



BY: Zhu_Ming

<CS1010S>

Tutorial 3

Recursion & Iteration
Order of Growth
Higher-Order Function



Lecture Recap

1. Recursion vs Iteration
2. Order of Growth
 - Space Complexity
 - Time Complexity
3. Higher-Order Function



1. Order of Growth

- Describe the limiting behavior of a function when the input value approach to a infinity
- Describe proportions of growth of time/space with respect to the growth of input value
- In another words, measure the efficiency of code

How to measure the efficiency??

- By space/memory and time required



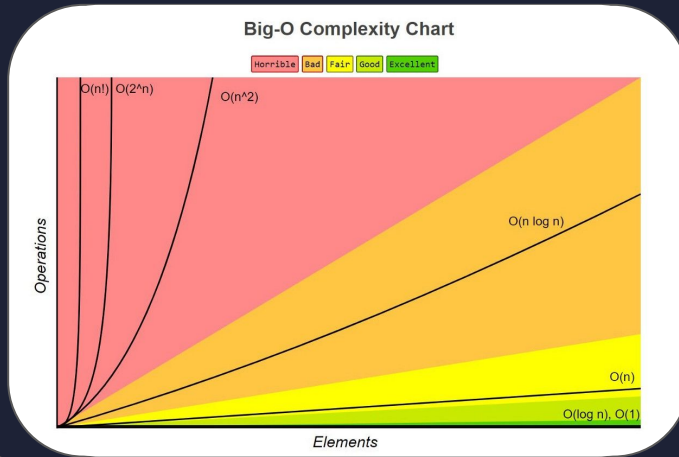
Why not numerical analysis??

Computer	code_iter(i)			code_rec(i)		
	i = 1	i = 10	i = 100	i = 1	i = 10	i = 100
A	1s	10s	100s	1s	1Day	1Year
B	10s	100s	1000s	10s	10Day	10Year

- Is the difference due to computer hardware???
- Is 1s fast or slow??
- How to normalize the data for analysis??
- Is the computer running in optimum condition??
- ...

Why Big-O Notation...

- Because Big-O has NO dependence on the hardware of computer
- Intrinsic properties of the code



Notation	Name
$O(1)$	Constant
$O(\log(n))$	Logarithmic
$O(n)$	Linear
$O(n \log(n))$	Linearithmic
$O(n^{**k})$, where $k \geq 1$	Polynomial
$O(k^{**n})$	Exponential
$O(n!)$	Factorial
$O(n^{**n})$	Tetration

Time complexity

- Depends on the number of operation / evaluation / leaves

Space complexity

- Depends on the number of pending operations / variables / arrays / string

Important!!

Always evaluate the **worst case scenario!!**

Always take note what is the n in $O(n)$ we referring to!!

Take the **most significant term**
 $O(n^2 + 2n) \gg O(n^2)$

IMPORTANT SKILL: CODE TRACING

<https://pythontutor.com/visualize.html#mode=edit>

```
def f(n):  
    return 2 ** n
```

Time Complexity: $O(1)$
Space Complexity: $O(1)$

Computer can do arithmetic
operation very fast

```
def f(n):  
    result = 0  
    for i in range(100):  
        result += 1  
    return result
```

Time Complexity: $O(1)$
Space Complexity: $O(1)$

For loop is independent to
the input!!

```
def f(n):  
    result = 0  
    for i in range(n):  
        result += 1  
    return result
```

Time Complexity: $O(n)$
Space Complexity: $O(1)$

```
def f(n):  
    if n == 0:  
        return 1  
    return 1 + f(n - 1)
```

Time Complexity: $O(n)$
Space Complexity: $O(n)$

```
def f(n):  
    result = 0  
    for i in range(0, n - 1):  
        result += 1  
    return result
```

Time Complexity: $O(n-1) \gg O(n)$
Space Complexity: $O(1)$

Simplify by taking the most significant term

```
def f(n):  
    if n == 0:  
        return 1  
    else:  
        return f(n - 1) + f(n - 1)
```

