

BRANCH-AND-BOUND: BASIC IDEA

Integer Linear Program: Example

$$\begin{array}{llll} \max & -x_1 - 2x_2 - 0.5x_3 - 0.2x_4 - x_5 + 0.6x_6 & & \\ \text{s.t.} & x_1 + 2x_2 & \geq & 1 \\ & x_1 + x_2 + 3x_6 & \geq & 1 \\ & x_1 + x_2 + x_6 & \geq & 1 \\ & x_3 - 3x_4 & \geq & 1 \\ & x_3 - 2x_4 - 5x_5 & \geq & 1 \\ & x_4 + 3x_5 - 4x_6 & \geq & 1 \\ & x_2 + x_5 + x_6 & \geq & 1 \\ 0 \leq & x_1, x_2, \dots, x_6 & \leq & 10 \\ & x_1, \dots, x_6 & \in & \mathbb{Z} \end{array}$$

Step #1: Solve the LP relaxation

Solving the LP relaxation.

```
x1.val = 0.3333333333333333  
x2.val = 0.6666666666666667  
x3.val = 2.6666666666666667  
x4.val = 0  
x5.val = 0.3333333333333333  
x6.val = 0
```

Optimal Value (LP relaxation): -3.333333

Step #2: Choose a branch variable

x1.val = 0.3333333333333333

x2.val = 0.6666666666666667

x3.val = 2.6666666666666667

x4.val = 0

x5.val = 0.3333333333333333

x6.val = 0

Choose a
variable that
should be
integer

Optimal Value (LP relaxation): -3.333333

Step #3: Branching Constraints

$$\begin{array}{llll} \max & -x_1 - 2x_2 - 0.5x_3 - 0.2x_4 - x_5 + 0.6x_6 & & \\ \text{s.t.} & x_1 + 2x_2 & \geq & 1 \\ & x_1 + x_2 + 3x_6 & \geq & 1 \\ & x_1 + x_2 + x_6 & \geq & 1 \\ & x_3 - 3x_4 & \geq & 1 \\ & x_3 - 2x_4 - 5x_5 & \geq & 1 \\ & x_4 + 3x_5 - 4x_6 & \geq & 1 \\ & x_2 + x_5 + x_6 & \geq & 1 \\ & 0 \leq x_1, x_2, \dots, x_6 & \leq & 10 \end{array}$$

Original Problem

b1

b2

Original Problem Constr.

+

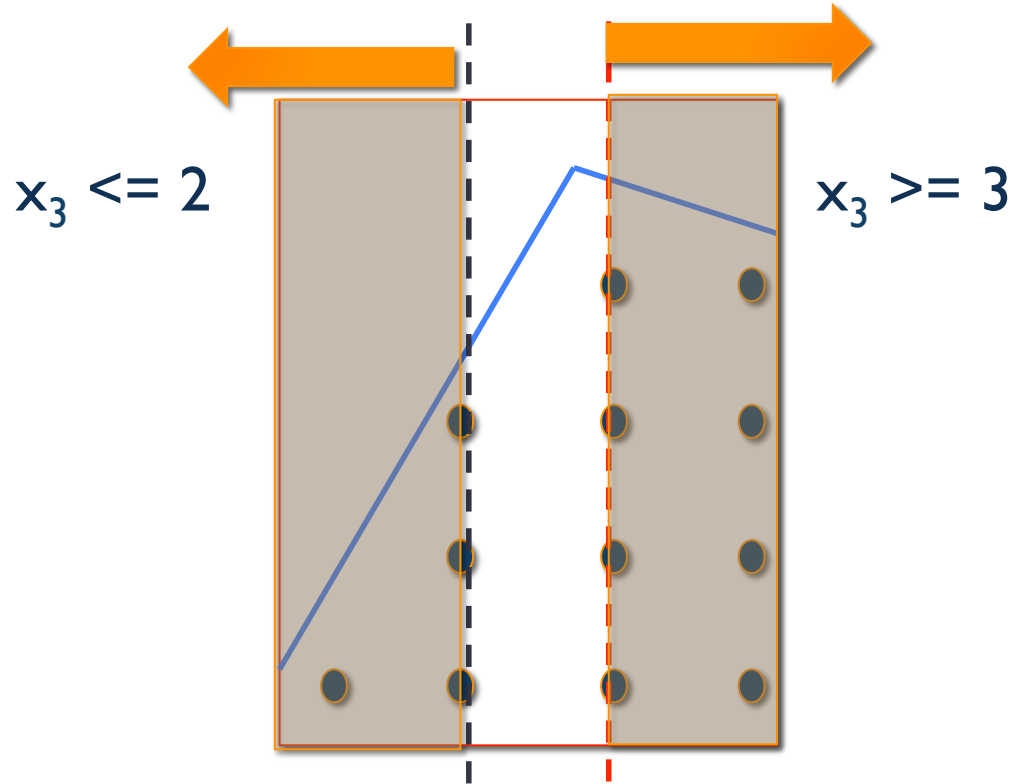
($x_3 \leq 2$)

Original Problem Constr.

