## Math 449: Numerical Applied Mathematics Lecture 07

09/15/2017

Today's topic: Intro to Numerical Linear Algebra

Intro to Numerical Linear Algebra

1— rariable :

linear systems are trivial

$$ax = b$$
  $(a \neq 0)$  has solution  $x = \frac{b}{a} = a^{-1}b$ 

division appears in 1-D newton's method.  $z_{k+1} = z_k - \frac{f(x)}{f'(x)}$ 

Focus: algorithms that are fast/efficient-for very large systems (nxn matrices with n large)

Google's Page Rank eigenvector of matrix of size (# of webpages). \* (# of webpages).

Motivations.

A 
$$\overrightarrow{z} = \overrightarrow{b}$$
 known
$$\uparrow$$

$$A \in \mathbb{R}^n \times \mathbb{R}^n \qquad \overrightarrow{z}, \overrightarrow{b} \in \mathbb{R}^n$$
A is nonsingular.

an inverse.

unique soln 
$$\vec{z} = A^{-1}\vec{b}$$

Ex. 
$$\begin{cases} \alpha x_1 + \beta x_2 = b_1 \\ \gamma x_2 + \delta x_2 = b_2 \end{cases} \begin{pmatrix} \alpha & \beta \\ \gamma & \delta \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \end{pmatrix} = \begin{pmatrix} b_1 \\ b_2 \end{pmatrix}$$

if 
$$det(A) = \alpha \cdot 8 - \beta \cdot r \neq 0$$
 then  $A^{-1} = \frac{1}{\alpha \cdot 8 - \beta \cdot r} \begin{pmatrix} 8 & -\beta \\ -r & \alpha \end{pmatrix}$ 

Soln: 
$$\vec{z} = A^{-1}\vec{b} = \frac{1}{\alpha \delta - \beta r} \begin{pmatrix} \delta - \beta \\ -r & \alpha \end{pmatrix} \begin{pmatrix} b_1 \\ b_2 \end{pmatrix}$$

$$= \frac{1}{\alpha \delta - \beta r} \begin{pmatrix} \delta b_2 - \beta b_1 \\ -xb_1 + \alpha b_2 \end{pmatrix}$$

Check 
$$Ax = \frac{1}{\alpha s - \beta r} \begin{pmatrix} \alpha & \beta \\ \gamma & s \end{pmatrix} \begin{pmatrix} sb_2 - \beta b_1 \\ -\gamma b_1 + \alpha b_2 \end{pmatrix}$$

$$= \frac{1}{\alpha \delta - \beta \tau} \begin{pmatrix} (\alpha \delta - \beta \tau) b_1 + O b_2 \\ O b_1 + (\alpha \delta - \beta \tau) b_2 \end{pmatrix} = \begin{pmatrix} b_1 \\ b_2 \end{pmatrix}$$

Cramer's Rule (works for nxn matrix)

from 2x2 -> nxn

Soln to 
$$\overrightarrow{A}\overrightarrow{\alpha} = \overrightarrow{b}$$
 is  $x_i = \frac{\det(A_i)}{\det(A)}$   $\overrightarrow{X} = \begin{pmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{pmatrix}$ 

$$\vec{X} = \begin{pmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{pmatrix}$$

where Ai is the matrix obtained by replacing the  $i^{th}$  column of A by  $\vec{b}$  (i=1,...,n)

Ex.  $A = \begin{pmatrix} \alpha & \beta \\ \delta & \delta \end{pmatrix}$   $A_1 = \begin{pmatrix} b_1 & \beta \\ b_2 & \delta \end{pmatrix}$   $A_2 = \begin{pmatrix} \alpha & b_1 \\ \delta & b_2 \end{pmatrix}$ 

$$A_{i} = \begin{pmatrix} b_{i} & \beta \\ b_{2} & \delta \end{pmatrix}$$

$$A_2 = \begin{pmatrix} \alpha & b_1 \\ \gamma & b_2 \end{pmatrix}$$

det (A) = a8-Bi

$$det(A_1) = \begin{vmatrix} b_1 & \beta \\ b_2 & \delta \end{vmatrix} = \delta b_1 - \beta b_2$$

$$\det(A_1) = \begin{vmatrix} b_1 & \beta \\ b_2 & \delta \end{vmatrix} = \delta b_1 - \beta b_2 \qquad \det(A_2) = \begin{vmatrix} \alpha & b_1 \\ \tau & b_2 \end{vmatrix} = \alpha b_2 - rb_1$$

$$x_1 = \frac{\delta b_1 - \beta b_2}{\alpha \delta - \beta r}$$
  $x_2 = \frac{\alpha b_2 - r b_1}{\alpha \delta - \beta r}$ 

$$\chi_2 = \frac{\alpha b_2 - \gamma b_1}{\alpha \delta - \beta r}$$

Expansion by rows / columns.

$$\begin{pmatrix} \cdot \cdot \cdot \\ & \end{pmatrix} \qquad det A = \sum_{j=1}^{n} (-1)^{j+1} a_{ij} \det(A_{ij})$$

Aij: (n-1) x (n-1) matrix excluding row I arel colj.

To compute  $det(n \times n)$ , compute det of n different  $(n-1) \times (n-1)$  matrices.

N=1:  $det(a_{11}) = a_{11}$ 

n=2: det (

n=3:

n = 4

24

n=5

in general, det(A) has n! terms.

describing the comparative growth rates of different sequences.

How many steps closes it take to solve an  $n \times n$  system?

How quickly does this grow with n?

direction methods

Method that returns the answer after a finite # of steps.

Def. let  $(x_n)$  and  $(y_n)$  be requere of real #.

"big 0"  $z = O(y_n)$  if  $|z_n| \le c|y_n|$  for some c > 0,  $N \in \mathbb{N}$  whenever  $n \ge N$  for some  $N \in \mathbb{N}$ 

Ex. 
$$x_n = O(1)$$
  
 $x_n = \sin(n) = O(1)$   
 $x_n = \cos(n)$   
 $x_n$ 

2f 
$$n \ge 1$$
, then  $1 + \frac{1}{n} \ge 2$   
Ex.  $2n+1 = O(n)$   
if  $n \ge 1$ , then  $2n+1 \le 2n+n = 3n$   
So take  $N=1$ ,  $C=3$ .  
 $n+1000,000 = O(n)$ .

## General idea

look at dominant term that grows fastest as  $n\to\infty$  ignore constant factors and lower-order terms.