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**School of Computing Science  
Simon Fraser University**

# **CMPT 471: Computer Networking II**

## **Introduction**

**Instructor: Mohamed Hefeeda**

# Course Objectives

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## ❑ Understand

- ❖ principles of designing and operating computer networks
- ❖ structure and protocols of the Internet
- ❖ services that can/cannot be offered by the Internet

## ❑ Know how to

- ❖ implement network protocols and applications

## ❑ Be informed about

- ❖ recent/hot topics in networking research and industry
- ❖ top technical conferences/journals in networking research

# Course Info: Textbooks and References

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## ❑ Textbooks

- ❖ Kurose and Rose, Computer Networking: A top-down Approach, 7<sup>th</sup> edition, 2016
- ❖ Peterson and Davie, Computer Networks: A Systems Approach,, 5th edition, 2012.
  - Available Online through SFU Library.

## ❑ References

- ❖ Posted on the course web page

## ❑ Course web page

- ❖ Check: [courses.cs.sfu.ca](http://courses.cs.sfu.ca)

# Course Info: Grading (Tentative)

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## ❑ Assignments, Projects: 50%

- ❖ Several programming projects, mostly in C and Java
- ❖ Assignments may include problems sets, researching topics, conducting experiments, presentations, ...
- ❖ **Must read Assignment Policy (on course web page)**

## ❑ Exams: 50%

# Course Info: Topics

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- ❑ Review of Networking Basics:
  - ❖ Internet Architecture and TCP/IP Stack
- ❑ IP Multicast
- ❑ Multimedia Networking
- ❑ Wireless Networks
- ❑ Selected topics from
  - ❖ Virtual Networks and Overlays
  - ❖ Network Security
  - ❖ Software Defined Networks
  - ❖ Cloud Computing
  - ❖ Data Center Networking

# Quick Survey: did you cover ...

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- ☐ Socket programming?
- ☐ Wireshark experiments?
- ☐ IP Multicast?
- ☐ Multimedia Networking?
- ☐ Wireless Networks?
- ☐ Network Security?

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# Basic Networking Concepts

# Review of Basic Networking Concepts

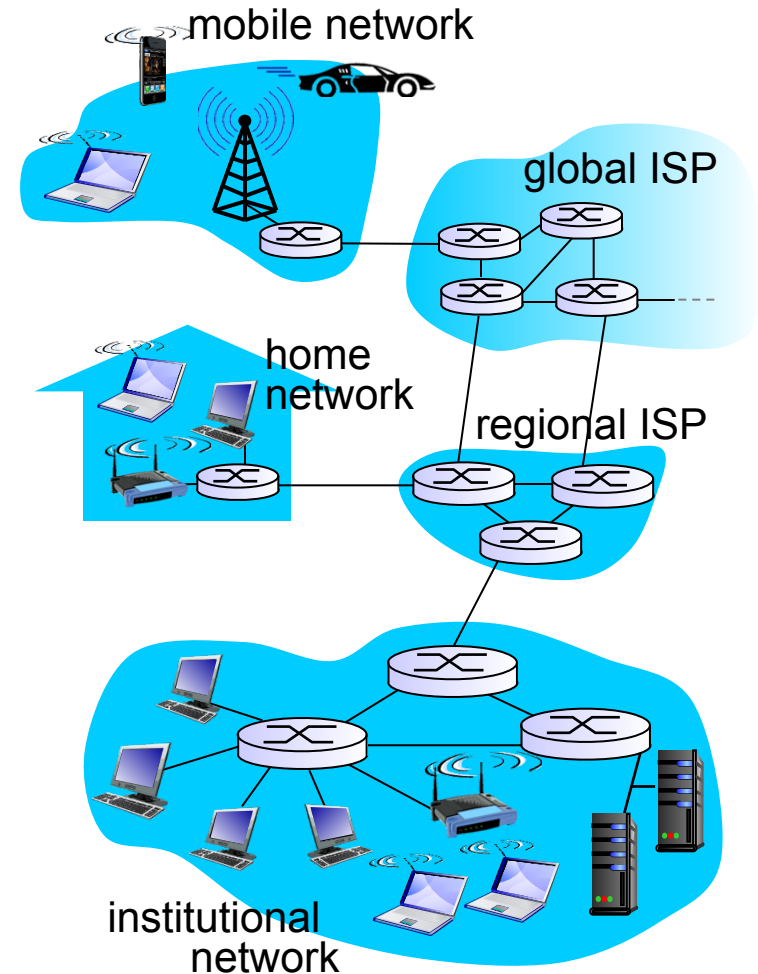
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- ❑ Internet structure
- ❑ Protocol layering and encapsulation
- ❑ Socket programming
- ❑ Transport layer
  - ❖ Reliability and congestion control
  - ❖ Performance modeling of TCP
- ❑ Network Layer
  - ❖ Addressing, Forwarding, Routing
  - ❖ IP Multicast