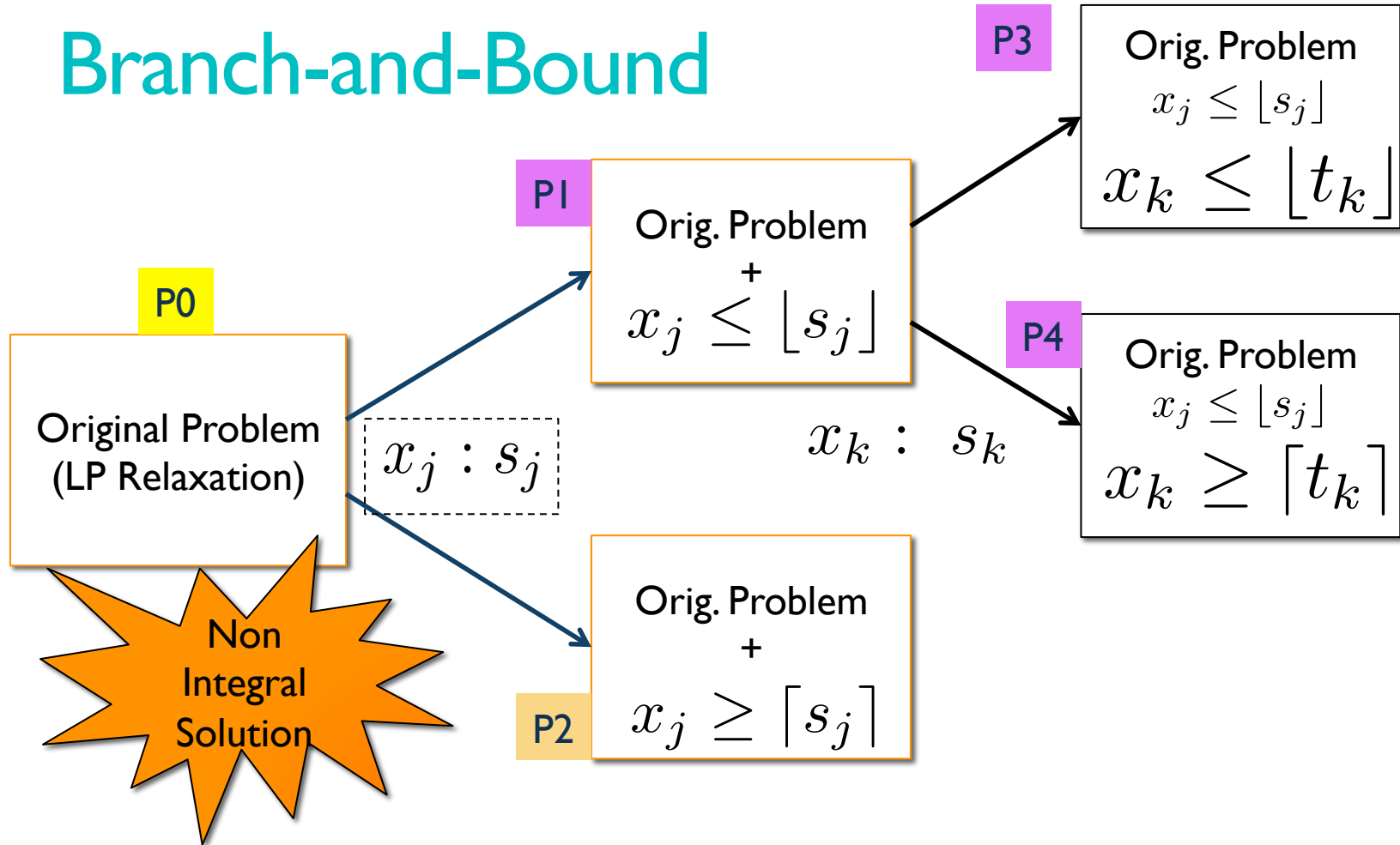
+

($x_3 \geq 3$)