Time Complexity: $O(2^n)$
Space Complexity: $O(n)$

It is a binary tree!!


```
def f(n):  
    if n == 0:  
        return 1  
    else:  
        temp = f(n - 1) 0(n)  
        return temp + temp 0(1)
```

Time Complexity: $O(n)$
Space Complexity: $O(n)$

Notice that this is a linear tree!!

```
def count_string(s):  
    result = 0  
    for i in range(len(s)):  
        result += 1  
    return result
```

Time Complexity: $O(n)$, where n is $\text{len}(s)$
Space Complexity: $O(1)$

```
def concatenation(n):  
    str_ = ""  
    lst_ = []  
    tup_ = ()  
    for i in range(n):  
        str_ += "1"  
        lst_ += [1]  
        tup_ += (1,)
```

Time Complexity: $O(n * 3n) \ggg O(3*n**2)$
 $\ggg O(n**2)$
Space Complexity: $O(n + n + n) \ggg O(3n)$
 $\ggg O(n)$

Concatenation is not efficient!!

```
def concatenation(n):  
    lst_ = []  
    for i in range(n):  
        lst_.append(1)
```

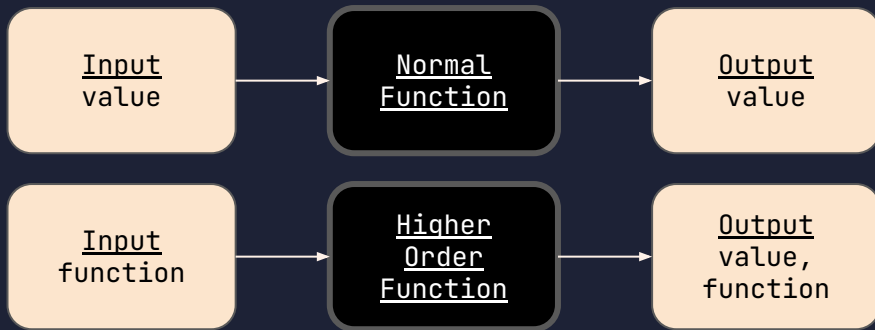
Time Complexity: $O(n * 1) \ggg O(n)$
Space Complexity: $O(1 + 1 \dots n \text{ times}) \ggg O(n)$

List.append is more time efficient than
concatenation +=!!

More information

- 1) [Python Website](#)
- 2) [Other resources](#)

2. Higher-Order Function



- Extract common pattern
- Make the code more generic

HOF return value

```
def general(func):  
    def helper():  
        return func(1)  
    return helper
```

```
general(add_one) # 2
```

HOF return function

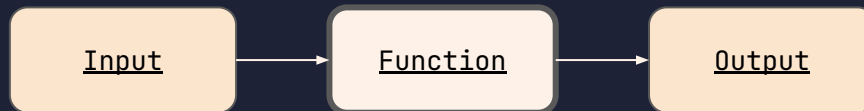
```
def general(func):  
    def helper(x):  
        return x  
    return helper(func)
```

```
general(func) # return func
```

3. Functional Abstraction & Abstraction Barrier

- Client & Implementer relationship/contract

What implementer see:



What client see:



- Once Functional Abstraction is established, **EVERYONE** should follow the abstraction barrier

“Let the function do its job”

