

CDB: A Crowd-Powered Database System

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Crowd-Based Database

Crowd-Based database can execute some queries which are hard for traditional database

Select Affiliation From Professor, Paper
Where Professor CROWDJOIN Paper AND
Paper.Conf. CROWDEQUAL "SIGMOD"

Do Guo. Li and G.Li refer to the same person?

Yes

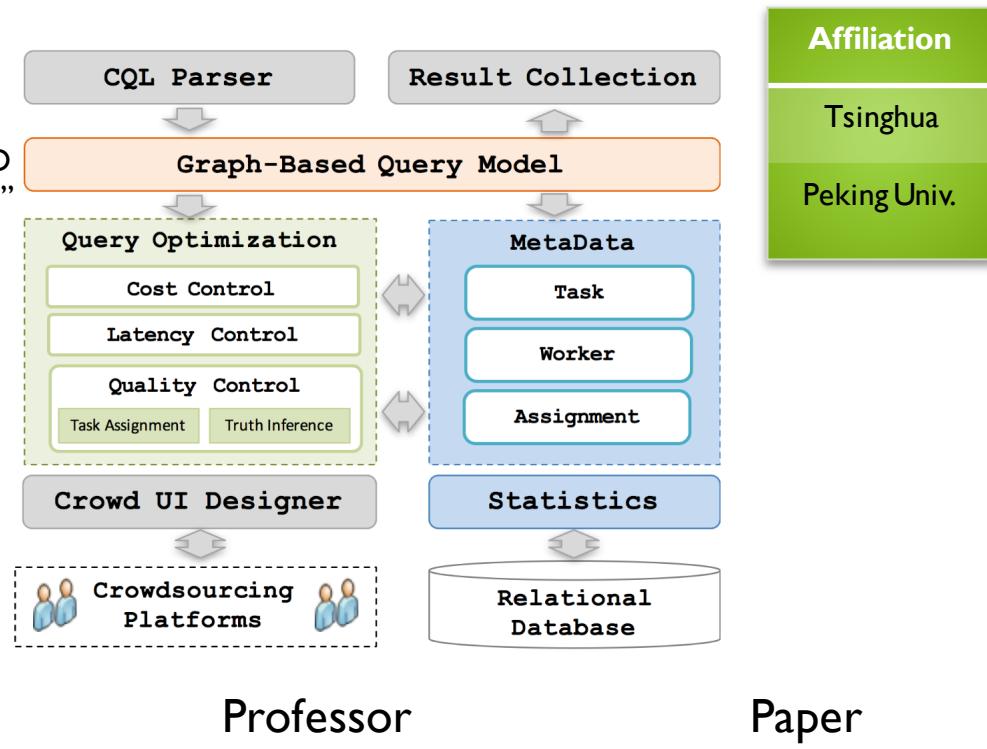
No

Do G. Li and J.Li refer to the same person?

Yes

No

.....



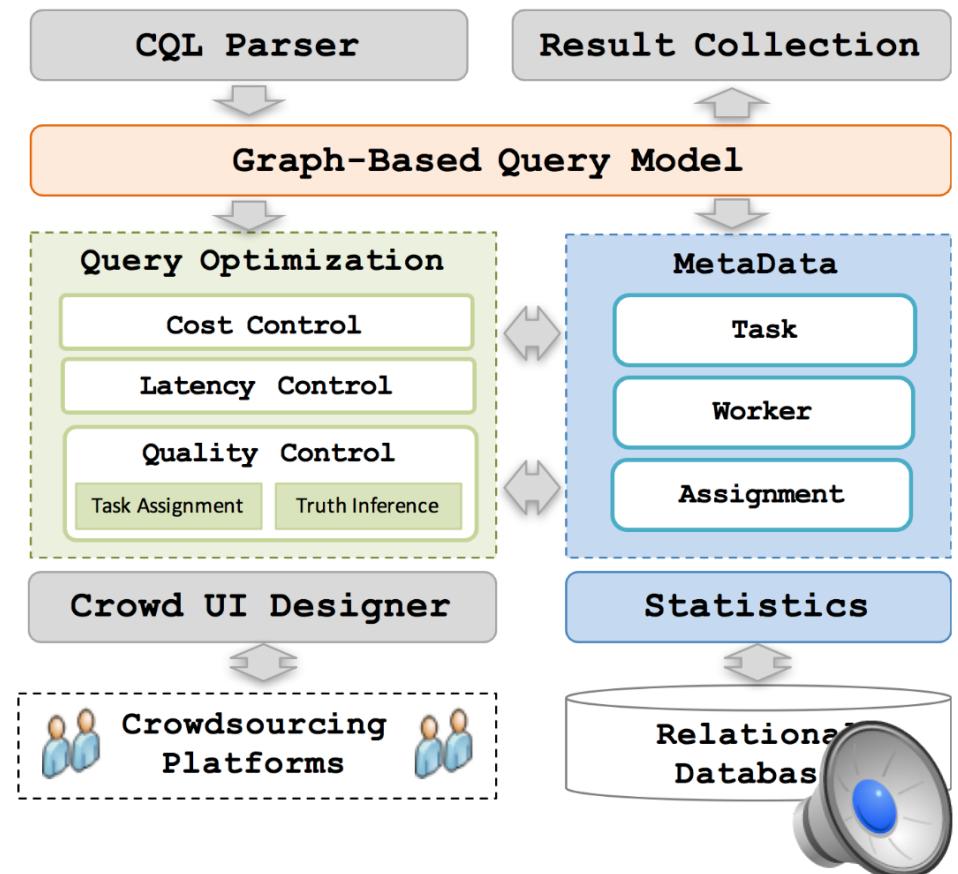
Prof.	Affiliation
Guo. Li	Tsinghua
J. Li	Peking Univ.

