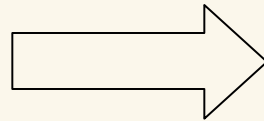
Entity to attribute



Convert TA entity again?



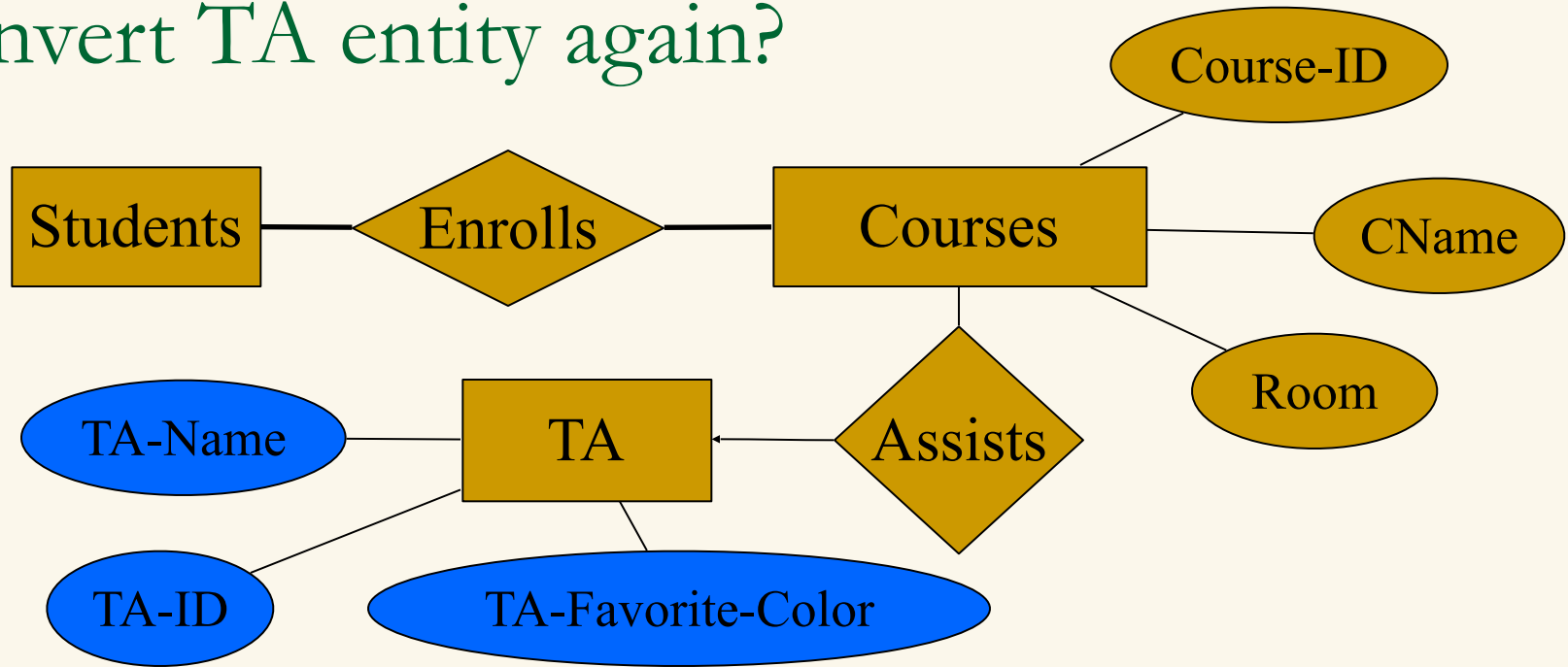
- No! Multiple TAs allowed — redundant course data
- Violates *condition (1)*



CName	CID	Room	TA-Name
DBMS	46	123	Howard
DBMS	46	123	Wesley
...			

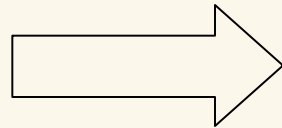


Convert TA entity again?



- No! TA has *dependent fields* — redundant TA data
- Violates *condition (2)*

□ How can it tell?



CName	TA-Name	TA-ID	TA-Color
DBMS	Ralph	678	Green
A.Soft.	Ralph	678	Green
...			