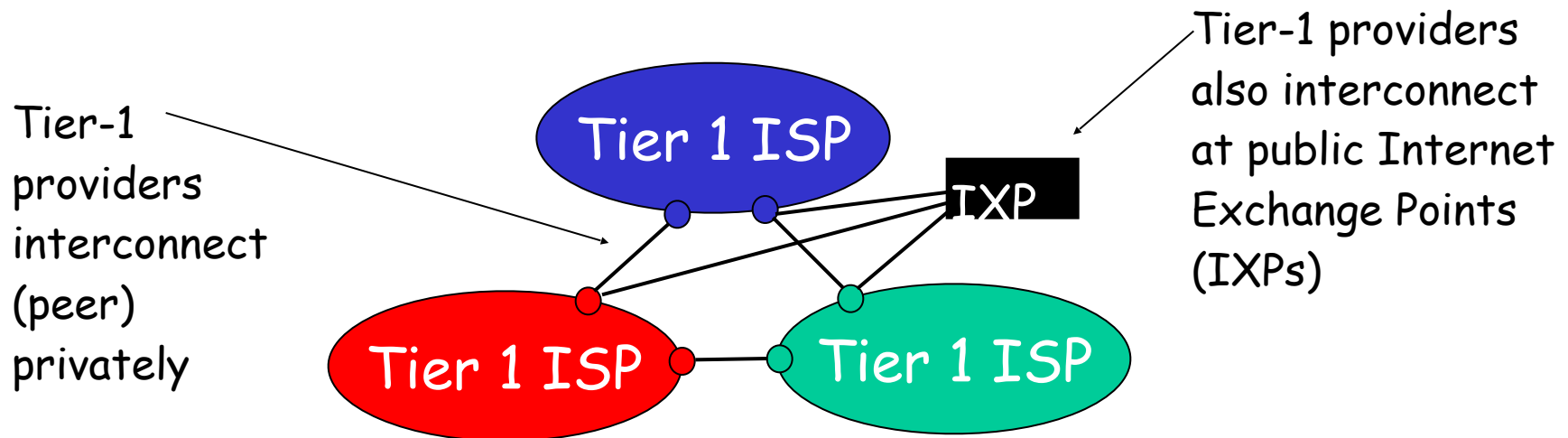
# The Internet: Network of Networks

- ❑ *Internet: “network of networks”*
  - ❖ Interconnected ISPs
- ❑ *protocols* control sending, receiving of messages
  - ❖ e.g., TCP, IP, HTTP, Skype, 802.11
- ❑ *Internet standards*
  - ❖ RFC: Request for comments
  - ❖ IETF: Internet Engineering Task Force

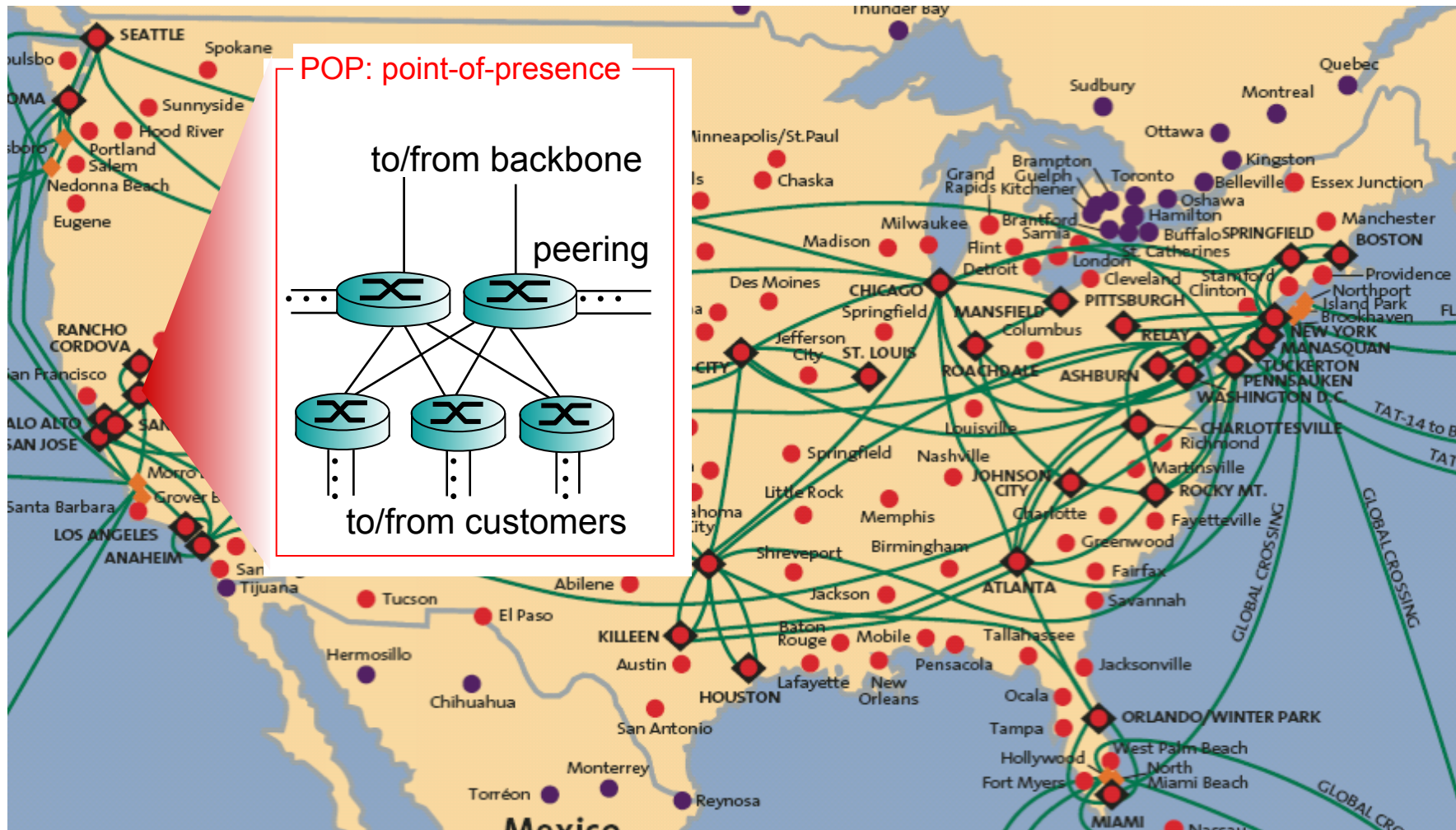


# Internet structure: network of networks

- ❑ roughly hierarchical
- ❑ **at center: “tier-1” ISPs** (e.g., MCI, Sprint, and AT&T), national/international coverage
  - ❖ treat each other as equals