Author	Title	Conf.
G. Li	xxx	SIGMOD
Jian. Li	xxx	SIGMOD



Workflow

- A requester submits her query using CQL, which will be parsed by CQL Parser.
- Graph-based query model builds a graph model based on the parsed result.
- Query optimization generates an optimized query plan
- Crowd UI Designer designs various interfaces and interacts with underlying crowdsourcing platforms.



Motivation

- Optimization Models

Existing works: Tree-based model (table-level)

CDB: Graph-based model (tuple-level).

- Optimizing Goals:

Existing works: Mainly on cost.

CDB: Focus on multiple goals (cost, quality and latency).



Graph Model

Weight(Jaccard, ED) $w(e) > threshold$

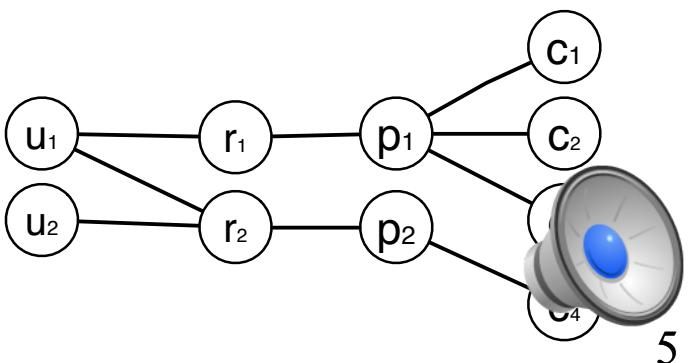
	Country	Name
u ₁	UK	Univ. of Cambridge
u ₂	US	Microsoft

	Affiliation	Name
r ₁	University of Cambridge	Nandan Parameswaran
r ₂	Microsoft Cambridge	S. Chaudhuri

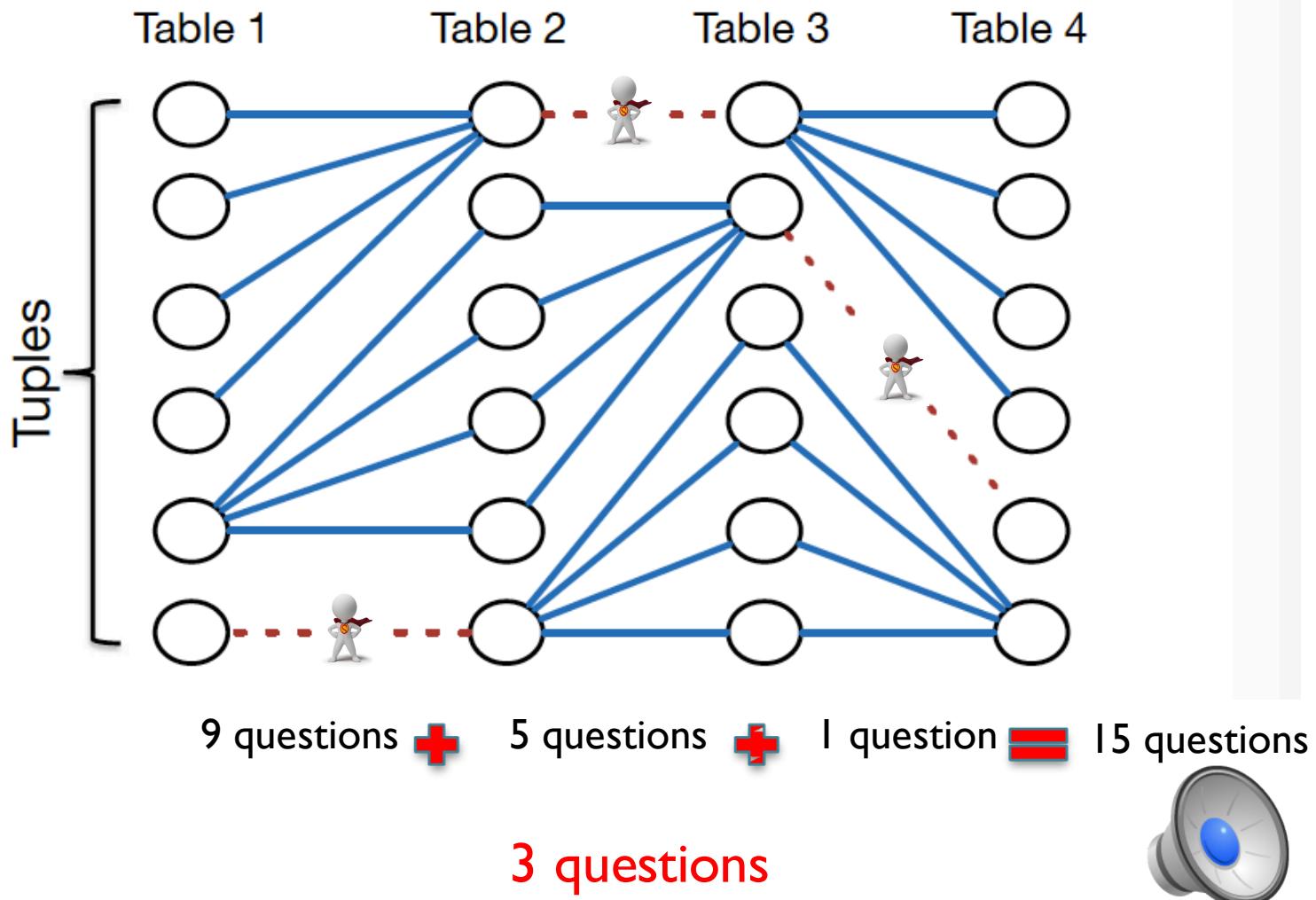
	Number	Title
c ₁	16	DataSift: An Expressive and Accurate Crowd-Powered Search Toolkit.
c ₂	4	A crowd powered search toolkit
c ₃	0	A Crowd Powered System for Similarity Search
c ₄	1	Query portals: dynamically generating portals for entity-oriented web queries.