A case Study

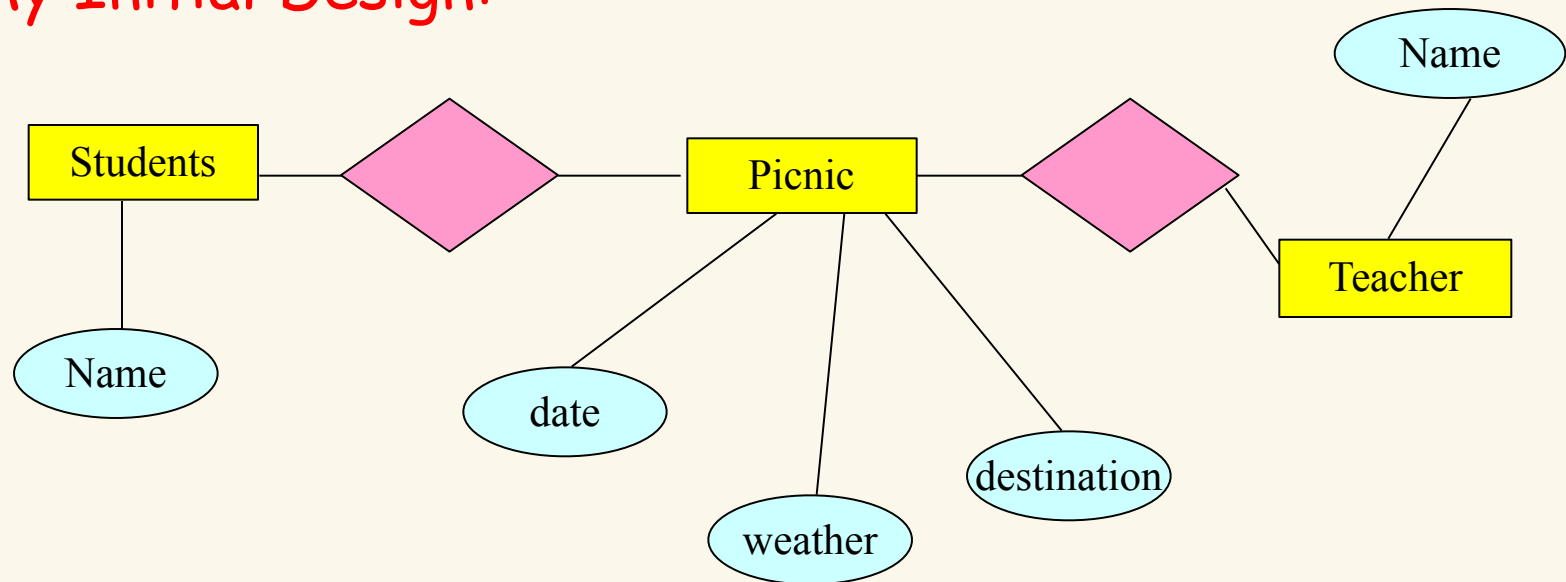
A primary school student writes a composition about a picnic:

Today is Sep 9, the weather is fine.

My classmates, John, Mary and I go to a picnic in Sai Kung.

Our teacher is Ms Wong

My Initial Design:

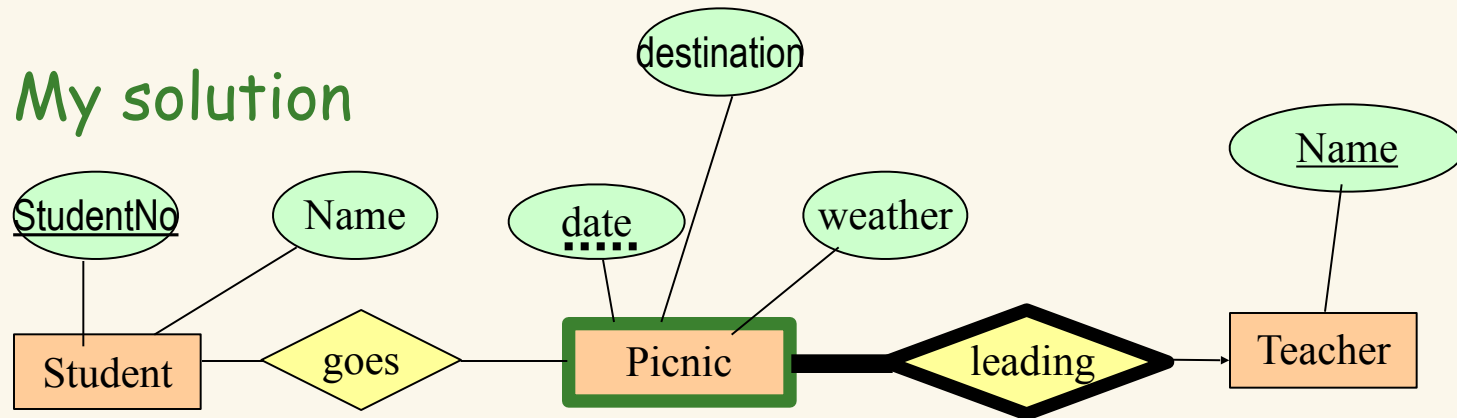


Questions ?

- Why “John”, “Mary”, “Miss Wong” are not in the ER diagram ?
- What do these names tell us ?
- What are the keys of Student, Picnic & Teacher ?
- What are the cardinalities of the relationships ?



