

## ON TWO CLASSICAL RAMSEY NUMBERS OF THE FORM $R(3, n)^*$

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**Abstract.** New lower bounds are given for the classical Ramsey numbers  $R(3, 10)$  and  $R(3, 12)$ . Both constructions were made using a variant of the Metropolis Algorithm and were built on smaller cyclic constructions.

**Key words.** Ramsey numbers, heuristics

**AMS(MOS) subject classification.** 05C55

Two new lower bounds for Ramsey numbers of the form  $R(3, n)$  are proved below. Both proofs are completed by extending cyclic constructions for  $R(3, n - 1)$ .

In general our notation follows that of Harary [2]. We use  $R(s, t)$  to denote the classical Ramsey number of  $K_s$  versus  $K_t$ , defined to be the smallest integer  $n$  such that in any 2-coloring of the edges of  $K_n$  there is a monochromatic copy of  $K_s$  in color 1 or a monochromatic copy of  $K_t$  in color 2. A coloring of a complete graph is called an  $(s, t)$ -coloring if there are no monochromatic copies of  $K_s$  in color 1 or of  $K_t$  in color 2. A graph of order  $n$  is called a *cyclic*  $n(a_1, \dots, a_k)$  graph if its vertices can be labeled with the integers from 0 to  $n - 1$  so that two vertices are adjacent if and only if their difference is  $a_i$ , for some  $i$ ,  $1 \leq i \leq k$ .

The underlying algorithm we used to make these constructions is a procedure that has been called *simulated annealing* [3], and is based on an algorithm devised by Metropolis et al. [4] for application to statistical mechanics. We offer a brief description. Let  $f$  be an integer-valued function of integer variables  $x_1, \dots, x_n$ , and suppose we wish to find the minimum value of  $f$ . At each step of the algorithm we have a *current vector*  $(x_1, \dots, x_n)$  that may initially be chosen at random. We consider a small random change in one of the variables  $x_i$ , yielding a new vector  $(x_1, \dots, x'_i, \dots, x_n)$ . The values  $y = f(x_1, \dots, x_i, \dots, x_n)$  and  $y' = f(x_1, \dots, x'_i, \dots, x_n)$  are compared, and  $\Delta Y = y' - y$  is computed. If  $\Delta Y \leq 0$ , then the new vector is accepted as the current vector, otherwise the new vector is accepted with probability  $\exp(-\Delta Y/k_B T)$ , where  $k_B$  is the analogue of the Boltzmann constant and  $T$  is an analogue of temperature. In the course of running the algorithm we usually begin with a relatively large value for  $T$  (i.e., a high temperature), and gradually lower the value of  $T$  (i.e., allow the system to cool).

The first problem that arises when applying this procedure to Ramsey numbers is that of determining  $f$ . This issue has been discussed in some detail in [1]. New issues arise when dealing with far off-diagonal cases. Specifically, with  $R(3, t)$ , we must decide how much weight to give to a  $K_3$  in color 1, as opposed to a  $K_t$  in color 2. In the context of the Metropolis Algorithm, the random change in the current vector corresponds to recoloring one edge in a given 2-coloring. Suppose that coloring a given edge in color 1 yields  $m_1$  monochromatic  $K_3$ 's in color 1, while coloring it with color 2 yields  $m_2$  monochromatic  $K_t$ 's in color 2. Let  $\rho = m_2/m_1$ . The question that must be answered is: For what values of  $\rho$  do we prefer color 1 and for what values do we prefer color 2? Let  $\rho_0$

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be the value of  $\rho$  for which we are indifferent. In other words, for values of  $\rho > \rho_0$  we choose color 1, for values of  $\rho < \rho_0$  we choose color 2, and when  $\rho = \rho_0$  we make a random choice. In practice we have found that choosing

$$\rho_0 = \binom{t}{2} / 3$$

is a good choice, so that the weights are inversely proportional to the number of edges in the graphs we are trying to avoid.

The table of Ramsey numbers given in [5] seems to be the most recently published. The values listed there for numbers of the form  $R(3, n)$  are as follows:

$$\begin{array}{ll} R(3, 3) = 6, & R(3, 4) = 9, \\ R(3, 5) = 14, & R(3, 6) = 18, \\ R(3, 7) = 23, & 28 \leq R(3, 8) \leq 29, \\ R(3, 9) = 35, & 39 \leq R(3, 10) \leq 44, \\ 46 \leq R(3, 11) \leq 54, & 49 \leq R(3, 12) \leq 63. \end{array}$$

We improve the lower bounds for  $R(3, 10)$  and  $R(3, 12)$  by one.

**THEOREM 1.**  $R(3, 10) \geq 40$ .

*Proof.* Begin with the cyclic  $(3, 9)$ -coloring of  $K_{35}$  given by having the edges of the cyclic graph  $35(1, 7, 11, 16, 19, 24, 28, 34)$  colored in color 1, and the edge of the complement colored in color 2. To this graph we add four vertices labeled  $a, b, c$ , and  $d$ . The edges joining  $a$  to  $c$  and  $b$  to  $d$  are colored in color 2. The remaining edges among these four vertices are colored in color 1. In addition, the four new vertices are joined in color 1 to those of the original 35 as listed below:

$$\begin{array}{lllll} a: & 2 & 15 & 19 & 27 & 32 \\ b: & 11 & 17 & 25 & 29 & \\ c: & 8 & 16 & 26 & 28 & 34 \\ d: & 1 & 4 & 10 & 18 & 22 & 24 & 30. \end{array}$$

The remaining edges are in color 2.

**THEOREM 2.**  $R(3, 12) \geq 50$ .

*Proof.* The construction proceeds just as in Theorem 1. We begin with the cyclic coloring derived from the graph  $45(3, 10, 11, 12, 16, 29, 33, 34, 35, 42)$ . Again we add four vertices,  $a, b, c$ , and  $d$ , with  $a$  adjacent to  $c$  in color 2 and  $b$  adjacent to  $d$  in color 2. All other edges among these four vertices are in color 1. The color 1 edges joining the new vertices to the original 45 are given below:

$$\begin{array}{lllll} a: & 1 & 22 & 24 & 31 & 39 \\ b: & 6 & 7 & 14 & 28 & 33 & 37 \\ c: & 10 & 12 & 27 & 34 & 35 & 36 & 40 \\ d: & 2 & 3 & 4 & 17 & 21 & 25 & 30 & 43. \end{array}$$

We note that evidence seems to be accumulating for the conjecture that it is the exception, rather than the rule, for Ramsey colorings to be cyclic.

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