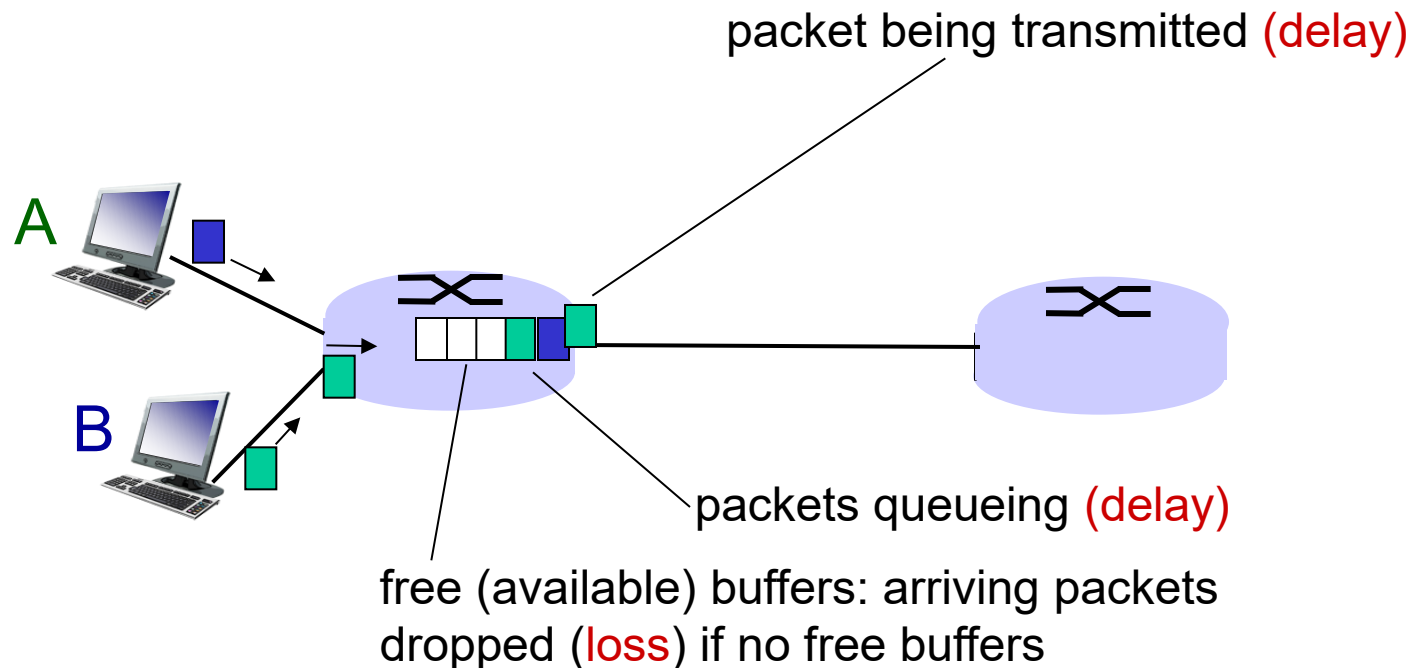


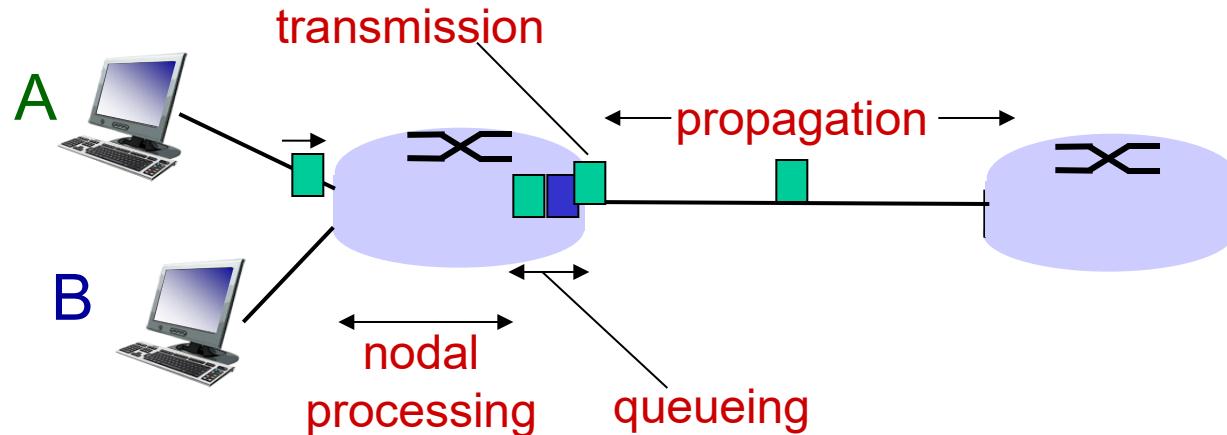
How do loss and delay occur?

packets *queue* in router buffers

- ❖ packet arrival rate to link (temporarily) exceeds output link capacity
- ❖ packets queue, wait for turn



Four sources of packet delay



$$d_{\text{nodal}} = d_{\text{proc}} + d_{\text{queue}} + d_{\text{trans}} + d_{\text{prop}}$$