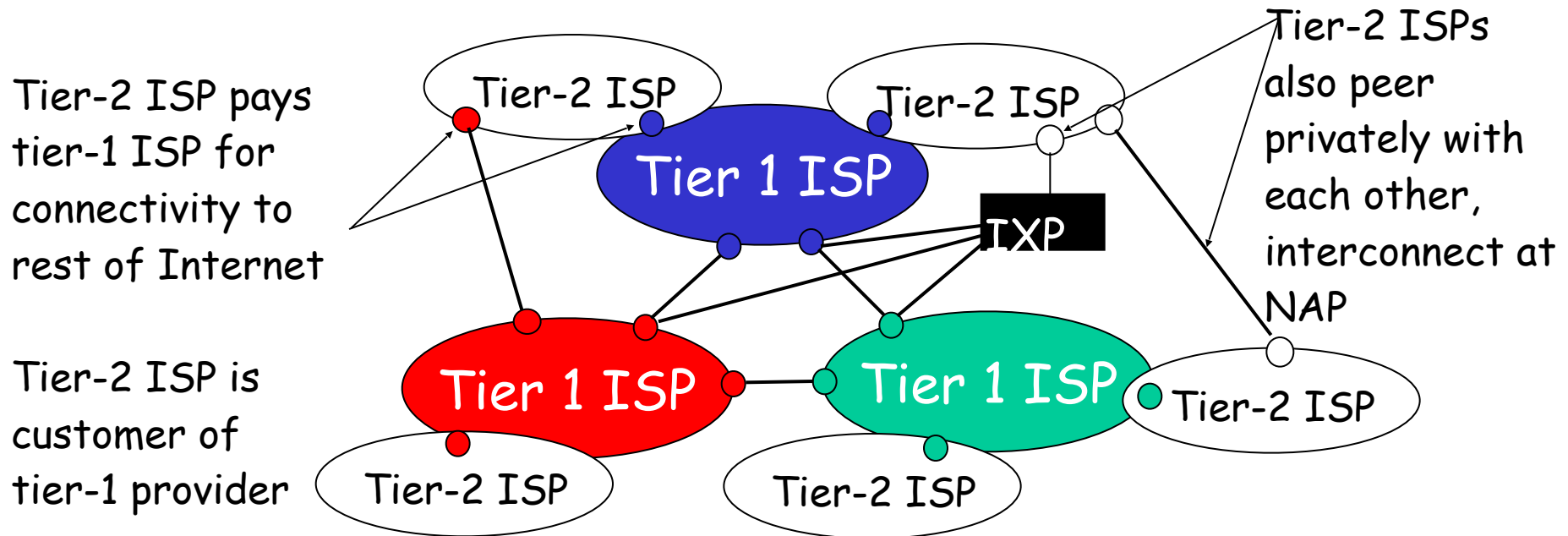


# Tier-1 ISP: e.g., Sprint



# Internet structure: Tier-2 ISPs

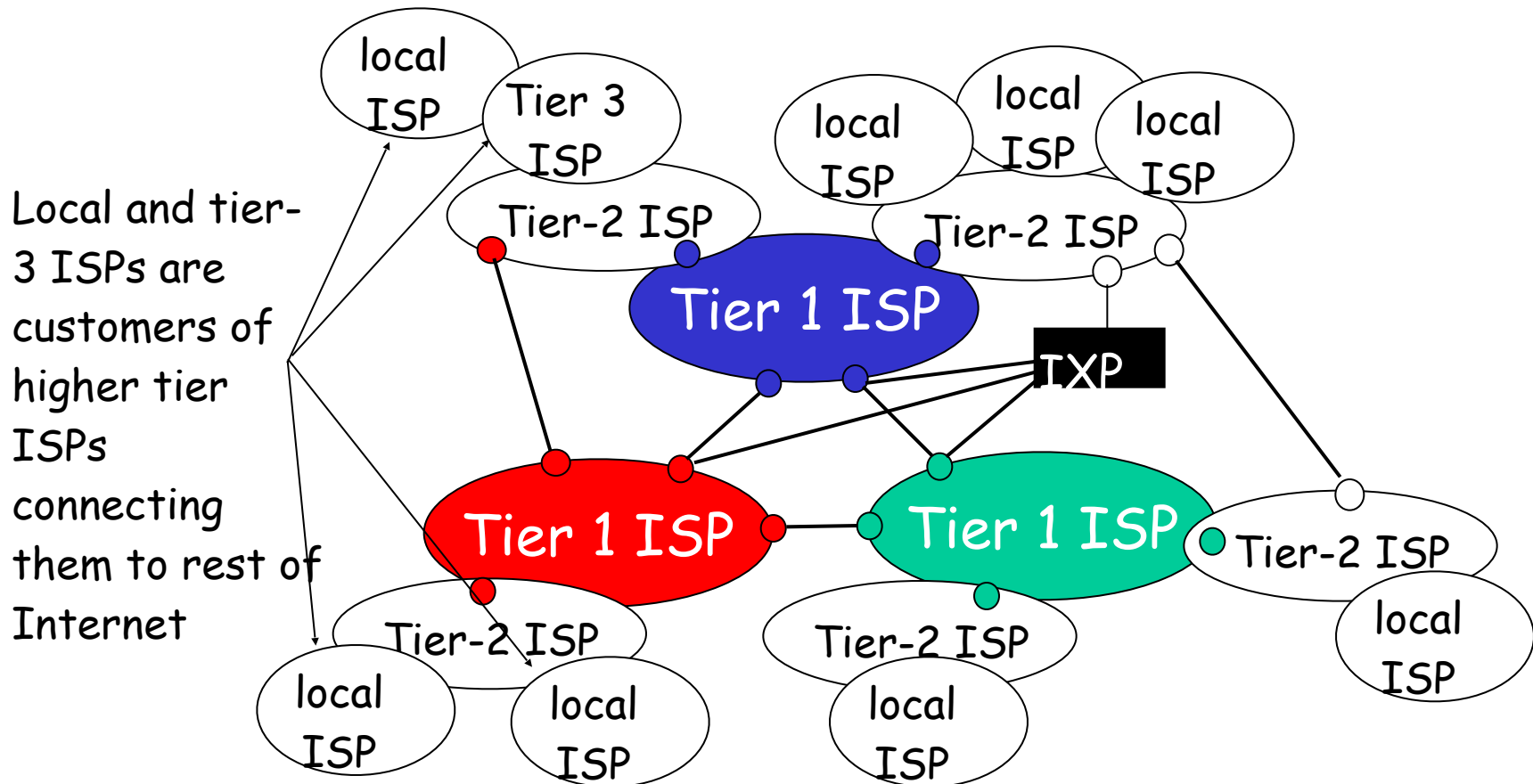
- ❑ **“Tier-2” ISPs: smaller (often regional) ISPs**
  - ❖ Connect to one or more tier-1 ISPs, possibly other tier-2 ISPs



