- Part I

 2. Let R be a relation schema, R, and R2 form a decomposition of R. Decomposition is a loss respect for all legal dotabases instences r of R, $\Pi_{R_1}(r) \bowtie \Pi_{R_2}(r) = r$
- 6. The time it takes for a disk I/O includes block-transfer time and access time, The latter consists of Stell time and rotational latency.
- 8. A sequence of primitive operations due can be used to evaluate a quary is a plan.
- 10. In query-processing, the query-execution engine takes the execution plan which cortains detailed information on how a particular query or a set of queries will be executed.
 - 11. In deferred database modification scheme, 1200? operation is the only operation used in the recovery procedure.
 - 12. Cascading rollbacks can be avoided by applying strict two phase locking protocol to transactions in currency protocols.
- 13. log is the most widely used structure for recording the modification of clostabase. Since a failure may occur while a update is taking place, it must be written out to nonvolatile storage before the actual update to database to be alone.

14. A (possibly concurrent) schedule is serializable if it is equal to a serial schedule.

2. Database design II: Consider a relation schema R(A,B,C,D,E) Given the functional dependencies ABC -> D, D->E, E>A. cas please find all candidate keys for R ABC, BCD. GUB cb, please de compose R in 3NF (we only decompose when there is violation of 3NF). If R is already in 3FN just write R(A, B, C, D, E) instead of decomposition. To={ABC>D,D>D=E>A} (ABCD) U(DE) U(EA) (C) Please 3 clecompose of P in SPENTSTAJ SABOJ 4. Relation R vontains 10,000 tuples and has 10 tuples per page. Relation S contains 2000 tuples and also has 10 tuples per page. If only 15 buffers were available, what would be the cost of a sore-merge join?
/ love + 200) - (2+ log 14 15 + log 14 15) + 1000 + 200 J. Describe the process of Cost-Based optimization. Is the following schedule S conflict serializable? Why? (T) ---> (T) write (A) read (A) 165 write(B)