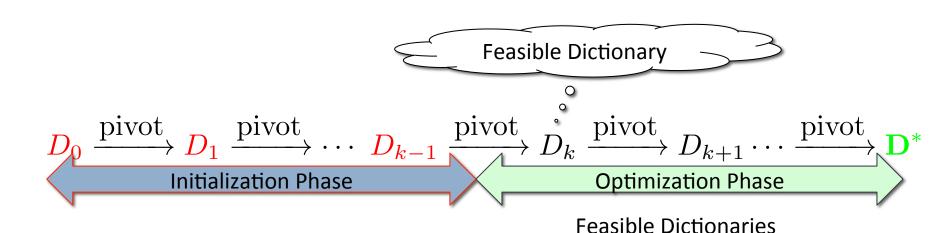
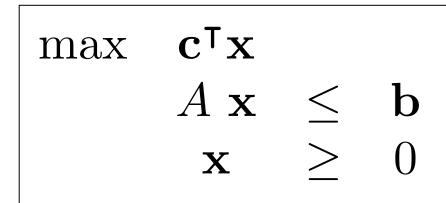
# INITIALIZATION USING THE DUAL

## Simplex Algorithm



#### Initialization

• Goal: Find a feasible dictionary for the problem.



Observation:

Objective does not matter.

#### Initialization

$$\begin{array}{ccc}
\mathbf{max} & \mathbf{c}^{\mathsf{T}} \mathbf{x} \\
A \mathbf{x} & \leq \mathbf{b} \\
\mathbf{x} & \geq 0
\end{array}$$

If all entries in  $b \ge 0$ 

- Initialization not needed.

If any entry in **b <0** 

- initialization needed.

**Original Problem** 

## Linear Program (Dual)

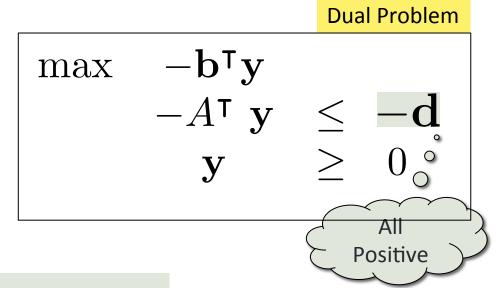
 $egin{array}{ccccc} oldsymbol{ ext{Dual Problem}} & oldsymbol{ ext{Dual Problem}} \ & oldsymbol{ ext{b}}^{\mathsf{T}} \mathbf{y} & & \geq \mathbf{c} \ & \mathbf{y} & & \geq 0 \end{array}$ 

$$\begin{array}{cccc}
\mathbf{max} & -\mathbf{b}^{\mathsf{T}}\mathbf{y} \\
-A^{\mathsf{T}} & \mathbf{y} & \leq & -\mathbf{c} \\
\mathbf{y} & \geq & 0
\end{array}$$

**Standard Form Converted Dual** 

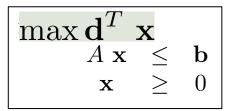
#### Initialization Using Dual: Basic Idea

1. Change the objective function



Idea: choose **d** to be all negative entries.

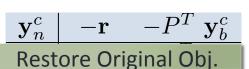
#### **Initialization Using Dual**



d has all negative entries

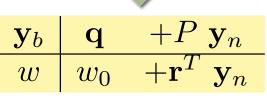
Dualize

 $\begin{array}{rcl}
\max & -\mathbf{b}^{\mathsf{T}}\mathbf{y} \\
-A^{\mathsf{T}}\mathbf{y} & \leq & -\mathbf{d} \\
\mathbf{y} & \geq & 0
\end{array}$ 



