# **CPT\_S 580 HW2**

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# 2.3.21

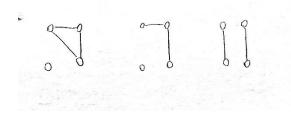
- (a) 1 clique (u, v, y)
- (b) w = 3
- (c) uxz, xyz, xyw, xvz, xuw

# 2.3.22

- (a) 3 cliques (u, x, y), (u, v, y), (y, v, w)
- (b) w = 3
- (c) xvz, xw, uw, uz, yz

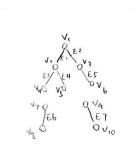
# 2.3.30

All possible isomorphism types of a simple 4-vertex graph with exactly 2 components



# 2.3.36

A forest with 10 vertices, 7 edges, and 3 components



### 2.4.8

Cut-vertices: {u, v, z}

Cut-edges: {uv}

### 2.4.9

Cut-vertices: {u, v, x, z, s}

Cut-edges: {uv, xz, zs, st}

## 2.4.30

Since the  $\mathbf{e}$  is the cutting-edge, removing it will separate parts that connected by  $\mathbf{e}$ , and since any vertex is not removed, which means the two separate parts are connected. Therefore, connected graph G – cutting edge  $\mathbf{e}$  has exactly two components.

## 2.5.3

The left-side graph has 1 4-edge circle, while the right-side graph has none.

# 2.5.5

Yes, they are in the pair of isomorphic

Cause we can map one graph to the other one.

Map: {a->1, b->4, c->2, d->3, e->5}

## 3.1.15

Suppose graph G has n vertices v<sub>1</sub> to v<sub>n</sub>. And suppose the average degree of its vertices is 2, then,

$$deg(v_1) + ... + deg(v_n) = 2n$$

which implies that the number of edges in G is n. The graph G has n vertices; thus it cannot be a tree. Thus G has at least one cycle, say C. Let e be an edge in C. Then the graph G - e is connected, and has n - 1 edges, hence it is a tree. By a theorem in the notes, adding one edge e to the tree G - e creates exactly one cycle. Thus G has exactly one cycle.

### 3.1.17

Let G have n vertices and n edges. Since G is a connected graph, it has a spanning tree T with n vertices and n-1 edges. Let e be the edge not in T, with its endpoints u and v. There is a unique path  $\gamma$  between u and v in T (since T is a tree). The union of e and  $\gamma$  is a cycle. Suppose that there is some other cycle  $\delta$ . If  $\delta$  does not contain e, then it is contained in T, contradicting that T has no cycles.

#### 3.1.18

We know that n-vertex simple graph with n edge has exactly 1 cycles. Then add 1 edge to the graph would connect 2 non-adjacent vertices (since simple graph), which make the 1 cycle into 2 cycles.

### 2.2.10

For every vertex in hypercube graph  $Q_n$ , the vertex-degree is identical, which means any vertex has the same local environment. Thus, for any two vertices of  $Q_n$ , the vertex-bijection that swaps them and fixes all the others is adjacency-preserving. It is its own inverse, so the inverse bijection is also adjacency-preserving.

### 2.6.24

When n = 1, the entry  $a_{ij}$  is 1 if {i, j}  $\in$  E. By definition, i {i, j} j is then an i – j walk of length 1 and this is the only one. So the statement is true for n = 1. Now, we assume the statement is true for n and then prove the statement is also true for n + 1. Since  $A_{ij}^{n+1} = A_{ij}^n \cdot A_{ij}$ , therefore,  $a_{ij}^{n+1} = \sum_{k=1}^m a_{ik}^n \cdot a_{kj}$ . Because  $a_{ki} = 0$  whenever {k, i}  $\in$ / E and  $a_{ki} = 1$  if {k, i}  $\in$  E, it follows that a n  $a_{ik}^n \cdot a_{kj}$  represents the number of those i–j walks that are i–k walks of length n joined by the edge {k, j}. In particular, all walks from i to j of length n + 1 are of this form for some vertex k.

Thus  $a_{ij}^{n+1} = \sum_{k=1}^m a_{ik}^n \bullet a_{kj}$  indeed represents the total number of i – j walks of length n + 1. This proves the statement for n + 1. Then by the principle of induction, we prove the statement for all natural numbers n.