Visualization

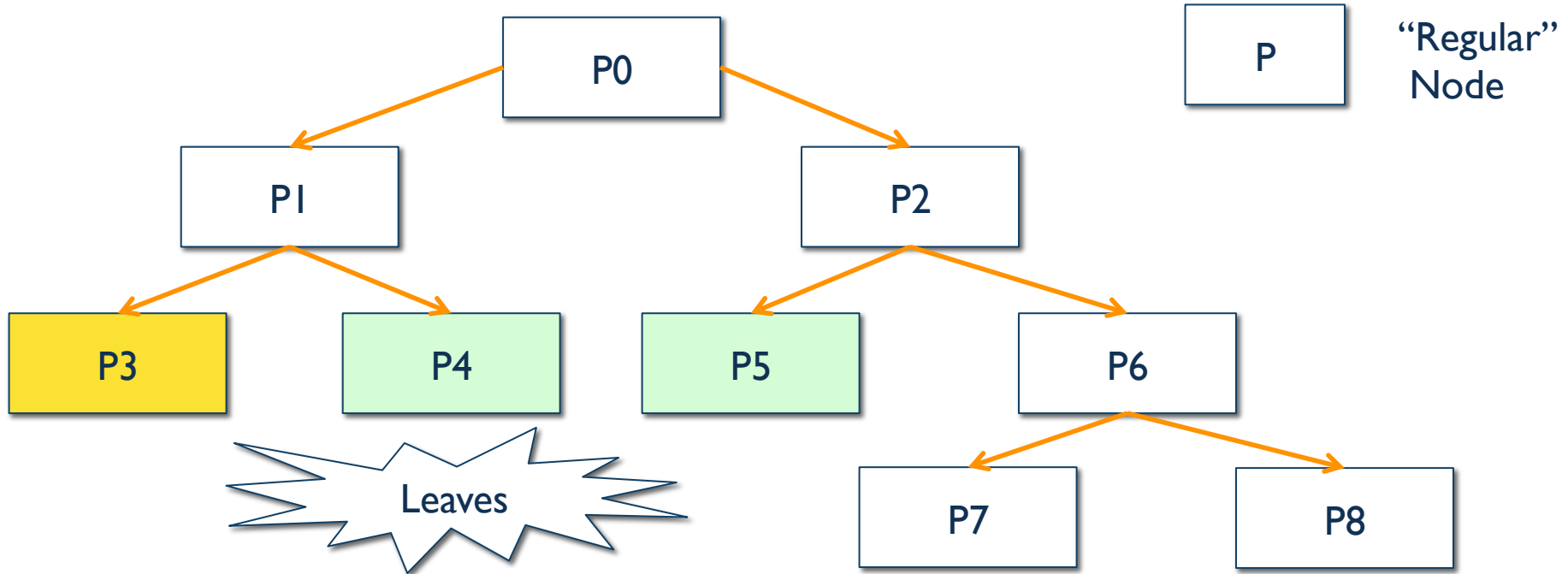


Branch-and-Bound



Branch-And-Bound: Operations

Branch-And-Bound Tree: Nodes are ILP instances.



