

# Enhancing Efficiency in Microchip Distribution Strategic Supply Chain Route Optimization

*Optimization Methods Final Project*

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# PROBLEM DESCRIPTION



Allocate 1000 orders to a warehouse and assign freight paths



# DATA OVERVIEW | RELATIONAL DATABASE



Order
Order ID : integer
Order Date : date
Origin Port : string
Carrier : string
TPT : integer
Service Level : string
Ship ahead day count : integer
Ship Late Day count : integer
Customer : string
Product ID : integer
Plant Code : string
Destination Port : string
Unit quantity : integer
Weight : float

FreightRates
Carrier : string
orig_port_cd : string
dest_port_cd : string
minm_wgh_qty : float
max_wgh_qty : float
svc_cd : string
minimum cost : float
rate : float
mode_dsc : string
tpt_day_cnt : integer
Carrier type : string

WhCosts
WH : string
Cost/unit : float

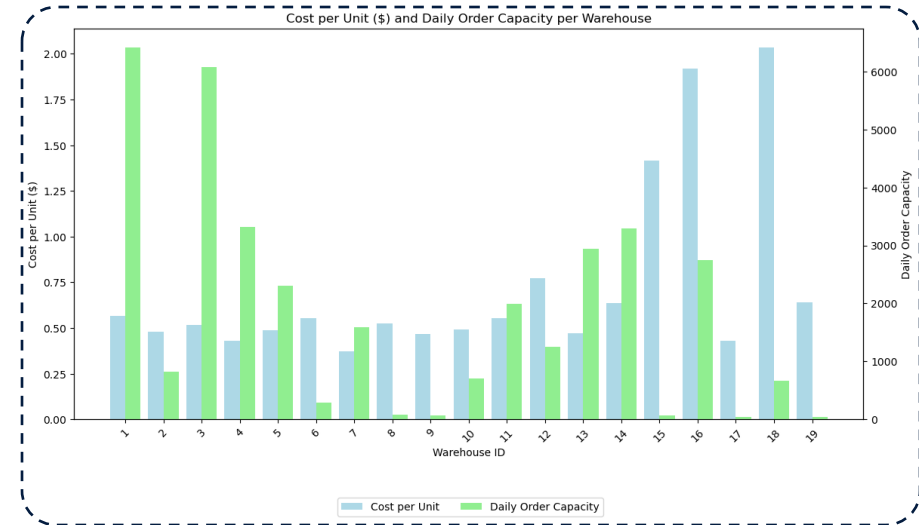
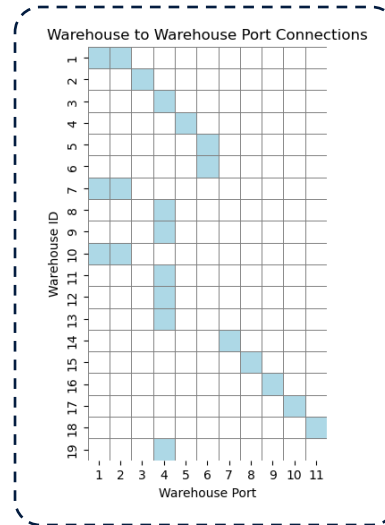
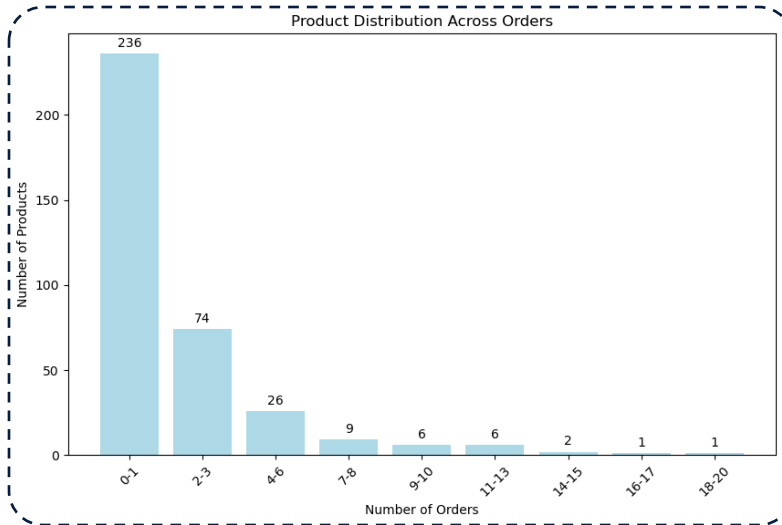
WhCapacities
Plant ID : string
Daily Capacity : integer