	Title	Author
p ₁	DataSift: a crowd-powered search toolkit	Aditya G. Parameswaran
p ₂	Dynamically generating portals for entity-oriented web queries.	Surajit Chaudhuri

- For each table T in the CQL query, there is a vertex for each tuple in this table.
- For each crowd join predicate $T.C_i \text{ CROWDJOIN } T'.C_i$ in the CQL query, there is an edge e between $t \in T$ and $t' \in T'$ with $w(e) > threshold$.



Tuple-level VS Table-level



Differences with Existing Systems

Reduce



		CrowdDB	Qurk	Deco	CrowdOP	CDB
Optimized Crowd Operators	COLLECT	✓	✗	✓	✗	✓
	FILL	✓	✗	✓	✓	✓
	SELECT	✓	✓	✓	✓	✓
	JOIN	✓	✓	✓	✓	✓
Optimization Objectives	Cost	✓	✓	✓	✓	✓
	Latency	✗	✗	✗	✓	✓
	Quality	MV	MV	MV	MV	✓
Optimization Strategies	Cost-Model	✗	✓	✓	✓	✓
	Tuple-Level	✗	✗	✗	✗	✓
Budget-Supported		✗	✗	✗	✗	✓
Task Deployment	Cross-Market	✗	✗	✗	✗	✓

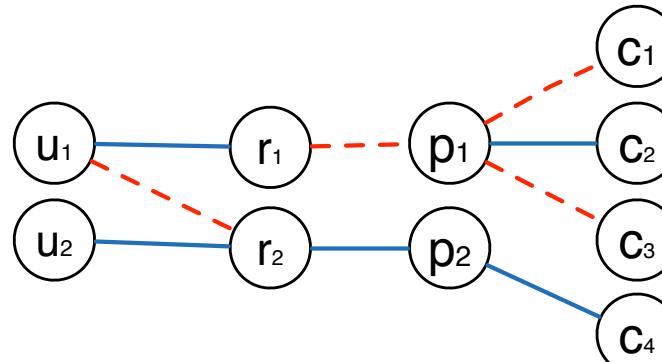


Contributions

- Optimization Models: Graph-based model (tuple-level).
- Optimizing Goals: Focus on multiple goals (cost, quality and latency).
- Many commonly used crowd-powered operators.
- Cross-market HITs deployment.



Cost Control



CQL Query Candidate:

(u_1, r_1, p_1, c_1) , (u_1, r_1, p_1, c_2) , (u_1, r_1, p_1, c_3) , (u_1, r_2, p_2, c_4) , (u_2, r_2, p_2, c_4)

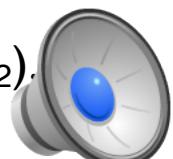
CQL Query Answer:

(u_2, r_2, p_2, c_4)

Given the colors of every edge, how to select the minimum number of edges to find all the answers ?

