

Frama-C Training Session

Browsing your code dependencies

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How to better understand a C code within Frama-C by extracting semantic information from this code

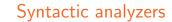
For what purpose

- helping to start verification of an unknown code
- helping to understand results of heavier analyses
- helping heavier analyses to give better results
- helping audit activities
- helping reverse-engineering activities

In what way

 using a battery of Frama-C plug-ins, either syntactic or semantic







Only deduce information from a direct use of the AST

Warnings

- ► those here-presented use the normalised program, not the original one
- does not use advanced semantical information (for instance, the value of a variable at some statement)
- ▶ in particular, does not handle pointers
- some may provide incorrect results in some cases

Syntactic analyzers within Frama-C

- analysing code using program syntax only is not the main goal of Frama-C
- only few syntactic analyzers in Frama-C





Syntactic analyzers what they (do not) provide

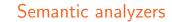
What their are good for

getting information quickly

What their are *not* good for

- providing a big amount of useful information
- providing confidence if they may provide incorrect results







In this session, all semantic analyzers are based on abtract interpretation and the value analysis plug-in

Features

- ► theoretically sound: always provide correct results, as long as there are no soudness implementation bugs
- ► handle pointers correctly

Semantic analyzers within Frama-C

most Frama-C plug-in are semantic analyzers



Semantic analyzers warnings

Warnings

- ► run the value analysis first
- may take a long time
- over-approximate the results
- all the ways to improve the efficiency/precision of the value analysis apply
- ▶ all the limitations of the value analysis also apply
- all the alarms emitted by the value analysis should be carefully examined







- 1. Lightweight analyzers
 - Metrics
 - Callgraphs
 - Constant foldings
 - Occurrence
- 2. Dependencies and effects
 - ► Functional dependencies and effects
 - ► Imperative effects
 - Operational effect
 - ► Data scoping
- 3. Reducing code to analyse
 - Slicing
 - Sparecode
 - ► Impact

syntactic both both semantic

semantic semantic semantic semantic

semantic semantic semantic







They are either:

- syntactic analyzers; or
- semantic analyzers remaining quite precise even if the value analysis does not give so precise results







Give some syntactic metrics about the analyzed code.

Features

- defined and undefined functions
- number of calls to each function
- potential entry points (the never-called functions)
- number of loc
- number of conditionals, assignments, loops, calls, gotos, pointer access

Warnings

- measures are done on the normalised code, not on the original one
- does not take function pointers into account





What is it good for

- helping to measure how difficult the analyses will be
- helping to identify whether some file is missing
- helping to identify which functions have to be stubbed or specified
- helping to identify entry points of the analyzed code

How to use

- ▶ -metrics dumps metrics on stdout
- -metrics-dump <f> dumps metrics on file f
- also (partially) available from the GUI





Indicate the callers of each function

Features

- representation as graphs into dot files
- notion of service, a group of related functions which seems to provide common functionalities

Warning

does not take function pointers into account

What is it good for

- helping to identify entry points of the analyzed code
- helping to discover services provided by an application
- grasping the code architecture





Syntactic callgraph (continuing)

How to use

- -cg <f> dumps callgraph in dot file f
- -cg-init-func <f> adds function f as a root service
- ► from the GUI: menu View, then Show Call Graph (still experimental)





Same as the syntactic callgraph... But using the program semantics

Features

correctly deal with function pointers

Warnings

run the value analysis first: may take a long time

What is it good for

computing the callgraph for codes with function pointers





Semantic callgraph (continuing)

How to use

- ▶ -scg-dump dumps the callgraph to stdout into dot format
- -cg-init-func <f> uses function f as a root service
- not available from the GUI

Warnings

- currently not the same interface as the syntactic callgraph (will be fixed soon)
- currently not exactly the same notion of service as the syntactic callgraph (will be fixed soon)





Same as the semantic callgraph... But not represented as a graph

Feature

display the callees of each functions

Warning

no service computed

What is it good for

extracting information with some external automatic tools (like grep)

