PS6: Assignment on ILP solvers

Help

The due date for this homework is Tue 16 Dec 2014 2:59 PM CST.

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Question 1

This assignment takes you through the process of "binarizing" an ILP to convert it into a 0-1 ILP. Consider the ILP shown below:

We represent x_1 by a three bit binary number $b_{1,3}b_{1,2}b_{1,1}$ ($b_{1,3}$ is the most significant bit and $b_{1,1}$ is the least significant bit). Which of the following expressions replaces x_1 in the original ILP?

 $x_1: 8b_{1,1}+4b_{1,2}+2b_{1,1}$

 $\bigcirc \ \ x_1:\ 4b_{1,3}+2b_{1,2}+b_{1,1}$

 $\bigcirc \ x_1: \ b_{1,1}+b_{1,2}+b_{1,3}$

 $\bigcirc \ x_1: \ 1000b_{1,3}+100b_{1,2}+10b_{1,1}$

Question 2

Continuing with problem 1, what is the minimum number of bits needed to represent x_2 in the binarized problem? Hint: you may have to solve an LP to find out.

3

4

- 2
- \bigcirc 5

Question 3

In our lecture we considered pure ILPs where all the variables in the problem are integers. In this problem, we consider mixed integer programs (MIPs), wherein a subset of the decision variables are integer variables and the remaining variables are considered real-valued.

Which of the following modifications to branch-and-bound procedure covered will help is solve mixed integer programs? Select all the correct answers. Assume that the problem is maximizing the objective function.

- Branching is considered when the LP relaxation yields a fractional solution for some decision variable.
- lt is possible to branch on any of the decision variables in the problem.
- A node can be converted into a leaf whenever its solution satisfies the integrality constraints for the integral decision variables.
- We cannot prune a node simply because its LP relaxation yields an optimum that is less than or equal to the best objective.
- We can branch on any real-valued variable x_j whose LP relaxation optimal value is s_j , the branch constraints will be $x_j \leq s_j$ and $x_j \geq s_j$
 - Branching is considered when the LP relaxation yields a fractional solution for an integer decision variable.

Question 4

We apply branch-and-bound on a 0-1 (binary) ILP. Select all the true facts below. Assume that the problem is a maximization problem with n decision variables and m constraints.

- \square The maximum depth of the tree is given by the number of decision variables n.
- We should explore the branch $x_i \geq 1$ first and the $x_i \leq 0$ since it is guaranteed to yield a larger value.
- lacksquare For a branch variable x_i , the branch constraints are always $x_i \leq 0$ and $x_i \geq 1$.
- lacksquare The total number of nodes in the final branch-and-bound tree will be less than 2^{n+1}

Question 5

Consider the following final dictionary encountered while solving an ILP given below (we assume all problem and slack variables are integers).

$$x_{2} \begin{vmatrix} \frac{8}{3} & +\frac{2}{3}x_{6} & -\frac{10}{3}x_{3} \\ x_{1} & \frac{5}{3} & -\frac{5}{3}x_{6} & +\frac{1}{3}x_{3} \\ x_{4} & 6 & +\frac{14}{3}x_{6} & -\frac{4}{3}x_{3} \\ x_{5} & \frac{4}{3} & +\frac{5}{6}x_{6} & -\frac{1}{3}x_{3} \\ z & -\frac{17}{3} & -\frac{7}{3}x_{6} & -\frac{1}{3}x_{3} \end{vmatrix}$$

Select all the valid cutting plane constraints from the list below.

- lacksquare $\frac{2}{3}\,x_6+rac{2}{3}\,x_3\geqrac{2}{3}\,$ corresponding to x_1
- lacksquare $\frac{1}{3} \, x_6 + \frac{1}{3} \, x_3 \geq 0$ corresponding to x_4
- lacksquare $\frac{5}{3} x_6 \frac{1}{3} x_3 \geq \frac{5}{3}$ corr. to x_1
- \square $\frac{5}{6} x_6 \frac{1}{3} x_3 \geq \frac{4}{3}$ corr. to x_5
- lacksquare $\frac{1}{6}\,x_6+rac{1}{3}\,x_3\geqrac{1}{3}\,$ corresponding to x_5
- lacksquare $\frac{1}{3}\,x_6+rac{1}{3}\,x_3\geqrac{2}{3}\,$ corresponding to x_2
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