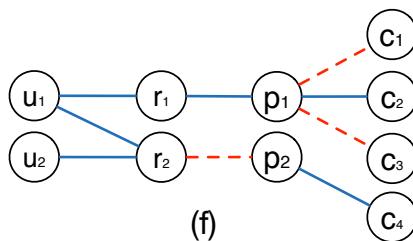
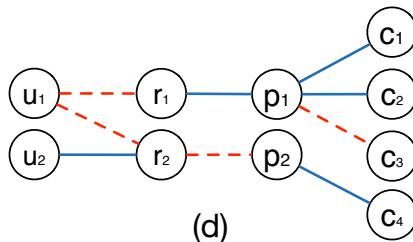
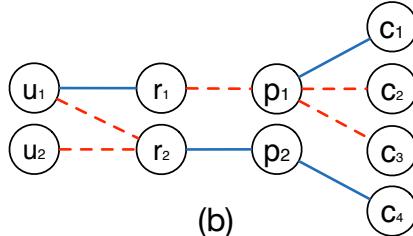
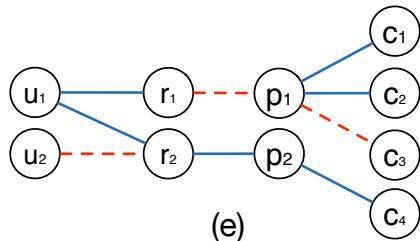
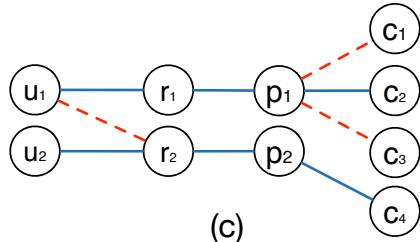
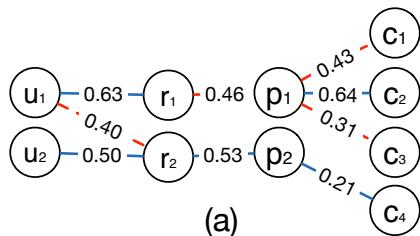
Min-Cut Based Algorithm (refer to the paper for detail)

In the example, the optimal edges are (u_2, r_2) (r_2, p_2) (p_2, c_4) (r_1, p_1) (u_1, r_2) .



Consider the case where the colors of edges are **unknown**. We aim to ask fewer edges to find all answers with high probability.

Sample Average



Given S sample graphs,
select the minimum number of
edges to resolve all samples

- (b) $(U_1, r_2), (U_2, r_1), (r_1, p_1), (r_1, p_2), (p_1, c_1), (p_1, c_2), (p_2, c_4)$
- (d) $(U_1, r_2), (U_2, r_1), (r_1, p_1), (r_1, p_2), (p_1, c_1), (p_1, c_2), (p_2, c_3)$
- (e) $(r_1, p_1), (U_2, r_2), (U_1, r_2), (U_2, r_1), (p_1, c_1), (p_1, c_2), (p_2, c_3)$

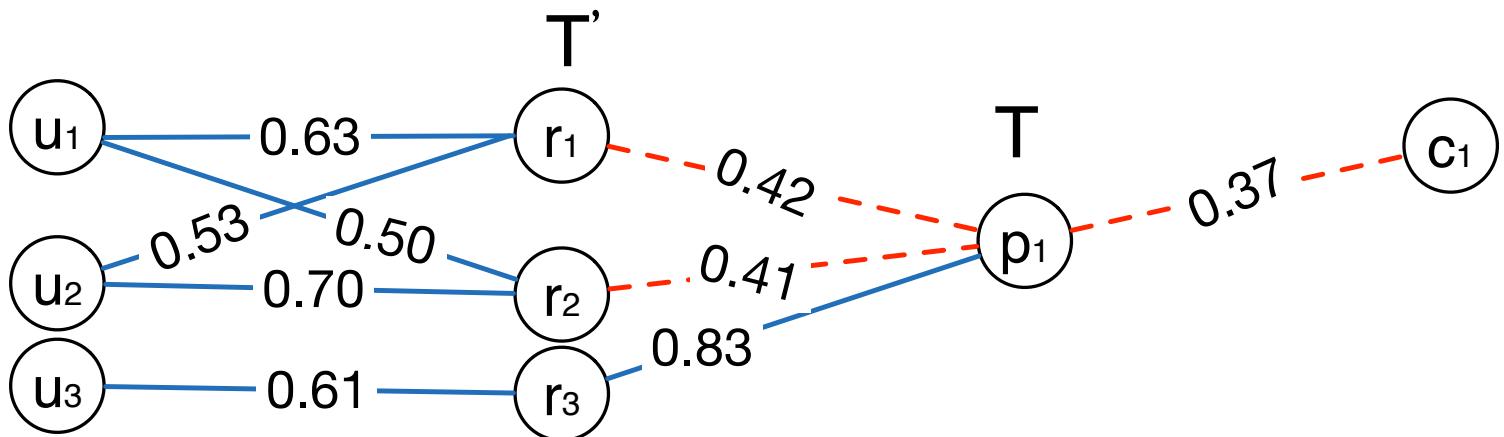
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NP-HARD