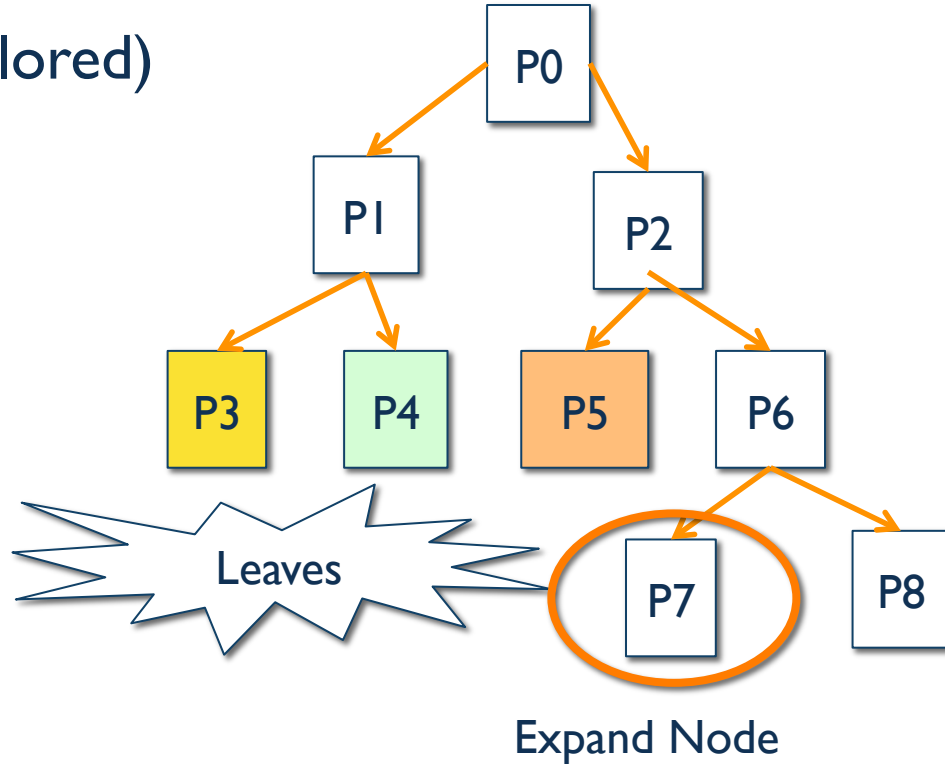
Branch-And-Bound Tree

P0 Regular Node (to be explored)

P3 Leaf node

P4 Leaf node

P5 Leaf node



Expanding a Node

1. Solve the LP relaxation for the node ILP.

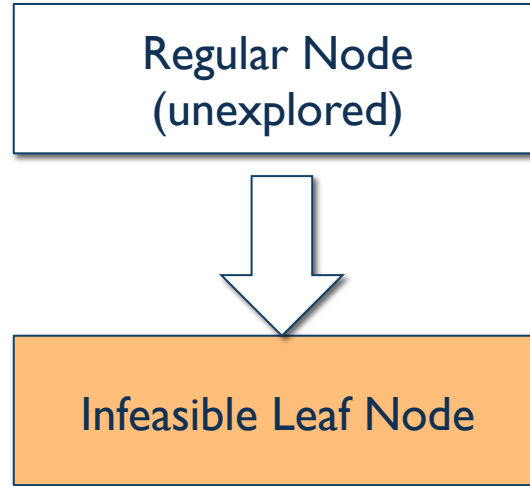
Regular Node
(unexplored)

LP relaxation solution

2. Three cases:
 1. Infeasible.
 2. Integral Solution.
 3. Fractional Solution.

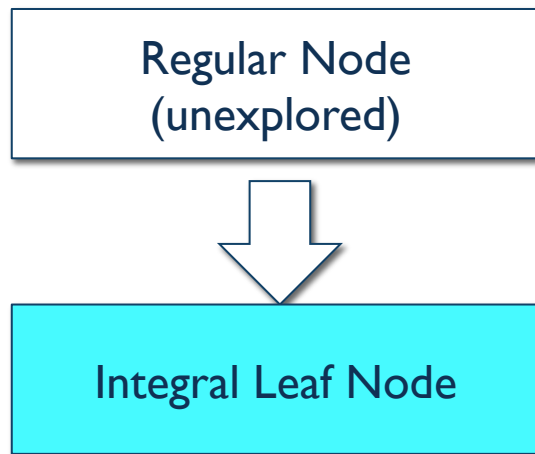
Case-I: Infeasible LP relaxation.

- LP relaxation is infeasible.



Case-2: LP relaxation yields integral solution

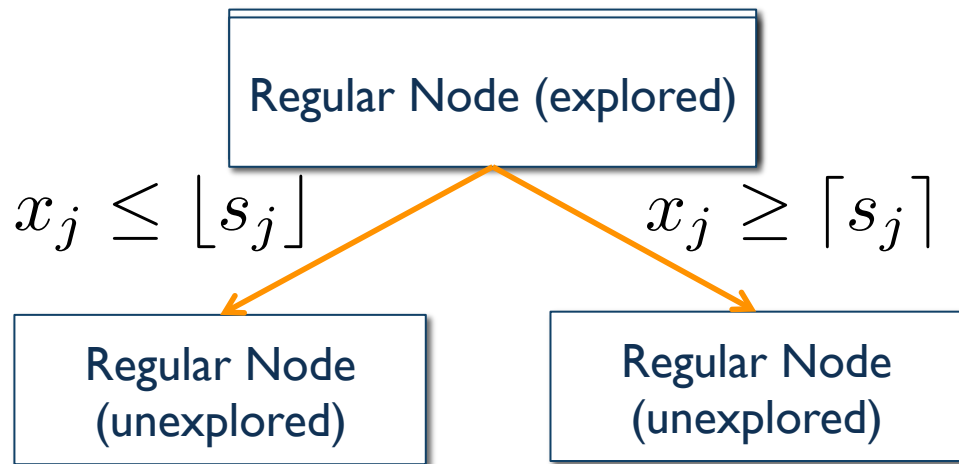
- LP relaxation solution is integral: ILP solution = LP solution.