Feasible
Primal Dictionary

Dualize



Final Dual Dictionary

**Optimization** 

**Phase Simplex** 

#### Example

```
max. x_1 + 2x_2

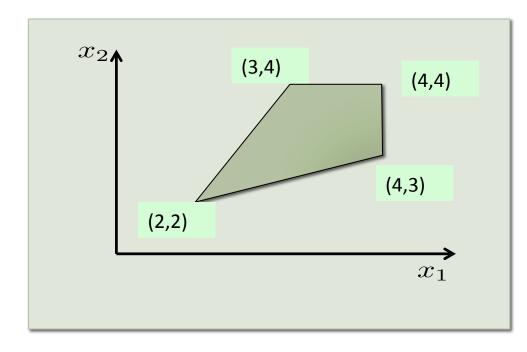
s.t. -2x_1 + x_2 \le -2

x_2 \le 4

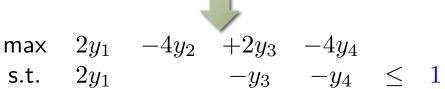
x_1 - 2x_2 \le -2

x_1 \le 4

x_1, x_2 \ge 0
```



#### Example #1: Problem Transformation





#### Solving the modified dual

 $y_1$  enters and  $y_5$  leaves

## Solving the modified dual (step 2)

 $y_3$  enters and  $y_6$  leaves

## Solving the modified dual (step 3)

**Final Dual Dictionary** 

#### Convert Dual back to Primal



$x_1$	$y_5$
$x_2$	$y_6$
$x_3$	$y_1$
$x_4$	$\mid y_2 \mid$
$x_5$	$y_3$
$x_6$	$y_4$

#### Conversion to Primal Dictionary

$x_1$	$y_5$
$x_2$	$y_6$
$x_3$	$y_1$
$x_4$	$y_2$
$x_5$	$y_3$
$x_6$	$y_4$

### Initialization: Restoring Objective

$$z = x_1 + 2x_2$$

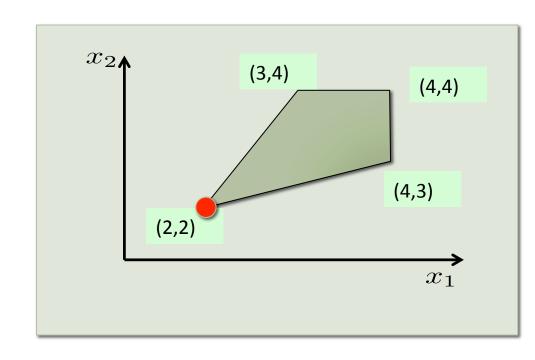
$$= 2 + \frac{2}{3}x_3 + \frac{1}{3}x_5 + 2(2 + \frac{1}{3}x_3 + \frac{2}{3}x_5)$$

$$= 6 + \frac{4}{3}x_3 + \frac{5}{3}x_5$$

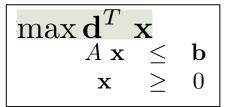
Initial Feasible Dictionary Found. We can now proceed to optimize!!

#### Example

max. 
$$x_1 + 2x_2$$
  
s.t.  $-2x_1 + x_2 \le -2$   
 $x_2 \le 4$   
 $x_1 - 2x_2 \le -2$   
 $x_1 \le 4$   
 $x_1, x_2 \ge 0$   
 $x_1 \mid 2 + \frac{2}{3}x_3 + \frac{1}{3}x_5$   
 $x_4 \mid 2 - \frac{1}{3}x_3 - \frac{2}{3}x_5$   
 $x_2 \mid 2 + \frac{1}{3}x_3 + \frac{2}{3}x_5$   
 $x_6 \mid 2 - \frac{2}{3}x_3 - \frac{1}{3}x_5$   
 $x_1 \mid 6 + \frac{4}{3}x_3 + \frac{5}{3}x_5$ 



#### **Initialization Using Dual**



Dualize

$$\begin{array}{rcl}
 \text{max} & -\mathbf{b}^{\mathsf{T}}\mathbf{y} \\
 -A^{\mathsf{T}}\mathbf{y} & \leq & -\mathbf{d} \\
 \mathbf{y} & \geq & 0
\end{array}$$

**d** has all negative entries



Optimization
Phase Simplex

#### **Unbounded!**

Original primal problem is infeasible.

$\mathbf{y}_b$	$\mathbf{q}$	$+P \mathbf{y}_n$
w	$w_0$	$+\mathbf{r}^T \; \mathbf{y}_n$

Final Dual Dictionary