Greedy algorithm



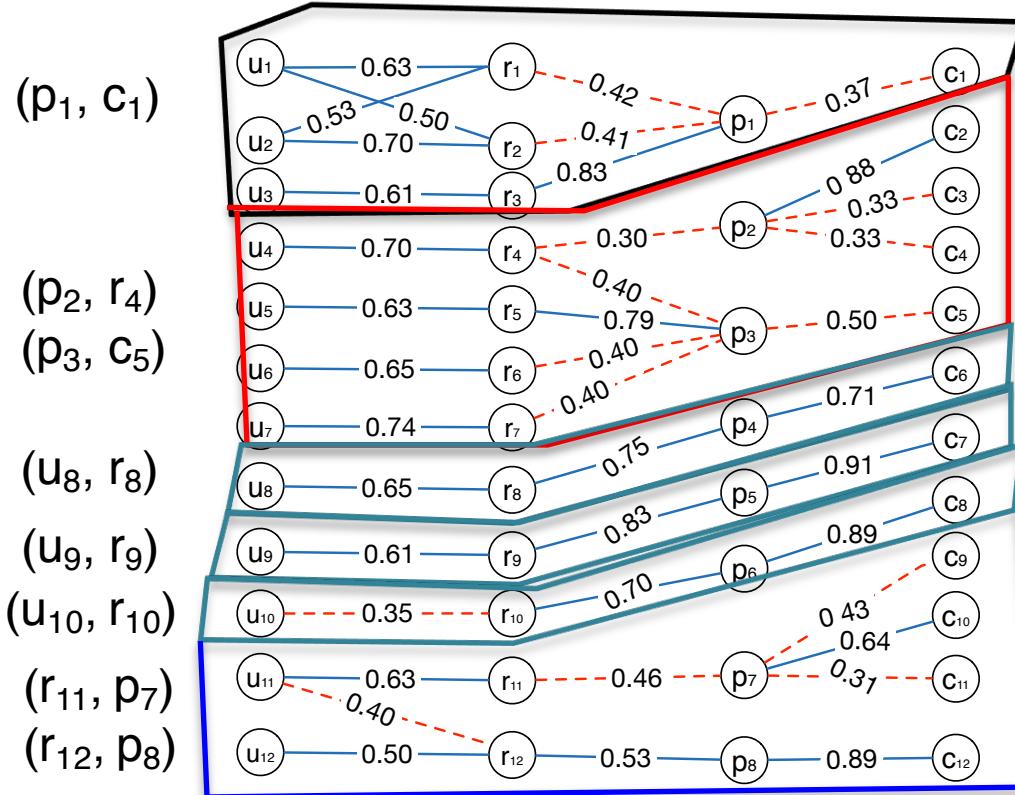
Expectation-based Method



$$\mathbb{E}(t, t') = \frac{\prod_{i=1}^x (1 - \omega(t, t_i))}{x} \alpha + \frac{\prod_{i=1}^y (1 - \omega(t_i, t'))}{y} \beta.$$

$$\mathbb{E}(r_1, p_1) = (1 - 0.42) * 2 + (1 - 0.42) * (1 - 0.41) * (1 - 0.83) * 6 / 3 = 1$$

Latency Control



Connected Components

e.g. (p_1, c_1) (p_2, c_2)

Edges Containing Tuples from the Same Table.

e.g. (p_1, r_1) (p_1, r_2)



Quality Control

Truth Inference

$W = \{w\}$: a set of workers

$T = \{t\}$: a set of tasks

$V_t = \{(w, a)\}$: worker w provides answer a for task t

The probability of the i -th choice being the truth for task t is computed as:

$$p_i = \frac{\prod_{(w, a) \in V_t} (q_w)^{\mathbb{1}_{\{i=a\}}} \cdot \left(\frac{1-q_w}{\ell-1}\right)^{\mathbb{1}_{\{i \neq a\}}}}{\sum_{j=1}^{\ell} \prod_{(w, a) \in V_t} (q_w)^{\mathbb{1}_{\{j=a\}}} \cdot \left(\frac{1-q_w}{\ell-1}\right)^{\mathbb{1}_{\{j \neq a\}}}}$$

Other types of tasks: refer to the paper for detail



Quality Control

Task Assignment

Assign a set of k tasks to worker w , such that the quality can be improved the most.

Two main problems:

- (i) unknown ground truth
- (ii) how the worker can answer each task.

Distribution of choices being true for each task t

$$P=(p_1, p_2, \dots p_{l-1})$$

Entropy function:

$$H(P)=-\sum p_i \log(p_i)$$

The lower $H(p)$ is, the more consistent P is, the higher quality will be achieved.



Quality Control

Task Assignment

Probability that the i -th choice will be answered by w :

$$p_i \cdot q_w + (1 - p_i) \cdot \frac{1 - q_w}{\ell - 1}$$

Then after worker w answers task t with the i -th choice, the distribution is as follows:

$$\vec{p}' = \left(\frac{p_1 \cdot \frac{1 - q_w}{\ell - 1}}{\Delta}, \dots, \frac{p_i \cdot q_w}{\Delta}, \dots, \frac{p_\ell \cdot \frac{1 - q_w}{\ell - 1}}{\Delta} \right)$$

The expected quality of improvement

$$\mathcal{I}(t) = \mathcal{H}(\vec{p}) - \sum_{i=1}^{\ell} \left[p_i \cdot q_w + (1 - p_i) \cdot \frac{1 - q_w}{\ell - 1} \right] \cdot \mathcal{H}(\vec{p}')$$

Other types of tasks: refer to the paper for detail



Task Type & UI Designer

Please choose the brand of the phone



- Apple
- Samsung
- Blackberry
- Other



Which ones are correct?