# Internet structure: Tier-3 ISPs

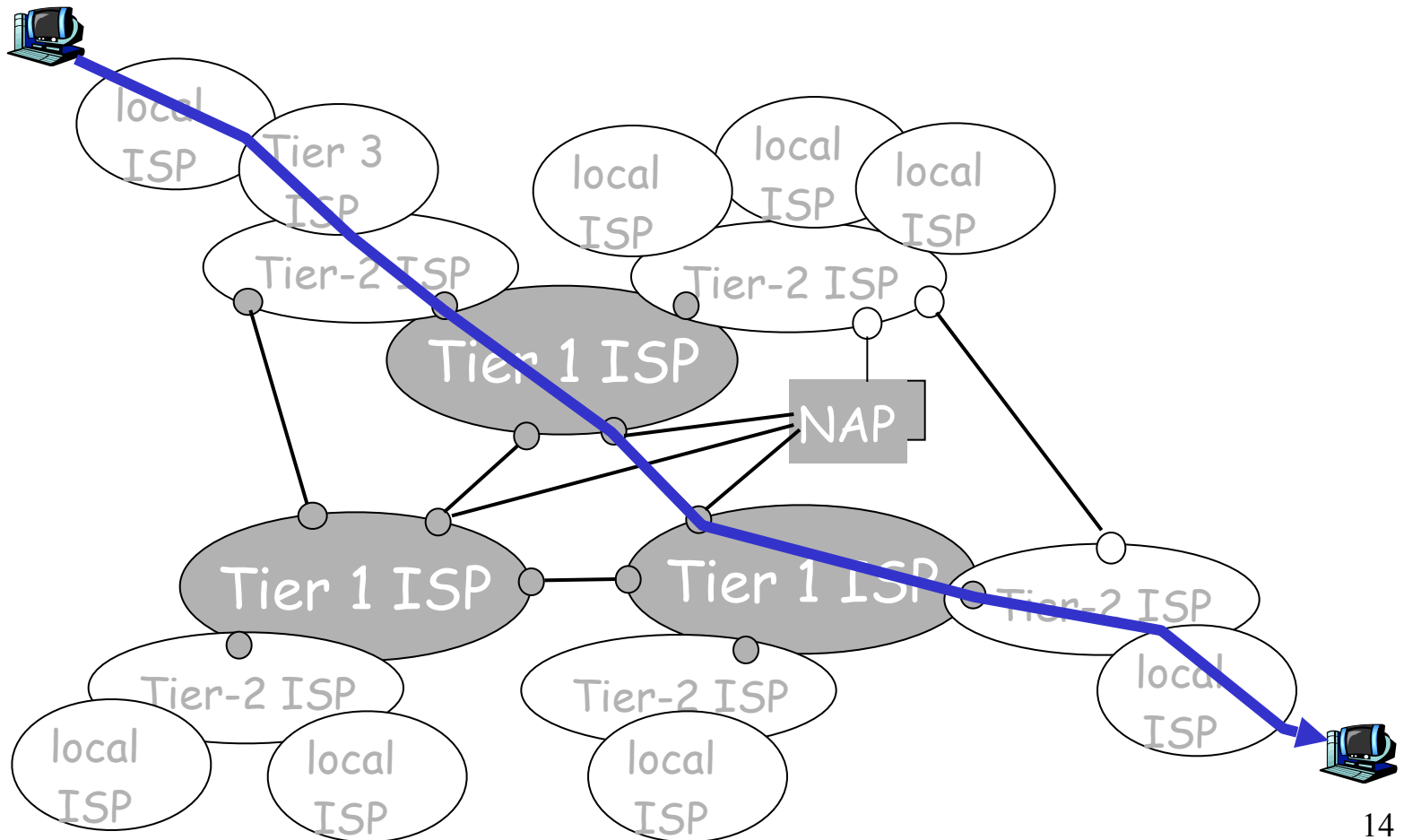
## ❑ “Tier-3” ISPs and local ISPs

❖ last hop (“access”) network (closest to end systems)



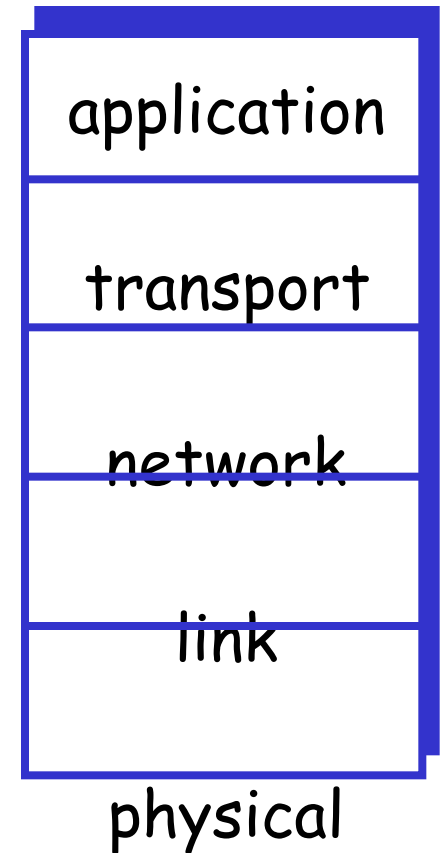
# Internet structure: packet journey

- ❑ a packet passes through many networks!

