

Math 449: Numerical Applied Mathematics

Lecture 20

10/20/2017 Wenzhen

Today's topic: Iterative Methods in Numerical Linear Algebra

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Today: Iterative methods in numerical linear algebra

Linear Systems $A\vec{x} = \vec{b}$

- direct methods

terminate w/ exact answer after finitely many steps

- iterative methods

start w/ guess $x^{(0)}$ iterate $x^{(0)}, x^{(1)}, x^{(2)}, \dots$, get answer in limit.

Why not Newton's method?

$$f(\vec{x}) = A\vec{x} - b$$

$$J_f(x) = A$$

$$x^{(k+1)} = x^{(k)} A^{-1} (Ax^{(k)} - b) = A^{-1} b$$

Requires solving the original linear system

Recall geometric series

if $|r| < 1$

$$\frac{1}{1-r} = 1 + r + r^2 + r^3 + r^4 + \dots$$

$$\approx 1 + r + r^2 + \dots + r^k$$

Note :

$$(1-r)(1+r+r^2+\dots+r^k) = 1 - \cancel{r} + \cancel{r} - \cancel{r^2} + \cancel{r^2} + \dots - \cancel{r^k} + \cancel{r^k} - r^{k+1} = 1 - r^{k+1}$$

$\rightarrow 1$ as $k \rightarrow \infty$

$$1 + r + \dots + r^k = \frac{1}{1-r} - \frac{r^{k+1}}{1-r} \rightarrow \frac{1}{1-r} \text{ as } k \rightarrow \infty$$

Generalization to matrices

$$(I-M)^{-1} \approx I + M + M^2 + \dots + M^k$$

$$(I-M)(I+M+M^2+\dots+M^k) = I - M^{k+1} \quad \text{when does this cvg} \rightarrow I$$

" " $M^k \rightarrow 0$ as $k \rightarrow \infty$.

Claim: If $\rho(M) < 1$, then $M^k \rightarrow 0$ as $k \rightarrow \infty$.

i.e. $\|M^k - 0\| \rightarrow 0$ so $M^k \rightarrow 0$ in $\|\cdot\|$

(and in every other norm,

Equivalently Contraction map $cy \rightarrow$ fixed pt. at 0)

Just proved

Prop: If $\rho(M) < 1$, then $I-M$ is invertible and

$$(I-M)^{-1} = I + M + M^2 + \dots$$

To solve, $(I-M)x = b$

$$x = (I-M)^{-1} \cdot b$$

$$= (I + M + M^2 + M^3 + \dots) b$$

Approx by $x^{(k)} = (I + M + M^2 + \dots + M^k) b \rightarrow x$ as $k \rightarrow \infty$

Observe:

$$x^{(k+1)} = (I + M + M^2 + \dots + M^k + M^{k+1}) b$$

$$= b + (M + M^2 + \dots + M^{k+1}) b$$

$$= b + M(I + M + \dots + M^k) b$$

$$= b + Mx^{(k)}$$

So take iteration :

$$\begin{cases} x^{(0)} = b \\ x^{(k+1)} = b + Mx^{(k)} \end{cases} \quad \text{only uses matrix-vector multiplication} \quad O(n^2)$$

If we can get a good approx with $\ll n$ steps,

this does better than $O(n^2)$

Especially true for "sparse" matrices, where most entries = 0.

i.e., $O(n^2)$ nonzero entries

Ex. diagonal matrices $\begin{pmatrix} \diagdown \end{pmatrix}$

tridiagonal matrix $\begin{pmatrix} \diagup \diagdown \diagdown \end{pmatrix}$

Prop Fixed pt ξ is a solution to $(I-M)\xi = b$

Proof. $\xi = b + M\xi$

$$\xi - M\xi = b$$

$$(I-M)\xi = b \quad \square$$

so if $\rho(M) < 1$, then $x^{(k)} \rightarrow \xi$ as $k \rightarrow \infty$.

To solve $A\vec{x} = \vec{b}$ split $A = I - M \Leftrightarrow M = I - A$ and if $\rho(M) < 1$, apply prev method.

What if we can't do this. i.e. what if $\rho(I-A) \geq 1$

Split $A = N - P$
 $\downarrow \quad \searrow$
 nice perturbation
 easy to solve

iteration : $Nx^{(k+1)} = b + Px^{(k)}$

Thm: If $\rho(N^{-1}P) < 1$, then $x^{(k)} \rightarrow \xi$ as $k \rightarrow \infty$

$$\Leftrightarrow x^{(k+1)} = N^{-1}b + N^{-1}Px^{(k)}$$

$$(N-P)\vec{x} = \vec{b} \quad \Leftrightarrow (I - N^{-1}P)\vec{x} = N^{-1}b$$