- The same band
- The same size
- Different bands
- Different sizes

Please fill the attributes of the product



Brand	
Price	
Size	
Whether has camera	



Please submit a picture of a phone, which is the same brand as the left one.



Submit



Experiment

Dataset Paper

Table	#Records	Attributes
Paper	676	author , <i>title</i> , conference
Citation	1239	<i>title</i> , number
Researcher	911	<u>affiliation</u> , name , gender
University	830	<u>name</u> , city, country

Award

Table	#Records	Attributes
Celebrity	1498	name , <i>birthplace</i> , <i>birthday</i>
City	3220	<i>birthplace</i> , country
Winner	2669	name , <u>award</u>
Award	1192	<u>name</u> , place



CQL Queries

Table 4: The 5 representative queries used on paper and award.

Query	Dataset paper	Dataset award
2 Joins (2J)	<pre>SELECT Paper.title, Researcher.affiliation, Citation.number FROM Paper, Citation, Researcher WHERE Paper.title CROWDJOIN Citation.title AND Paper.author CROWDJOIN Researcher.name</pre>	<pre>SELECT Winner.award, City.country FROM Winner, City, Celebrity WHERE Celebrity.name CROWDJOIN Winner.name AND Celebrity.birthplace CROWDJOIN City.name</pre>
2 Joins 1 Selection (2J1S)	<pre>SELECT Paper.title, Researcher.affiliation, Citation.number FROM Paper, Citation, Researcher WHERE Paper.title CROWDJOIN Citation.title AND Paper.author CROWDJOIN Researcher.name AND Paper.conference CROWDEQUAL "sigmod"</pre>	<pre>SELECT Winner.award, City.country FROM Winner, City, Celebrity WHERE Celebrity.name CROWDJOIN Winner.name AND Celebrity.birthplace CROWDJOIN City.name AND Celebrity.birthplace CROWDEQUAL "New York"</pre>
3 Joins (3J)	<pre>SELECT Paper.title, Citation.number, University.country FROM Paper, Citation, Researcher, University WHERE Paper.title CROWDJOIN Citation.title AND Paper.author CROWDJOIN Researcher.name AND University.name CROWDJOIN Researcher.affiliation</pre>	<pre>SELECT Winner.name, Award.place, City.country FROM Winner, City, Celebrity, Award WHERE Celebrity.name CROWDJOIN Winner.name AND Celebrity.birthplace CROWDJOIN City.name AND Winner.award CROWDJOIN Award.name</pre>
3 Joins 1 Selection (3J1S)	<pre>SELECT Paper.title, Citation.number FROM Paper, Citation, Researcher, University WHERE Paper.title CROWDJOIN Citation.title AND Paper.author CROWDJOIN Researcher.name AND University.name CROWDJOIN Researcher.affiliation AND University.country CROWDEQUAL "USA"</pre>	<pre>SELECT Winner.name, City.country FROM Winner, City, Celebrity, Award WHERE Celebrity.name CROWDJOIN Winner.name AND Celebrity.birthplace CROWDJOIN City.name AND Winner.award CROWDJOIN Award.name AND Award.place CROWDEQUAL "US"</pre>
3 Joins 2 Selections (3J2S)	<pre>SELECT Paper.title, Citation.number FROM Paper, Citation, Researcher, University WHERE Paper.title CROWDJOIN Citation.title AND Paper.author CROWDJOIN Researcher.name AND University.name CROWDJOIN Researcher.affiliation AND Paper.conference CROWDEQUAL "sigmod" AND University.country CROWDEQUAL "USA"</pre>	 <pre>SELECT Winner.name, City.country FROM Winner, City, Celebrity, Award WHERE Celebrity.name CROWDJOIN Winner.name AND Celebrity.birthplace CROWDJOIN City.name AND Winner.award CROWDJOIN Award.name AND Celebrity.birthplace CROWDEQUAL "New York" AND Award.place CROWDEQUAL "US"</pre>