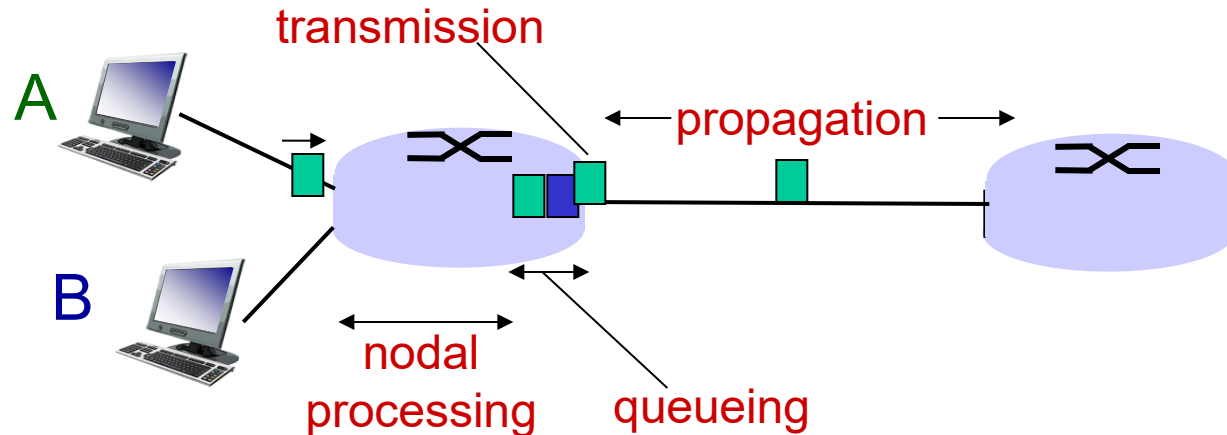
d_{proc} : nodal processing

- check bit errors
- determine output link
- typically < msec

d_{queue} : queueing delay

- time waiting at output link for transmission
- depends on congestion level of router

Four sources of packet delay



$$d_{\text{nodal}} = d_{\text{proc}} + d_{\text{queue}} + d_{\text{trans}} + d_{\text{prop}}$$

d_{trans} : transmission delay:

- L : packet length (bits)
- R : link bandwidth (bps)
- $d_{\text{trans}} = L/R$

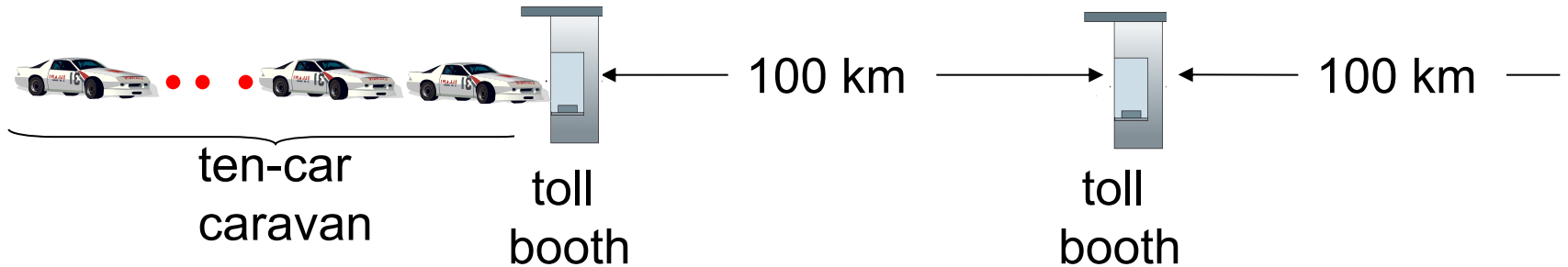
d_{prop} : propagation delay:

- d : length of physical link
- s : propagation speed in medium ($\sim 2 \times 10^8$ m/sec)
- $d_{\text{prop}} = d/s$

d_{trans} and d_{prop}
very different

* Check out the Java applet for an interactive animation on trans vs. prop delay

Caravan analogy



- ❖ cars “propagate” at 100 km/hr
 - ❖ toll booth takes 12 sec to service car (bit transmission time)
 - ❖ car ~ bit; caravan ~ packet
 - ❖ **Q: How long until caravan is lined up before 2nd toll booth?**
- time to “push” entire caravan through toll booth onto highway = $12 * 10 = 120$ sec
 - time for last car to propagate from 1st to 2nd toll booth:
 $100\text{km} / (100\text{km/hr}) = 1$ hr
 - **A: 62 minutes**

Q1

- Consider the first bit, before it can be transmitted, all of the bits in the same packet must be generated. This requires

$$\frac{64 \times 8}{128 \times 10^3} = 4 \text{ msec}$$

- The time required to transmit the packet is

$$\frac{64 \times 8}{4 \times 10^6} = 128 \text{ microsec}$$

- Propagation delay is 8 msec.
- The delay until decoding is 4 msec + 128 microsec + 8 msec = 12.128 msec.
- A similar analysis shows that all bits experience a delay of 12.128 msec.

Q2

- The first end system requires L/R_1 to transmit the packet onto the first link;
- the packet propagates over the first link in d_1/s_1 ; the packet switch adds a processing delay of d_{proc} ;
- after receiving the entire packet, the packet switch connecting the first and the second link requires L/R_2 to transmit the packet onto the second link;
- the packet propagates over the second link in d_2/s_2 . Similarly, we can find the delay caused by the second switch and the third link: L/R_3 , d_{proc} , and d_3/s_3 .

Q2

- Adding these eight delays gives

$$d_{end-end} = L/R_1 + L/R_2 + L/R_3 + d_1/s_1 + d_2/s_2 + d_3/s_3 + d_{proc} + d_{proc}$$

- To answer the second question, we simply plug the values into the equation to get $6 + 6 + 6 + 20 + 16 + 4 + 3 + 3 = 64$ msec.