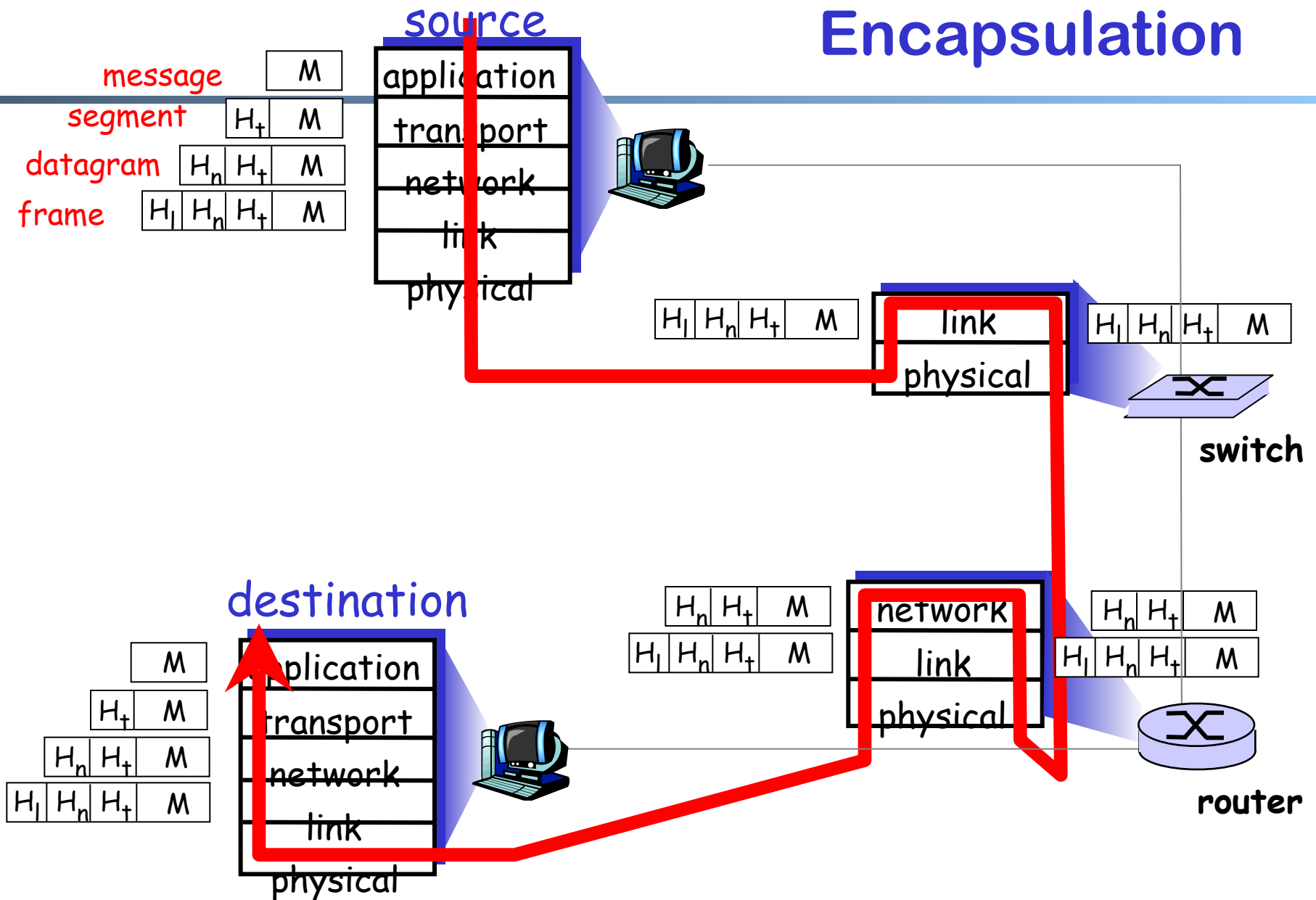


# Internet protocol stack

- ❑ **application:** supporting network applications
  - ❖ FTP, SMTP, HTTP
- ❑ **transport:** process-process data transfer
  - ❖ TCP, UDP
- ❑ **network:** routing of datagrams from source to destination
  - ❖ IP, routing protocols
- ❑ **link:** data transfer between neighboring network elements
  - ❖ PPP, Ethernet
- ❑ **physical:** bits “on the wire”



# Encapsulation





# Internet Services

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- ❑ View the Internet as a **communication infrastructure** that provides **services** to apps
  - ❖ Web, email, games, e-commerce, file sharing, ...
  
- ❑ Two communication services
  - ❖ Connectionless unreliable
  - ❖ Connection-oriented reliable

# Internet Services

## ❑ Connection-oriented

- ❖ Prepare for data transfer ahead of time
- ❖ establish connection → **set up state** in the two communicating hosts
- ❖ Usually comes with reliability, flow and congestion control
- ❖ **TCP**: Transmission Control Protocol

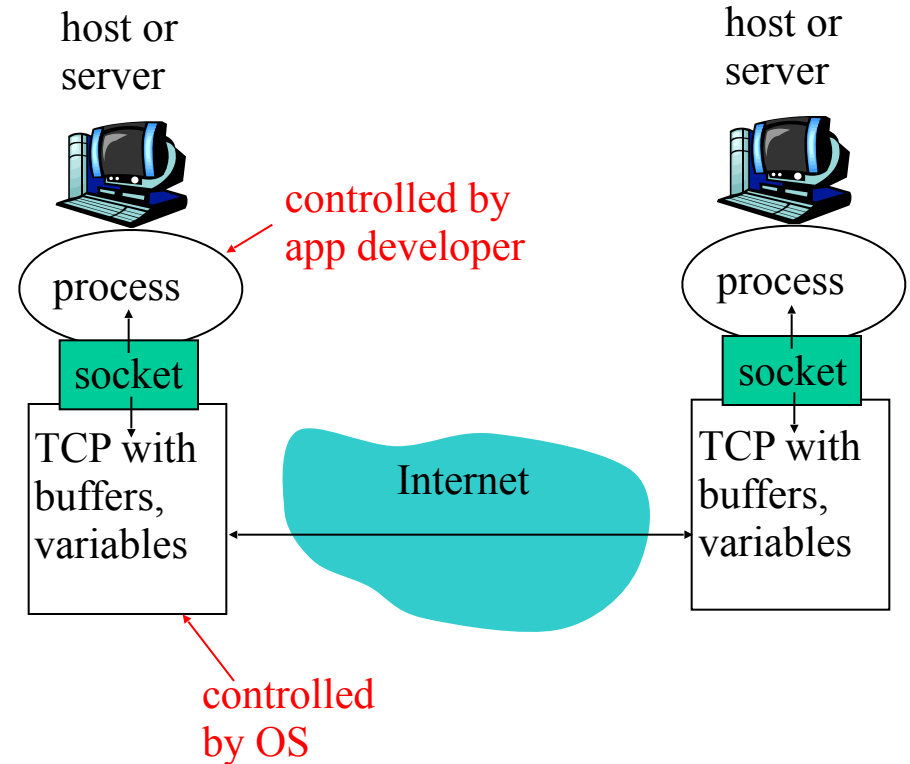
## ❑ Connectionless

- ❖ No connection set up, simply send
- ❖ Faster, less overhead
- ❖ No reliability, flow control, or congestion control
- ❖ **UDP**: User Datagram Protocol