Cost

Reduce 2-3 times cost

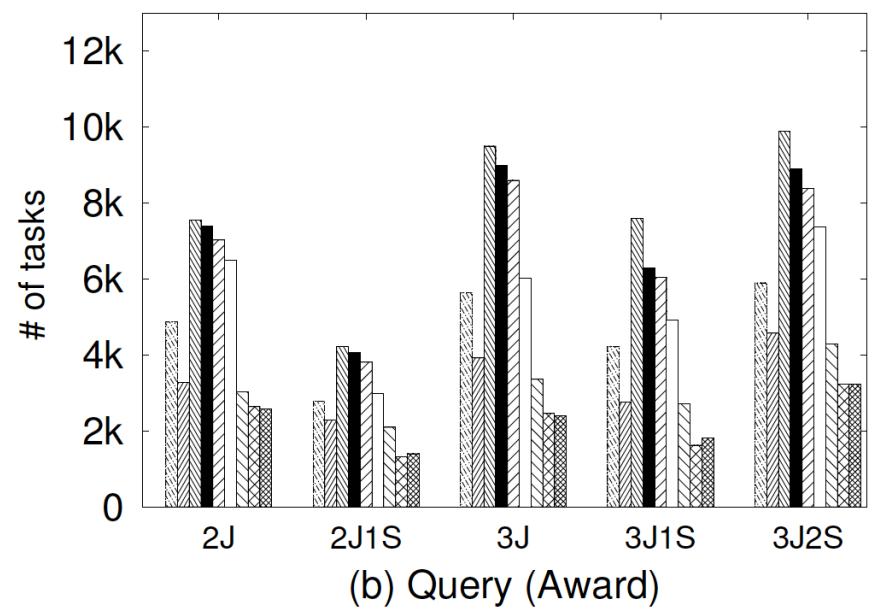
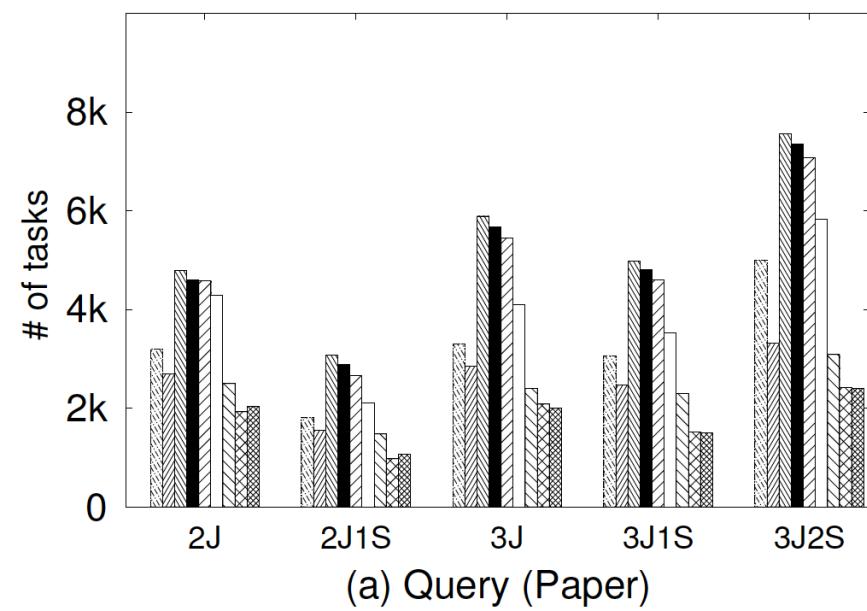
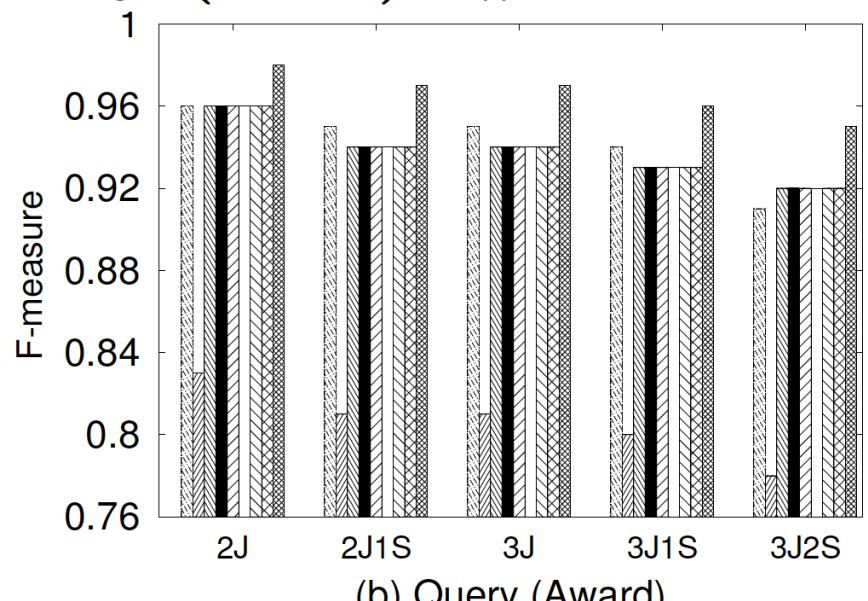
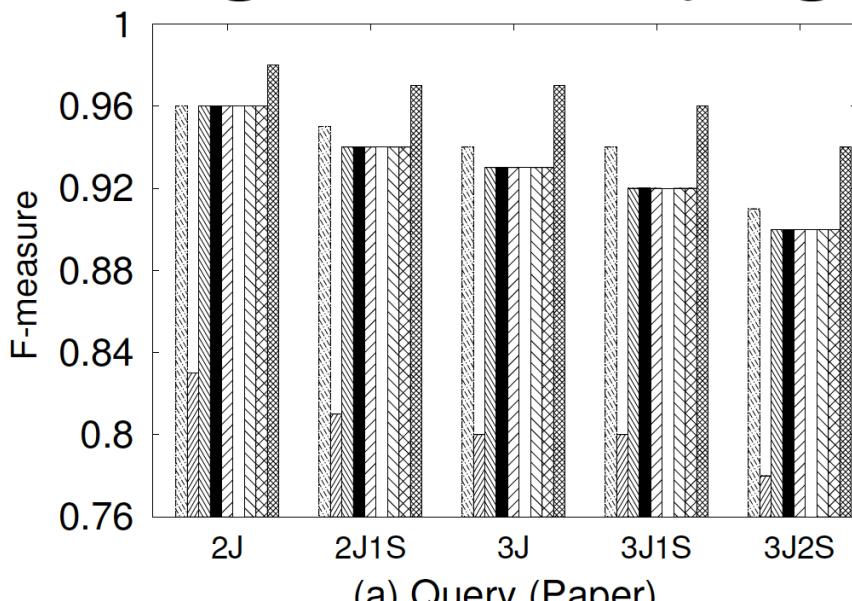