Q3

- Because bits are immediately transmitted, the packet switch does not introduce any delay; in particular, it does not introduce a transmission delay. Thus,

$$d_{end-end} = L/R + d_1/s_1 + d_2/s_2 + d_3/s_3$$

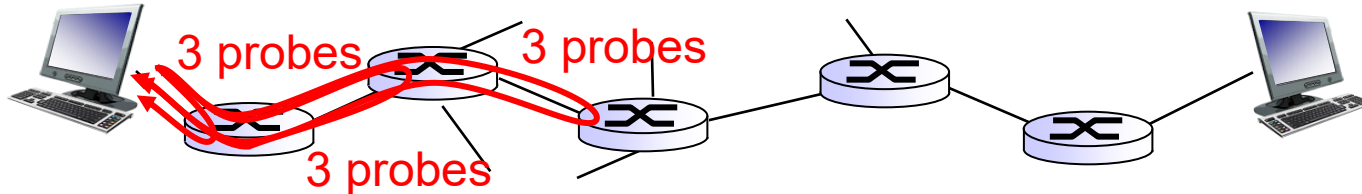
- For the values given, we get $6 + 20 + 16 + 4 = 46$ msec.

Q4

- The arriving packet must first wait for the link to transmit $4.5 * 1,500$ bytes = 6,750 bytes or 54,000 bits.
- Since these bits are transmitted at 2 Mbps, the queuing delay is 27 msec.
- Generally, the queuing delay is $(nL + (L - x))/R$.

“Real” Internet delays and routes


- ❖ what do “real” Internet delay & loss look like?
- ❖ `traceroute` program: provides **delay measurement** from source to router along end-end Internet path towards destination. For all i :
 - sends three packets that will reach router i on path towards destination
 - router i will return packets to sender
 - sender times interval between transmission and reply.



“Real” Internet delays, routes


traceroute: gaia.cs.umass.edu to www.eurecom.fr

3 delay measurements from
gaia.cs.umass.edu to cs-gw.cs.umass.edu



1 cs-gw (128.119.240.254) 1 ms 1 ms 2 ms
2 border1-rt-fa5-1-0.gw.umass.edu (128.119.3.145) 1 ms 1 ms 2 ms
3 cht-vbns.gw.umass.edu (128.119.3.130) 6 ms 5 ms 5 ms
4 jn1-at1-0-0-19.wor.vbns.net (204.147.132.129) 16 ms 11 ms 13 ms
5 jn1-so7-0-0-0.wae.vbns.net (204.147.136.136) 21 ms 18 ms 18 ms
6 abilene-vbns.abilene.ucaid.edu (198.32.11.9) 22 ms 18 ms 22 ms
7 nycm-wash.abilene.ucaid.edu (198.32.8.46) 22 ms 22 ms 22 ms
8 62.40.103.253 (62.40.103.253) 104 ms 109 ms 106 ms
9 de2-1.de1.de.geant.net (62.40.96.129) 109 ms 102 ms 104 ms
10 de.fr1.fr.geant.net (62.40.96.50) 113 ms 121 ms 114 ms
11 renater-gw.fr1.fr.geant.net (62.40.103.54) 112 ms 114 ms 112 ms
12 nio-n2.cssi.renater.fr (193.51.206.13) 111 ms 114 ms 116 ms
13 nice.cssi.renater.fr (195.220.98.102) 123 ms 125 ms 124 ms
14 r3t2-nice.cssi.renater.fr (195.220.98.110) 126 ms 126 ms 124 ms
15 eurecom-valbonne.r3t2.ft.net (193.48.50.54) 135 ms 128 ms 133 ms
16 194.214.211.25 (194.214.211.25) 126 ms 128 ms 126 ms
17 * * *
18 * * *
19 fantasia.eurecom.fr (193.55.113.142) 132 ms 128 ms 136 ms

trans-oceanic link



* means no response (probe lost, router not replying)

Q5

- IP address of the routers between the source and the destination, three round-trip time measurements for each router
- An asterisk is displayed if a packet is not acknowledged within the expected timeout (The default timeout is 5 sec). Timeout appears mainly because a firewall blocks the traffic.

Q5

- The round-trip delays include all of the delays, including transmission delays, propagation delays, router processing delays, and queuing delays. Because the queuing delay is varying with time, the round-trip delay of packet N sent to a router N can sometimes be longer than the round-trip delay of packet $N+1$ sent to router $N+1$.