

Chapter 3: Describing Syntax and Semantics

1. What is the difference between a sentence and a sentential form?

Ans:

- A sentence is a sentential form that has only terminal symbols. On the other hand, every string of symbols in a derivation is a sentential form.
- Every sentence is a sentential form but every sentential form is not a sentence.

2. Describe the basic concept of operational Semantics.

Ans:

- Operational Semantics describe the meaning of a program by executing its statements on a machine, either simulated or actual.
- The change in the state of the machine (memory, registers, etc.) defines the meaning of the statement.
- To use operational semantics for a high-level language, a virtual machine is needed.
- Uses: (i) Language manuals and textbooks (ii) Teaching programming languages.

3. Define syntax and semantics.

Ans:

Syntax: Syntax is the form or structure of the expressions, statements, and program units.

Semantics: Semantics is the meaning of the expressions, statements, and program units.

Syntax and semantics provide a language's definition

4. What is the difference between a synthesized and an inherited attribute?

Ans:

| Synthesized Attribute | Inherited Attribute |
|--|---|
| An attribute is said to be Synthesized attribute if its parse tree node value is determined by the attribute value at child nodes. | An attribute is said to be Inherited attribute if its parse tree node value is determined by the attribute value at parent and/or siblings' node. |
| If all attributes are synthesized, the tree can be decorated in bottom-up order. | If all attributes are inherited, the tree can be decorated in top-down order. |
| The production must have non-terminal as its head. | The production must have non-terminal as a symbol in its body. |

5. Is aabcdccc in the language generated by the following grammar? If so, draw the parse tree. If not, show why.

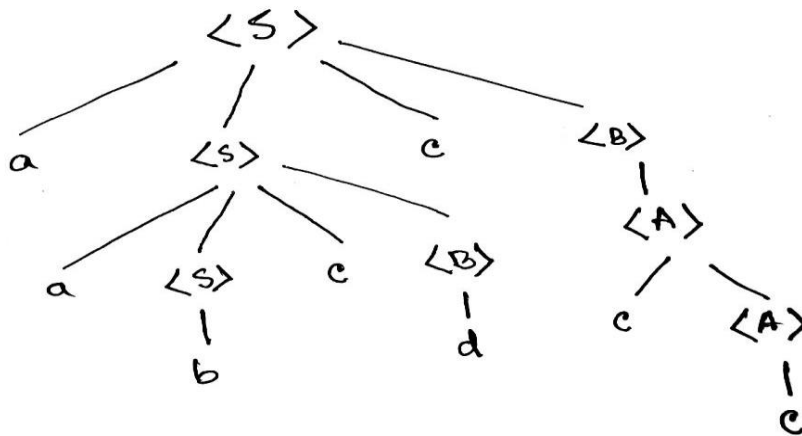
$\langle S \rangle \rightarrow a \langle S \rangle c \langle B \rangle \mid \langle A \rangle \mid b$

$\langle A \rangle \rightarrow c \langle A \rangle \mid c$

$\langle B \rangle \rightarrow d \mid \langle A \rangle$

Ans: Yes, aabcdccc in the language is generated by the following grammar.

Parse tree:



6. Is aacccdcc in the language generated by the following grammar? If so, draw the parse tree. If not, show why.

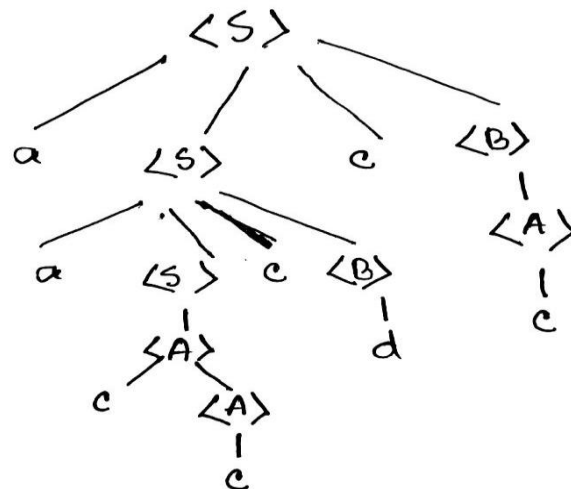
$\langle S \rangle \rightarrow a \langle S \rangle c \langle B \rangle \mid \langle A \rangle \mid b$

$\langle A \rangle \rightarrow c \langle A \rangle \mid c$

$\langle B \rangle \rightarrow d \mid \langle A \rangle$

Ans: Yes, aacccdcc in the language is generated by the following grammar.