How to use

-users dumps the function callees on stdout







Fold all constant expressions in the code before analysis

Feature

replace constant expressions by their results

Warning

▶ local propagation only: do not propagate the assignment of a constant to a left-value in the propram

What is it good for

- quickly simplifying programs with lots of constant expressions
- using analysis puzzled by big constant expressions

How to use

-constfold performs this analysis before all others







Propagate constant expressions in the whole program

More precisely

- generate a new program where expressions of the input program which are established as constant by the value analysis are
 - replaced by their value
 - propagated through the whole program

Features

- ▶ the output program is a compilable C code
- ▶ it has the same behaviour as the original one
- ▶ handle constant integers and pointers, even function pointers



Semantic constant folding (continuing)

Warning

does not handle floating-point values yet

What is it good for

- simplifying programs with lots of constant values
- using analysis puzzled by constant expressions

How to use

- -semantic-const-folding propagates constants and pretty print the new source code
- -semantic-const-fold <f1>, ..., <fn> propagates
 constants only into functions f1, ..., fn
- -cast-from-constant replaces expressions by constants even when doing so requires a pointer cast



Show the uses of a variable in a program

More precisely

highlight the left-values that may access a part of the location denoted by the selected variable

Features

- take aliasing into account
- also show uses of a C variable in logic annotations
- mainly a graphical plug-in

Warnings

- quite difficult to use in batch mode
- does not handle logic variable yet





What is it good for

- understanding a quite mysterious piece of code
- discovering some unknown aliases of the program

How to use

- -occurrence dumps the occurrences of each variable on stdout
- ▶ from the GUI: left panel and contextual menu







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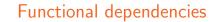




Features

- several notions of input/output for functions
- several kinds of dependencies







Dependencies between inputs and outputs of functions

Definitions

- functional output of a function f: left-value that may be modified in f when f terminates
- functional input of a function f: left-value which may impact the output value of a functional output of f

Features

- functional outputs and inputs
- dependencies between outputs and inputs
- indicate whether the analyzer knows that an output is always modified (when the function terminates)
- ignore local variables (from the next release)



Functional dependencies (continuing)

How to use

- mainly a batch plug-in
- ▶ -deps displays the functional dependencies for each function
- -calldeps displays the functional dependencies by callsite: if a function is called several times, results are not merged

What is it good for

- providing dataflow specifications of functions
- ▶ helping to understand relations between inputs and outputs of each function
- improving precision of other analyser through -calldeps





What is read, what is written, what is read before being written

Definitions

- imperative input of a function f: left-value that may be read in f
- imperative output of a function f: left-value that may be written in f
- operational input of a function f: left-value that is read without having been previously written to, when f terminates

Features

- ▶ imperative inputs and outputs
- operational inputs

Warnings

- ► mainly a batch plug-in
- operational inputs are still experimental: the specification may change
- operational outputs exist but are not yet documented





How to use

- -input displays the imperative inputs of each function; locals and function parameters are not displayed
- -input_with_formals same as -input, but displaying function parameters
- -out displays the imperative outputs of each function
- ▶ -inout displays the operational inputs of each function



Dependencies of a given left-value / at a given program point L

Features

- ▶ show defs: statements that may define the value of / at L
- zones: statements that may contribute to define the value of / at L
- ► data scope: statements where / is guaranteed to have the same value as at L

Warning

► still experimental





What is it good for

- locally better understand what the program does
 - relations between left-values
 - where the current value of a left-value comes from
 - scope of definition of a left-value

How to use

 only available from the GUI: sub-menu Dependencies of the contextual menu with three entries (Show defs, Zones, DataScope)







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Features

- generate a new program in a new project
- ▶ the new program is compilable
- the new program is usually shorter
- ▶ the new program is usually easier to analyze

Warning

usually is not always...







Specialize the program according to some user-provided criteria

Features

- generate a new program in a new project
- the new program is compilable
- ▶ the new program and the analysed one have the same behaviour according to the slicing criterion



Slicing available criteria

What are the available criteria?