Figure 14: Varying Query (Real): # Tasks

Quality

Higher quality by about 5%



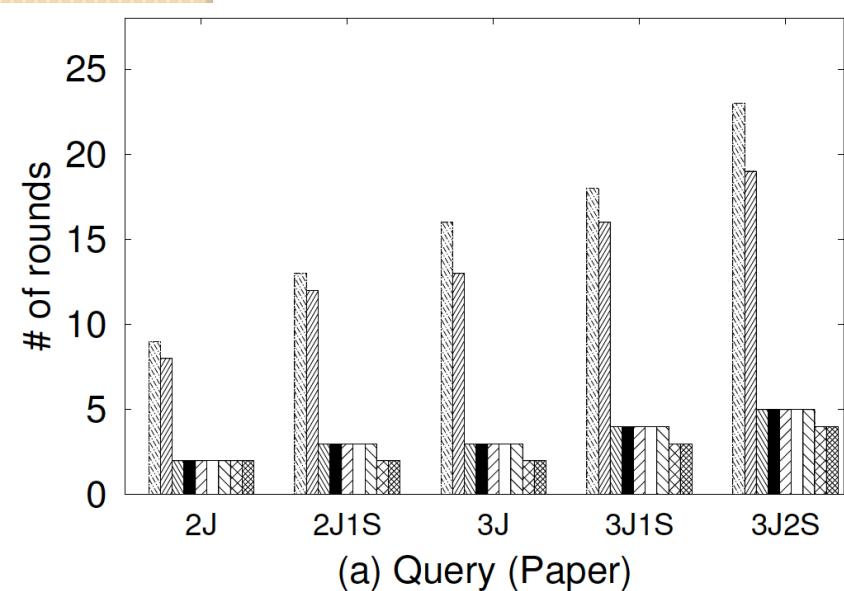
ACD Trans CrowdDB Qurk Deco OptTree MinCut CDB CDB+
(a) Query (Paper) (b) Query (Award)

Figure 15: Varying Query (Real): Fmeasure

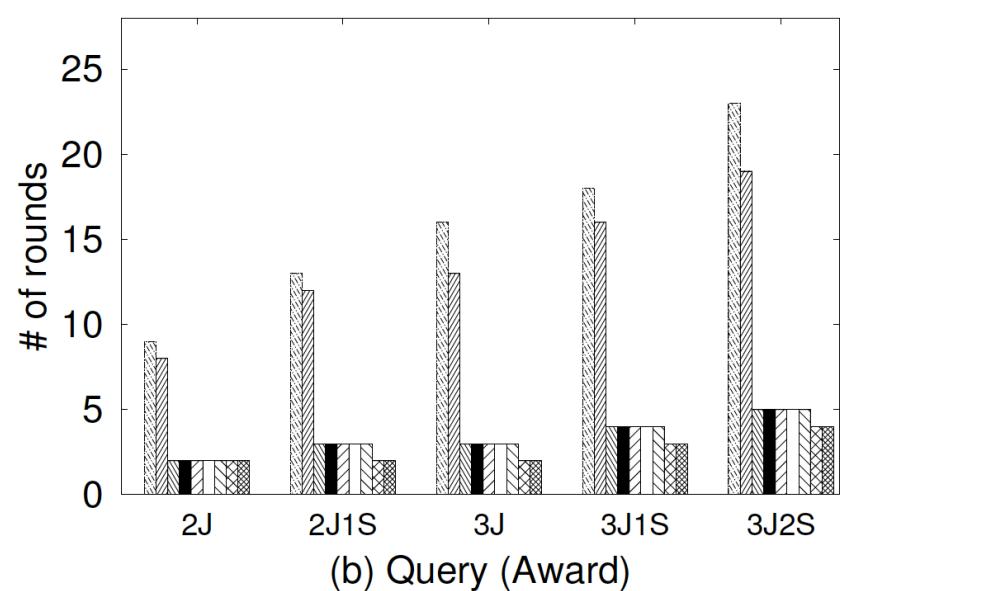


Latency

Lower Latency



(a) Query (Paper)



(b) Query (Award)

Figure 16: Varying Query (Real): # Rounds



Thank you!