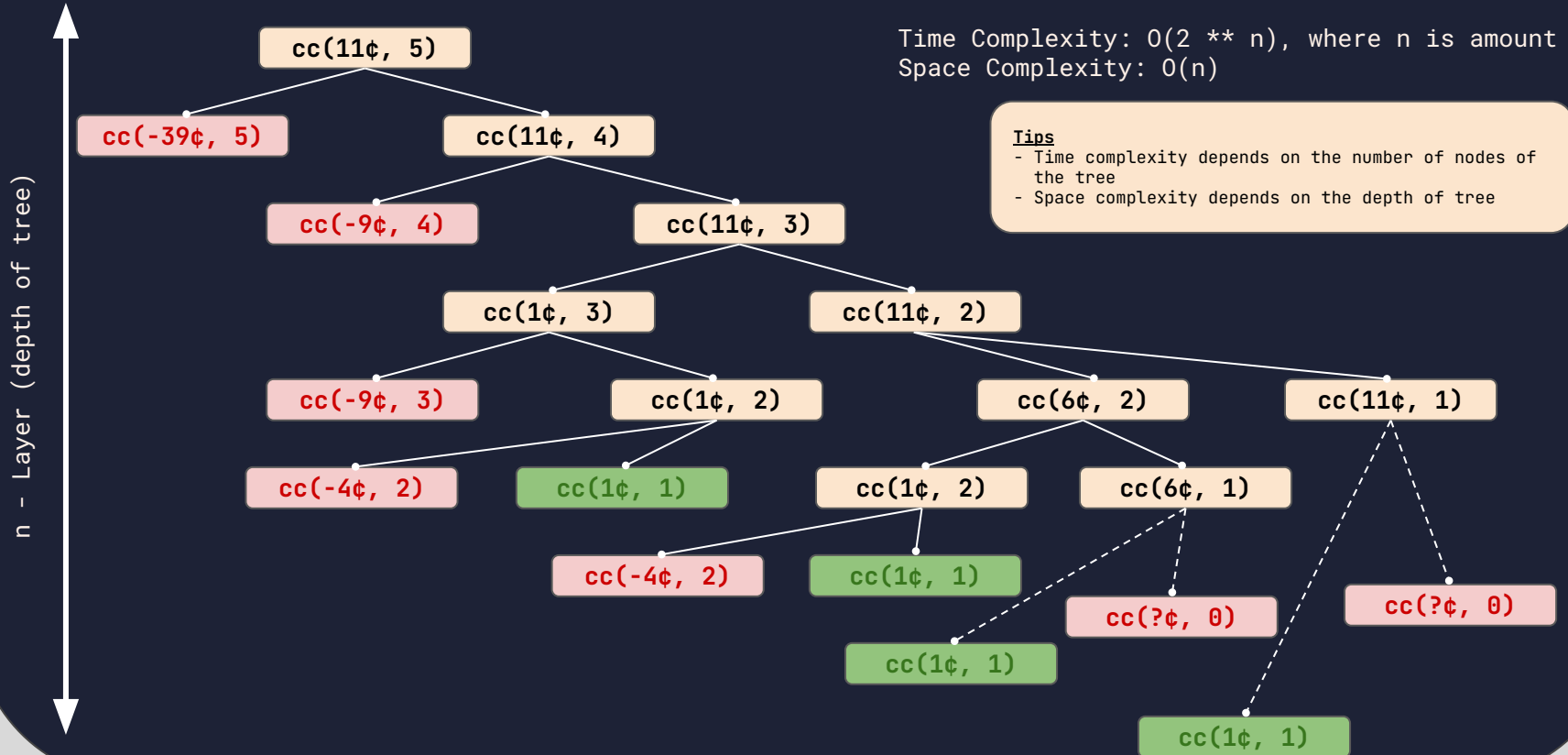


BY: Zhu_Ming

Any Questions?

Draw the tree illustrating the process generated by the `cc(amount,d)` function given in the lecture in making change for **11 cents**.

```
def cc(amount, d):  
    if amount == 0: # only one way to make 0¢  
        return 1  
    elif amount < 0 or d == 0: # cant make change for negative ¢  
        return 0 # cant make any change if no denoms  
    else:  
        return cc(amount - first_denomination(d), \  
                    d) + \  
                    cc(amount, d-1)  
  
def first_denomination(d):  
    # Functional Abstraction  
    . . .
```



$$f(n) = \begin{cases} n & n < 3 \\ f(n-1) + 2f(n-2) + 3f(n-3) & n \geq 3 \end{cases}$$