My solution

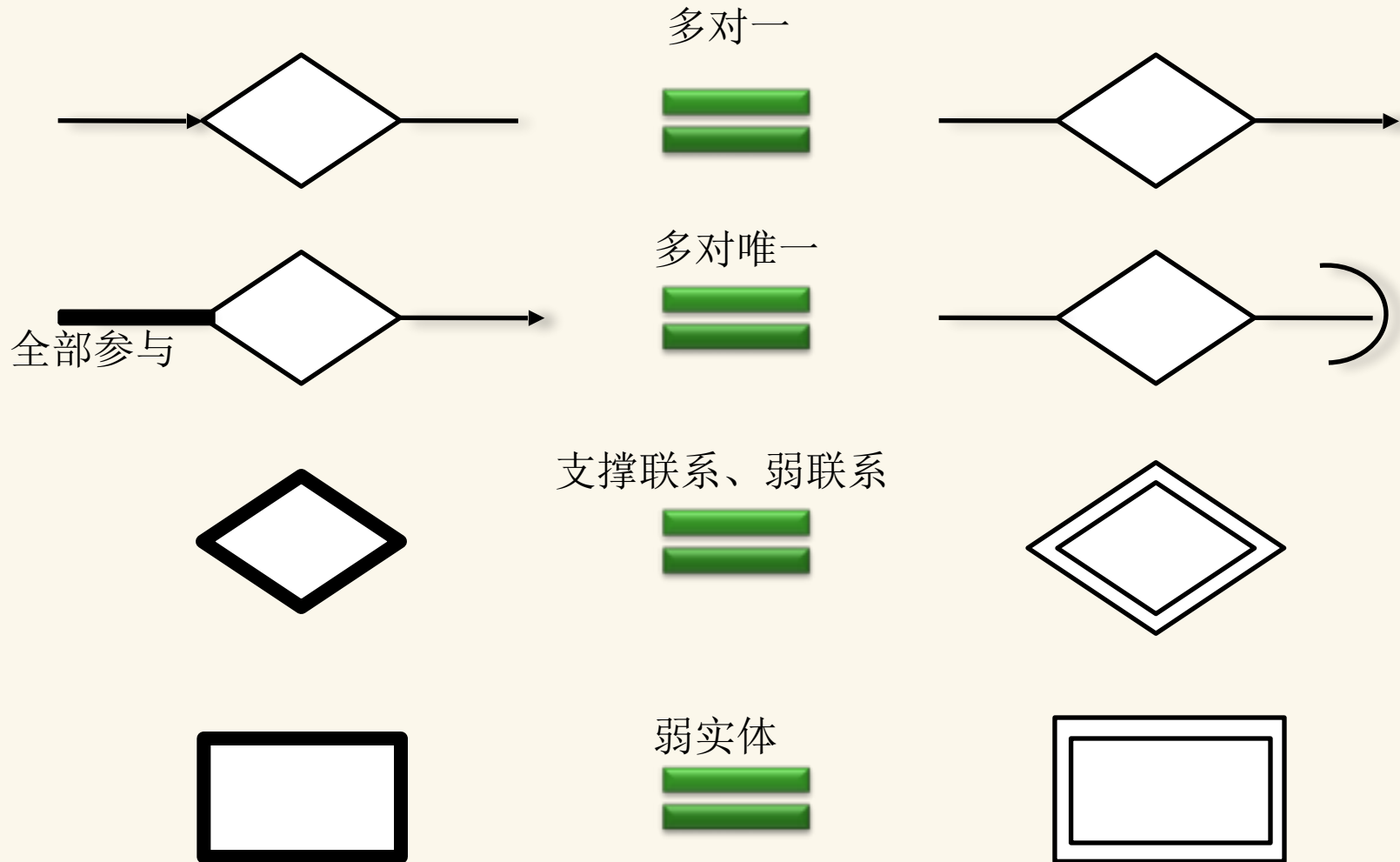


- Every student has an ID number, it is better to keep it in the database and use it as a key
- I bet that there won't be teachers with the same name; otherwise, I'll add employee number and use it as a key
- **goes** is N:M, why ? *A picnic has more than one student participating; also, a student can go to more than 1 picnic. However, this N:M relationship allows a student to go to more than one picnic on the same date*
- **leading** is N:1 , why? *Depends on your assumptions*
 - I assume a teacher can only lead 1 picnic on a certain date, so given the teacher name and the date, I can identify a picnic
- Picnic is made a weak entity. I *could* have added a PicnicNo, but it would be very awkward.

Question:
How to record number of students in a picnic?



Textbook vs. PPT



Review

The DB dev path

How to draw ER Diagram

Concepts of Keys

Multiplicity

Understand relations is a set

How to avoid redundancy

...