Criteria for code observation

- preserving effects of statements
- preserving the read/write accesses of/to left-values

Criteria for proving properties

- preserving behaviour of assertions
- preserving behaviour of loop invariants
- preserving behaviour of loop variants
- preserving behaviour of threats (emitted by the value analysis)





Pragmas

- /*@ slice pragma ctrl; */ preserves the reachability of this control-flow point
- /*@ slice pragma expr e; */ preserves the value of the ACSL expression e at this control-flow point
- /*@ slice pragma stmt; */ preserves the effects of the
 next statement

How to use

- from command line options
- ▶ from the GUI: left panel and contextual menu







Each option preserves the semantics of the input program according to a specific criterion

Options of the form -slice-criterion <f1>, ..., <fn>

- ▶ -slice-calls: calls to these functions
- ▶ -slice-return: the return of the these functions
- ▶ -slice-pragma: slicing pragmas in theses functions
- ▶ -slice-assert: assertions of these functions
- ► -slice-loop-inv: loop invariants in these functions
- ► -slice-loop-var: loop variants in these functions
- ▶ -slice-threat: threats in these functions





command line options (continuing)

Options of the form -slice-criterion <v1>, ..., <vn>

- -slice-value values of these left-values at the end of the entry point
- -slice-rd read access to these left-values
- ▶ -slice-wr write access to these left-values

Warning

addresses of the left-values are evaluated at the beginning of the entry point





Custom options

- -slicing-level <n> specifies how to slice the callees
 - ▶ 0: never slice the called functions
 - ▶ 1: slice the callees but preserves all their functional outputs
 - ▶ 2: slice the callees but create at most 1 slice by function
 - ▶ 3: most precise slices; create as many slices as necessary

Default level is 2

- -no-slice-undef-functions does not slice the prototype of undefined functions (default)
- -slice-undef-functions slices the prototype of undefined functions
- -slice-print pretty prints the sliced code

Warning

▶ the higher the slicing level is, the slower the slicing is





What is it good for

- helping to extract the signifiant parts of a program according to your own criteria
- helping to understand where a behavior comes from
- helping analyses to give better results
- helping audit activities





Remove useless code of the program

Features

- generate a new program in a new project
- ▶ the new program is compilable
- ▶ the values assigned to the output variables of the main function are preserved in the new program
- ▶ slicing pragmas may be used to keep some statements
 - /*@ slice pragma ctrl; */
 - /*@ slice pragma expr e; */
 - /*@ slice pragma stmt; */





Warnings

- ▶ still experimental
- ▶ partial support of ACSL: only the annotations inside function bodies (e.g. assertions) are processed at the moment; all the others are ignored and do not appear in the new program

What is it good for

- ▶ help to discover what is useless in a program
- may improve the results of others analyzers which are puzzled by some useless code



Sparecode

command line options

How to use

- -sparecode-analysis removes statements and functions that are not useful to compute the result of the program
- -rm-unused-globals removes unused types and global variables
- -sparecode-no-annot may remove some useless code even if it changes the validity of some ACSL properties





What could be discovered if the side effect of a statement would be revealed

More precisely

- ▶ a statement s is impacted by a statement s' iff modifying the effect of s' by another possible one may modify the effect of s
- ▶ an effect is possible iff there is an execution of the program that generates this effect. For instance, the possible effects of z=x+y; in x=c?0:1; y=c?0:1; z=x+y are z becomes equal to 0 or 2.

Warning

▶ still experimental



What is it good for

- helping to understand what a statement is useful for
- helping to apprehend code changes
- helping audit activities, in particular security audits

How to use

-impact-pragma <f1>, ..., <fn> computes the impact from the pragmas in functions f1, ..., fn. Only the following pragma is yet usable.

```
/*@ impact pragma stmt; */
```

- -impact-print dumps the result of the analysis on stdout
- ▶ from the GUI: left panel and contextual menu



Battery of Frama-C plug-ins presented

For what purpose

- helping to start verification of an unknown code
- helping to understand results of heavier analyses
- helping heavier analyses to give better results
- helping audit activities
- helping reverse-engineering activities

Browse your code dependencies more easily!