Write a function **f(n)** that computes f by a **recursive** process.

```
def f(n):  
    if n < 3: # only one way to make 0¢  
        return n  
    else:  
        return f(n - 1) + 2 * f(n - 2) + 3 * f(n - 3)
```

Time Complexity: $O(3^{**} n)$
Space Complexity: $O(n)$

$$f(n) = \begin{cases} n & n < 3 \\ f(n-1) + 2f(n-2) + 3f(n-3) & n \geq 3 \end{cases}$$

Write a function **f(n)** that computes *f* by an **iterative** process.

```
def f(n):  
    if n < 0: # only one way to make 0¢  
        return n  
    else:  
        f0, f1, f2 = 2, 1, 0  
        for i in range(n - 2):  
            fn = f0 + (2 * f1) + (3 * f2)  
            f0, f1, f2 = fn, f0, f1  
        return fn
```

Time Complexity: $O(n-2) \ggg O(n)$
Space Complexity: $O(1)$

Write a function `is_fib(n)` that returns `True` if `n` is a Fibonacci number, and `False` otherwise.

What is the order of growth in terms of time and space for the function that you wrote? Explain.

```
def is_fib(n):  
    if n < 0:  
        return False #not fib  
    elif n == 0 or n == 1:  
        return True  #is fib  
    a, b = 0, 1  
    while b < n: #generate fib  
        a, b = b, a + b  
        if b == n:  
            return True  
    return False
```

Time Complexity: $O(\log(n))$
Space Complexity: $O(1)$

Define a function **make_fare** that takes as arguments `stage1`, `stage2`, `start_fare`, `increment`, `block1` and `block2` and returns a function that calculates the taxi fare using those values.

```
def make_fare(stage1, stage2, start_fare, increment, block1, block2):
    def taxi_fare(distance):
        if distance <= stage1:
            return start_fare
        elif distance <= stage2:
            return (start_fare +
                    (increment * ceil((distance - stage1) / block1)))
        else:
            return (taxi_fare(stage2) + \
                    (increment * ceil((distance - stage2) / block2)))
    return taxi_fare
```

```
make_fare(1000, 10000, 3.0, 0.22, 400, 350)(3500)
>>> taxi_fare(3500)
>>> 4.54
```

Important!!

`stage1`, `stage2`, `start_fare`, `increment`, `block1`, `block2` is GLOBAL with respect to `taxi_fare`, but LOCAL to `make_fare`

`distance` is LOCAL to `taxi_fare`, but UNDEFINED to `make_fare`



Extra_Question.tut02

EXTRA Practices

(EXTRA)

Order of Growth

QUESTION 1

```
def isprime(n):  
    if n <= 1:  
        return False  
    for i in range(2, int((n**0.5)+1)):  
        return bool(n % i)  
    return True
```

Time complexity: $O(n^{0.5})$
Space complexity: $O(1)$

(EXTRA)

Order of Growth

QUESTION 2

```
def weird(n):  
    b = 0  
    for i in range(n):  
        for j in range(i, n):  
            b += 1
```

Time complexity: $O(n^2)$
Space complexity: $O(1)$

(EXTRA)

Order of Growth

QUESTION 3

```
def weird(n):  
    b = []  
    for i in range(n):  
        b.append(1)  
  
    b = []  
    for i in range(n): 0(n)  
        b += [1]       0(n)
```

Time complexity: $O(n^2)$
Space complexity: $O(n)$

(EXTRA)

Order of Growth

QUESTION 4

```
def what_is_in(s, str_):  
    return s in str_
```

Time complexity: $O(n)$
Space complexity: $O(1)$

(EXTRA)

Nested Lambda Function

Question 3

```
def f(x):  
    return lambda y: (x, y(x))  
  
def g(y):  
    return lambda x: x(y)  
  
print(g(2)(f)(lambda x: x + 1))
```

OUTPUT:

(2, 3)



BY: Zhu_Ming

Any Questions?



Thank You!!

The End

See you next lesson