Parse tree:



7. Derive the string 'acabcedcd' using the following grammar (make leftmost derivation):

$$\langle S \rangle \rightarrow a \langle S \rangle c \langle B \rangle \mid \langle A \rangle \mid b$$

$$\langle A \rangle \rightarrow c \langle S \rangle \mid c$$

$$\langle B \rangle \rightarrow d \mid b \langle A \rangle$$

Ans: Derivation of the string 'acabcedcd':

$$\langle S \rangle \rightarrow a \langle S \rangle c \langle B \rangle$$

$$\rightarrow a \langle A \rangle c \langle B \rangle \quad [\text{Using } \langle S \rangle \rightarrow \langle A \rangle]$$

$$\rightarrow a c \langle S \rangle c \langle B \rangle \quad [\text{Using } \langle A \rangle \rightarrow c \langle S \rangle]$$

$$\rightarrow a c a \langle S \rangle c \langle B \rangle c \langle B \rangle \quad [\text{Using } \langle S \rangle \rightarrow a \langle S \rangle c \langle B \rangle]$$

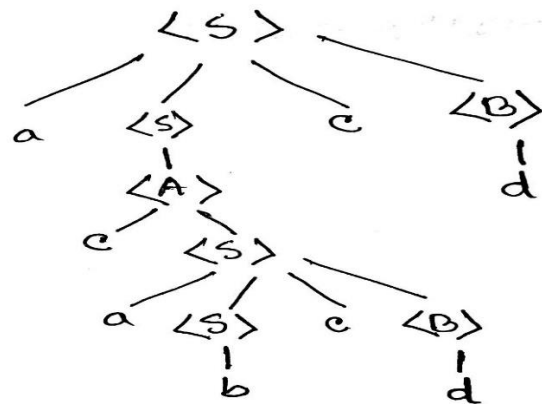
$$\rightarrow a c a b c \langle B \rangle c \langle B \rangle \quad [\text{Using } \langle S \rangle \rightarrow b]$$

$$\rightarrow a c a b c d c \langle B \rangle \quad [\text{Using } \langle B \rangle \rightarrow d]$$

$$\rightarrow a c a b c d e d \quad [\text{Using } \langle B \rangle \rightarrow d]$$

(derived)

using the Parse tree:



8. Prove that the following grammar is ambiguous:

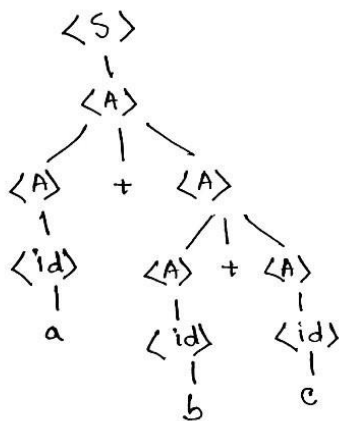
$$\langle S \rangle \rightarrow \langle A \rangle$$

$$\langle A \rangle \rightarrow \langle A \rangle + \langle A \rangle \mid \langle \text{id} \rangle$$

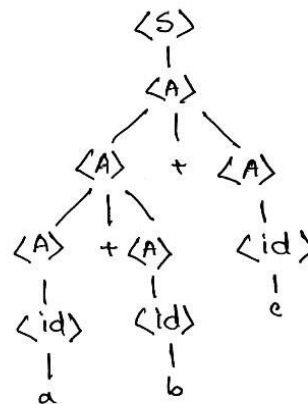
$$\langle \text{id} \rangle \rightarrow a \mid b \mid c$$

Ans: If a grammar generates different parse tree, then we can say that, the grammar is ambiguous means it is not a good grammar. The given grammar generates two different parse trees. The parse trees are shown below.

Parse Tree - I



Parse Tree - II



So, the given grammar is ambiguous.

9. Consider the following grammar:

$$\langle S \rangle \rightarrow a \langle S \rangle c \langle B \rangle \mid \langle A \rangle \mid b$$

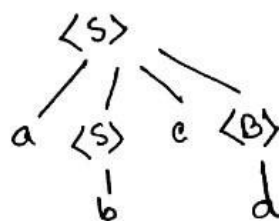
$$\langle A \rangle \rightarrow c \langle A \rangle \mid c$$

$$\langle B \rangle \rightarrow d \mid \langle A \rangle$$

Which of the following sentences are in the language generated by this grammar?

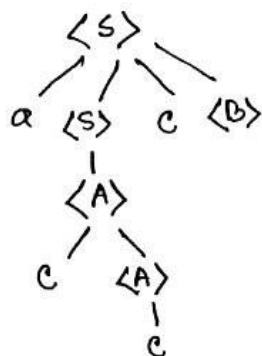
- a. abcd
- b. acccbd
- c. acccbcc
- d. acd
- e. accc

(a) abad:



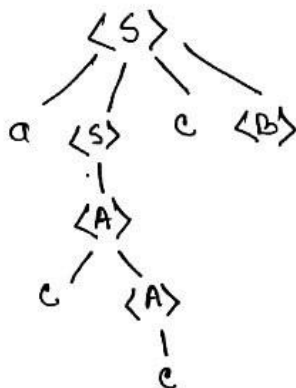
so, the sentence 'abad' ~~are~~^{is} in the language generated by this grammar.