**How can we access these services?**

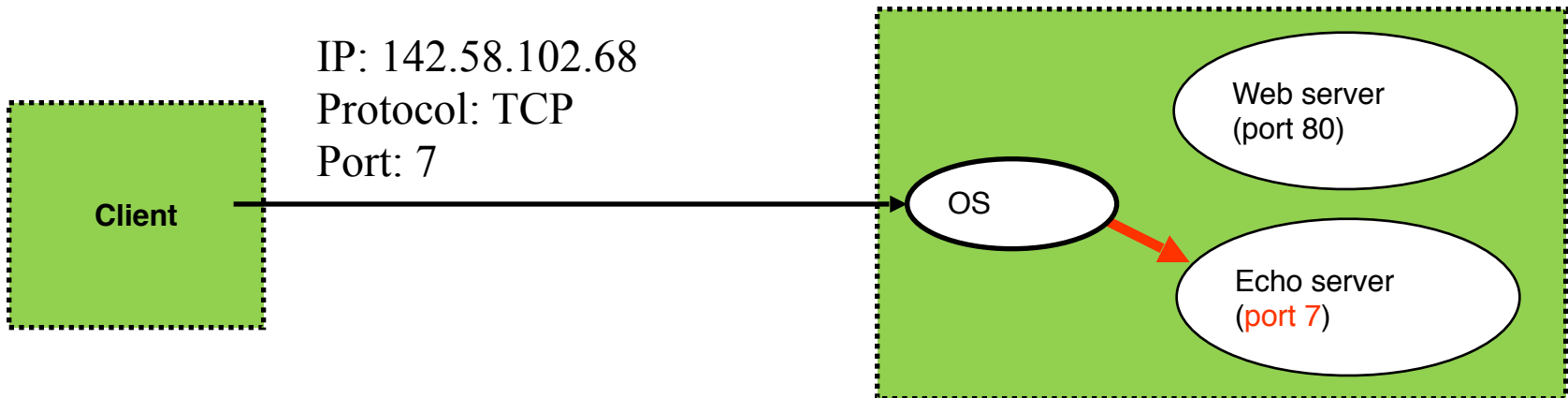
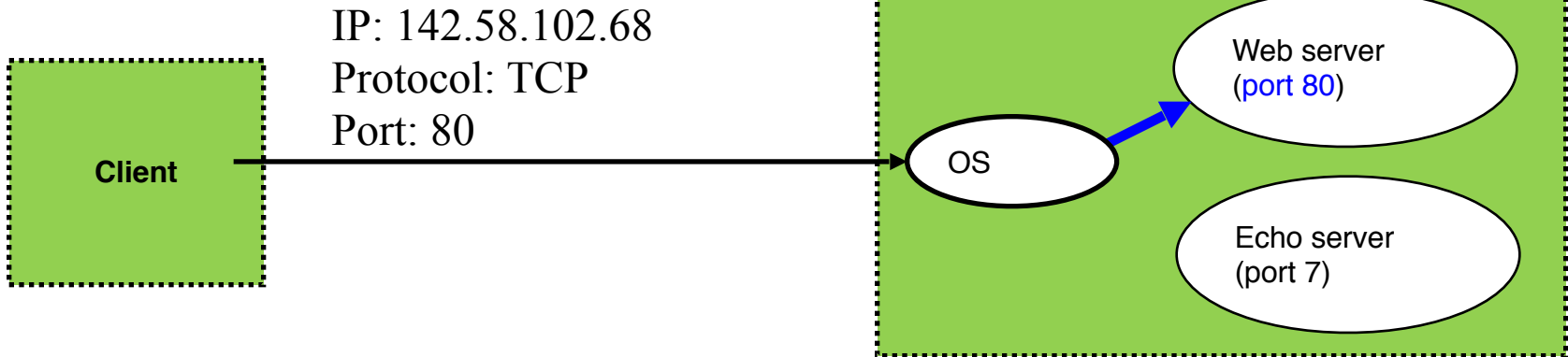
# Network (Socket) Programming

- ❑ Process sends/receives messages to/from its **socket**
- ❑ Socket is the **interface** (API) between application and transport layer
- ❑ Process is identified by:
  - ❖ IP address,
  - ❖ Transport protocol, and
  - ❖ Port number



# Identifying Processes

Server: 142.58.102.68



# Port Numbers

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- ❑ Popular applications have well-known ports
  - ❖ E.g., port 80 for Web and port 25 for e-mail
  - ❖ See <http://www.iana.org/assignments/port-numbers>
  
- ❑ Server port:
  - ❖ Known/fixed (e.g., port 80)
  - ❖ Ports between 0 and 1023 (require root to use)
  
- ❑ Client port:
  - ❖ Client chooses an unused ephemeral (i.e., temporary) port, Between 1024 and 65535

# UNIX Socket API

## ❑ Socket interface

- ❖ Originally provided in Berkeley UNIX
- ❖ Later adopted by all popular operating systems
- ❖ Simplifies porting applications to different OSes