ProductsPerPlant
Plant Code : string
Product ID : integer

VmiCustomers
Plant Code : string
Customers : string

PlantPorts
Plant Code : string
Port : string

# EXPLORATION DATA ANALYSIS



Various types of products

Sparse supply ability

Lack of low-cost warehouse



**COMPLICATED!!!**

# BASELINE MODEL | YAN-TIAN GREEDY ALGORITHM



- 1 Traverse through all orders
- 2 Check constraints one by one
- 3 Assign orders to available place
- 4 Calculate costs

## Algorithm 1 Yan-Tian Greedy Algorithm For Order Assignment

```
// Assign orders to warehouse and freight
for orderk from 1 to norder do
  for warehousei from 1 to nwarehouse do
    for freightj from 1 to nfreight do
      if warehousei can produce orderk's product then
        if warehousei can serve orderk's customer then
          if warehousei has capacity then
            if freightj has capacity then
              if warehousei can transport products to freightj's warehouse port then
                if freightj's transportation time satisfy orderk's demanding time then
                  Assign orderk to warehousei and freightj
                  BREAK
                end if
              end if
            end if
          end if
        end if
      end if
    end for
  end for
end for
// Calculate transportation cost TC and Warehouse Cost WC, Penalty Cost PC, and total cost C
Initialize TC = 0, WC = 0, PC = 0, and C = 0
for warehousei from 1 to nwarehouse do
  WC = WC + warehousei's cost
end for
for freightj from 1 to nfreight do
  TC = TC + freightj's cost
end for
for orderk from 1 to norder do
  if orderk is not assigned then
    PC = PC + unit penalty cost
  end if
end for
C = PC + TC + WC
// Output results
```