(b) accabd:



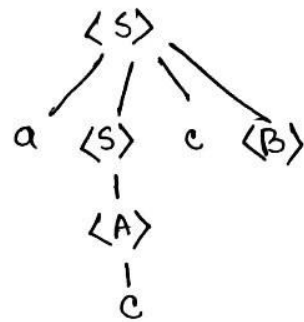
so, the sentence 'accabd' isn't in the language generated by this grammar. the parse tree can't generate the sentence

(c) accabcc:



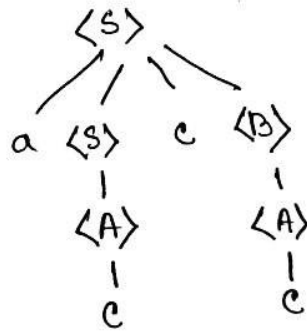
so, the sentence 'accabcc' isn't in the language generated by this grammar. the parse tree can't generate the sentence.

(d) acd:



So, the sentence 'acd' isn't in the language generated by this grammar. The parse tree can't generate the sentence.

(e) accc:



So, the sentence 'accc' is in the language generated by this grammar.

10. Convert BNF to EBNF

(i) BNF: $\langle \text{expr} \rangle \rightarrow \langle \text{expr} \rangle + \langle \text{term} \rangle$
 $\quad \quad \quad | \langle \text{expr} \rangle - \langle \text{term} \rangle$
 $\quad \quad \quad | \langle \text{term} \rangle$
EBNF: $\langle \text{expr} \rangle \rightarrow \langle \text{term} \rangle \{ (+ | -) \langle \text{term} \rangle \}$

(ii) BNF: $\langle \text{term} \rangle \rightarrow \langle \text{term} \rangle * \langle \text{factor} \rangle$
 $\quad \quad \quad | \langle \text{term} \rangle / \langle \text{factor} \rangle$
 $\quad \quad \quad | \langle \text{factor} \rangle$
EBNF: $\langle \text{term} \rangle \rightarrow \langle \text{factor} \rangle \{ (* | /) \langle \text{factor} \rangle \}$

(iii) BNF: ~~begin~~ $\langle \text{program} \rangle \rightarrow \text{begin} \langle \text{stmt_list} \rangle \text{end}$
EBNF: $\langle \text{program} \rangle \rightarrow \text{begin} \langle \text{stmt_list} \rangle \text{end}$

(iv) BNF: $\langle \text{stmt_list} \rangle \rightarrow \langle \text{stmt} \rangle$
 $\quad \quad \quad | \langle \text{stmt} \rangle ; \langle \text{stmt_list} \rangle$
EBNF: $\langle \text{stmt_list} \rangle \rightarrow \langle \text{stmt} \rangle \{ ; \langle \text{stmt_list} \rangle \}$

(v) ~~stmt~~ BNF: $\langle \text{stmt} \rangle \rightarrow \langle \text{var} \rangle = \langle \text{expression} \rangle$
EBNF: $\langle \text{stmt} \rangle \rightarrow \langle \text{var} \rangle = \langle \text{expression} \rangle$

(vi) BNF: $\langle \text{var} \rangle \rightarrow A | B | C$
EBNF: $\langle \text{var} \rangle \rightarrow A | B | C$

(vii) BNF: $\langle \text{expression} \rangle \rightarrow \langle \text{var} \rangle + \langle \text{var} \rangle$
 $\quad \quad \quad | \langle \text{var} \rangle - \langle \text{var} \rangle$
 $\quad \quad \quad | \langle \text{var} \rangle$
EBNF: $\langle \text{expression} \rangle \rightarrow \langle \text{var} \rangle \{ (+ | -) \langle \text{var} \rangle \}$

$\left\{ \begin{array}{l} \text{(viii) } \underline{\text{BNF:}} \quad \langle \text{assign} \rangle \rightarrow \langle \text{id} \rangle = \langle \text{expr} \rangle \\ \underline{\text{EBNF:}} \quad \langle \text{assign} \rangle \rightarrow \langle \text{id} \rangle = \langle \text{expr} \rangle \end{array} \right.$

$\left\{ \begin{array}{l} \text{(ix) } \underline{\text{BNF:}} \quad \langle \text{id} \rangle \rightarrow A|B|C \\ \underline{\text{EBNF:}} \quad \langle \text{id} \rangle \rightarrow A|B|C \end{array} \right.$

$\text{(x) } \underline{\text{BNF:}} \quad \langle \text{expr} \rangle \rightarrow \langle \text{expr} \rangle + \langle \text{expr} \rangle$
 $\qquad \qquad \qquad | \langle \text{expr} \rangle * \langle \text{expr} \rangle$
 $\qquad \qquad \qquad | \langle \text{expr} \rangle$
 $\qquad \qquad \qquad | \langle \text{id} \rangle$

$\underline{\text{EBNF:}} \quad \langle \text{expr} \rangle \rightarrow \langle \text{expr} \rangle \{ (+ | *) \langle \text{expr} \rangle \}$
 $\qquad \qquad \qquad | \langle \text{id} \rangle$