$\text{bestObjective} := \max(\text{lpOptimum}, \text{bestObjective})$

Case-3: LP relaxation yields a fractional solution

- LP relaxation yields a fractional solution.



LP relaxation solution:

x1: ...

x2: ...

x3: ...

...

Opt. Solution: **optSolution**

Optimal Pruning

Regular Node
(explored)

$\text{optSolution} \leq \text{bestObjective?}$

LP relaxation solution:

x_1 : ...

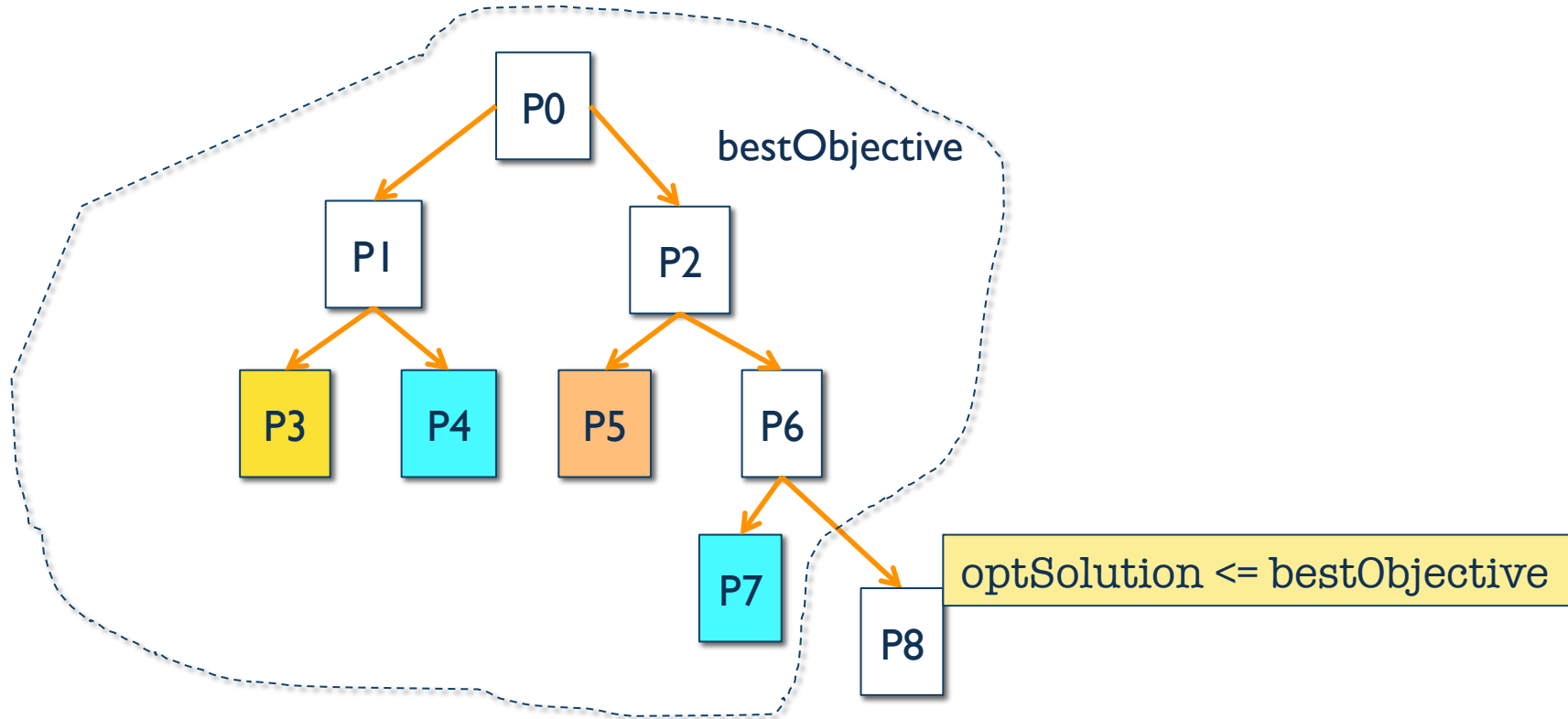
x_2 : ...

x_3 : ...

...

Opt. Solution: optSolution

Optimal Pruning



Branch-And-Bound Initial Node

Original ILP
(unexplored
node)

`bestObjective := -Infinity`