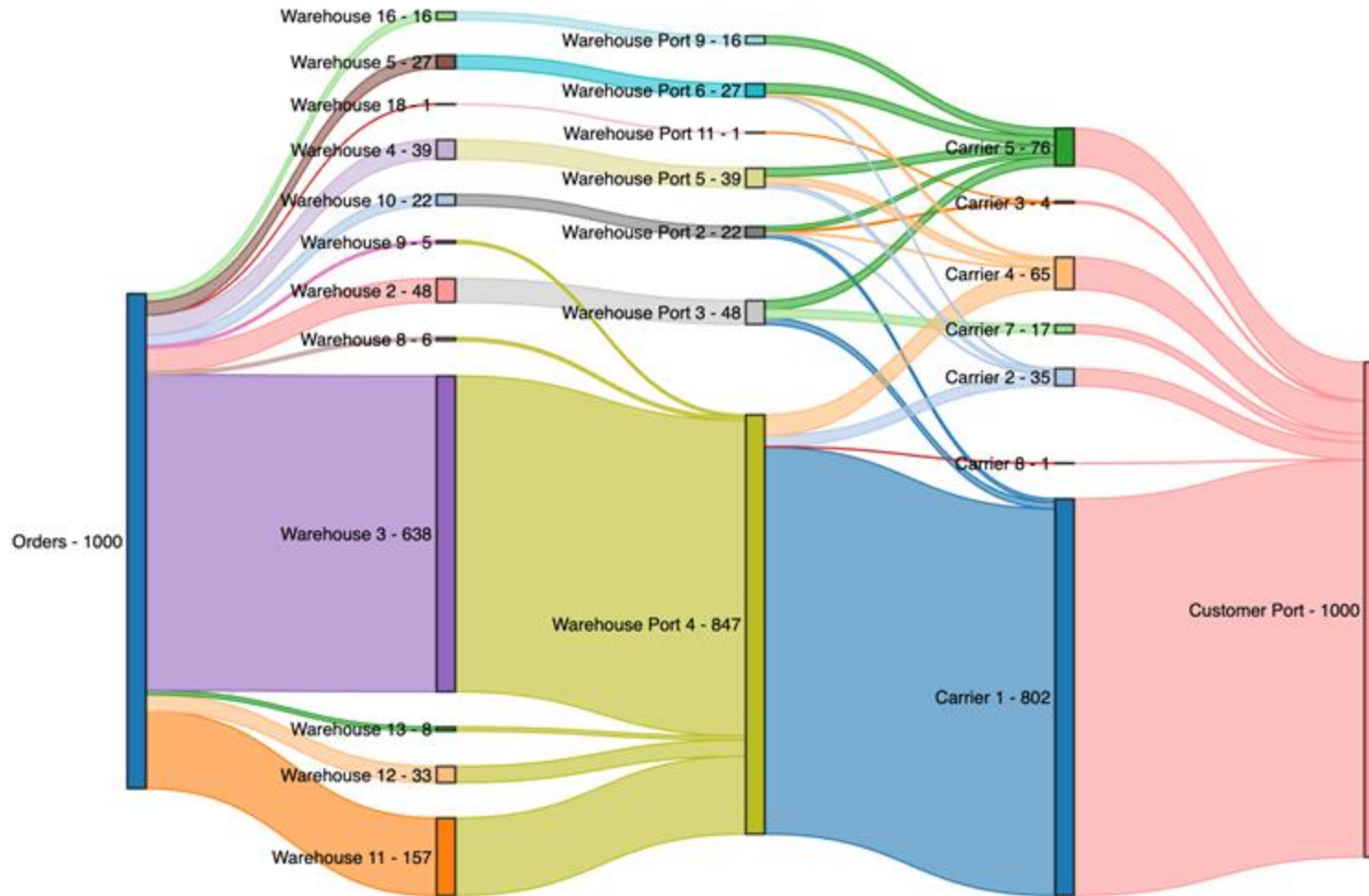
# FORMULATING OPTIMIZATION MODEL



Minimize Total Cost	→	min	$\sum_{k=1}^{n_{\text{order}}} \sum_{i=1}^{n_{\text{warehouse}}} X_{ki} \cdot p_i \cdot q_k + \sum_{k=1}^{n_{\text{order}}} \sum_{j=1}^{n_{\text{freight}}} Y_{kj} \cdot TC_{kj}$
Assignment Constraint	→	s.t.	$\sum_{i=1}^{n_{\text{warehouse}}} X_{ki} = 1, \forall k; \quad \sum_{j=1}^{n_{\text{freight}}} Y_{kj} = 1, \forall k$
Capacity Constraint	→		$\sum_{k=1}^{n_{\text{order}}} X_{ki} \leq \text{cap}_i, \forall i; \quad \sum_{k=1}^{n_{\text{order}}} Y_{kj} \cdot \text{ow}_k \leq \text{maxw}_j, \forall j$
Warehouse Constraint	→	{	$X_{ki} \leq M(PW_{\text{pro}_k i}), \forall k, i$ $X_{ki} \leq M(WC_{\text{icus}_k}), \forall k, i$ $Y_{kj} \leq M(1 - X_{ki} + WP_{\text{io}_j}), \forall k, i, j$
Delivery Time Constraint	→		$\sum_{j=1}^{n_{\text{freight}}} Y_{kj} \cdot t_j \leq \text{ot}_k, \forall k$
Cost Calculation	→	{	$TC_{kj} \geq s_k[(1 - m_j) \cdot TCA_{kj} + m_j \cdot TCG_j], \forall k, j$ $\text{minc}_j \leq TCA_{kj}, \forall k, j; \quad \text{ow}_k \cdot r_j \leq TCA_{kj}, \forall k, j$ $z_j r_j \leq TCG_j, \forall j; \quad z_j \leq \sum_{k=1}^{n_{\text{order}}} Y_{kj}, \forall j; \quad Y_{kj} \leq z_j, \forall k, j$