## ❑ In UNIX, everything is like a file

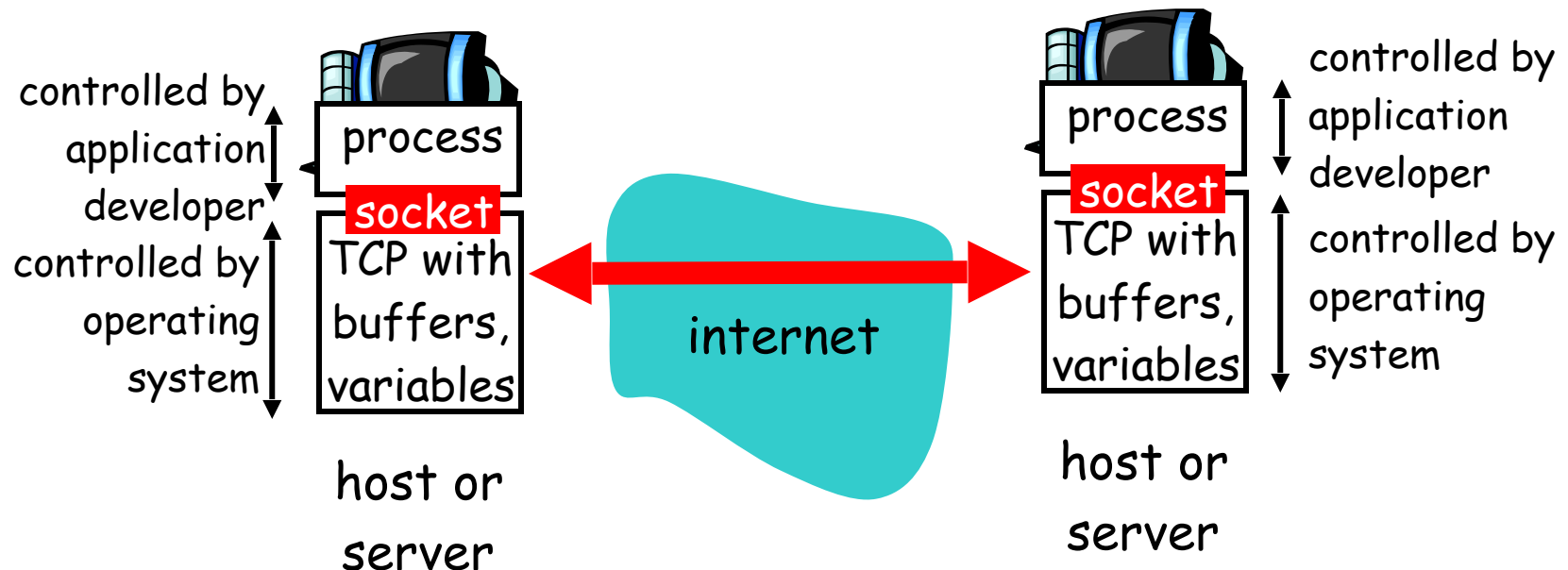
- ❖ All input is like reading a file
- ❖ All output is like writing a file
- ❖ File is represented by an integer file descriptor

## ❑ API implemented as system calls

- ❖ E.g., connect, read, write, close, ...

# Socket Programming using TCP

- ❑ **TCP service:** reliable transfer of **bytes** from one process to another
  - ❖ **virtual pipe** between sender and receiver



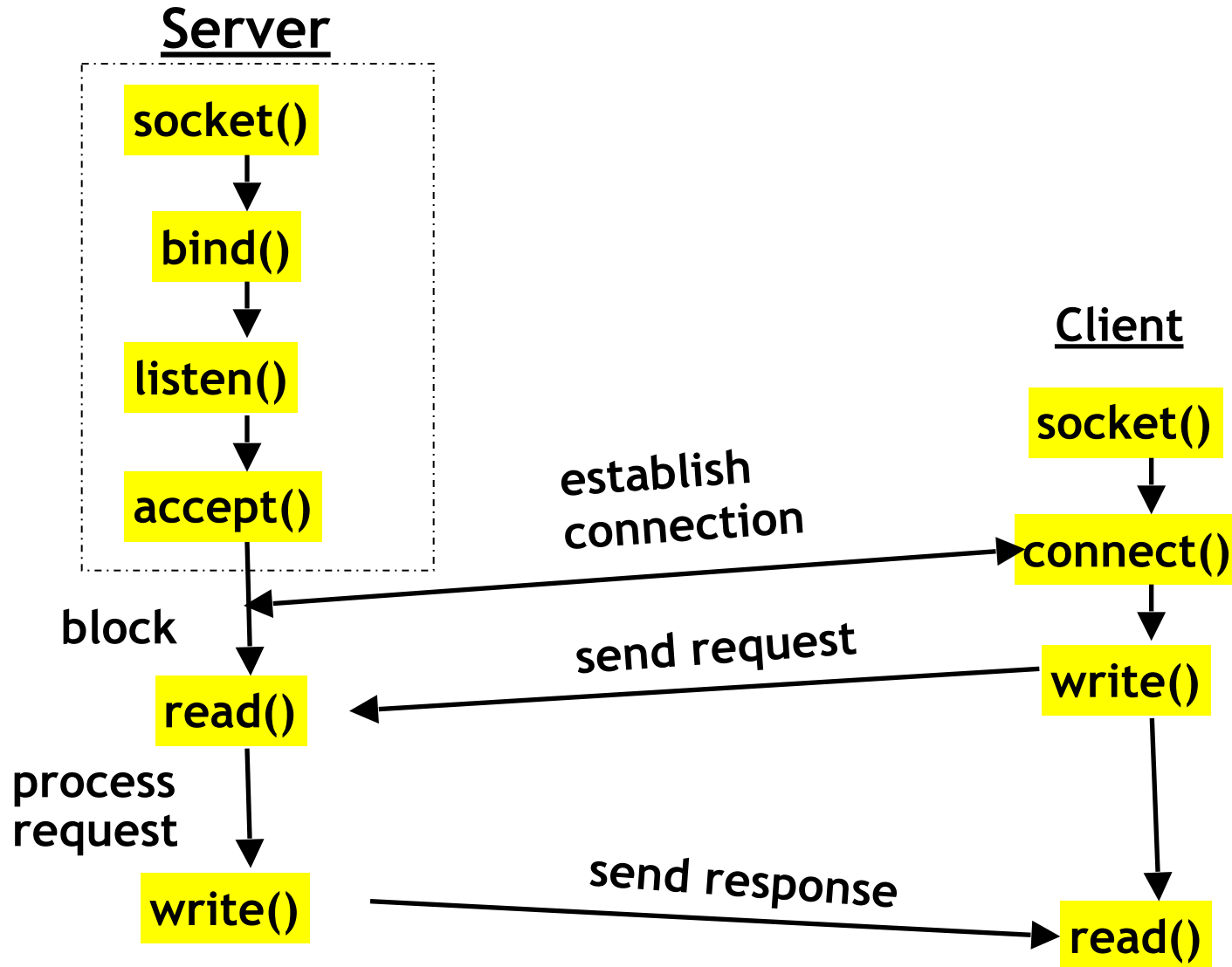
# Socket Programming using TCP

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- ❑ Server process must be running **first**, and
  - ❖ creates a socket (door) that accepts client's contact, then wait
- ❑ Client contacts server by creating local TCP socket using IP address, port number of server process
- ❑ When client creates socket
  - ❖ client TCP (in OS kernel) establishes connection to server TCP
  - ❖ Then data start to flow



# TCP Socket: Basic Structure (Unix/C)

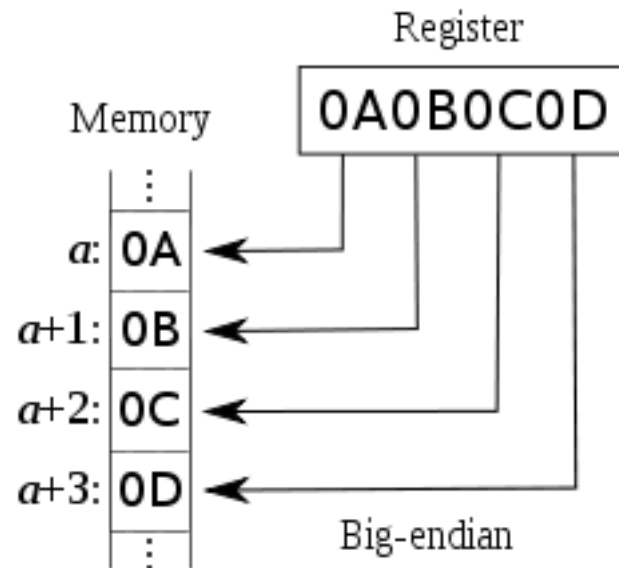
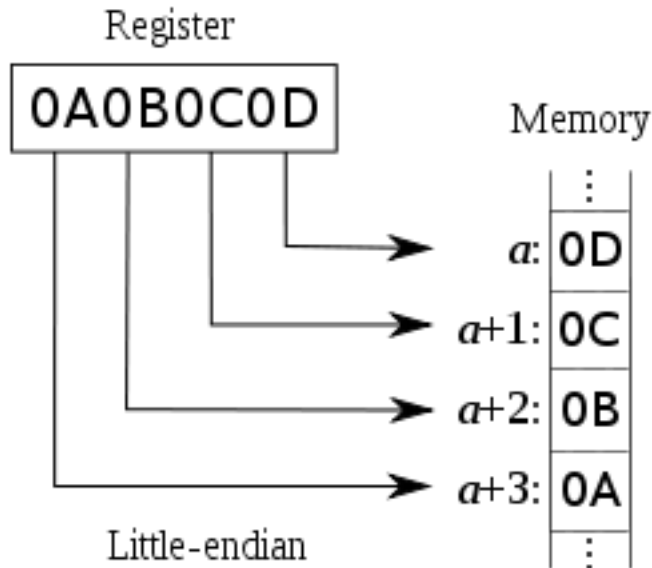


# TCP Daytime Server

```
int main (int argc, char **argv) {  
    int listenfd, connfd;  
    struct sockaddr_in servaddr;  
    char buff[MAXLINE];  
    time_t ticks;  
    listenfd = socket(AF_INET, SOCK_STREAM, 0);  
    bzero(&servaddr, sizeof(servaddr)).
```

# htonX and ntohs macros: Important

- ❑ Some machines use “big endian” and others use “little endian” to store numbers
  - ❖ Whenever sending numbers to network use htonl
  - ❖ Whenever receiving numbers from network use ntohs
    - Replace X with l for long integer (4 bytes), and s for short (2 bytes)



# Creating Socket: socket()

❑ **int socket(int domain, int type, int protocol)**

❖ Returns a file descriptor (or handle) for the socket

❑ **domain: protocol family**

❖ PF\_INET for the Internet (IPv4)

❑ **type: semantics of the communication**

❖ SOCK\_STREAM: reliable byte stream (TCP)

❖ SOCK\_DGRAM: message-oriented service (UDP)

❑ **protocol: specific protocol**

❖ UNSPEC: unspecified

❖ (PF\_INET and SOCK\_STREAM already implies TCP)

# TCP Daytime Client

```
int main(int argc, char **argv) {  
    ...  
    if ( (sockfd = socket(AF_INET, SOCK_STREAM, 0)) < 0) {  
        printf("socket error\n"); exit(1); }  
}
```

# Concurrent TCP Servers

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- ❑ Daytime server accepts one connection at a time
  - ❖ Not good for other servers, e.g., Web servers
- ❑ How would you make it handle multiple connections concurrently?
- ❑ We need some parallelism!
  - ❖ But where?

# TCP Daytime Server

```
int main (int argc, char **argv) {
    int listenfd, connfd;
    struct sockaddr_in servaddr;
    char buff[MAXLINE];
    time_t ticks;
    listenfd = socket(AF_INET, SOCK_STREAM, 0);
    bzero(&servaddr, sizeof(servaddr));
    servaddr.sin_family = AF_INET;
    servaddr.sin_addr.s_addr = htonl(INADDR_ANY);
    servaddr.sin_port = htons(DAYTIME_PORT); /* daytime server */
    bind(listenfd, (struct sockaddr *) &servaddr, sizeof(servaddr));
    listen(listenfd, LISTENQ);

    for ( ; ; ) {
        connfd = accept(listenfd, (struct sockaddr *) NULL, NULL);
        ticks = time(NULL);
        snprintf(buff, sizeof(buff), "%.24s\r\n", ctime(& here
        write(connfd, buff, strlen(buff));
        printf("Sending response: %s", buff);
        close(connfd);
    }
```

# Concurrent Server

```
for ( ; ; ) {  
    connfd = accept(listenfd, ...);  
  
    if ( (pid = fork()) == 0 ) {  
        close(listenfd); /*child closes listening socket */  
  
        doit(connfd); /*process the request */  
  
        close(connfd); /*done with this client */  
  
        exit(0); /*child terminates */  
    }  
  
    close(connfd); /*parent closes connected socket */  
}
```

❑ Fork: duplicates the entire process that called it

❑ Fork returns twice!

- ❖ One to the child process, return value = 0
- ❖ Second to parent with non zero (pid of the created child)



# Summary

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- ❑ Internet structure
- ❑ Protocol layering
- ❑ Socket programming