# RESULT | COST EFFECTIVE STORAGE & WELL-CONNECTED FREIGHT PATHS

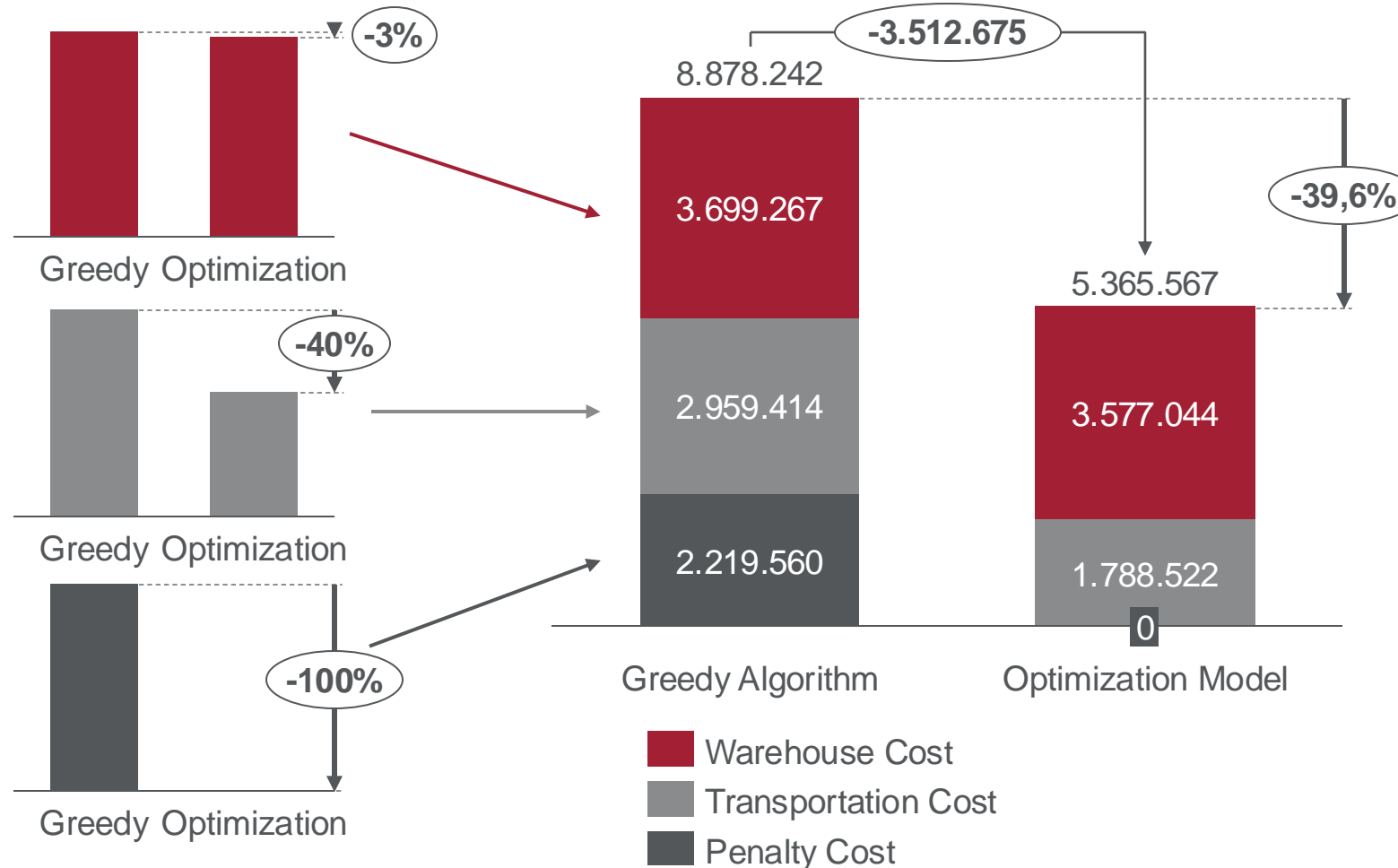


Economical Storage  
**80% Warehouse 3**

Dual Transportation  
**80% Carrier 1**

High Connectivity  
**85% WPort 4**

# BUSINESS IMPACT | NOTABLE COST REDUCTION & TRACTABLE SOLVING TIME



**Cost Reduction**  
**~ \$3.5M**

**Cost Saving**  
**~ 40%**

**Solving Time**  
**< 2 Minutes**



For a customer order:

Decide

If  $X_{ki} = 1$ ,

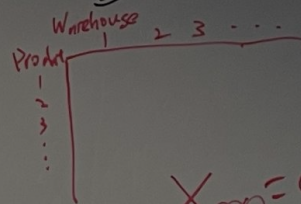
then  $Y_{kj} = 0$  if  $C_{idj} = 0$

$\Rightarrow$  If  $X_{ki} = 1$  and  $C_{idj} = 0$   
then  $Y_{kj} = 0$

$X_{ki} = 0$ , for  $i \notin \text{Set} \rightarrow Y_{kj} = 0$

②, for  $i \in \text{Set}$

$Y_{kj} TC_{kj}$



$X_{00} = 0$   $X_{00} = 0$

$\rightarrow$  get from "orderList"

$i \leftarrow \text{port}$   
 $X_{ki} + Y_{kj} = 0$

we know warehouse.

$X_{ki} + Y_{kj} = 0$

$X_{ki} + Y_{kj} = 0$

# THANK YOU

$C_{idj}$  = if we connect to port otherwise 0

Original port:  $O_j$   
Destination port:  $D_j$

Carrier:  $G_j$

Decision variable:

Objective:

① Warehouse:  $X_{ki} \ i \in [h], k \in [L]$

② Route assignment:  $Y_{kj}, j \in [t], k \in [L]$

$\sum_{k=1}^L \sum_{i=1}^h X_{ki} P_{ik} - \text{Warehouse cost}$

$\sum_{k=1}^L \sum_{j=1}^t Y_{kj} TC_{kj}$

$TC_{kpica}$

Min weight:  $\min W_j$   
Rate:  $R_j$   
Min Cost:  $\min C_j$

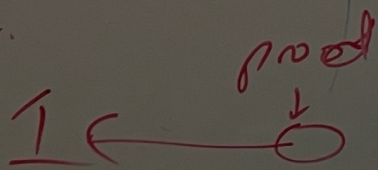
for order  $k$ , the  $TC_{kj}$ ,  $\in \{f\}$

$\max (MING_j, W_k R_j) \rightarrow TC_{kj}$

$Z_{kj} \geq MING_j$

$Z_{kj} \geq W_k R_j$

$TC_{kj} = (1 - M_j) Z_{kj} + M_j \left( \frac{R_j W_k Y_{kj}}{\sum_{k,j} W_k Y_{kj}} \right) S_k$



$\{0, 1\}$